Algorithms & Data Structure
Day 6: Tree

Kiran Waghmare

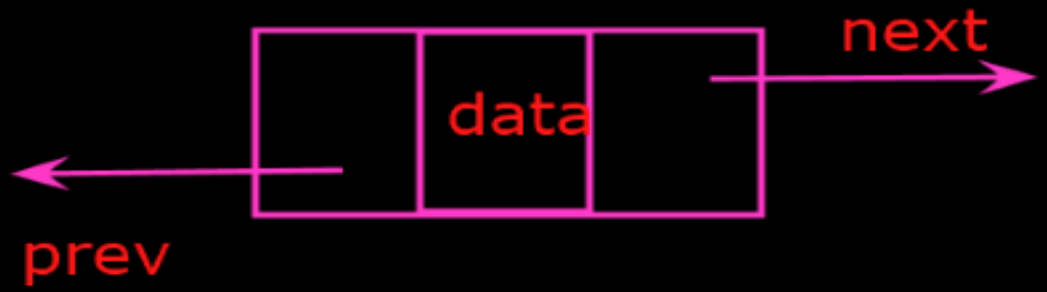
Imagination is more important than knowledge

– Albert Einstein

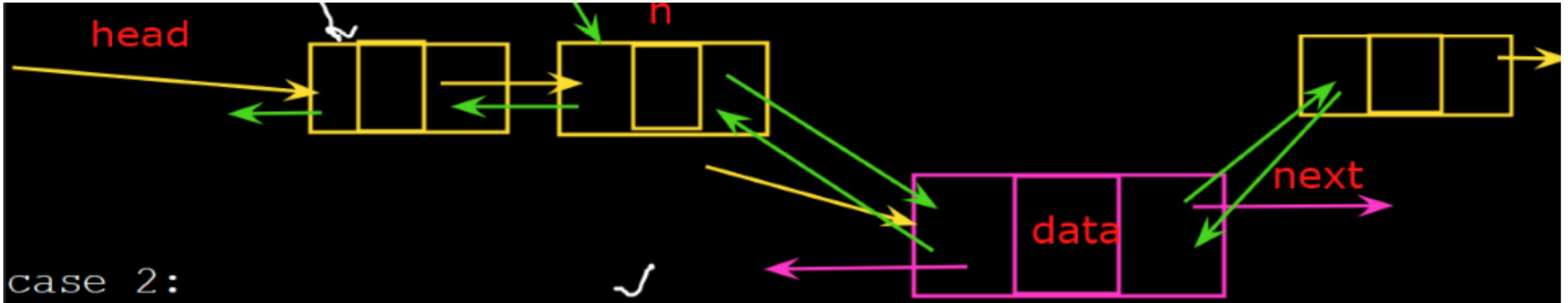
Node structure:

```
class Node
{
    int data;
    Node next;
    Node prev;

    Node(int d)
    {
        data = d;
        next = null;
        prev = null;
    }
}
```



this.data = d;

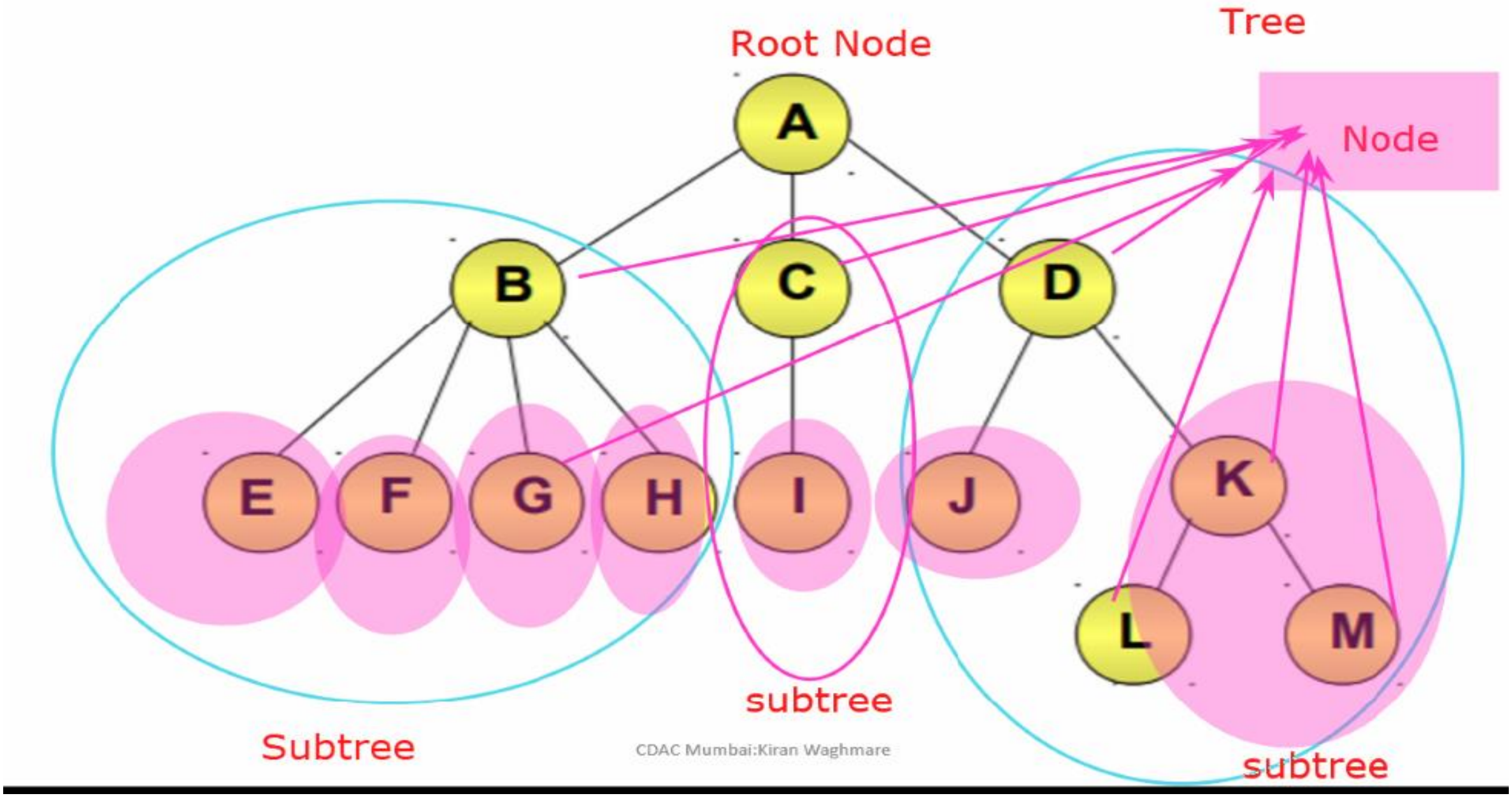


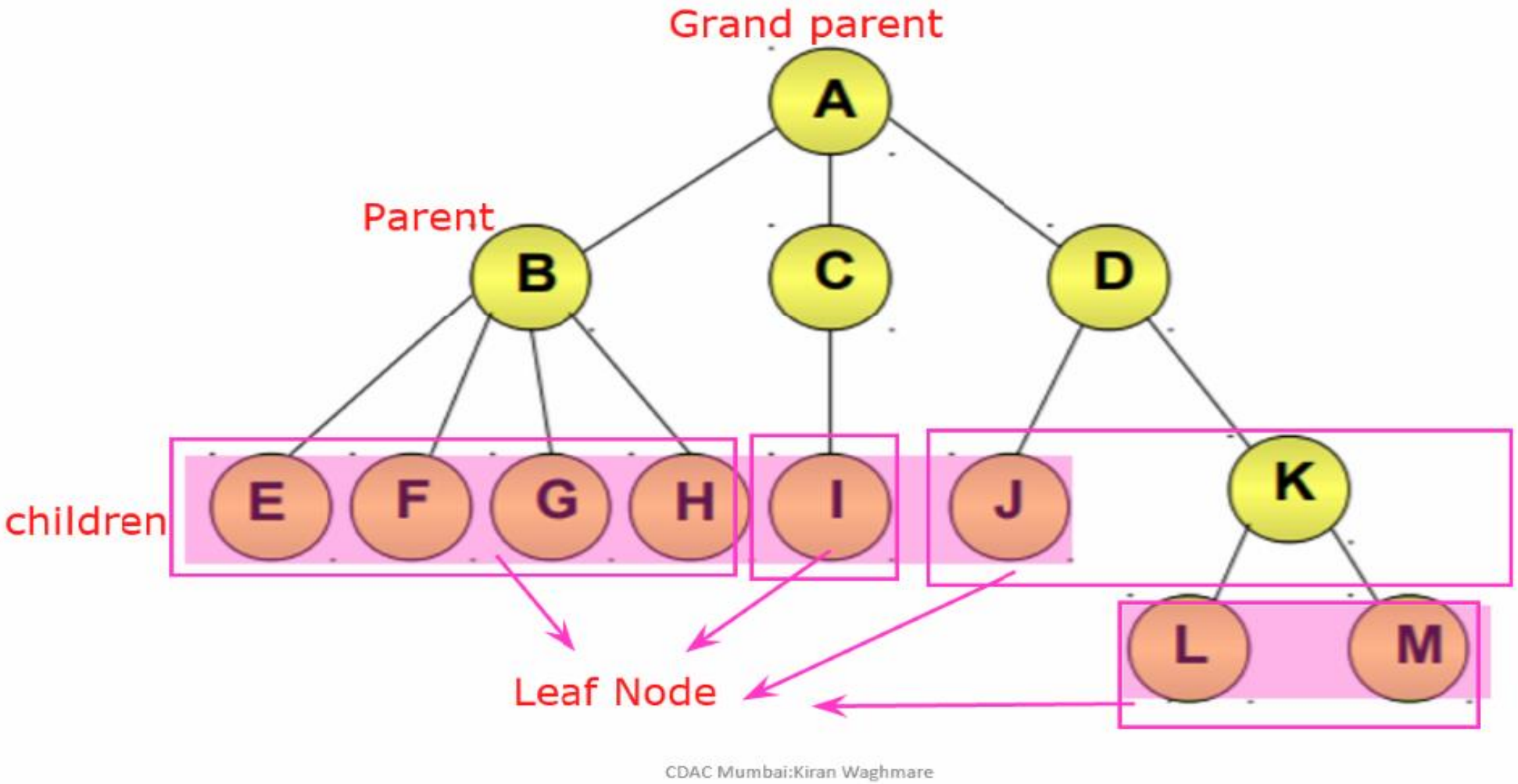

```
case 2:
void insertAfter(Node n,int new_data)
{
    if(n = null)
        {return;}
    Node new_node = new Node(new_data);
    new_node.next = n.next;
    n.next = new_node;
    new_node.prev = n;
    new_node.next.prev = new_node;

}
```

this.data = d;


```
void delete(Node n)
{
    if(head == null)
        return;
    if(head == n)
    {
        head = n.next
    }
    //It is not a last node
    if(n.next != null)
    {
        n.next.prev = n.prev;
    }
    if(n.prev != null)
    {
        n.prev.next = n.next;
    }
    return;
```


Root Node
Tree
Node
A
B
C
D
E
F
G
H
I
J
K
L
M
subtree
Subtree
CDAC Mumbai:Kiran Waghmare
subtree



