

## Lecture 8: Quantum Computing (18 July - End)

*Lecturer: Valentine Kabanets**Scribe: Lily Li*

## 8.1 Quantum Computing

### 8.1.1 Qubit

Let  $|0\rangle$  measures 0 and  $|1\rangle$  measures 1. You should think of each state as an energy level of the electron. However, in reality these states are continuous —rather than discrete. Consider a superposition of  $|0\rangle$  and  $|1\rangle$ . We model quantum mechanics as  $\alpha \cdot |0\rangle + \beta \cdot |1\rangle$  where  $\alpha, \beta \in \mathcal{C}$  such that  $|\alpha|^2 + |\beta|^2 = 1$ . When we measure the qubit, the duality collapses and we just get classical bits with probability:

$$\begin{cases} 0 & \text{with probability } |\alpha|^2 \\ 1 & \text{with probability } |\beta|^2. \end{cases}$$

Observe that  $\alpha, \beta$  are *not* the probabilities, but their squares *are*. Thus even though the probabilities are non-negative, the actual values  $\alpha$  and  $\beta$  are complex (and can be negative).

### 8.1.2 Quantum Registers

For  $m$  qubits, we can describe a quantum register as

$$\sum_{x \in \{0,1\}^m} \alpha_x \cdot x$$

where  $\alpha_x$  is the amplitude of bit string  $x$  and  $\sum |\alpha_x|^2 = 1$ . Observe that for  $m$  qubits, the quantum system can be described by a vector in  $\mathcal{C}^{2^m}$  (since the coefficient of each  $\alpha_x$  is one variable). This gets exceedingly big exceedingly fast.

### 8.1.3 Quantum Operations

First let us describe the operation  $F$  on the quantum system equivalent to negation on the classical bits. As it turns out that these operations  $F$  are exactly unary matrices. Generally, it is too much to define  $F$  for a quantum system with  $m$  qubits even if  $m$  is modest. Thus we must restrict the number of qubits we operate upon at one time. For us, each operation will be applied to at most three qubits. Thus  $F$  will be a  $2^3 \times 2^3$  or  $8 \times 8$  matrix. It is unclear if each operator  $F$  is realizable.

For bits  $|0\rangle$  and  $|1\rangle$  observe that

### 8.1.4 Complexity BQP

**Definition 8.1** Let  $f : \{0,1\}^n \rightarrow \{0,1\}^n$

Consider next the  $\wedge$  operation. Classically, we have  $x, y \in \{0, 1\}$  such that  $\text{AND}(x, y) = x \wedge y$ . In quantum computation we need the Trefoli gate which takes in three qubits  $x, y, z$  and outputs  $\text{QAND}(x, y, z) = (x, y, z \oplus (x \wedge y))$ .