

Evaluation of Topology-based OpenFlow Controller for Multipath TCP on PRAGMA-ENT

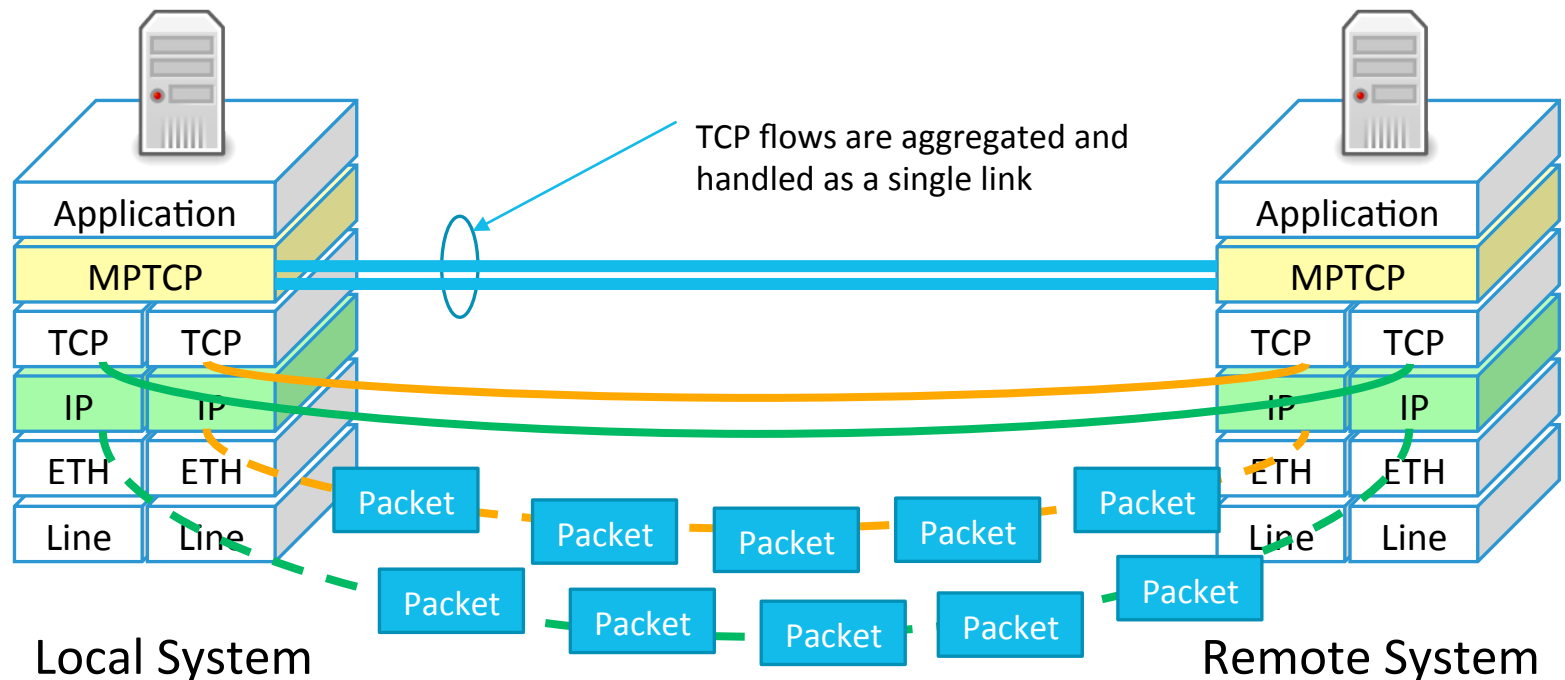
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2015/10/07, PRAGMA 29 Lightning Talk @Universitas Indonesia

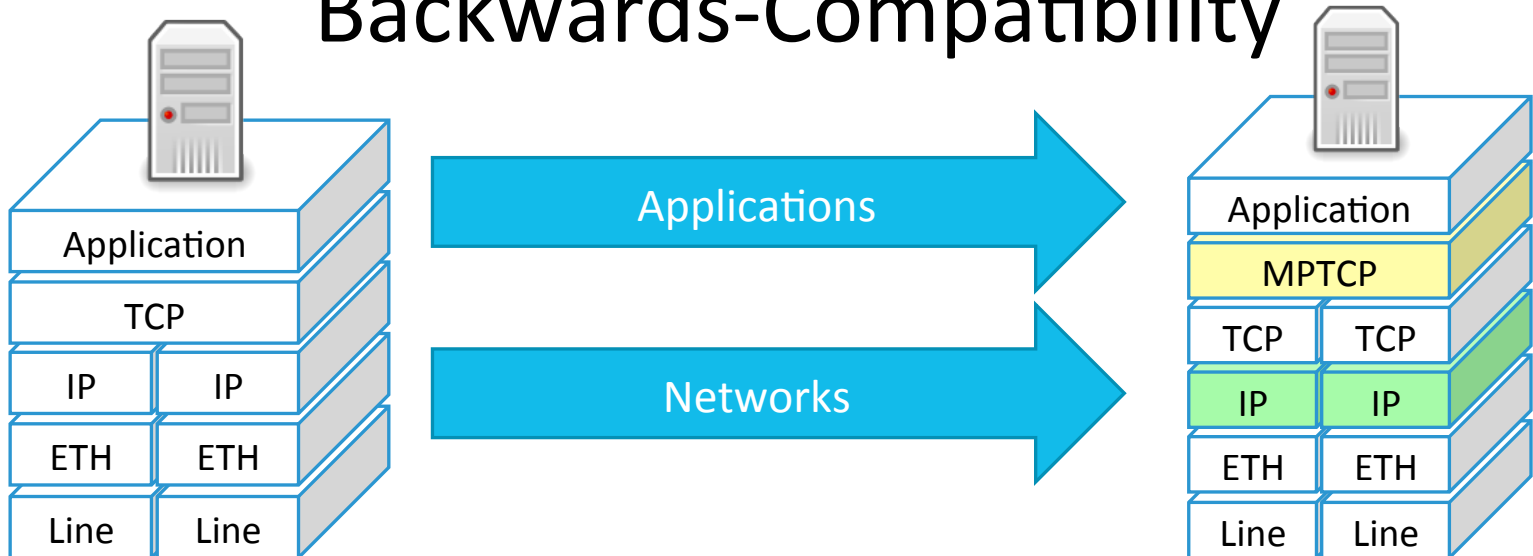
MPTCP aggregates TCP so application needs one socket.