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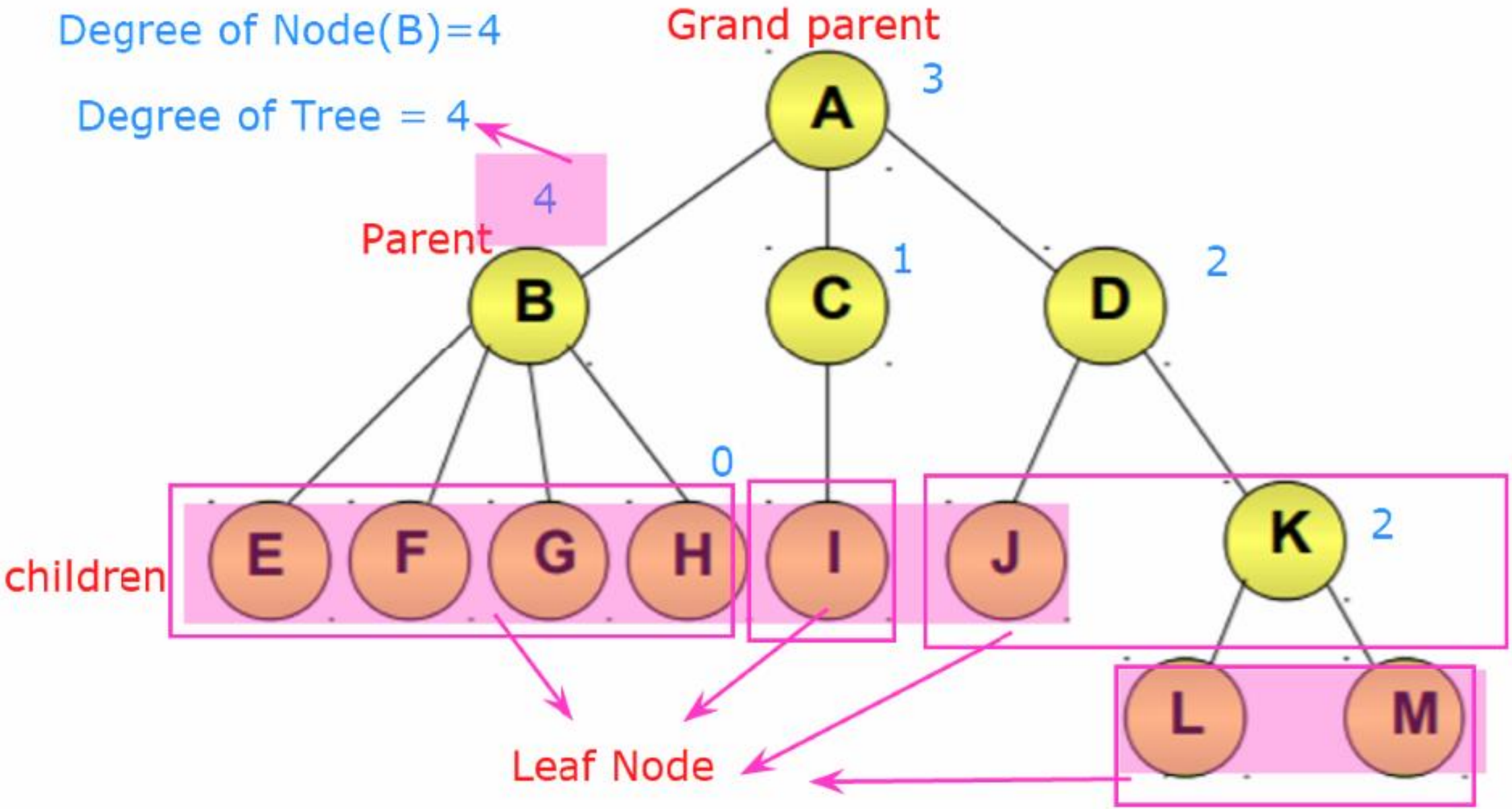
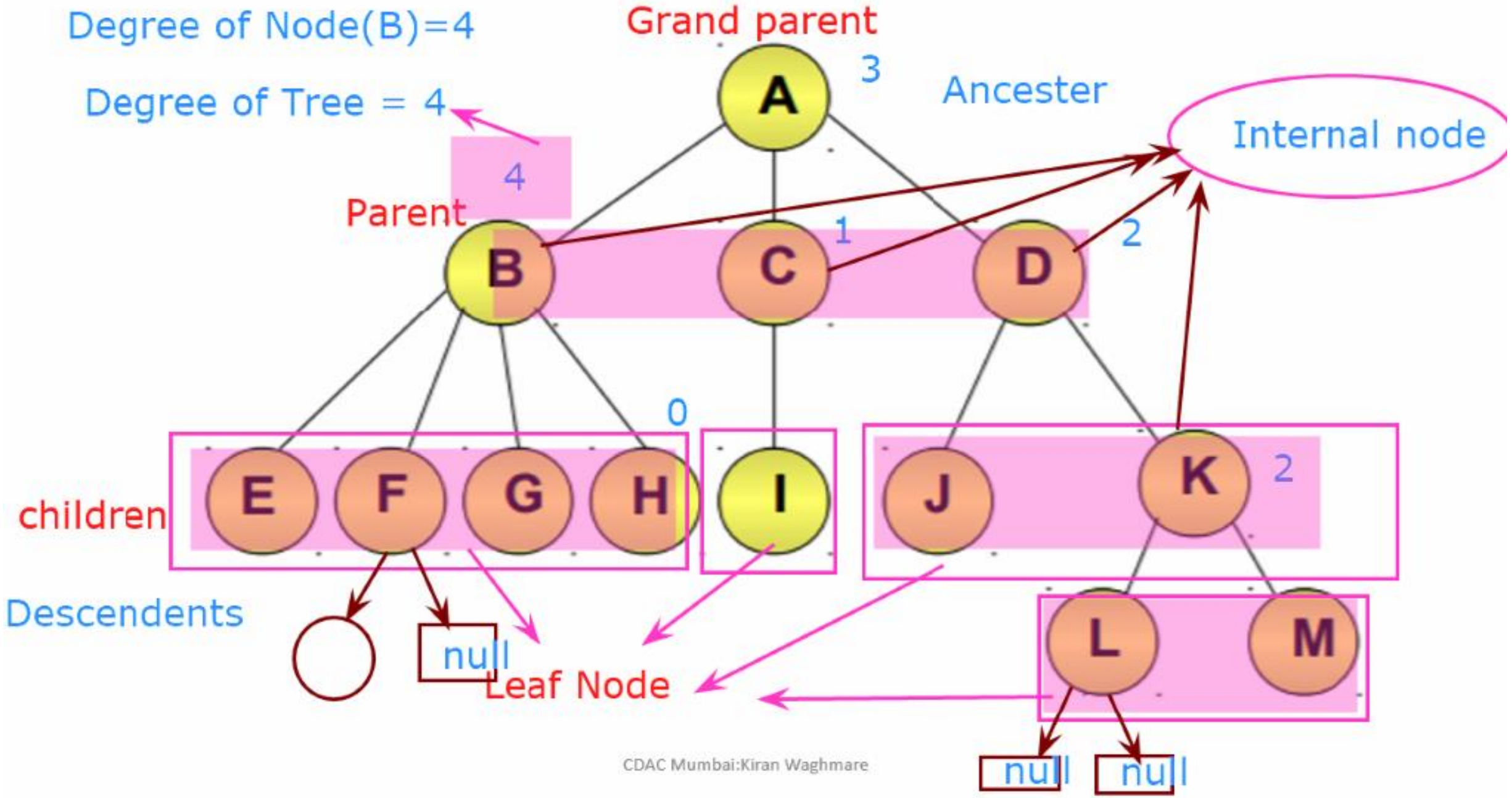

Degree of Node(B)=4
Degree of Tree = 4
4
Grand parent
3
A
Parent
B
C
1
D
2
0
children
E
F
G
H
I
J
K
2
L
M
Leaf Node


Degree of Node(B)=4
Degree of Tree = 4
Grand parent
3
A
Ancester
Internal node
4
Parent
B
C
1
D
2
0
children
E
F
G
H
I
J
K
2
Descendents
null
Leaf Node
L
M
null
null
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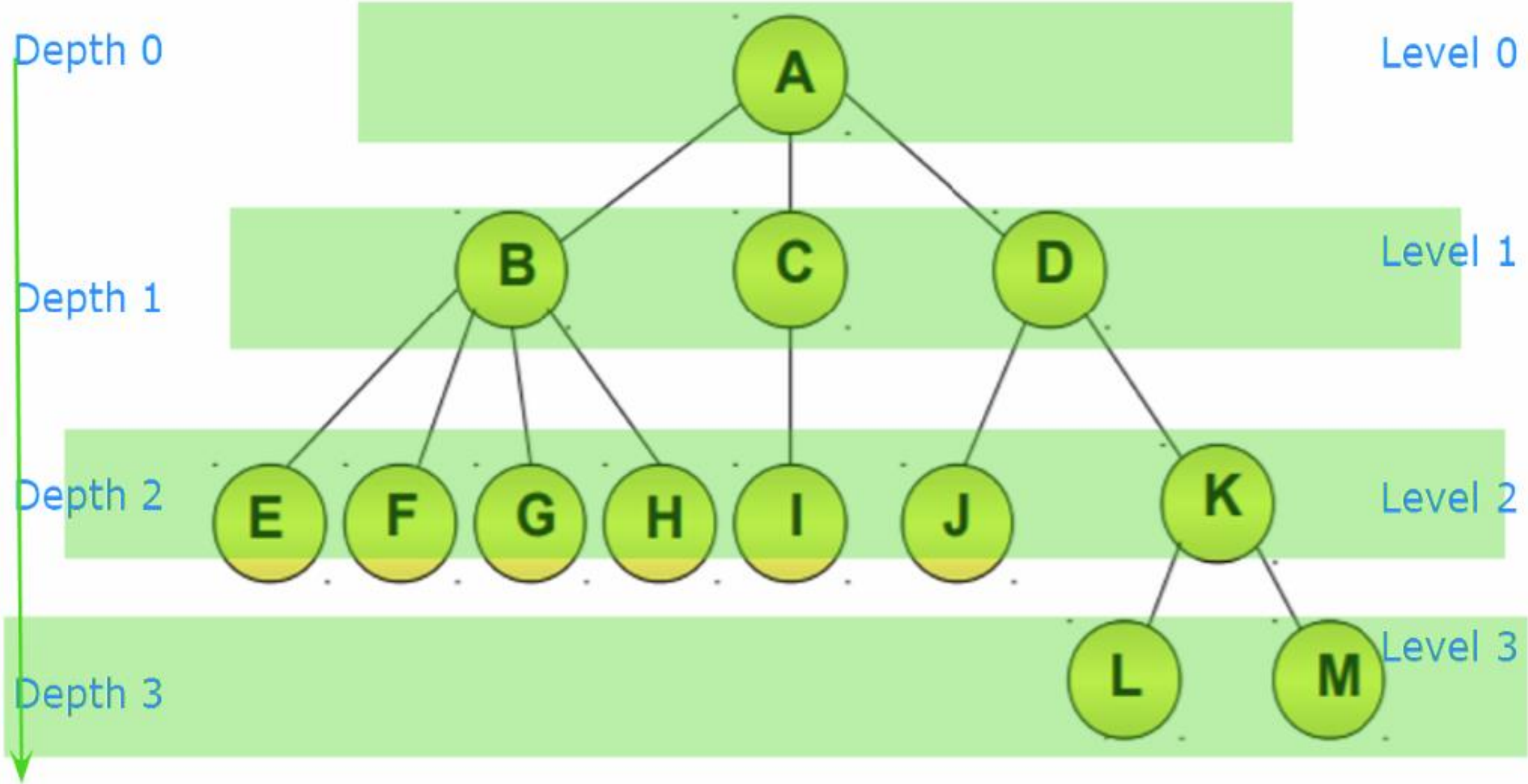
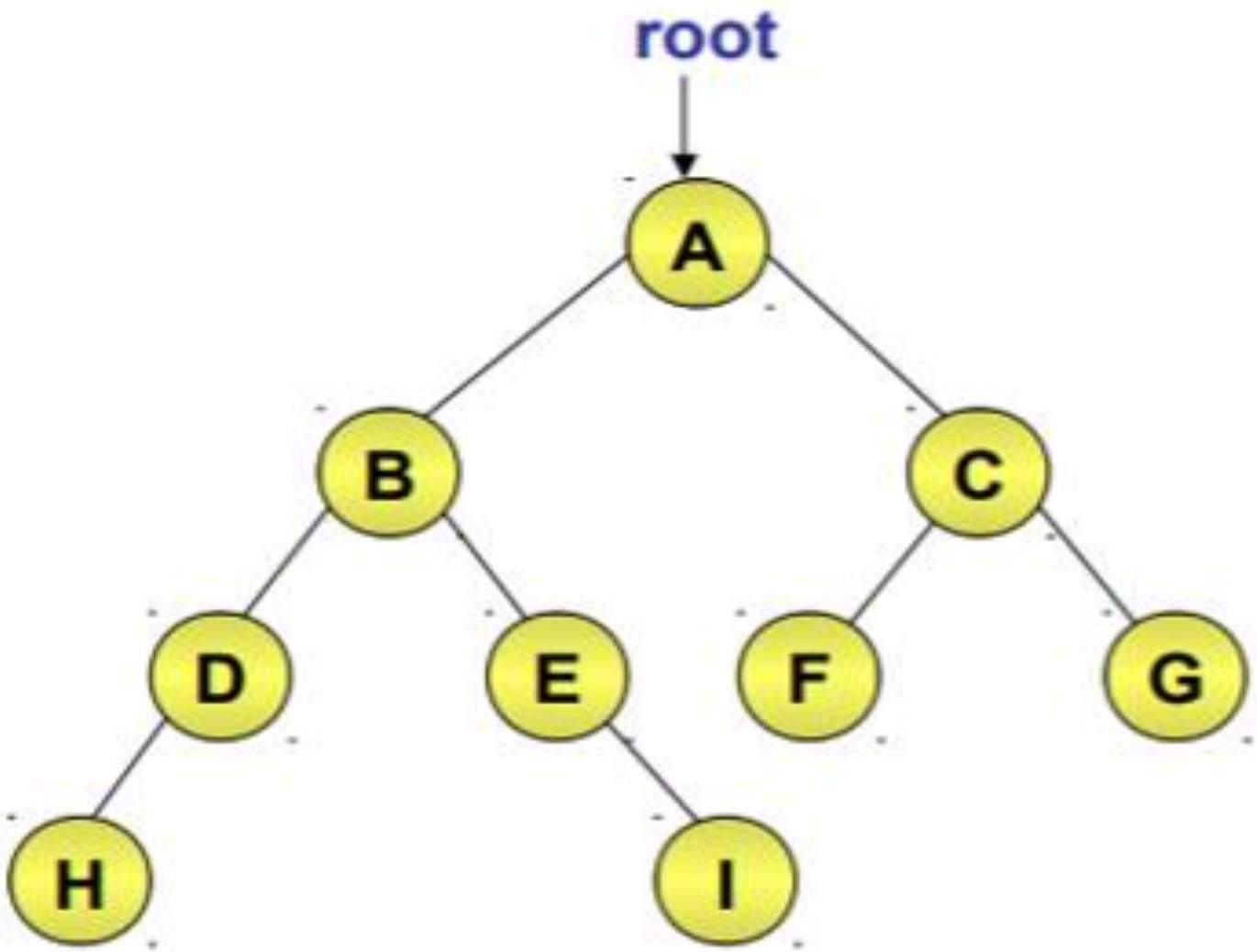

