Longest Harmonious Subsequence

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Problem Statement

We define a harmonious array as an array where the difference between its maximum value and its minimum value is **exactly** 1.

Given an integer array nums, return the length of its longest harmonious subsequence among all its possible subsequences.

Example 1:

Input: nums = [1,3,2,2,5,2,3,7]

Output: 5

Explanation: The longest harmonious subsequence is [3,2,2,2,3].

Example 2:

Input: nums = [1,2,3,4]

Output: 2

Explanation: The longest harmonious subsequences are [1,2], [2,3], and [3,4], all of which have a length of

2.

Example 3:

Input: nums = [1,1,1,1]

Output: 0

Explanation: No harmonic subsequence exists.

Editorial

Brute Force Approach

Approach:

- Generate all possible subsequences using recursion or bitmasking.
- For each subsequence:
 - Find max and min values.
 - If max min == 1, check and update maximum length.
- Return the longest length found.

Code

C++:

```
#include <bits/stdc++.h>
using namespace std;
class Solution {
public:
  int findLHS(vector<int>& nums) {
     int n = nums.size(), maxLen = 0;
     int total = 1 << n; // 2^n possible subsequences
     for (int mask = 1; mask < total; ++mask) {
       vector<int> subseq;
       for (int i = 0; i < n; ++i) {
         if (mask & (1 << i)) subseq.push_back(nums[i]);
       }
       int mini = *min_element(subseq.begin(), subseq.end());
       int maxi = *max_element(subseq.begin(), subseq.end());
       if (maxi - mini == 1) {
         maxLen = max(maxLen, (int)subseq.size());
       }
    }
     return maxLen;
  }
};
int main() {
  vector<int> nums = {1,3,2,2,5,2,3,7};
  Solution s1;
  cout << s1.findLHS(nums) << endl;
  return 0;
}
```

Python: Usually not implemented — impractical due to exponential growth.

Complexity Analysis:

- Time Complexity: O(2^n * n) exponential due to all subsequences.
- Space Complexity: O(n) for temporary subsequence storage.

Better Approach (Sorting + Scanning)

Approach

- · Sort the array.
- Use two pointers to maintain a window:
 - Move start when the difference between current and start > 1.
 - If max min == 1, update the result.

• This avoids generating subsequences.

Code

C++:

```
#include <bits/stdc++.h>
using namespace std;
class Solution {
public:
  int findLHS(vector<int>& nums) {
     sort(nums.begin(), nums.end());
    int start = 0, maxLen = 0;
    for (int end = 0; end < nums.size(); ++end) {
       while (nums[end] - nums[start] > 1) ++start;
       if (nums[end] - nums[start] == 1)
         maxLen = max(maxLen, end - start + 1);
    }
    return maxLen;
  }
};
int main() {
  vector<int> nums = {1,3,2,2,5,2,3,7};
  Solution s1;
  cout << s1.findLHS(nums) << endl;
  return 0;
}
```

Python:

```
def findLHS(nums):
    nums.sort()
    start, res = 0, 0
    for end in range(len(nums)):
        while nums[end] - nums[start] > 1:
            start += 1
        if nums[end] - nums[start] == 1:
            res = max(res, end - start + 1)
    return res
```

Complexity Analysis:

- Time Complexity: O(n log n) for sorting + O(n) scanning.
- Space Complexity: O(1) (no extra space apart from sorting).

Optimal Approach (Hashmap Frequency Count)

Approach

- Use a unordered_map to store frequency of each element.
- For each element x, check if x+1 exists.
- If yes, freq+ freq[x+1] is a valid subsequence.
- Keep track of max length.

Code

C++:

```
#include <bits/stdc++.h>
using namespace std;
class Solution {
public:
  int findLHS(vector<int>& nums) {
     unordered_map<int, int> freq;
     for (int num: nums) freq[num]++;
     int maxLen = 0;
     for (auto& [num, count] : freq) {
       if (freq.count(num + 1)) {
         maxLen = max(maxLen, count + freq[num + 1]);
       }
     return maxLen;
};
int main() {
  vector<int> nums = {1,3,2,2,5,2,3,7};
  Solution s1;
  cout << s1.findLHS(nums) << endl;
  return 0;
}
```

Python:

```
def findLHS(nums):
    from collections import Counter
    freq = Counter(nums)
    res = 0
    for key in freq:
        if key + 1 in freq:
        res = max(res, freq[key] + freq[key+1])
    return res
```

Complexity Analysis:

- Time Complexity: O(n) for building hashmap and iterating over it.
- Space Complexity: O(n) for the hashmap.

THE END