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Learning Objectives

By the end of this module, you should have a better understanding of:



- Ways computers talk to each other
- Sniffing and capturing the network traffic











Protocols







2.1 Protocols

How does this support my pentesting career?

- Protocols are used in every computer network communication
- You need to know how things work to exploit them

2.1 Protocols

In a computer network, machines talk to each other by means of **protocols**.

These protocols ensure that different computers, using different hardware and software, can communicate.









2.1 Protocols

There is a large variety of networking protocols on the Internet, each one with its own purpose.

We are going to discuss a few of them in detail and point you towards free online resources for others.









The primary goal of networking is to exchange information between networked computers; this information is carried by packets.

Packets are nothing but streams of bits running as electric signals on physical media used for data transmission. Such media can be a **wire** in a LAN or **the air** in a WiFi network.

These electrical signals are then interpreted as bits (zeros and ones) that make up the information.









Every packet in every protocol has the following structure.









The **header** has a protocol-specific structure: this ensures that the receiving host can correctly interpret the payload and handle the overall communication.













The **payload** is the actual information. It could be something like part of an email message or the content of a file during a download.

Header

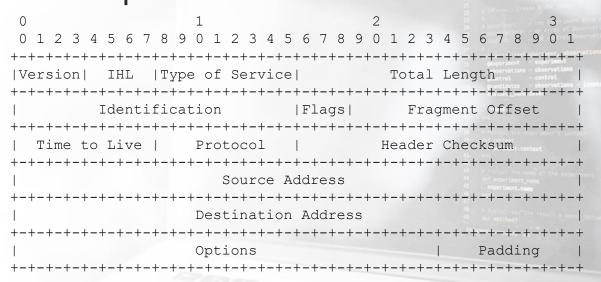
Payload







For example, the IP protocol header is at least 160 bits (20 bytes) long, and it includes information to interpret the content of the IP packet.



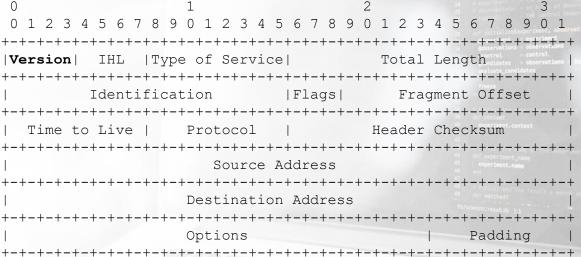






The first four bits identify the **IP version**. Today they can be used to represent IP version 4 or 6.













The 32 bits starting at position 96 represent the **source address**.

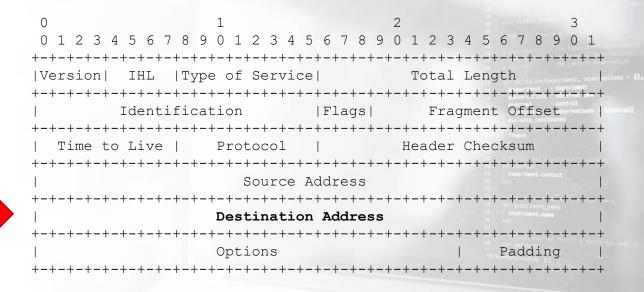
0 1	2
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4	5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-	+-
Version IHL Type of Servi	ce Total Length
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-	+-
Identification	Flags Fragment Offset
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-	+-
Time to Live Protocol	Header Checksum
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-	+-
Source Address	
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-	+-
Destination Address	
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-	+-
Options	Padding
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-	+-







The following four bytes represent the **destination address**.



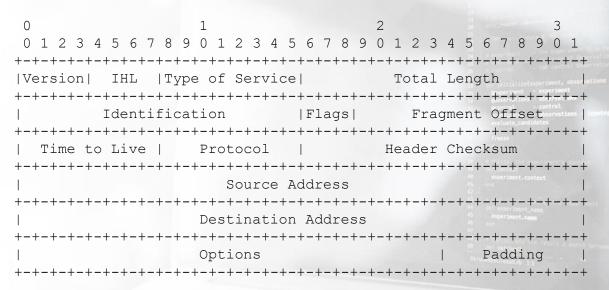








Using the information in the header, the nodes involved in the communication can understand and use IP packets.