Everything else is transparently handled by MPTCP in the OS kernel.



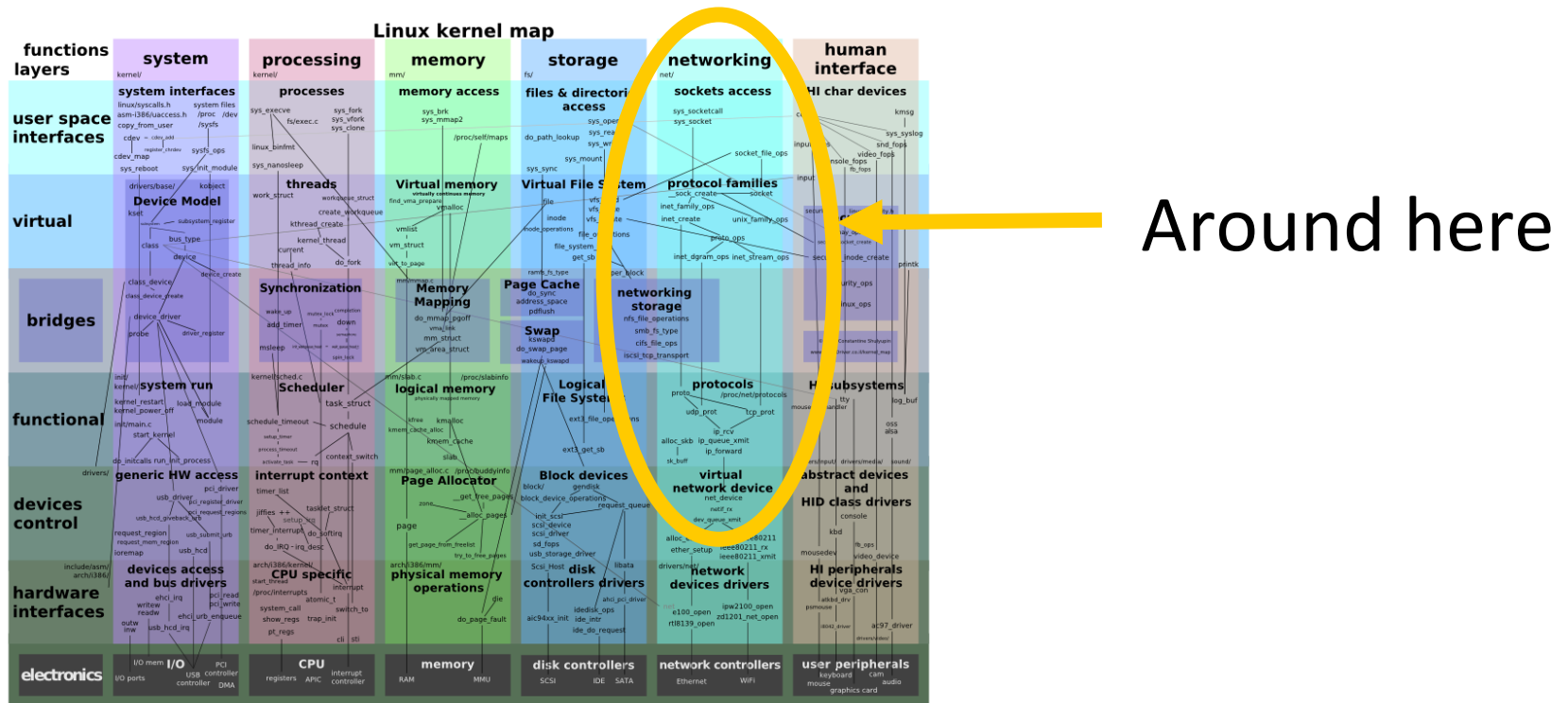
Benefits of MPTCP

Backwards-Compatibility



Limitation:
MPTCP needs OS modification

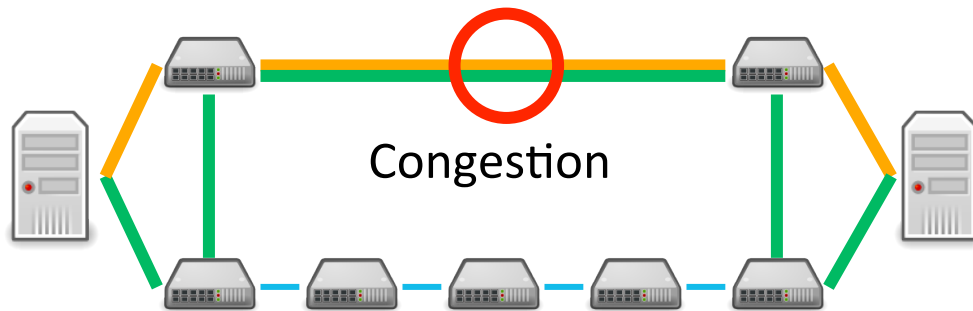
MPTCP must be implemented as a separate kernel, not even a module.
Support is confined usually into the active development community.



