

Data Structures and Algorithms

SORTING ALGORITHMS HANDOUT

Bubble Sort

ALGORITHM *BubbleSort*($A[0..n-1]$)
// Sorts a given array using bubble sort
// Input: An array $A[0..n-1]$ of orderable
// elements
// Output: Array $A[0..n-1]$ sorted in
// ascending order
for $i \leftarrow 0$ **to** $n-2$ **do**
 for $j \leftarrow 0$ **to** $n-2-i$ **do**
 if $A[j+1] < A[j]$
 swap $A[j]$ and $A[j+1]$

Selection Sort

ALGORITHM *SelectionSort*($A[0..n-1]$)
// Sorts a given array using selection sort
// Input: An array $A[0..n-1]$ of orderable
// elements
// Output: Array $A[0..n-1]$ sorted in
// ascending order
for $i \leftarrow 0$ **to** $n-2$ **do**
 $\text{min} \leftarrow i$
 for $j \leftarrow i+1$ **to** $n-1$ **do**
 if $A[j] < A[\text{min}]$
 $\text{min} \leftarrow j$
 swap $A[i]$ and $A[\text{min}]$

Insertion Sort

ALGORITHM *InsertionSort*($A[0..n-1]$)
// Sorts a given array using insertion sort
// Input: An array $A[0..n-1]$ of orderable
// elements
// Output: Array $A[0..n-1]$ sorted in
// ascending order
for $i \leftarrow 1$ **to** $n-1$ **do**
 $v \leftarrow A[i]$
 $j \leftarrow i-1$
 while $j \geq 0$ **and** $A[j] > v$ **do**
 $A[j+1] \leftarrow A[j]$
 $j \leftarrow j-1$
 $A[j+1] \leftarrow v$

Merge Sort

ALGORITHM *MergeSort*($A[0..n-1]$)
// Sorts a given $A[0..n-1]$ by recursive mergesort
// Input: An array $A[0..n-1]$ of orderable elements
// Output: Array $A[0..n-1]$ sorted in nondecreasing order
if $n > 1$
 copy $A[0 \dots \lfloor n/2 \rfloor - 1]$ to $B[0 \dots \lfloor n/2 \rfloor - 1]$
 copy $A[\lfloor n/2 \rfloor \dots n-1]$ to $C[0 \dots \lceil n/2 \rceil - 1]$
 MergeSort($B[0 \dots \lfloor n/2 \rfloor - 1]$)
 MergeSort($C[0 \dots \lceil n/2 \rceil - 1]$)
 Merge(B, C, A)

ALGORITHM *Merge*($B[0 \dots p-1], C[0 \dots q-1], A[0 \dots p+q-1]$)
// Merges two sorted arrays into one sorted array
// Input: Arrays $B[0 \dots p-1]$ and $C[0 \dots q-1]$ both sorted
// Output: Sorted array $A[0 \dots p+q-1]$ of the elements of B and C
 $i \leftarrow 0$
 $j \leftarrow 0$
 $k \leftarrow 0$
while $i < p$ **and** $j < q$ **do**
 if $B[i] \leq C[j]$
 $A[k] \leftarrow B[i]$
 $i \leftarrow i+1$
 else
 $A[k] \leftarrow C[j]$
 $j \leftarrow j+1$
 $k \leftarrow k+1$
if $i = p$
 copy $C[j \dots q-1]$ to $A[k \dots p+q-1]$
else
 copy $B[i \dots p-1]$ to $A[k \dots p+q-1]$

Quick Sort

ALGORITHM *QuickSort*($A[l \dots r]$)
// Sorts a subarray by quicksort
// Input: A subarray $A[l \dots r]$ of $A[0 \dots n-1]$, defined by its left and right indices l and r
// Output: Subarray $A[l \dots r]$ sorted in nondecreasing order
if $l < r$
 $s \leftarrow \text{Partition}(A[l \dots r])$
 QuickSort($A[l \dots s-1]$)
 QuickSort($A[s+1 \dots r]$)

ALGORITHM Partition($A[l \dots r]$)

// Partitions a subarray by using its first element as a pivot

// Input: A subarray $A[l \dots r]$ of $A[0 \dots n-1]$, defined by its left and right indices l and r ($l < r$)

// Output: Subarray $A[l \dots r]$, with split position returned as this functions value

$p \leftarrow A[l]$

$i \leftarrow l$

$j \leftarrow r + 1$

repeat

repeat $i \leftarrow i + 1$ **until** $A[i] \geq p$

repeat $j \leftarrow j - 1$ **until** $A[j] \leq p$

 swap($A[i]$ and $A[j]$)

until $i \geq j$

swap ($A[i]$, $A[j]$)

swap ($A[l]$, $A[j]$)

return j

Heap Sort

Step 01: Heap construction – construct a heap for a given array

Step 02: Maximum Deletions – Apply the root-deletion operation $n-1$ times to the remaining heap

Algorithm for Heap Construction**ALGORITHM** HeapBottomUp($H[1..n]$)

// Constructs a heap from the elements of a given array by the bottom up algorithm

// Input: An array $H[1..n]$ of orderable items

// Output: A heap $H[1..n]$

for $i \leftarrow \lfloor n/2 \rfloor$ **downto** 1 **do**

$k \leftarrow i$

$v \leftarrow H[k]$

 heap \leftarrow false

while not heap **and** $2 * k \leq n$ **do**

$j \leftarrow 2 * k$

if $j < n$

if $H[j] < H[j + 1]$

$j \leftarrow j + 1$

if $v \geq H[j]$

 heap \leftarrow true

else

$H[k] \leftarrow H[j]$

$k \leftarrow j$

$H[k] \leftarrow v$

Algorithm for maximum key deletion

Step 01: Exchange the root's key with the last key K of the Heap

Step 02: Decrease the heap's size by 1

Step 03: "Heapify" the smaller tree