In the previous example, we saw the header of the IP (Internet Protocol). There are many protocols out there, each one for a specific purpose. Purposes such as:

- Exchanging emails, files or performing VoIP calls
- Establishing a communication between a server and a client
- Identifying computers on a network
- Transmitting data









Instead of using specific examples, let's rewrite the previous list focusing on the features that a protocol provides:

- Make an application (email client, FTP, browser, ...) work
- Transport data between processes (the server and the client programs)
- Identify hosts
- Use the physical media to send packets









Moreover, we can rewrite the list again as:

- Application layer
- Transport layer
- Network layer
- Physical layer

These **layers** work on top of one another, and every layer has its own **protocol**.









Each layer serves the one above it.

The application layer does not need to know how to identify a process on a host, how to reach it and how to use the copper wire to establish a communication.

It just uses its underlying layers.







2.1.3 ISO/OSI

In 1984, the International Organization for Standardization (ISO) published a theoretical model for network systems communication: the Open System Interconnection (OSI) model.

The **ISO/OSI** model was never implemented, but it is widely used in literature or when talking about IT networks.









2.1.3 ISO/OSI

ISO/OSI consists of seven layers and is used as a reference for the implementation of actual protocols.

You can find more information about ISO/OSI here.



7

But how do protocols work together? If every protocol has a header and a payload, how can a protocol use the one on its lower layer?

The idea is simple. The **entire upper protocol packet** (header plus payload) is the **payload** of the lower one; this is called **encapsulation**.









In the following slides, you will see how encapsulation is used by the IP protocol suite, or TCP/IP.

TCP/IP is a real-world implementation of a networking stack and is the protocol stack used on the Internet.









TCP/IP has four layers:

Application

Transport

Network

Data Link







You will learn how TCP/IP works in the remainder of this module.

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Application

Transport

Network

Data Link

Header **Payload** a Header **Payload** \Box × Header **Payload** Header **Payload**

Application

Transport

Network

Data Link

0 W e a

Payload Header Header Header **Payload**

The application layer gives its packet to the transport layer, which adds its own header.



a





Application

Transport

Network

Data Link

0 W e a

Header Payload

Header Payload

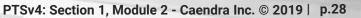
The application packet is now the transport protocol's payload.



a







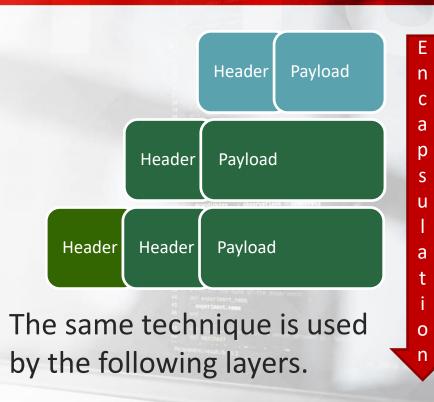
Application

Transport

Network

Data Link

0 W e a



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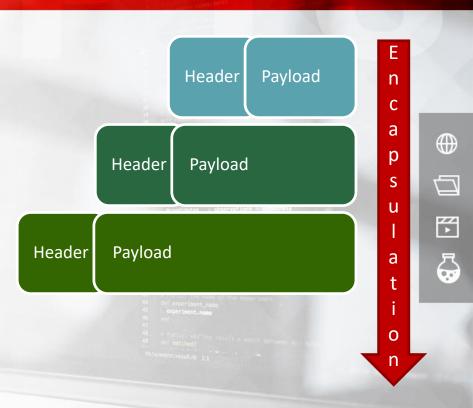
Application

Transport

Network

Data Link

0 W е a e



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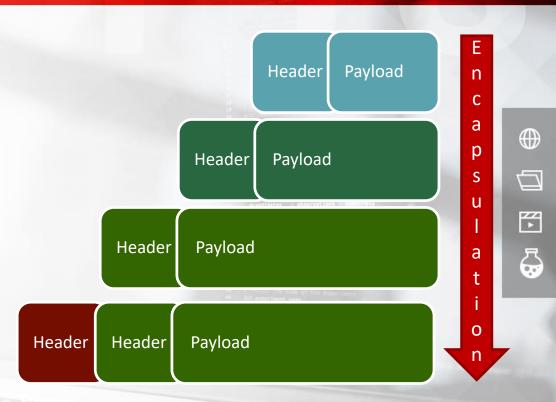
e