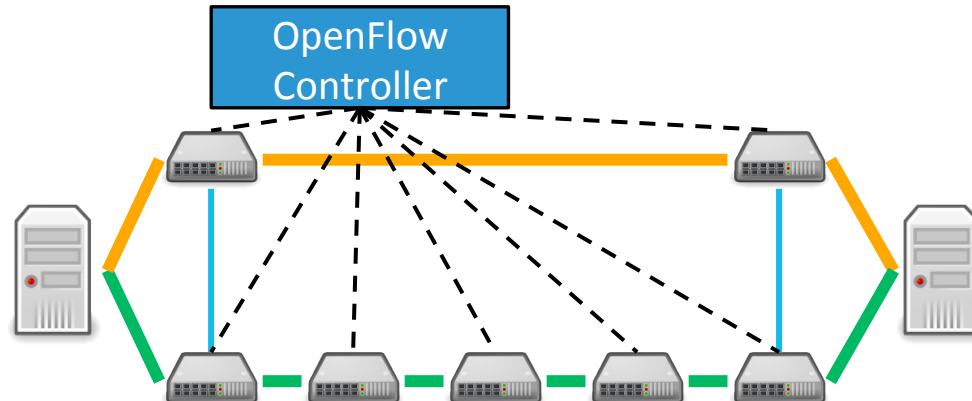
https://upload.wikimedia.org/wikipedia/commons/5/5b/Linux_kernel_map.png

Limitation:

Cannot route itself, may not be fully efficient with some networks



Congestion caused by traditional shortest-path routing



Multipath routing can avoid congestion

Algorithms

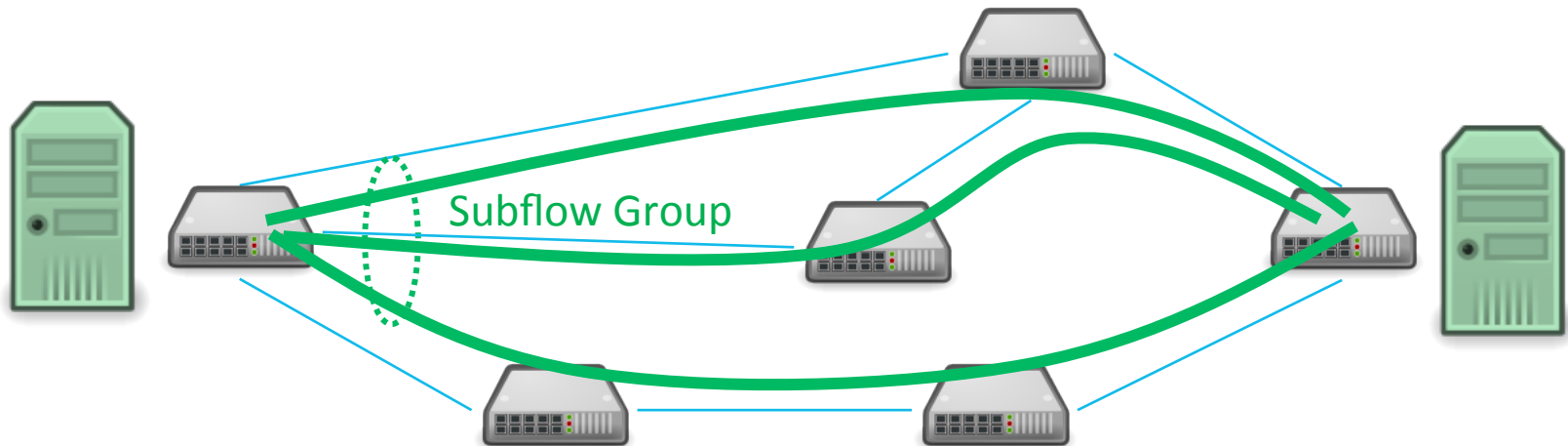
Routing

Flow
Management

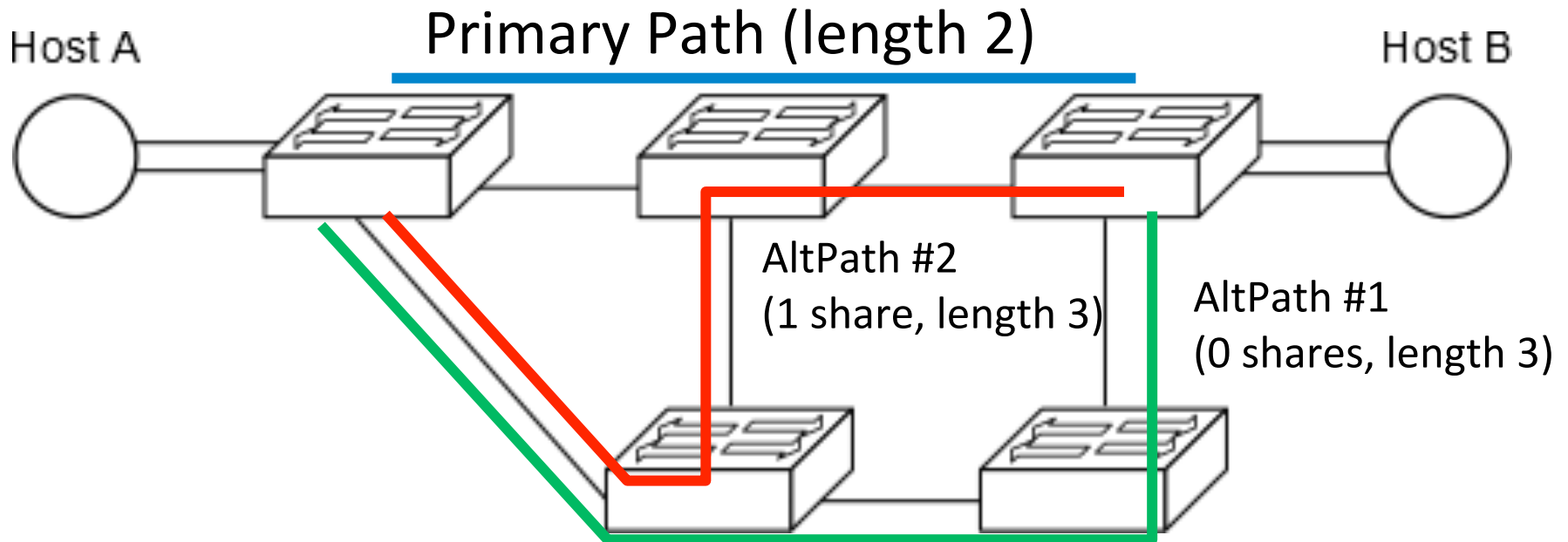
Goals

Flow management determines which subflow belongs to which MPTCP session

Routing algorithm distribute flows to different paths, minimizing path conflict that leads to bottlenecks.



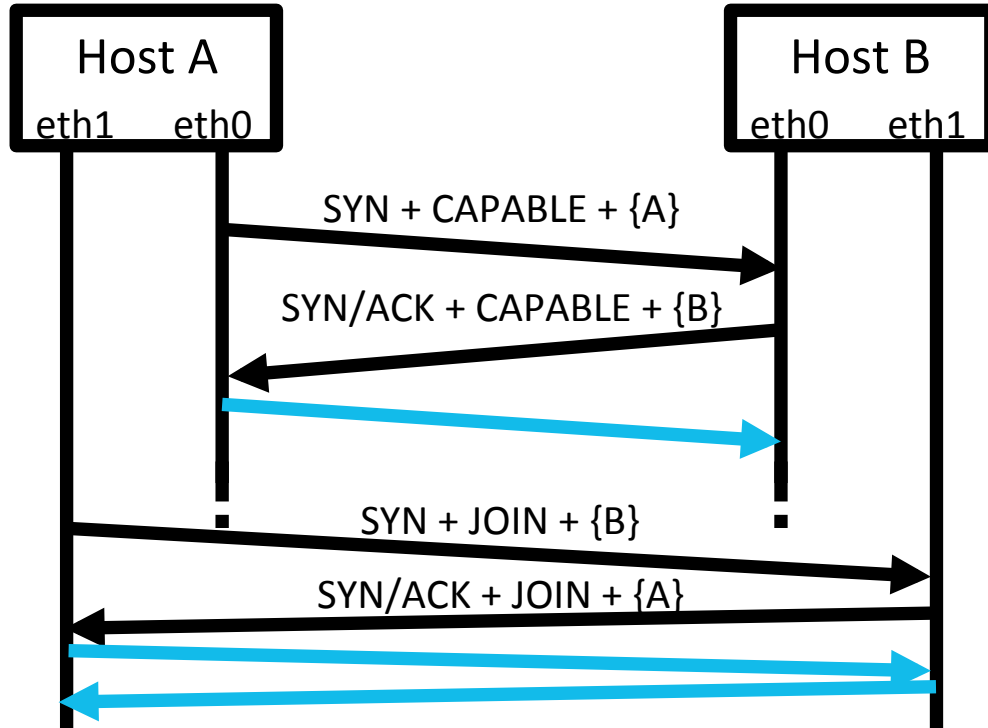
Routing Mechanism is based on shortest path



Sample routing results from our algorithm

The TCP Handshake

One of the most well-known topics about TCP is handshake. MPTCP expands on that by adding more information.



(1) Host introduce themselves by giving out its own keys.

(2) When hosts want to create another subflow, it sends the other party's key to indicate which MPTCP session it wants to talk to.

