

ANALYSIS OF LOAD BALANCING CLUSTER COMPUTING PERFORMANCE ON PROXMOX VE

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Motivation

This research aims to examine the performance of load balancing on a cluster. Load balancing is a solution for large access load and minimizes downtime in serving requests from users. Load balancing distributes loads of traffic evenly to servers with particular algorithms.

Our Works



Scenario 1

Implemented using Linux Virtual Server (LVS) topology via direct routing on Proxmox VE and use round-robin algorithm for load balancing

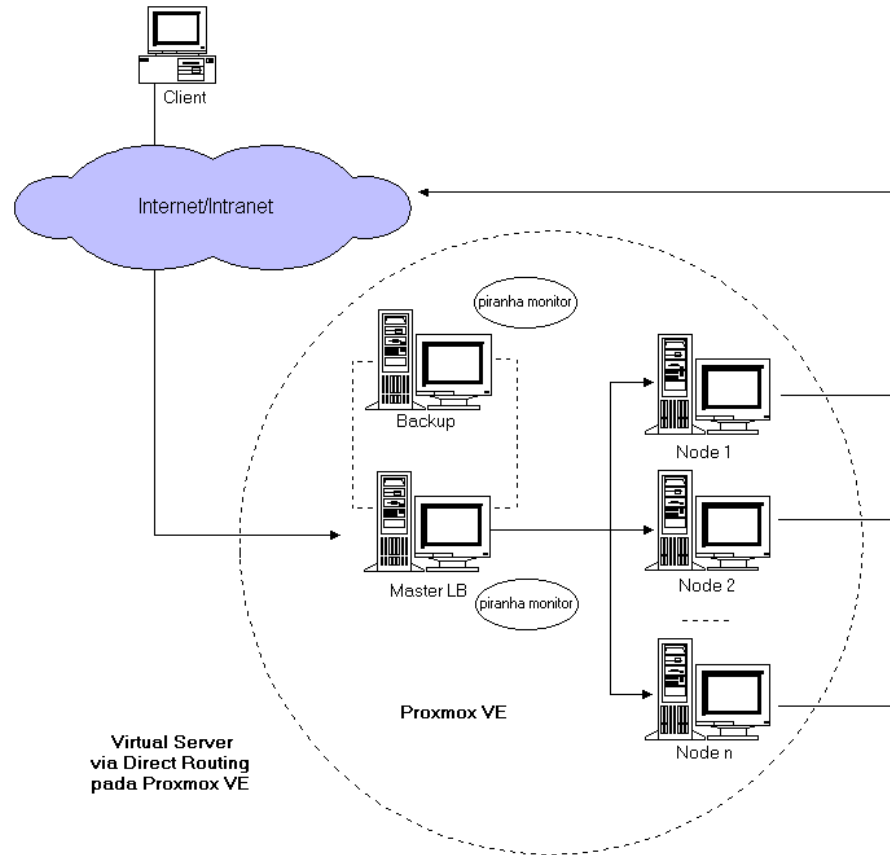
Scenario 2

Performed to find out the number of requests per second and packet loss with load of 250, 500, and 1000 users as many as 3 times for 1 minute between single server with load balancing on cluster computing.

Scenario 3

Intends to prove that the failover on load balancing can manage to requests copying of the process and configuration from Master to Backup with an average value of downtime.

Topology



Hardware & Software

Host

CPU Octa Core @ 2.5 GHz
Memory 4 GB DDR3
Storage 300 GB (IDE)
Operating System Proxmox VE 4.1
Network Adapter 1 : Custom (PCnet –FAST III)

CPU Single Core @ 2.5 GHz
Memory 500 MB DDR3
Storage 50 GB (Local-LVM)
Operating System CentOS 6.0 Final (i386-minimal)
Network Adapter 1-4 : Bridge (vmb)br
Ipsadm (Piranha)
LAMP (Linux-Apache-MySQL-PHP)