Depth 0
Level 0
A
Depth 1
Level 1
B
C
D
Depth 2
Level 2
E
F
G
H
I
J
K
Level 3
L
M
Depth 3

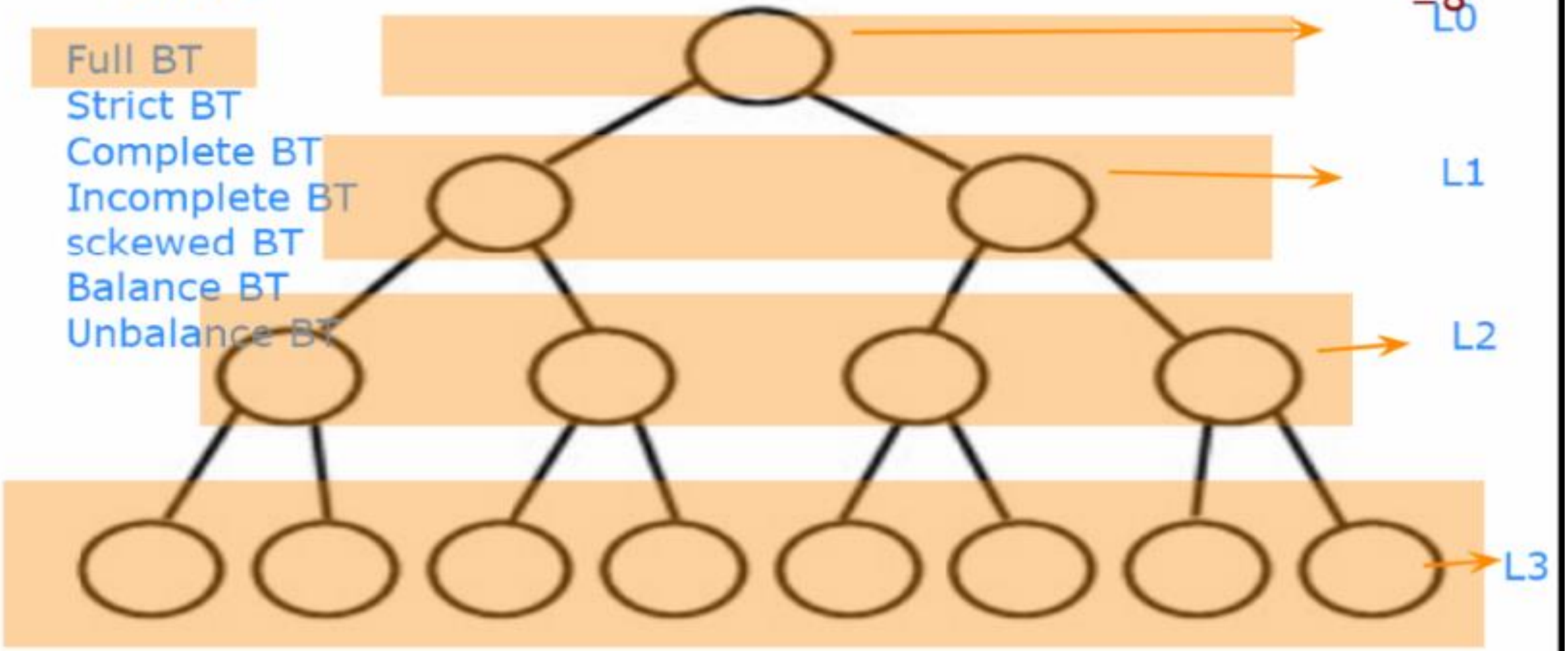

root
A
B
C
D
E
F
G
H
I

The max no. o f nodes on level i of a binary tree is 2^i, i>=0

(0,1,2)

2^i
=2^3
=8

Full BT
Strict BT
Complete BT
Incomplete BT
sckewed BT
Balance BT
Unbalance BT



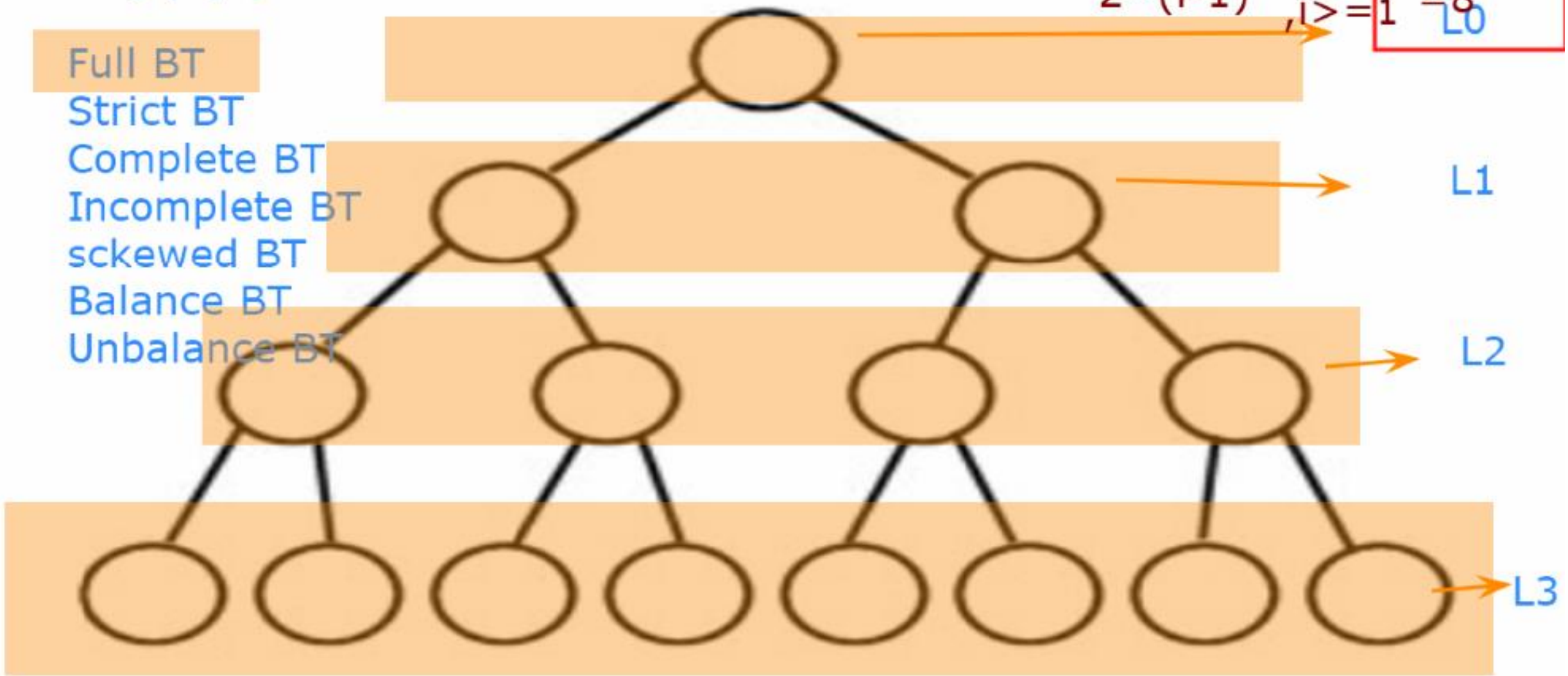
The max no. o f nodes on level i of a binary tree is 2^i, i>=0

(0,1,2)