$\{X\}$ = X's MPTCP session key

pending_capable

From	To	key	pathset

pending_join

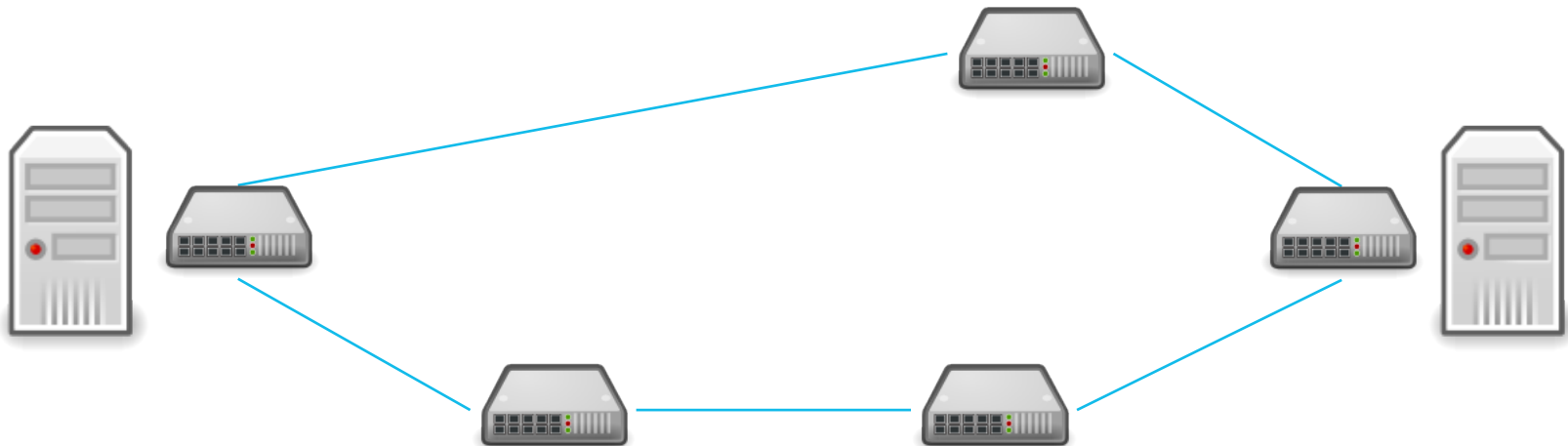
From	To	key	pathset

established

From	To		

X = X's IP address & Port

{X} = X's MPTCP Hash/Key



Host A

Host B

Routing

Flow
Management

pending_capable

From	To	key	pathset

pending_join

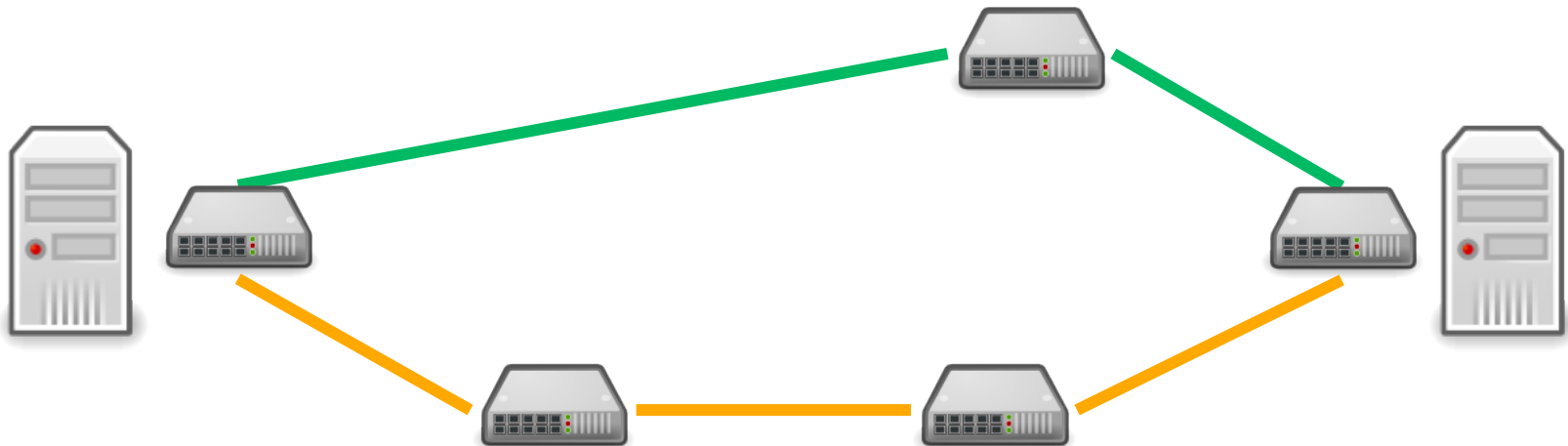
From	To	key	pathset
A	B	{B}	A=>B[1]
B	A	{A}	B=>A[1]

established

From	To	Pathset
{A}	{B}	A=>B[0]
{B}	{A}	B=>A[0]