Application

Transport

Network

Data Link



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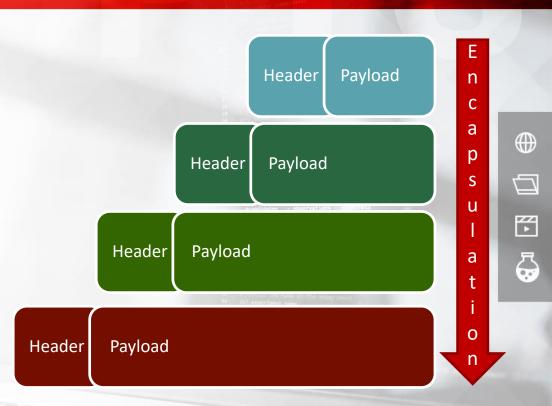
e

Application

Transport

Network

Data Link



During encapsulation every protocol adds its own header to the packet, treating it as a payload.

This happens to every packet sent by a host.







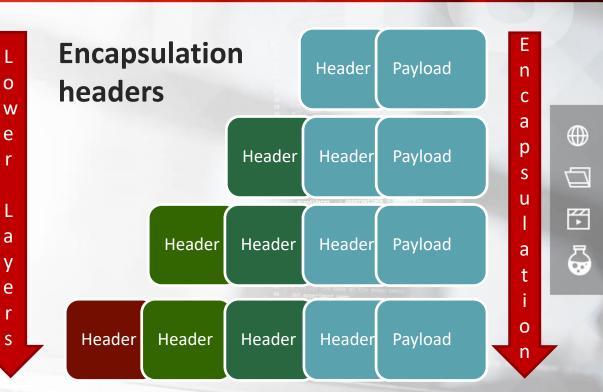


Application

Transport

Network

Data Link



The receiving host does the same operation in reverse order. Using this method, the application does not need to worry about how the transport, network and link layers work. It just hands in the packet to the transport layer.

You will see encapsulation in practice later, during the Wireshark section.

















2.2 IP



- Understanding network attacks
- Using network attack tools at their maximum
- Studying other networking protocols

2.2 IP

The Internet Protocol (IP) is the protocol that runs on the Internet layer of the Internet Protocol suite, also known as TCP/IP.

IP is in charge of delivering the **datagrams** (IP packets are called datagrams) to the hosts involved in a communication, and it uses **IP addresses** to identify a host.







When you write a letter, you have to specify the recipient's address on the envelope before sending it. Similarly, the Internet uses its addressing scheme to deliver packets to the right destination.

Any host on a computer network, be it a private network or the Internet, is identified by a **unique IP address**.









The vast majority of networks run IP version 4 (IPv4).

An IPv4 address consists of four bytes, or octets; a byte consists of 8 bits.

73.5.12.132







A dot delimits every octet in the address.









As you may recall from the introduction module, with 8 bits, you can represent **up to 2**⁸ **different values** from 0 to 255.

This does not mean that you can **assign** any address starting from 0.0.0.0 to 255.255.255.255 to a host. Some addresses are **reserved** for special purposes.









2.2.2 Reserved IPv4 Addresses

For example, some reserved intervals are:

- 0.0.0.0 0.255.255.255 representing "this" network.
- 127.0.0.0 127.255.255.255 representing the local host (e.g., your computer).
- 192.168.0.0 192.168.255.255 is reserved for private networks.

You can find the details about the special use of IPv4 addresses in RFC5735.









2.2.3 IP/Mask

To fully identify a host, you also need to know its **network**. To do that, you will need an IP address and a **netmask**, or subnet mask.

With an IP/netmask pair, you can identify the network part and the host part of an IP address.

IP address: 192.168.5.100

Subnet mask: 255.255.255.0







2.2.3 **IP/Mask**

To find the network part you have to perform a **bitwise AND operation** between the netmask and the IP address.