Guest

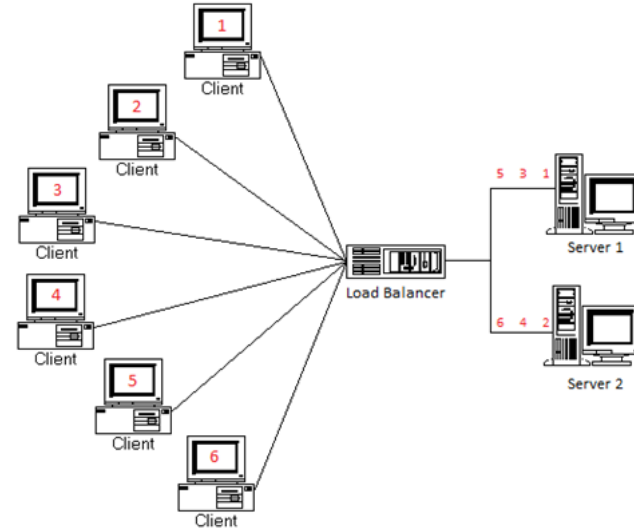
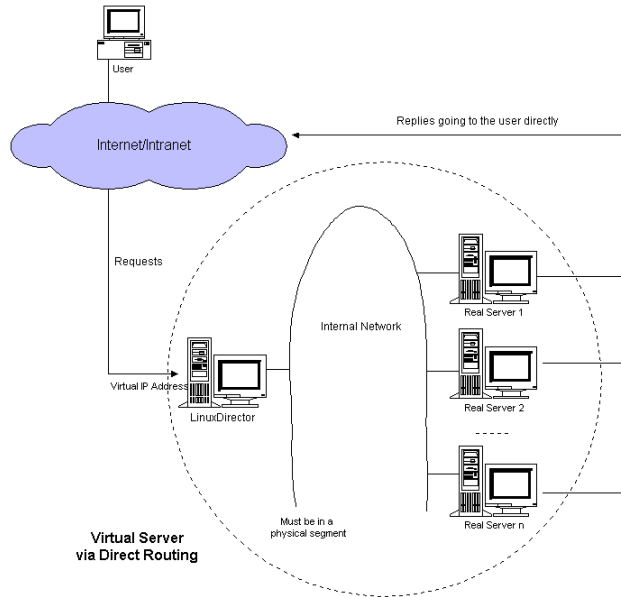
User

CPU Octa Core @ 1.9 GHz
Memory 8 GB DDR3
Storage 500 GB
Operating System Windows 10
Network Adapter : VB Ethernet Adapter
WebServer Stress Tool 8
Net Uptime Monitor

Conclusion

Scenario 1

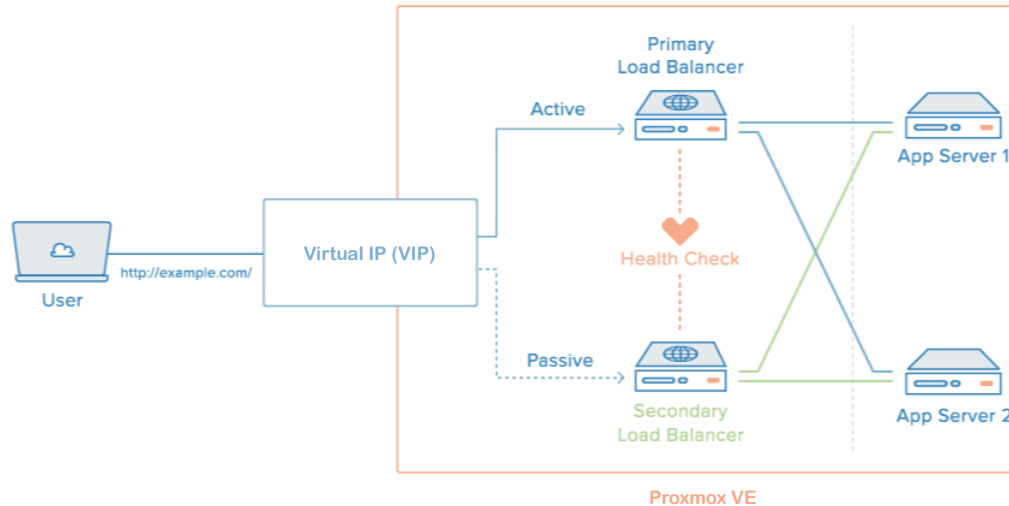
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Conclusion

Scenario 1

Implemented using Linux Virtual Server (LVS) topology via direct routing on Proxmox VE and use round-robin algorithm for load balancing