X = X's IP address & Port

{X} = X's MPTCP Hash/Key



Host A

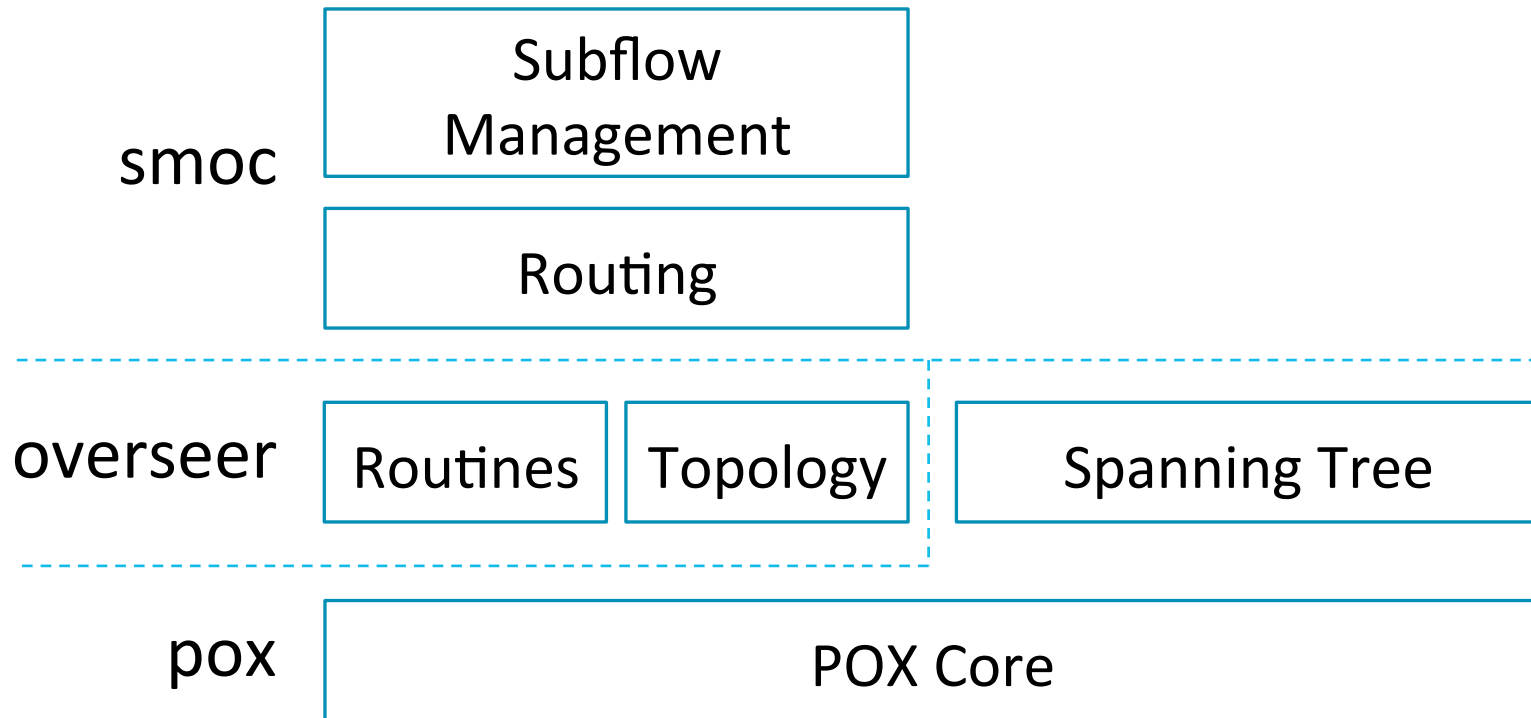
Host B

Routing

Flow
Management

smoc:

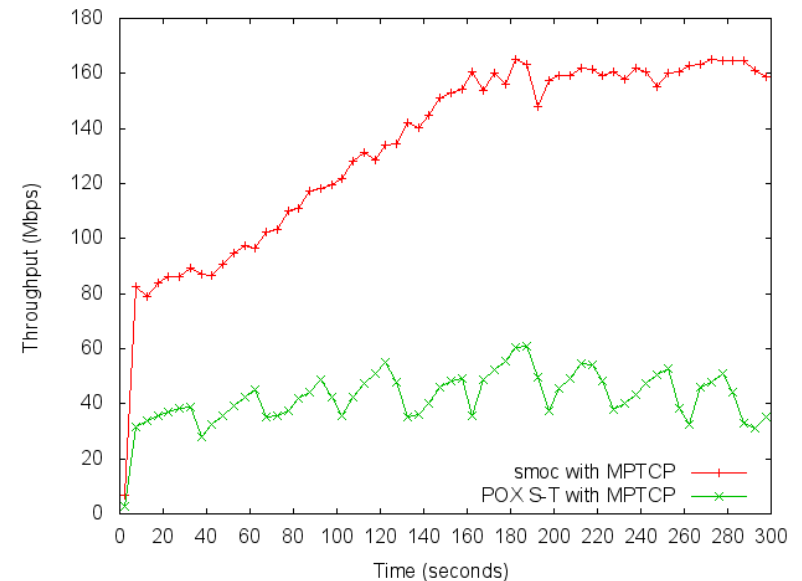
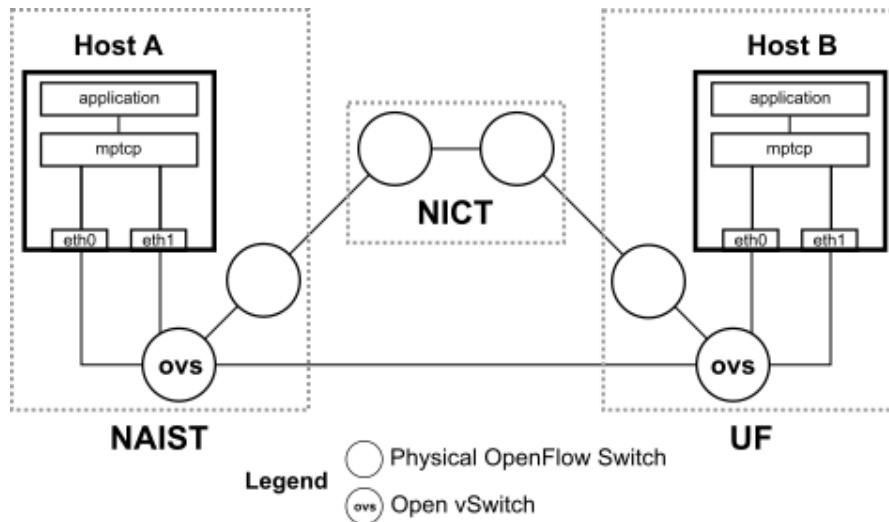
Simple Multipath OpenFlow Controller