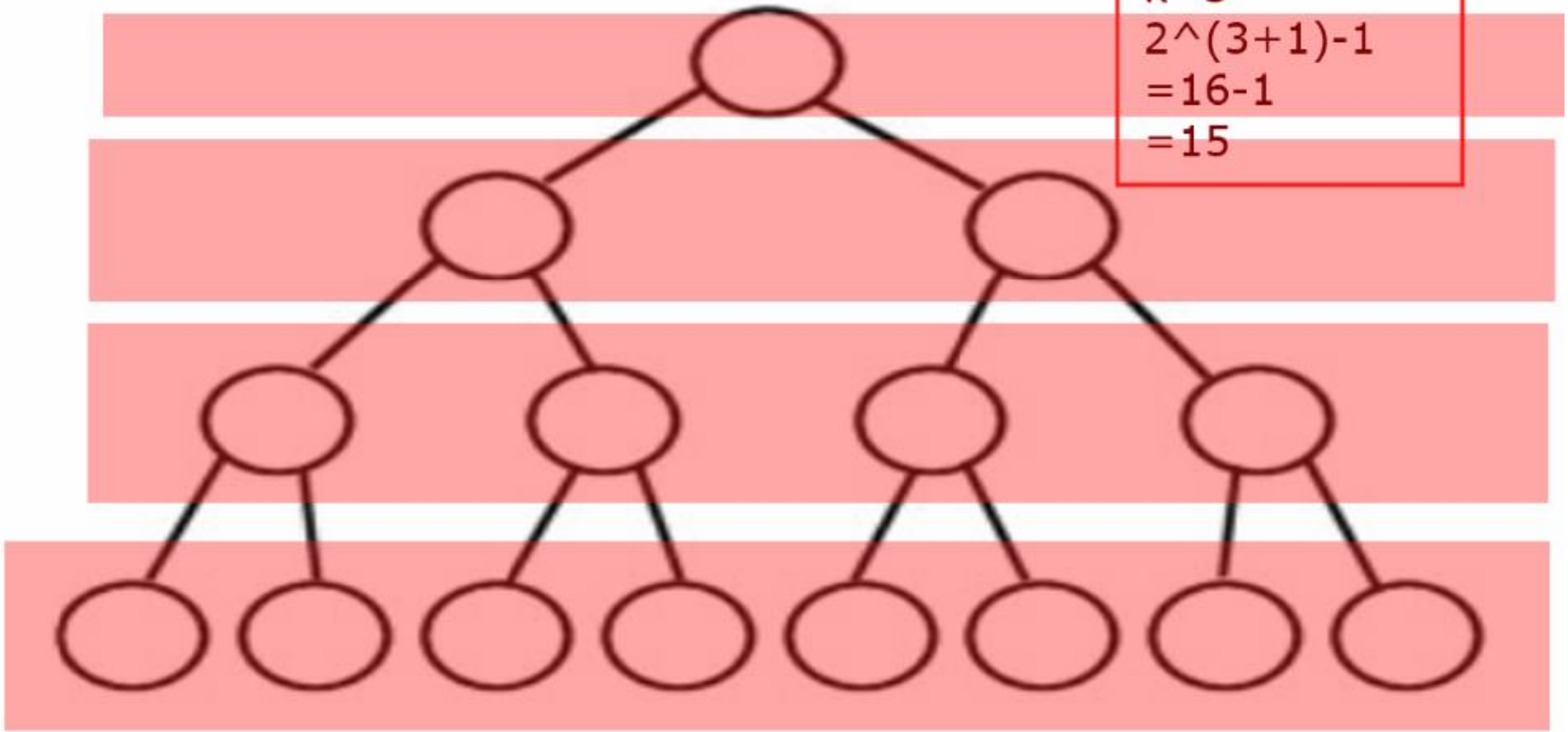
2^i
=2^3
=8

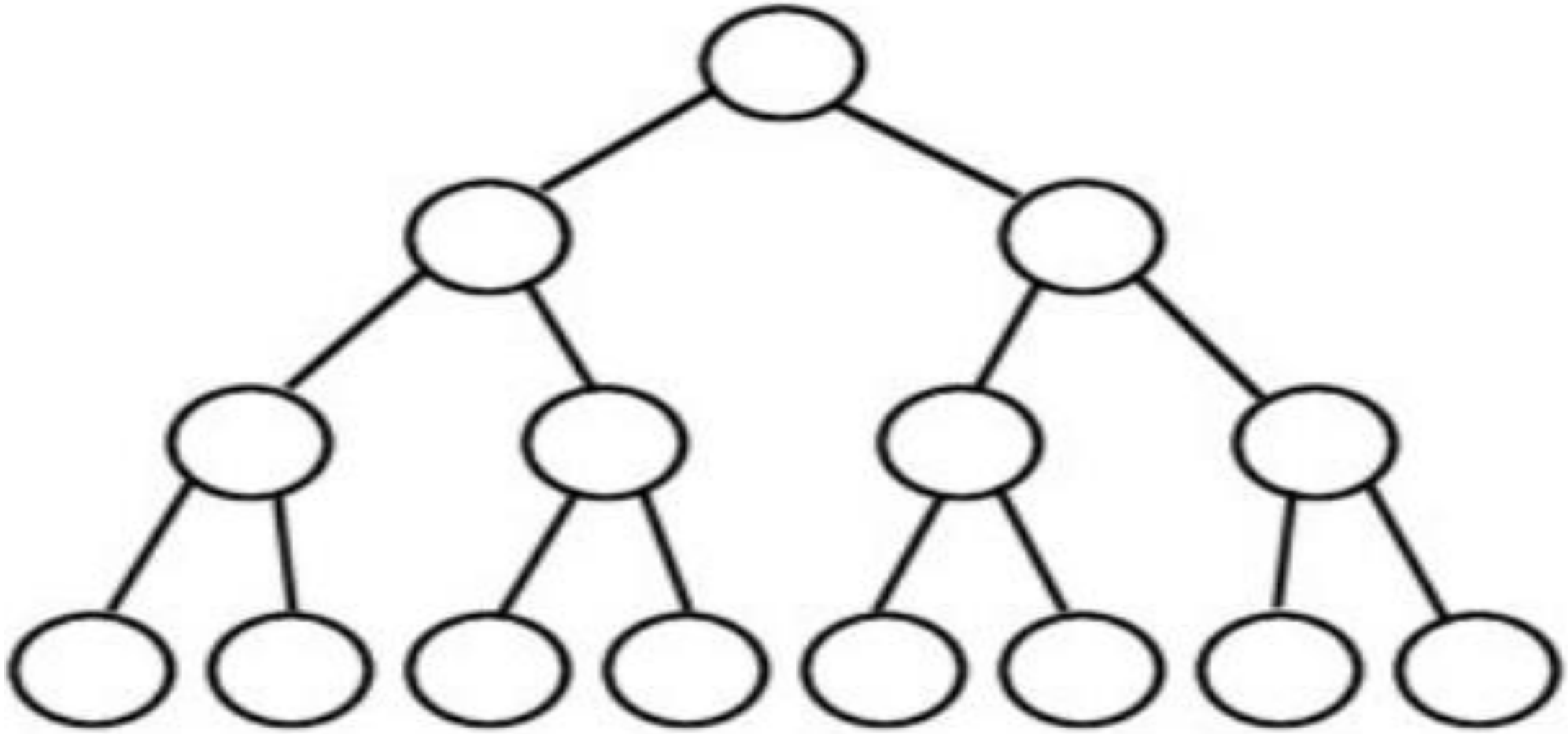
2^(i-1) ,i>=1

Full BT
Strict BT
Complete BT
Incomplete BT
sckewed BT
Balance BT
Unbalance BT



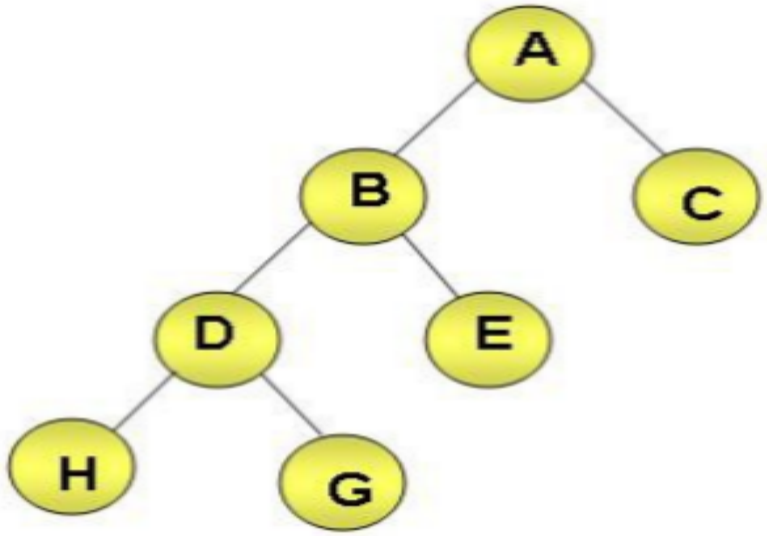
The max no nodes in a BT of depth k =>(2^(k+1))-1, k>=0 1

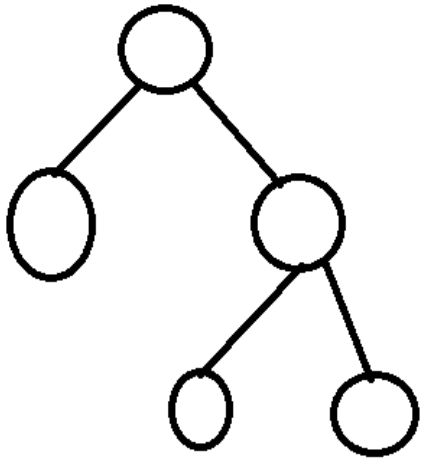




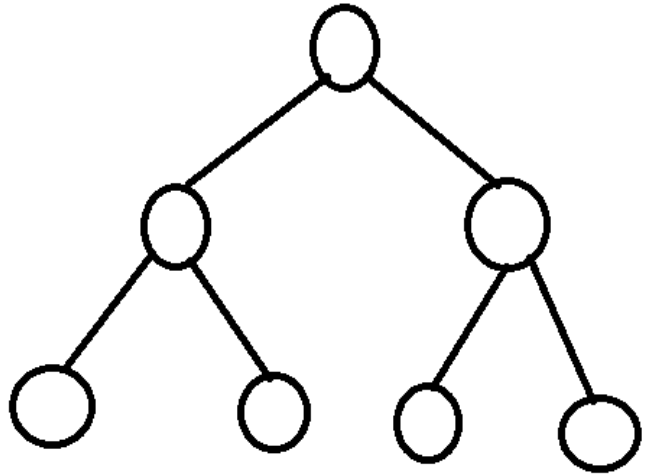
Defining Binary Trees

- Binary tree is a specific type of tree in which each node can have at most two children namely left child and right child.
- There are various types of binary trees:
  - Strictly binary tree
  - Full binary tree
  - Complete binary tree
- Strictly binary tree:
  - A binary tree in which every node, except for the leaf nodes, has non-empty left and right children.

