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# Find Median for Data streams

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1 | P  $\Rightarrow$  [1, 2, 3]  
 0 | P  $\Rightarrow$  [1, 5, 2]

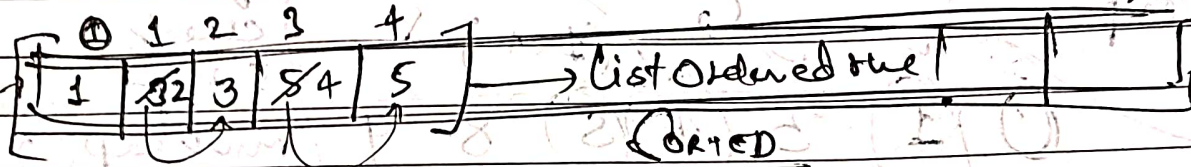
Median = Ordered list  $\rightarrow$  middle

any moment  
 you have to give  
 a median.

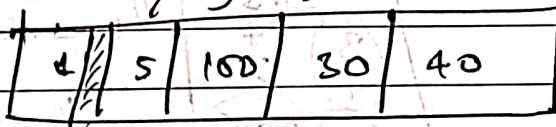
Odd  $\in$  1 2 (5) 7 8

Even 1 (2+7) 8  
 $= 3.5$

A1) Brute Force Given: Almost  $5 \times 10^4$  calls will be made to add new & find Median.



1, 3, 2, 5



TLE

Large no. of elements

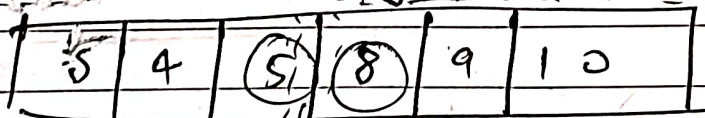
A2) Priority Queue  $\Rightarrow$  I don't want to sort the list everytime when short element comes!

(Heap)

~~$n \log n$~~

That's why we need max heap so that it give sorted way.

$n=6$



even case

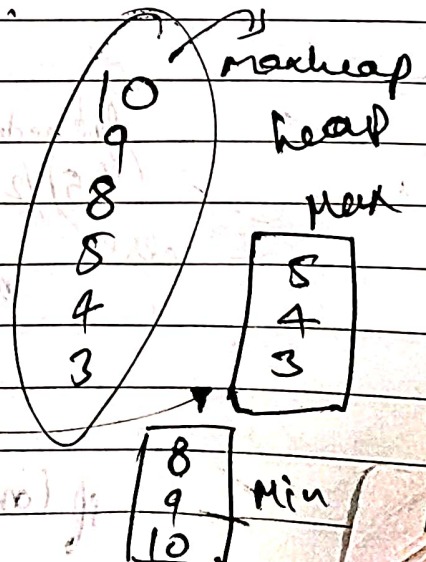
max heap PQ 1 - heap min heap PQ 2 -

Why we are doing this?

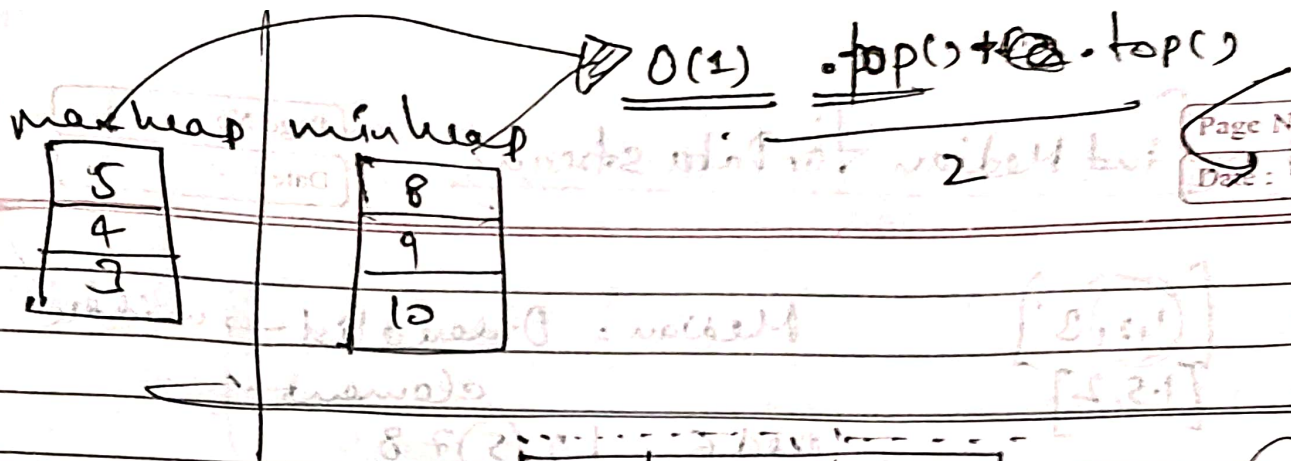
Because we need that in topmost!