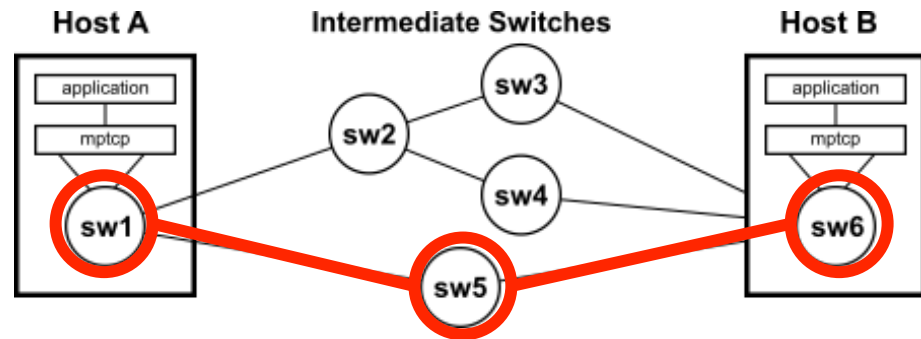
✂ networkx is used for additional path calculation

Experiment in PRAGMA

Experimental Network Testbed

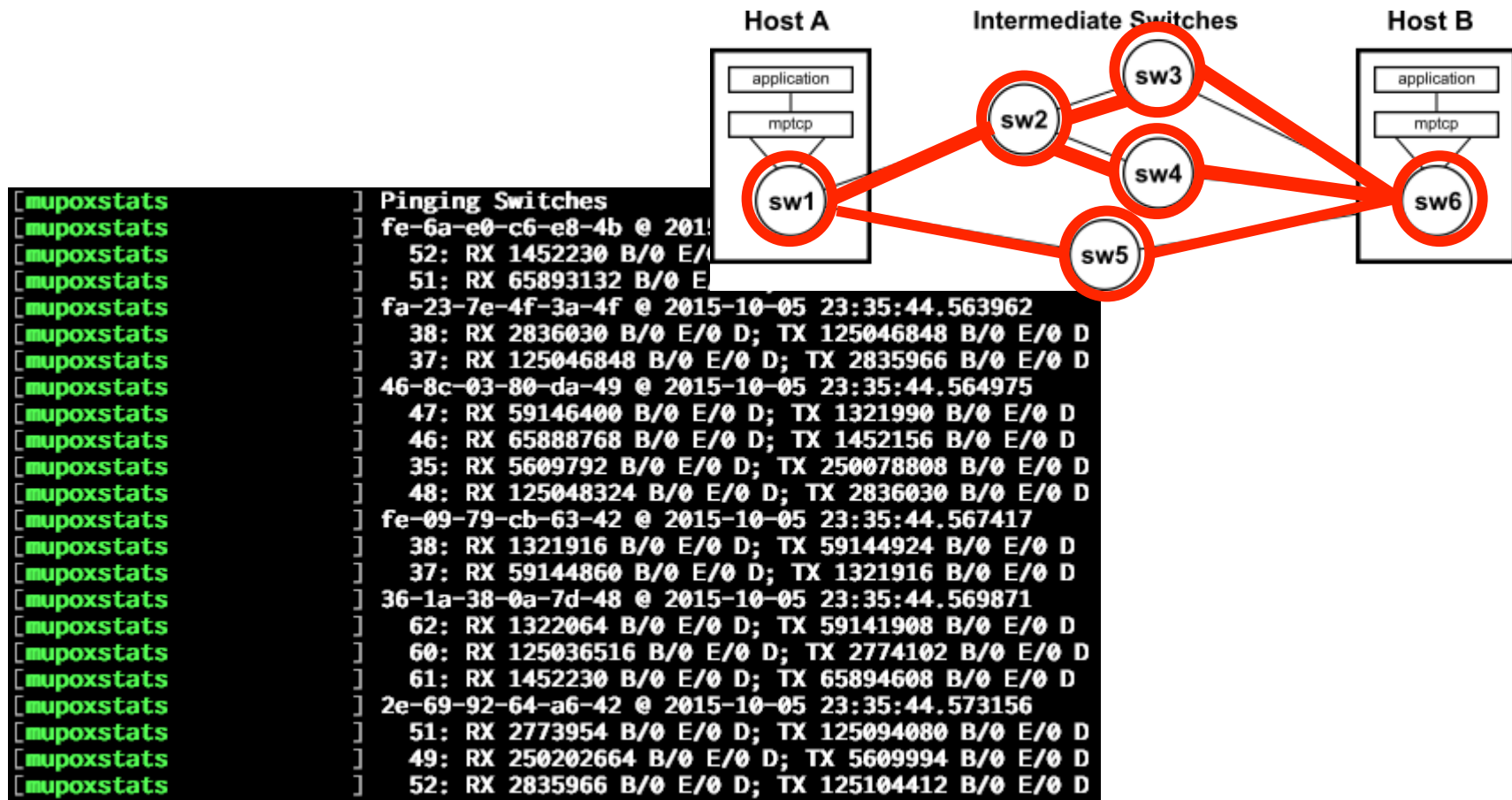


smoc w/o MPTCP

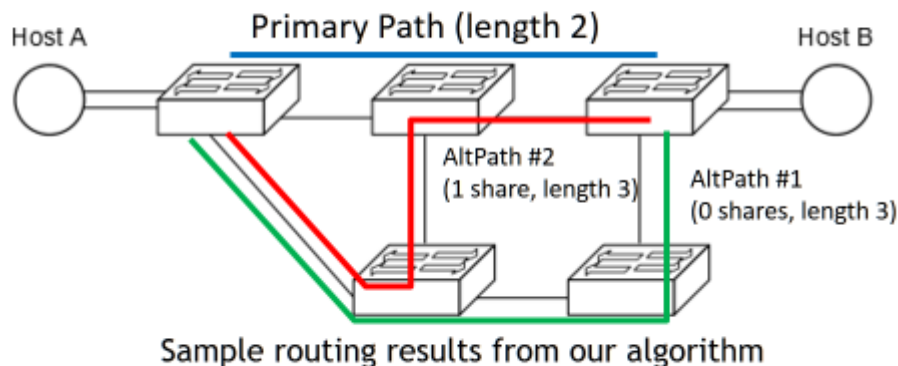


```
[mupoxstats] Pinging Switches
[mupoxstats] 46-8c-03-80-da-49 @ 2015-10-05 23:33:04.385336
[mupoxstats] 35: RX 816330 B/0 E/0 D; TX 103564262 B/0 E/0 D
[mupoxstats] 48: RX 103564220 B/0 E/0 D; TX 816416 B/0 E/0 D
[mupoxstats] fe-6a-e0-c6-e8-4b @ 2015-10-05 23:33:04.390581
[mupoxstats] 2e-69-92-64-a6-42 @ 2015-10-05 23:33:04.392020
[mupoxstats] 49: RX 102965694 B/0 E/0 D; TX 816352 B/0 E/0 D
[mupoxstats] 52: RX 816416 B/0 E/0 D; TX 102965758 B/0 E/0 D
[mupoxstats] fa-23-7e-4f-3a-4f @ 2015-10-05 23:33:04.395016