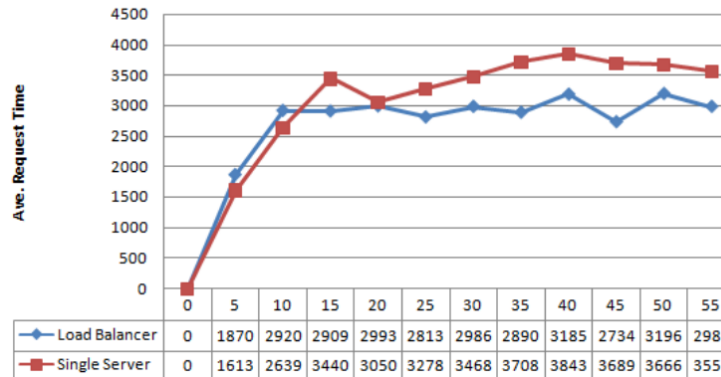
- 1 Active/Passive Cluster is healthy
- 2 Primary node fails
- 3 Floating IP is assigned to Secondary node

Conclusion

Scenario 2

Performed to find out the number of requests time and packet loss with load of 250, 500, and 1000 users as many as 3 times for 1 minute between single server with load balancing on cluster computing.

Request 250 User

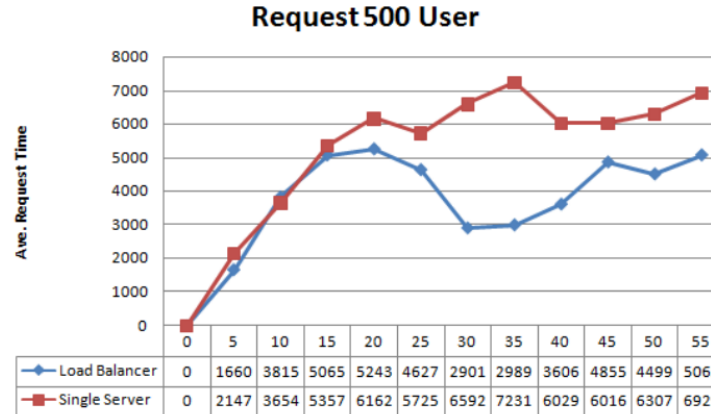


Name	Clicks	Errors	Errors [%]	Time Spent [ms]	Avg. Click Time [ms]
Load Balancer	4,291	0	0.00	12,325,289	2,872
Single Server	3,941	36	0.91	12,792,645	3,276

Conclusion

Scenario 2

Performed to find out the number of requests time and packet loss with load of 250, 500, and 1000 users as many as 3 times for 1 minute between single server with load balancing on cluster computing.

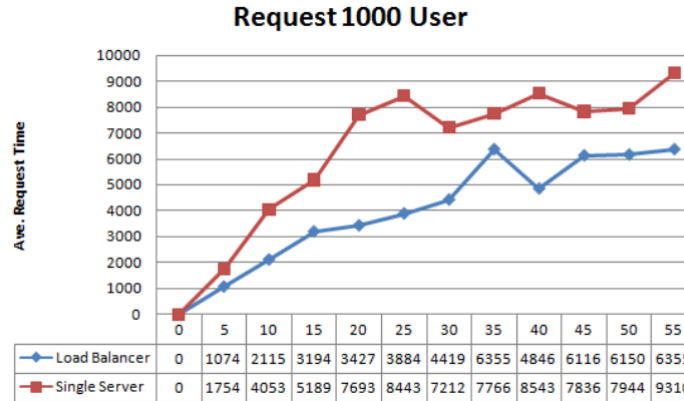


Name	Clicks	Errors	Errors [%]	Time Spent [ms]	Avg. Click Time [ms]
Load Balancer	3,784	0	0.00	15,220,957	4,022
Single Server	3,594	239	6.65	19,424,124	5,790

Conclusion

Scenario 2

Performed to find out the number of requests time and packet loss with load of 250, 500, and 1000 users as many as 3 times for 1 minute between single server with load balancing on cluster computing.



Name	Clicks	Errors	Errors [%]	Time Spent [ms]	Avg. Click Time [ms]
Load Balancer	3,395	0	0.00	14,138,671	4,165
Single Server	3,934	884	22.47	22,170,186	7,269

Conclusion

Scenario 3

Intends to prove that the failover on load balancing can manage to requests copying of the process and configuration from Master to Backup with an average value of downtime.

Failover	Downtime
1	16 s
2	17 s
3	17 s
4	19 s
5	16 s
6	13 s
7	18 s
8	19 s
9	17 s
10	14 s

Thank You

The background features a large, solid blue shape in the upper left corner. Below and to the right of this, there is a series of overlapping, angular shapes in shades of blue and light gray, creating a sense of depth and movement. The overall composition is clean and modern.