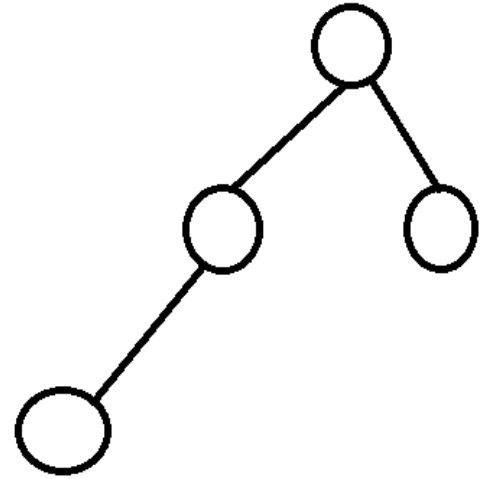




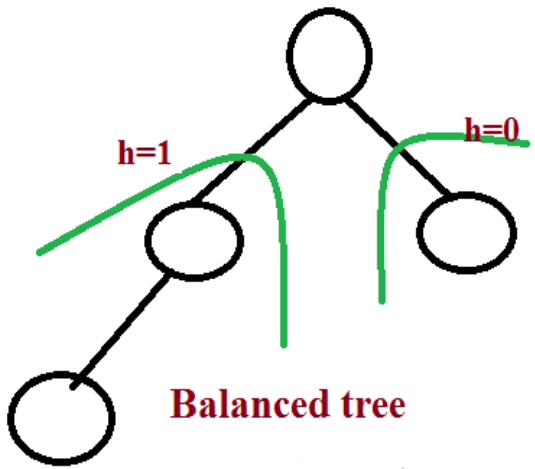
Binary Tree



Fully (perfect) Binary Tree

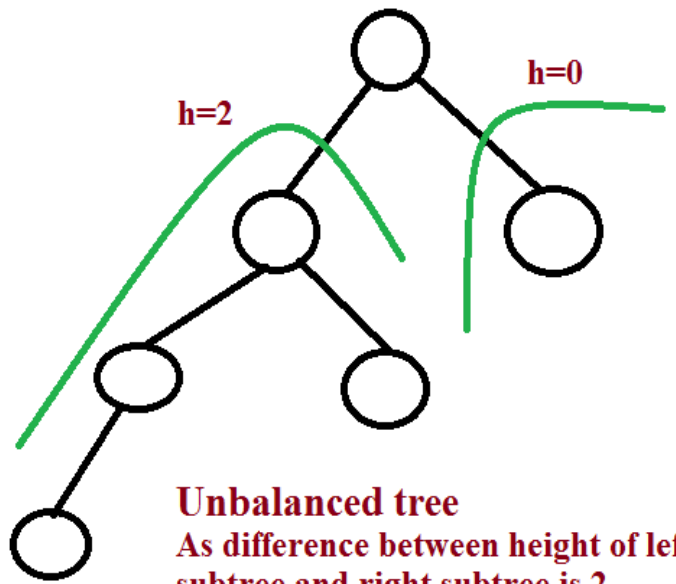


Complete Binary Tree



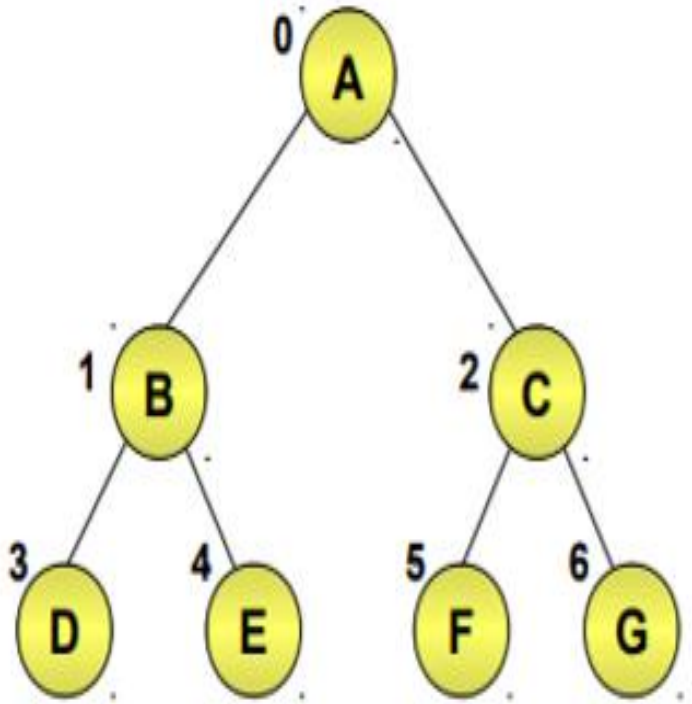
Balanced tree

As difference beween height of
left subtree and right subtree is 1

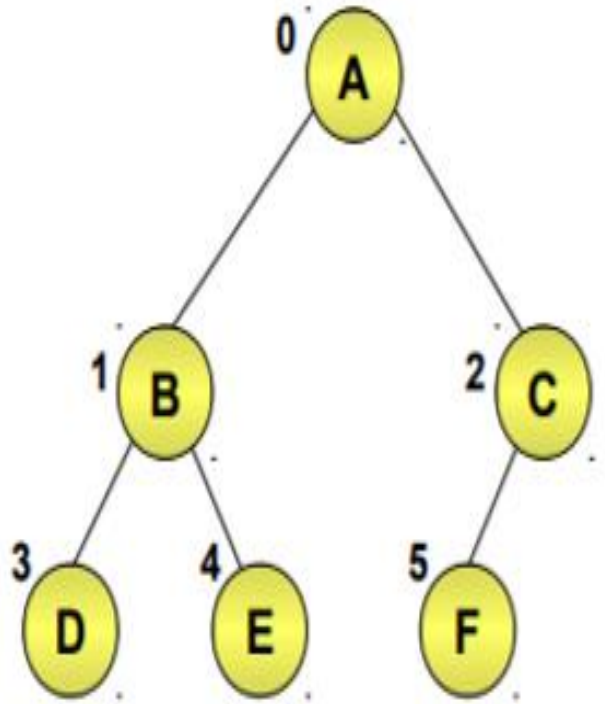


Unbalanced tree

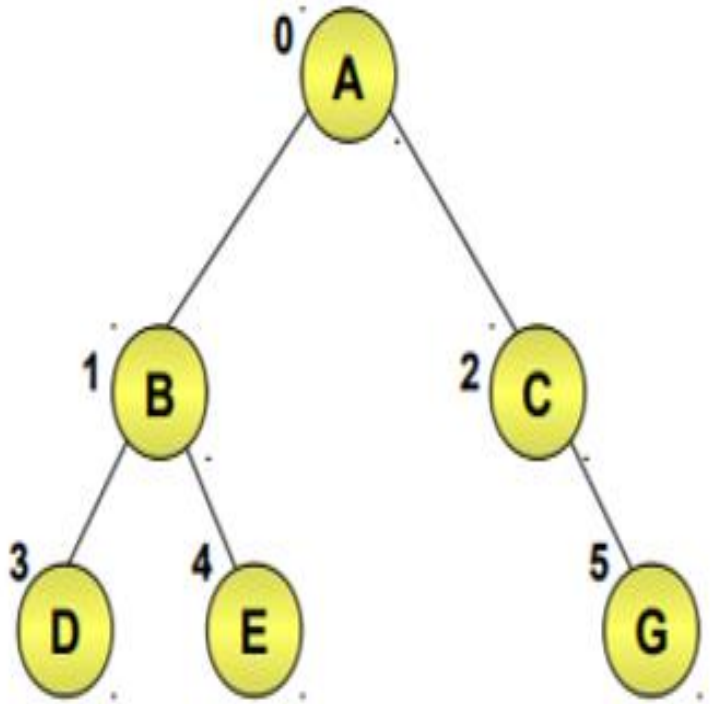
As difference between height of left
subtree and right subtree is 2