(5)

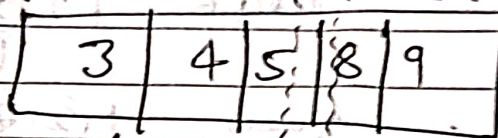
$O(1)$  😊



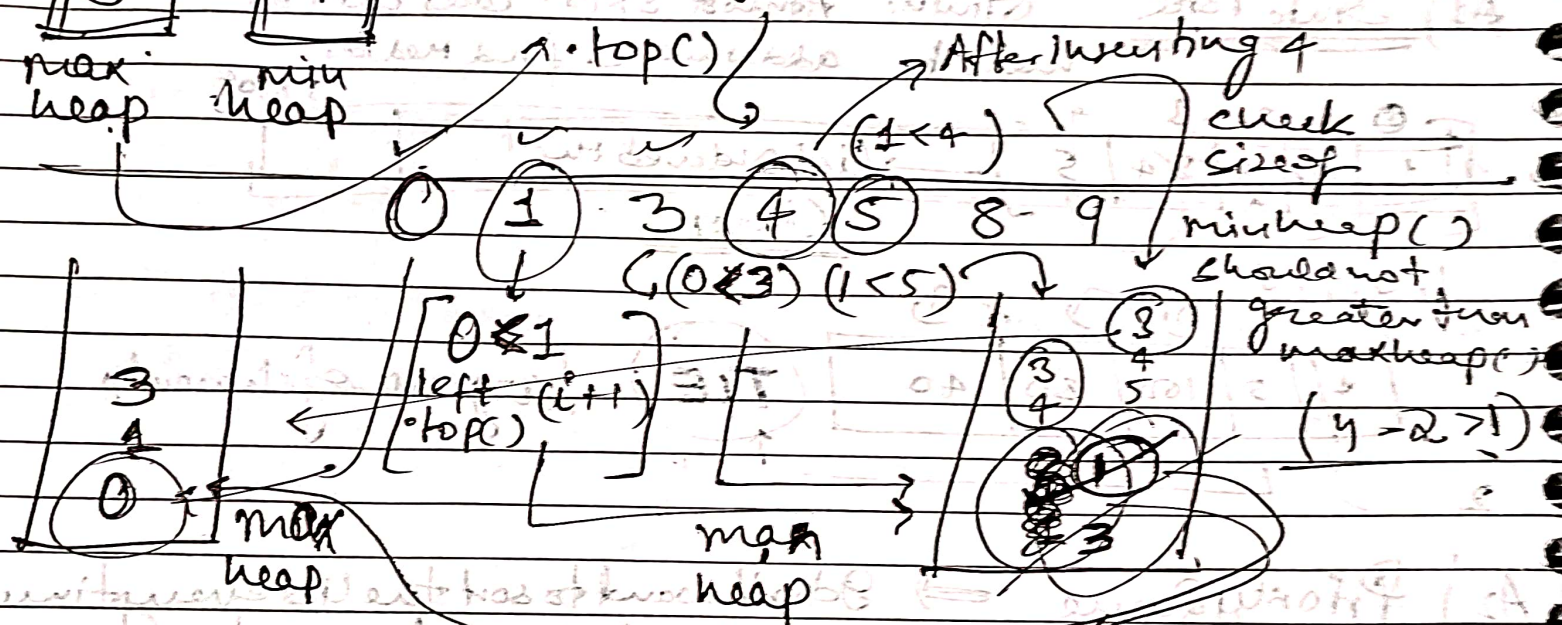
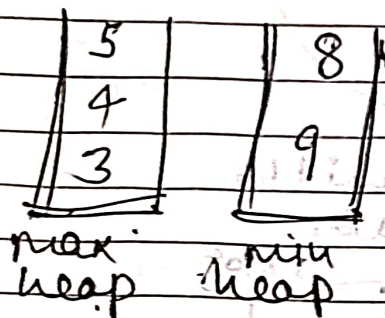




$n = \text{odd}$



Assuming that



find median

$(5/2) \neq 0$

return max heap =  $\text{sum} + 1$

else

int a = max heap()

b = min heap()

$(a+b)/2$

if (min heap > max heap)



P2 > L;  
P2 < R;

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Median Finder

```
// maxheap L = new P2<> (Collection<Integer>());
// minheap R = new P2<> ();
```

```
addNum(x) {
    num = L.peek();
    if (L.isEmpty() || num > x)
        L.offer(x);
    else
        R.offer(x);
}
```

Left heap  
// maxheap  
ke andar kab  
push karna hai

minheap and num > L.peek() may have  
At this time R has more elements

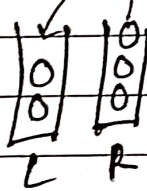
// Always maintain L.size() >= R.size()

$(L \leq R + 1)$

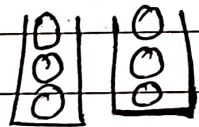
```
if (abs(L.size() - R.size()) > 1)
    Repush(L.top()); L.pop();
elseif (L.size() < R.size())
    L.push(R.pop());
```

$L - R \geq 1$   
 $L \geq R + 1$

that means  
L.size() must  
not be greater  
than R.size() + 1



double findMedian()



even no. of element

if (L.size() == R.size())

return (double) (L.top() + R.top()) / 2;

// odd case return L.top();

NOTE: Do remember to use # Offer, Peek, Poll

As this is P2

Remember initialisation of priority queue,