

# **DAA ASSIGNMENT 5**

## **Group No 20**

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# Problem Statement

- ▶ Given an  $N \times N$  chessboard and a Knight at position  $(x, y)$ . The Knight has to take exactly  $K$  steps, where at each step it chooses any of the 8 directions uniformly at random. What is the probability that the Knight remains in the chessboard after taking  $K$  steps, with the condition that it can't enter the board again once it leaves it? Solve using Dynamic programming

# Introduction

Dynamic Programming (DP) is an algorithmic technique for solving an optimization problem by breaking it down into simpler sub-problems and utilizing the fact that the optimal solution to the overall problem depends upon the optimal solution to its sub-problems.

# ALGORITHMIC DESIGN

One thing that we can observe is that at every step the Knight has 8 choices to choose from. Suppose, the Knight has to take  $k$  steps and after taking the  $K$ th step the knight reaches  $(x,y)$ . There are 8 different positions from where the Knight can reach to  $(x,y)$  in one step, and they are:  $(x+1,y+2)$ ,  $(x+2,y+1)$ ,  $(x+2,y-1)$ ,  $(x+1,y-2)$ ,  $(x-1,y-2)$ ,  $(x-2,y-1)$ ,  $(x-2,y+1)$ ,  $(x-1,y+2)$ .

The final probability after  $K$  steps will simply be equal to the ( probability of reaching each of these 8 positions after  $K-1$  steps)/8.

We are dividing by 8 because each of these 8 positions has 8 choices and position  $(x,y)$  is one of the choices.

For the positions that lie outside the board, we will either take their probabilities as 0 or simply neglect it.

Since we need to keep track of the probabilities at each position for every number of steps, we need Dynamic Programming to solve this problem.

# Time complexity

In this method, we are working on  $n \times n$  elements and there are  $k$  layers considered

Therefore, Time complexity would be  $O(n^2 \times k)$

# Space complexity

The size of dp array is  $n \times n$  and some constant variables are used.

Therefore, Space complexity is  $O(n^2)$



# Conclusion

Solutions in which some steps are repeating again and again can be made efficient using dynamic programming

# References

- ▶ [1] Wikipedia: Dynamic Programming,  
[https://en.wikipedia.org/wiki/Dynamic\\_programming](https://en.wikipedia.org/wiki/Dynamic_programming)
- ▶ [2] GeeksforGeeks: Dynamic Programming,  
<https://www.geeksforgeeks.org/dynamic-programming>