Full Binary Tree



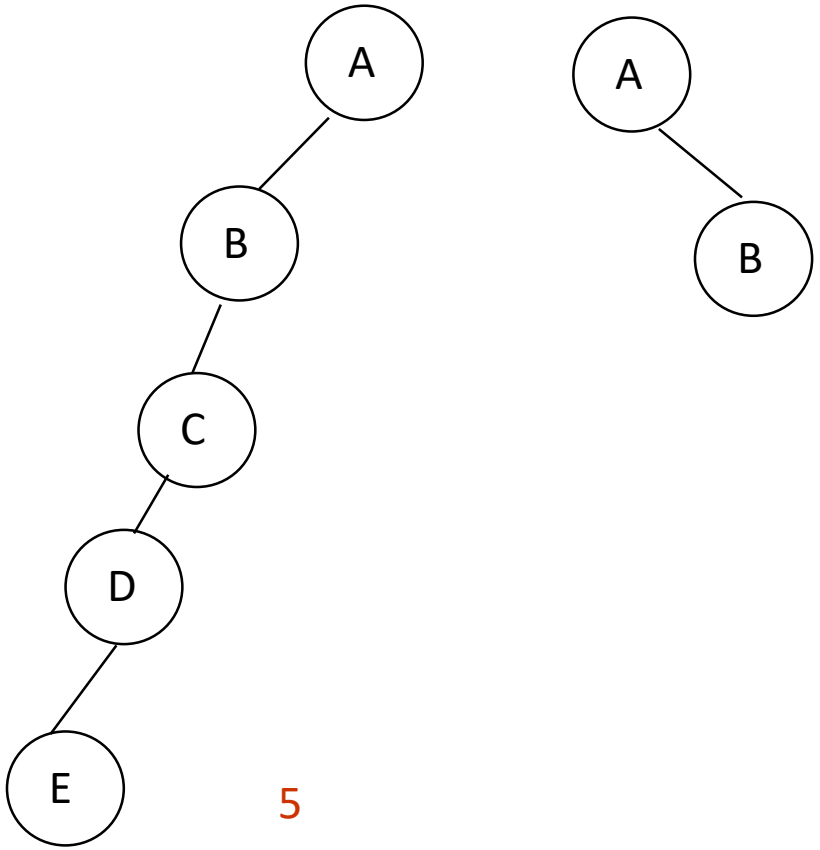
Complete Binary Tree



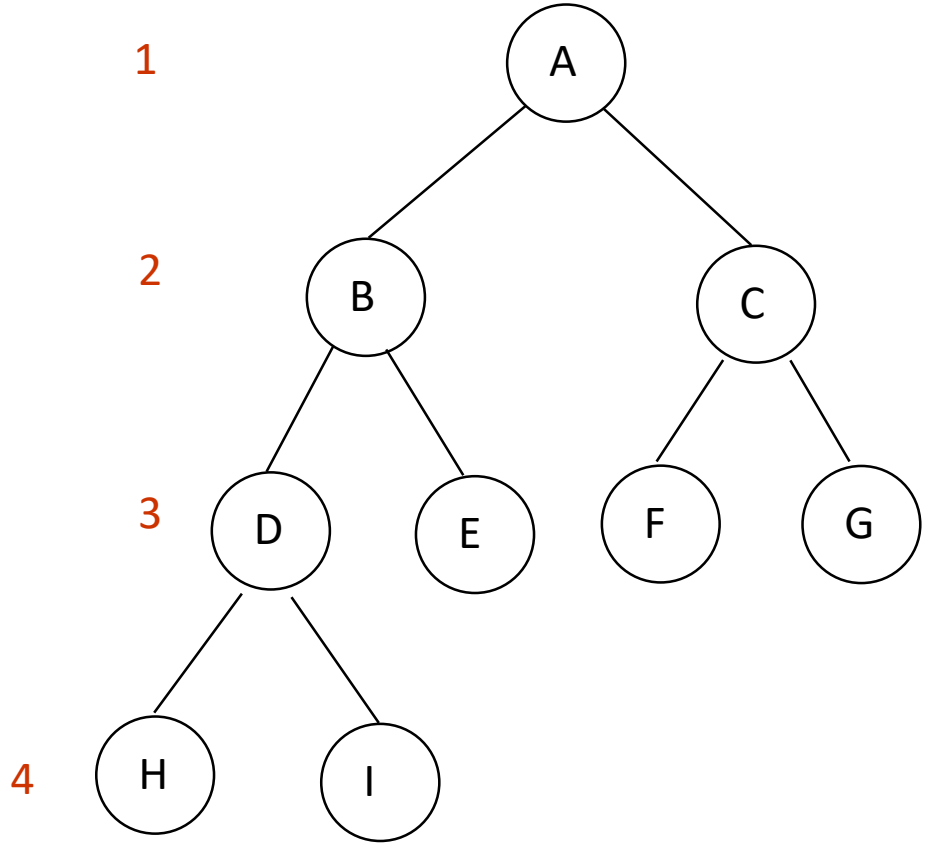
Incomplete Binary Tree

Examples of the Binary Tree

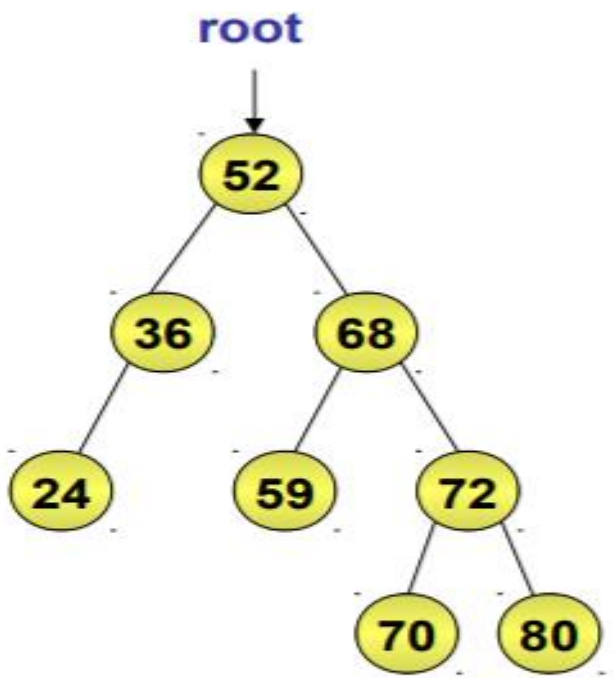
Skewed Binary Tree



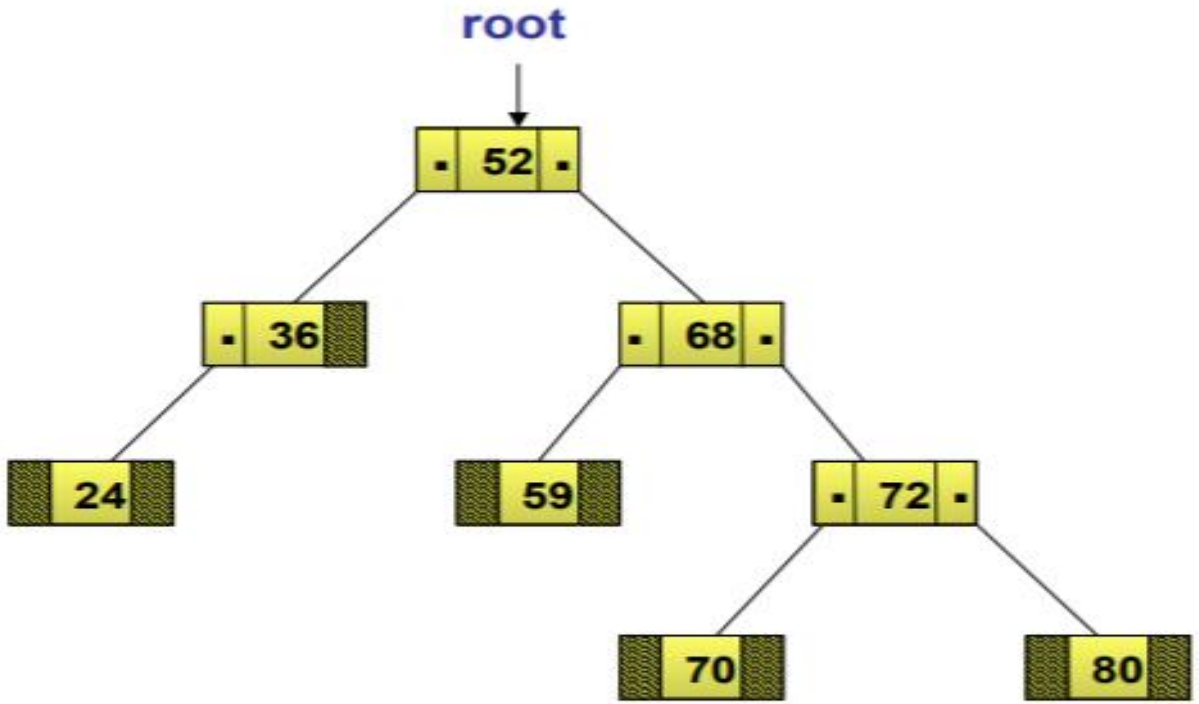
Complete Binary Tree



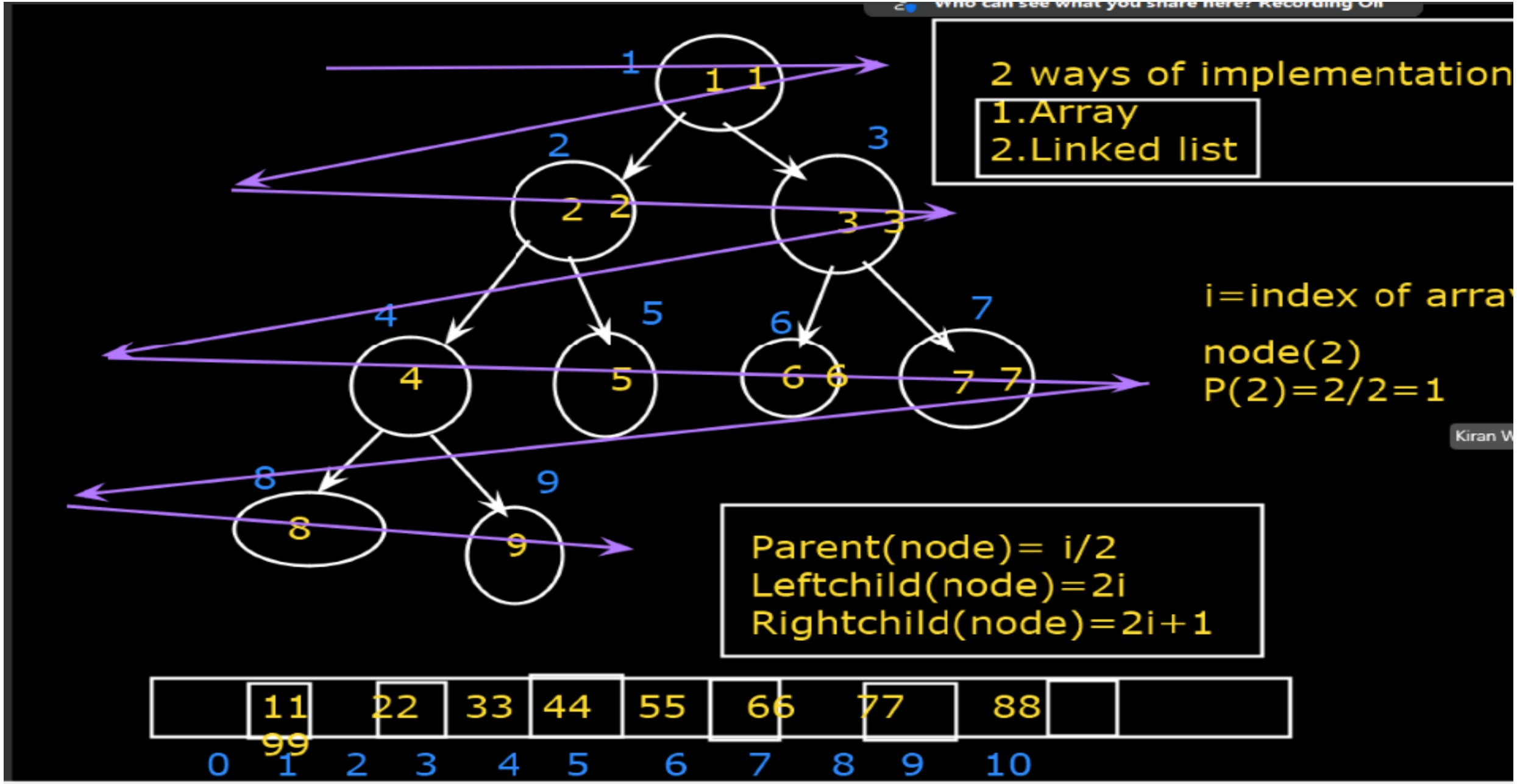
Representing a Binary Tree (Contd.)



Binary Tree



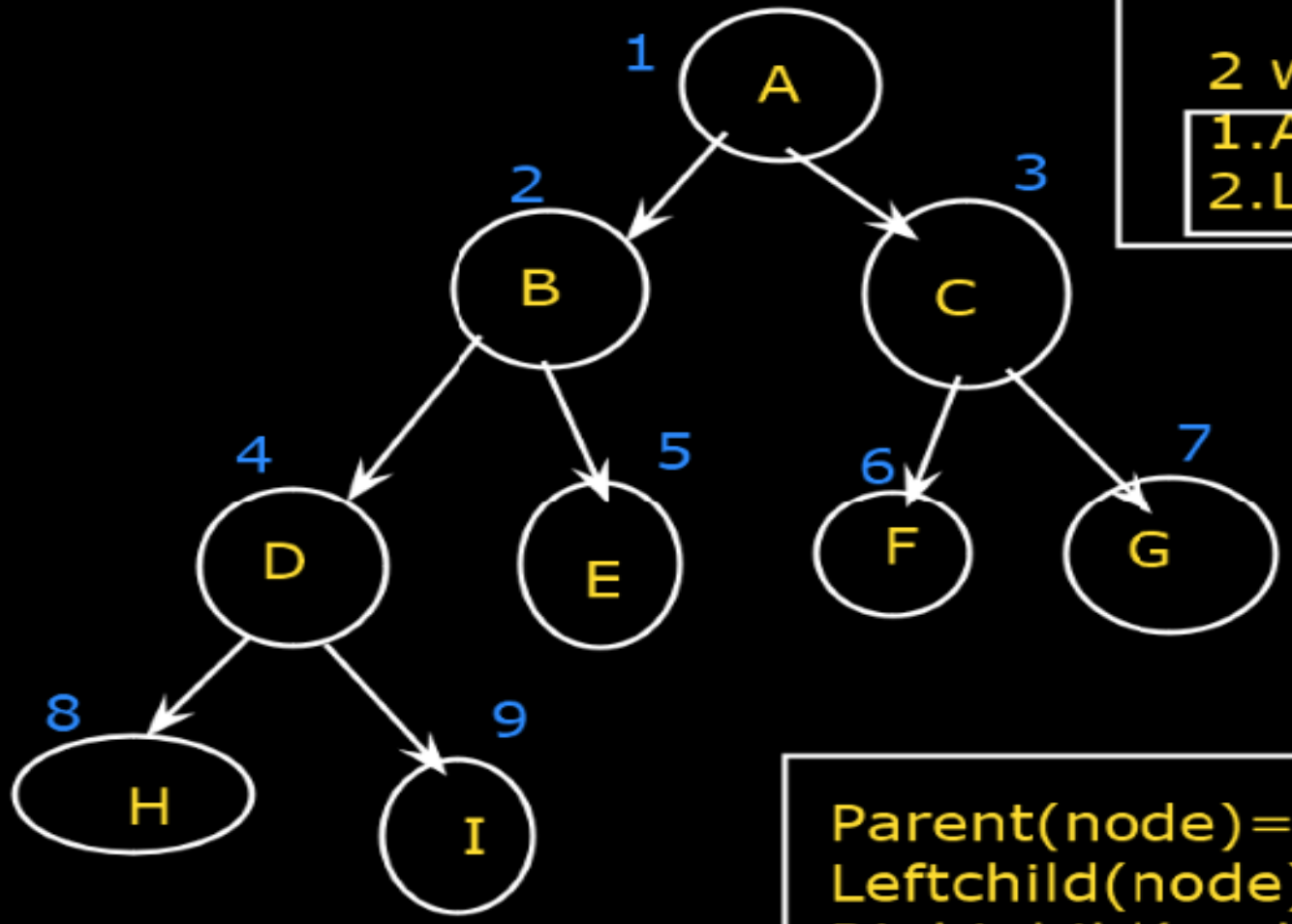
Linked Representation


2 ways of implementation
1.Array
2.Linked list
i=index of array
node(2)
P(2)=2/2=1
Parent(node)= i/2
Leftchild(node)=2i
Rightchild(node)=2i+1
11
22
33
44
55
66
77
88
99
0 1 2 3 4 5 6 7 8 9 10

in the full binary tree.

Mouse Select Text Draw Stamp Spotlight Eraser Format Undo Redo

Who can see what you share here? Recording On



2 ways of implementation

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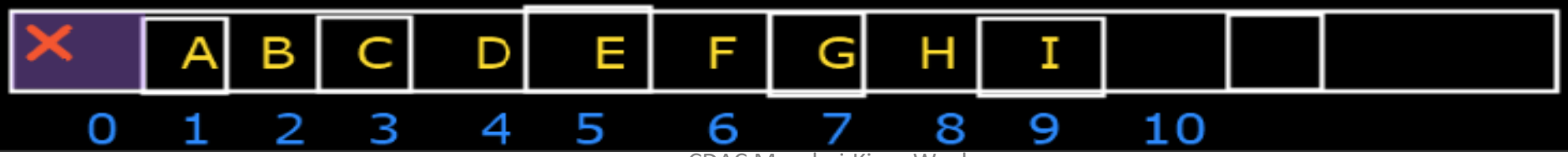
2.Linked list

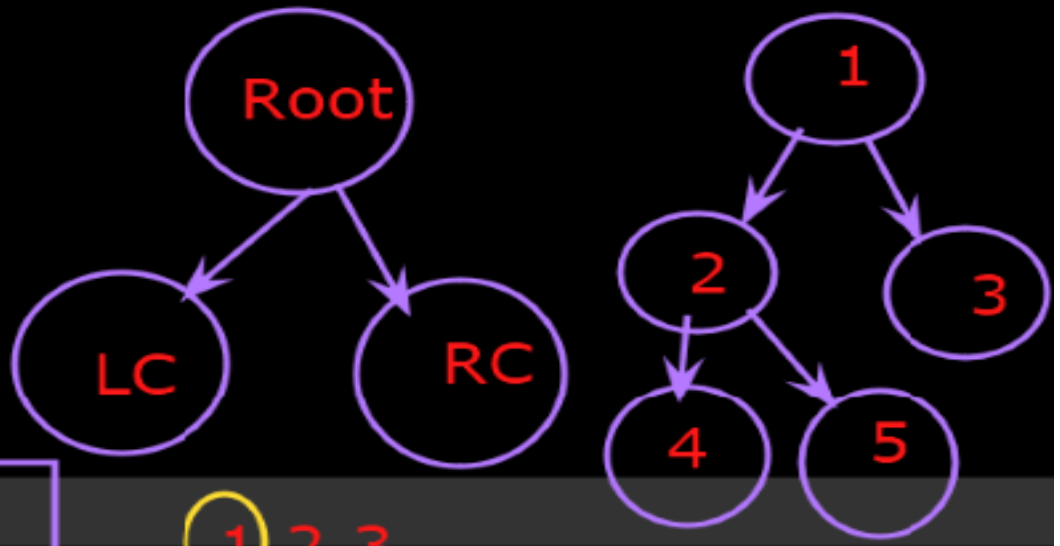
i=index of array

node(2)
P(2)=2/2=1

Parent(node)= i/2
Leftchild(node)=2i
Rightchild(node)=2i+1



```
Node
{
    int data;
    Node left, right;
    Node(int d)
    {
        data = d;
        left=right=null;
    }
}
```