[mupoxstats] 38: RX 816416 B/0 E/0 D; TX 103331834 B/0 E/0 D
[mupoxstats] 37: RX 103331770 B/0 E/0 D; TX 816416 B/0 E/0 D
[mupoxstats] fe-09-79-cb-63-42 @ 2015-10-05 23:33:04.397678
[mupoxstats] 36-1a-38-0a-7d-48 @ 2015-10-05 23:33:04.399173
```

smoc with MPTCP

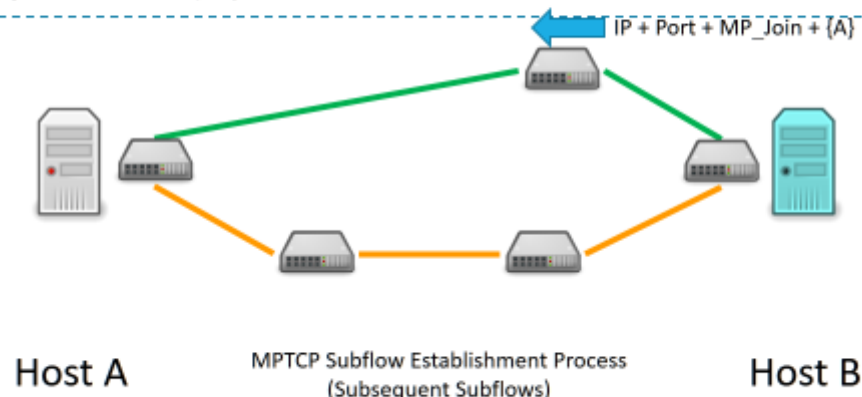


Routing Mechanism is based on shortest path

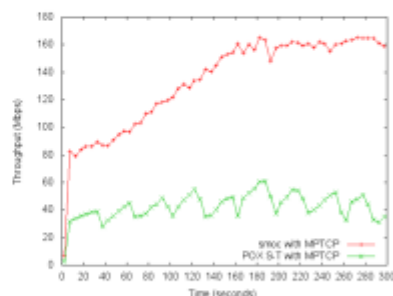
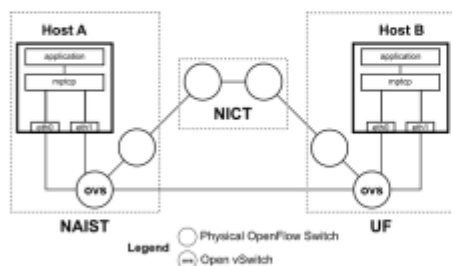


pending_capable				pending_join				established		
From	To	key	pathset	From	To	key	pathset	From	To	Pathset
				A	B	{B}	A=>B[1]	{A}	{B}	A=>B[0]
				B	A	{A}	B=>A[1]	{B}	{A}	B=>A[0]

X = X's IP address & Port
 {X} = X's MPTCP Hash/Key



Experiment in PRAGMA Experimental Network Testbed



Canned Results:
smoc with MPTCP