Preorder: Root, LC, RC

Inorder: LC, Root, RC

Postorder: LC, RC, Root

1,2,3

2,1,3

2,3,1

1,2,4,5,3

4,2,5,1,3

4,5,2,3,1

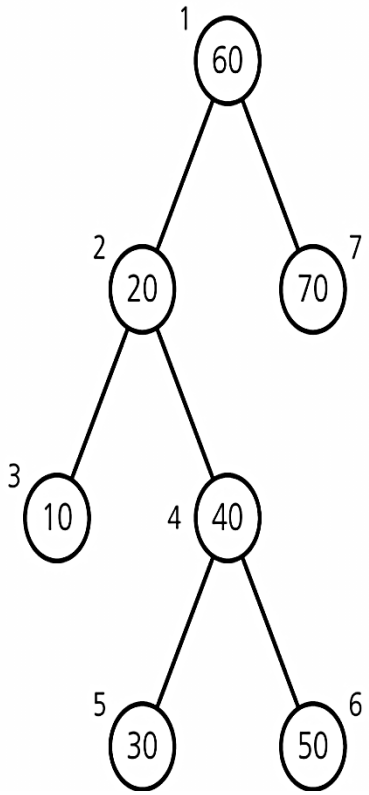
OPERATIONS ON TREES

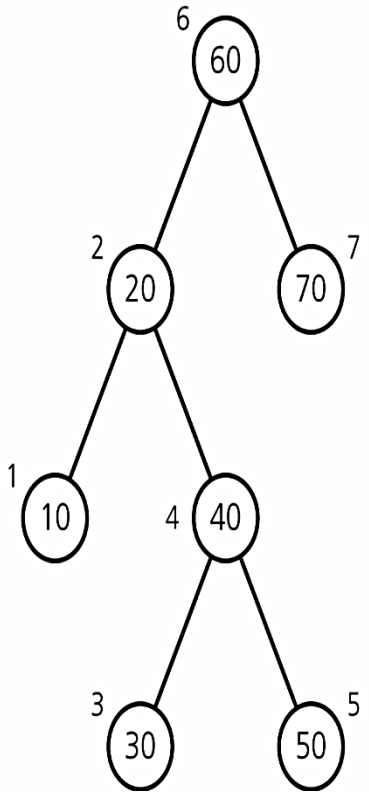
Traversing a Binary Tree

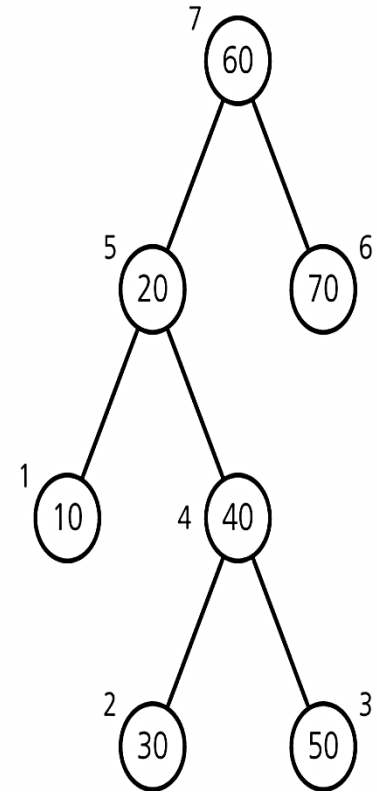
1)TRAVERSING

- You can implement various operations on a binary tree.
- A common operation on a binary tree is traversal.
- Traversal refers to the process of visiting all the nodes of a binary tree once.
- There are three ways for traversing a binary tree:
  - Inorder traversal
  - Preorder traversal
  - Postorder traversal

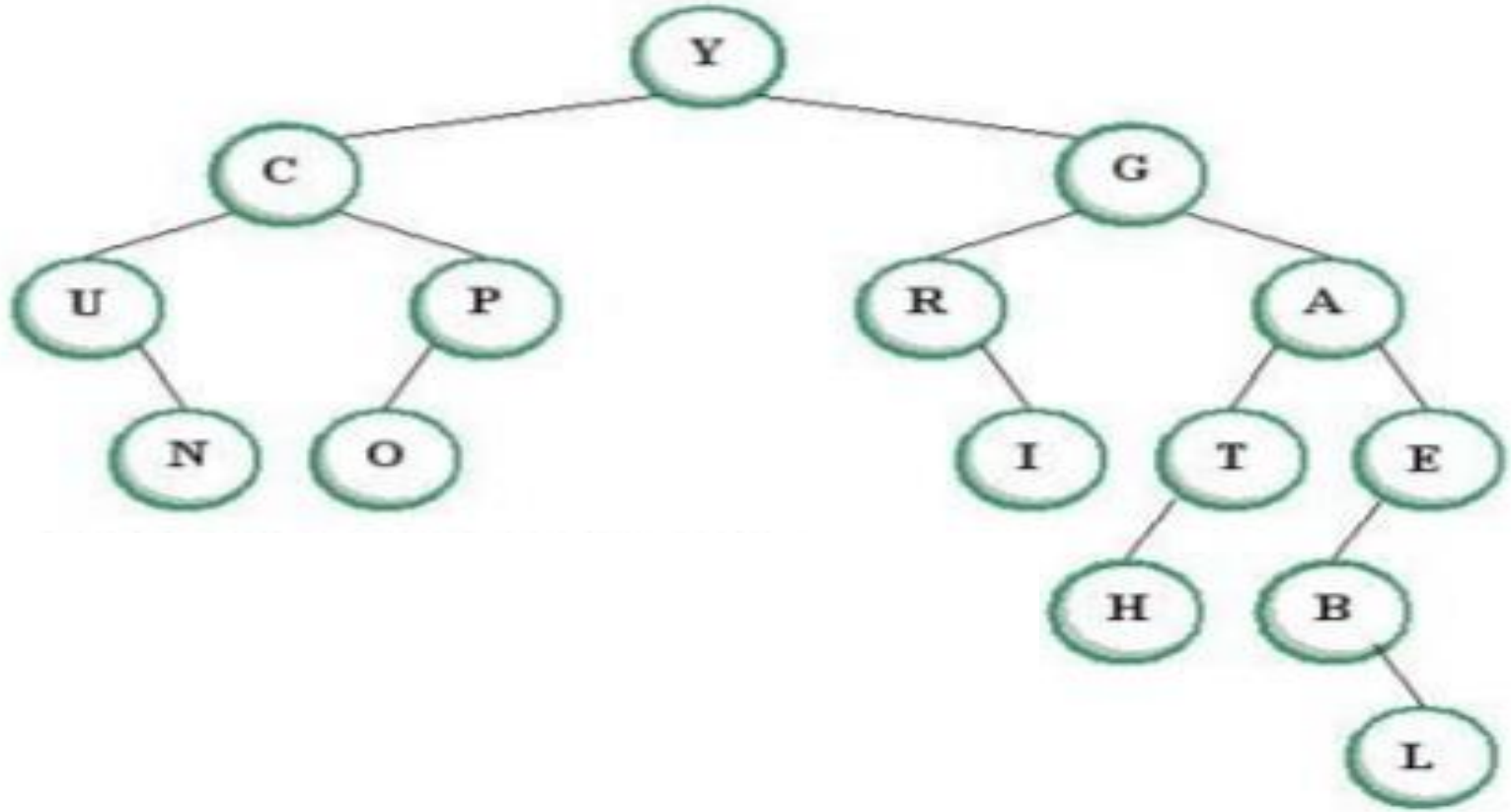
Binary Tree Traversals

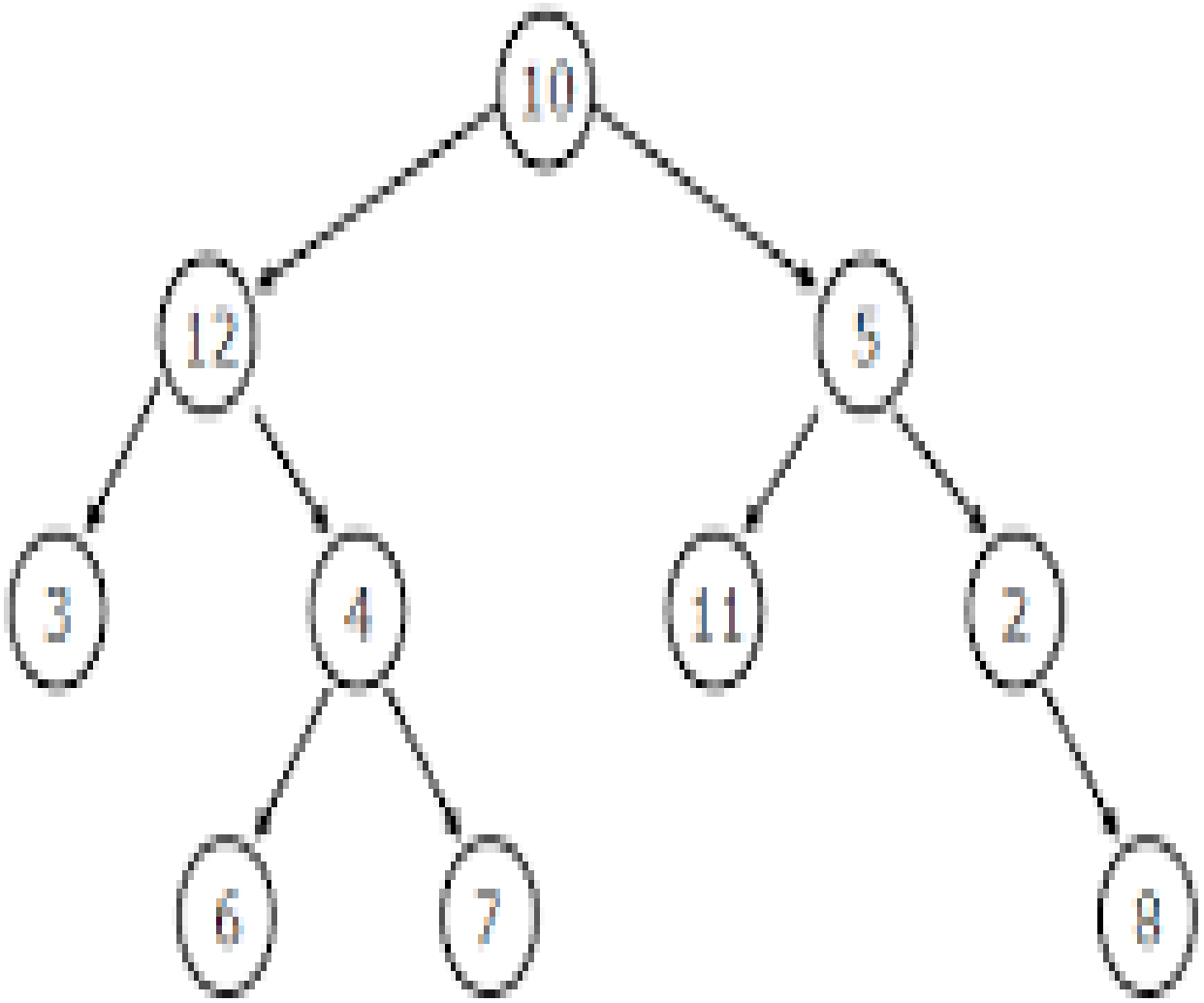


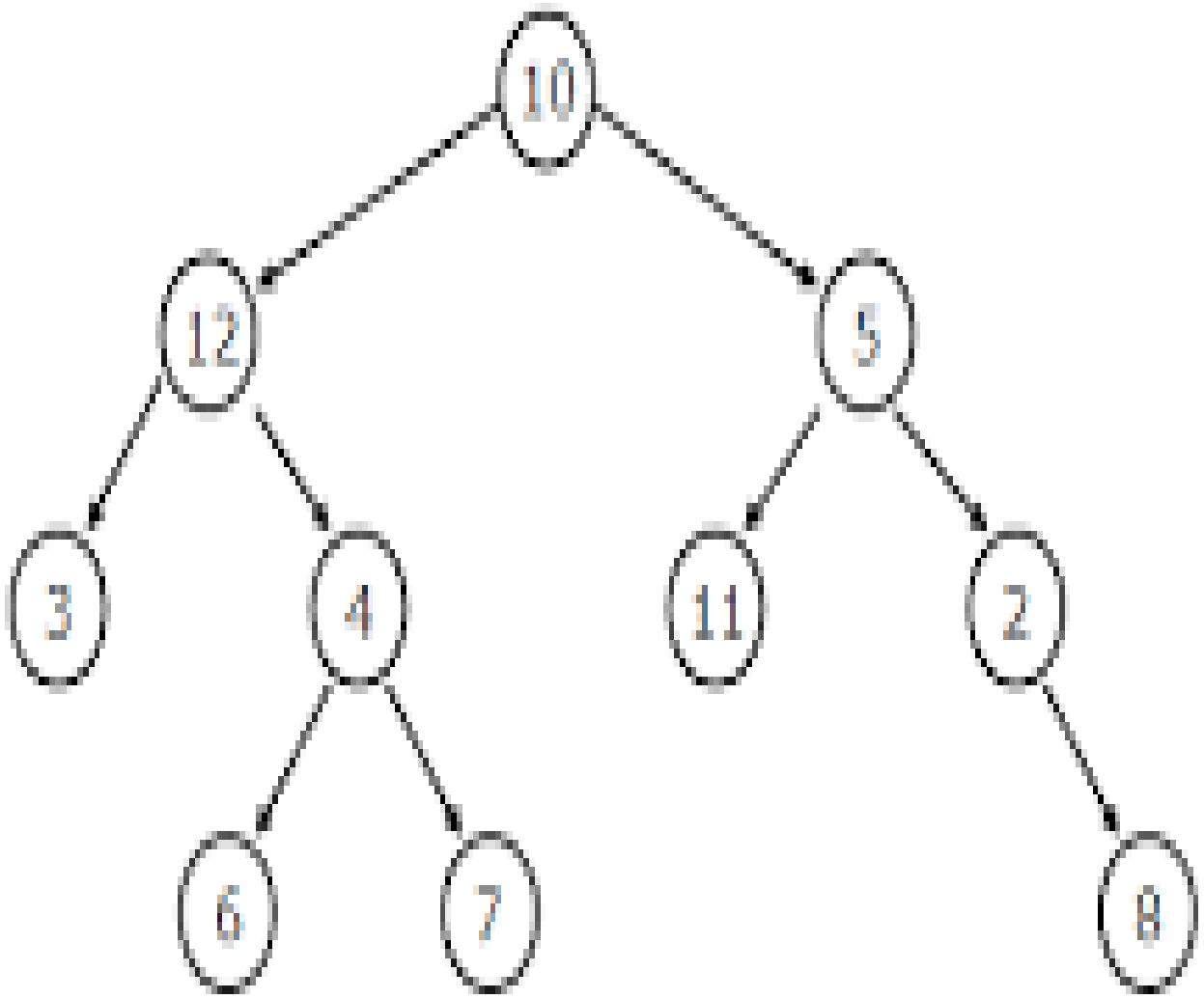




Traversal Examples




10
12
5
3
4
11
2
6
7
8


10
12
5
3
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11
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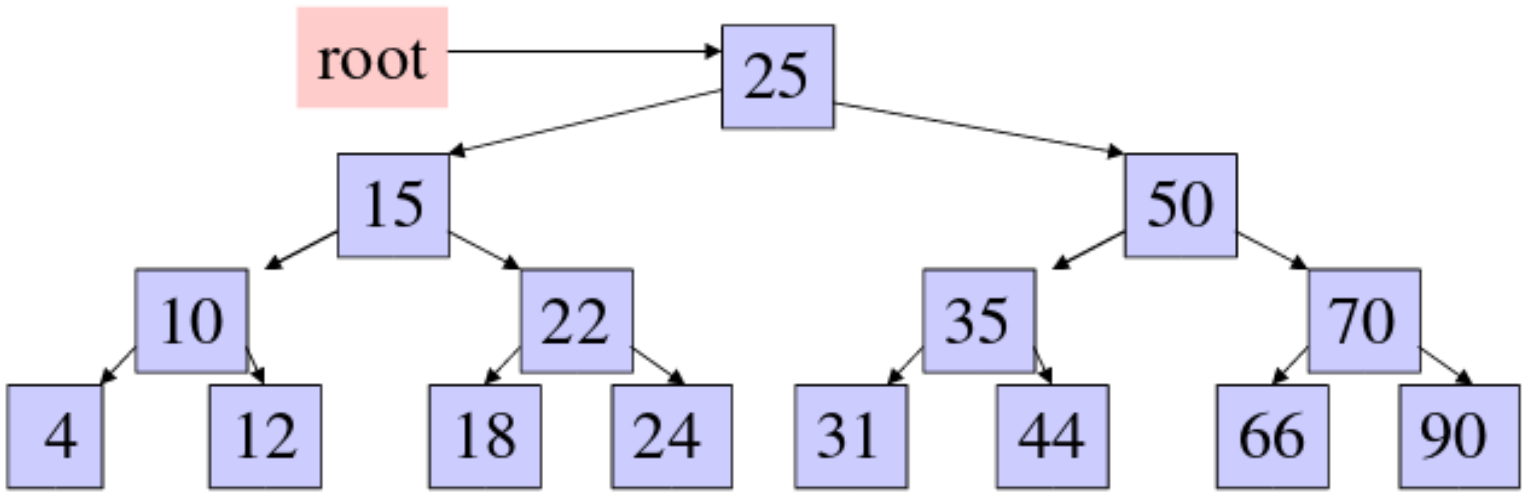
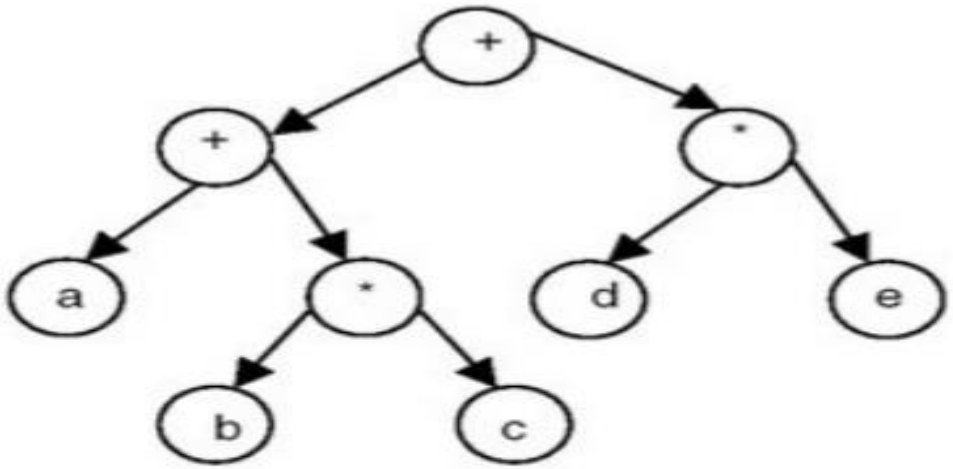
6
7
8


root
25
15
50
10
22
35
70
4
12
18
24
31
44
66
90

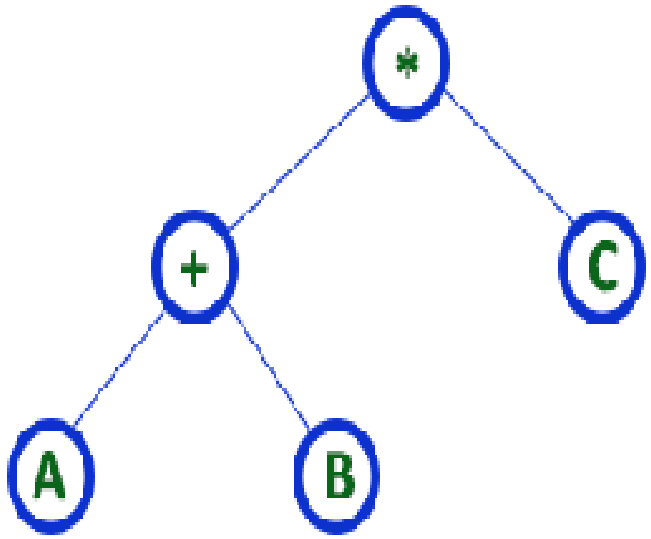
Expression Binary Tree Traversal

If an expression is represented as a binary tree, the inorder traversal of the tree gives us an infix expression, whereas the postorder traversal gives us a postfix expression as shown in Figure.

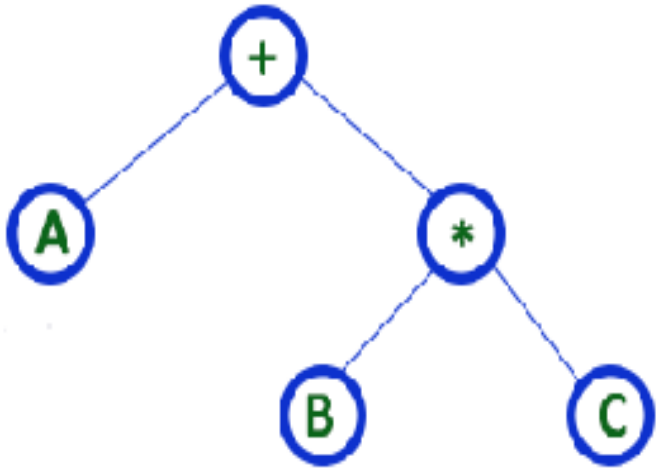


Inorder : a + b * c + d * e
postorder : abc*+de*+

Strictly binary tree data structure is used to represent mathematical expressions.



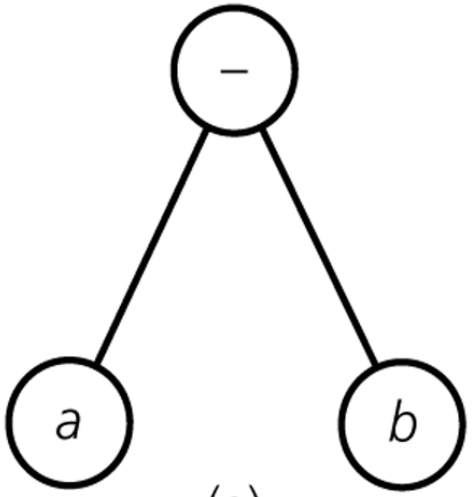
(A + B) * C



A + B * C

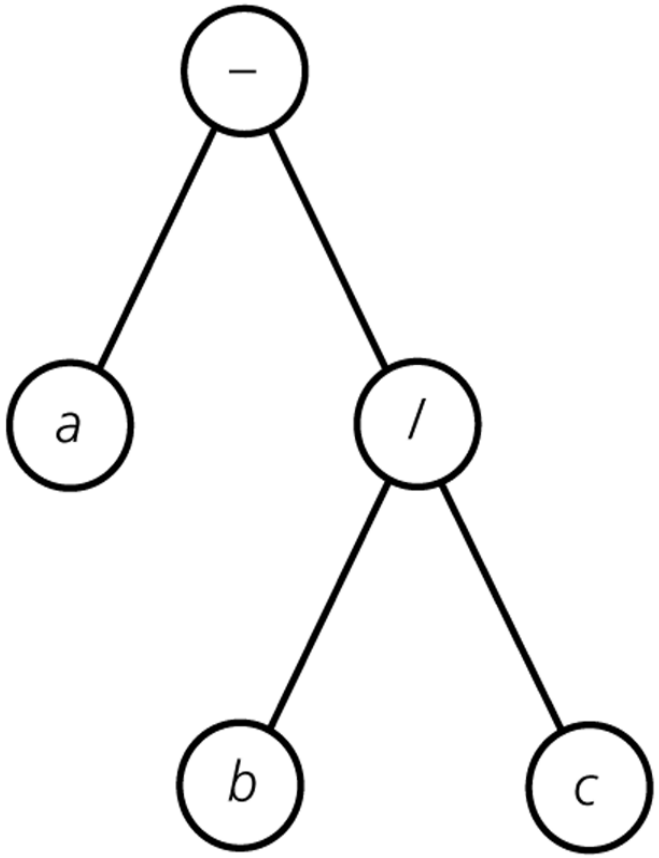
Binary Tree – Representing Algebraic Expressions

$a - b$



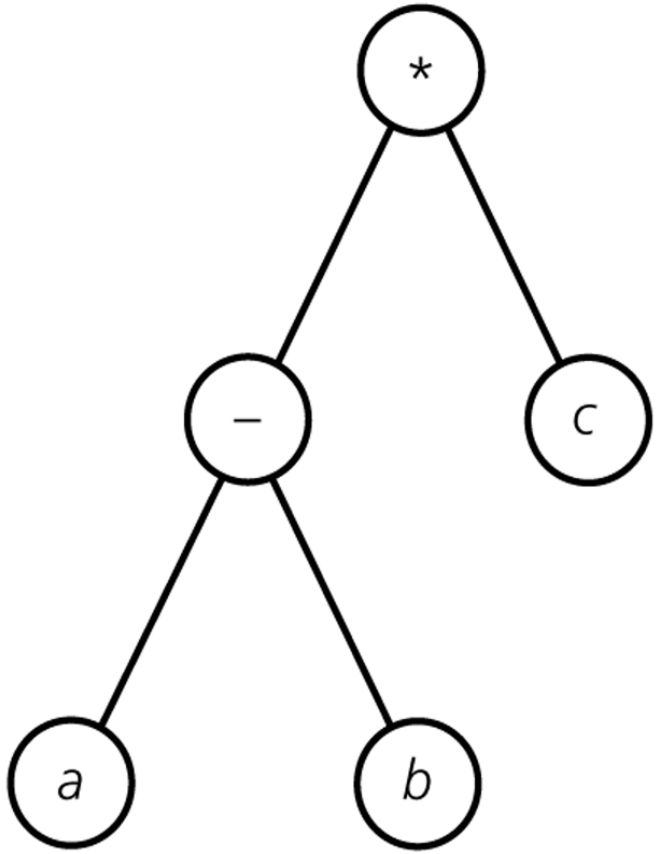
(a)

$a - b / c$



(b)

$(a - b) * c$



(c)

Thanks