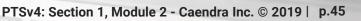
In the following example, we are going to see how to find the network part of this IP address/mask pair:

192.168.33.12/255.255.224.0



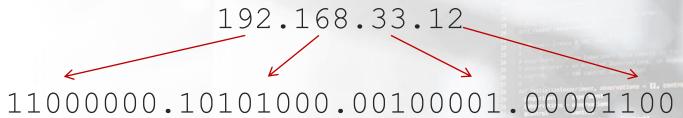


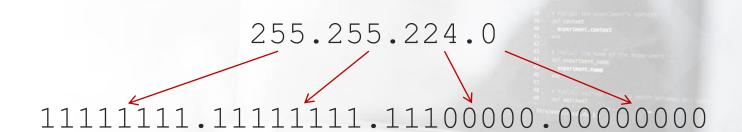




2.2.3.1 IP/Mask CIDR Example

Convert the octets in binary form:











2.2.3.1 IP/Mask CIDR Example

2 Perform the bitwise AND:

```
IP: 11000000.10101000.00100001.00001100
```

1111111.1111111.11100000.0000000

Network: 11000000.10101000.00100000.00000000

Network prefix in decimal notation:

192.168.32.0







2.2.3.1 IP/Mask CIDR Example

192.168.32.0 is the **network prefix**. You can identify the network by using the following notation:

192.168.32.0/255.255.224.0

Or, as the netmask is made by 19 consecutive "1" bits:

192.168.32.0/19

The latter is the **Classless Inter-Domain Routing (CIDR)** notation.









The address part not covered by the netmask is the **host** part of the IP address. You can find it by performing a bitwise *AND* with the **inverse of the netmask**.

Let's look at an example with the same IP/mask.









1

Convert the octets in binary form:

192.168.33.12 11000000.10101000.00100001.00001100









Invert the netmask by performing a bitwise NOT:

 \neg (11111111.1111111.11100000.00000000)

_

0000000.00000000.00011111.11111111







3

Perform the final bitwise AND:

Host: 00000000.0000000.0000001.00001100

Host part in decimal notation:

0.0.1.12









Moreover, the inverse of the netmask lets you know how many hosts a network can contain.

In our example, we have 13 bits to represent the hosts; this means that the network can contain 2^{13} = **8192 different** addresses.









2.2.4 Network and Broadcast Addresses

There are two special addresses:

- One with the host part made by all zeros.
- Another with the host part made by all ones.

These special addresses were used as the **network** and **broadcast** addresses, thus reducing by 2 the number of hosts on a given network. This technical limitation should be extinct (RFC1878) but is still used to keep compatibility with old equipment.









Let's recap by going over some IP examples.









- 10.54.12.0/24 (10.54.12.0/255.255.255.0)
 - Contains 2⁸ = 256 addresses

- 10.54.12.0 is the network address according to the pre-CIDR standard
- 10.54.12.255 is the broadcast address according to the pre-CIDR standard









- 192.168.114.32/27 (192.168.114.32/255.255.255.224)
 - Contains 2⁵ = 32 addresses

192.168.114.32 is the pre-CIDR network address

192.168.114.63 is the pre-CIDR broadcast address









- Given the network 172.16.2.0/23
 - 172.16.3.12 and 172.16.2.66 **are** in the same network
 - 172.16.3.240 and 172.16.4.2 **are not** in the same network

- The network 192.168.1.0/16
 - Does not make sense; a bitwise *AND* between 192.168.1.0 and 255.255.0.0 leads to 192.168.0.0 as network address
 - Could be a valid IP address in the 192.168.1.0/16 network









2.2.6 Subnet Calculators

You can practice more on this topic by using a subnet calculator.

Here are two subnet calculators you can check out:

- A classful calculator
- A CIDR calculator









IPv4 addresses are being consumed rapidly due to a large number of new devices connecting to the internet every day.

One day IPv4 addresses might be exhausted.







As a 32-bit address, **IPv4** has 2^32 = **4.294.967.296** possible addresses.

While a 128-bit **IPv6** address has **2^128 = 2^32 * 2^96** possible addresses.

2⁹⁶ is equal to **79 octillion addresses**









An **IPv6** address consists of **16-bit hexadecimal numbers** separated by a **colon** (:). Hexadecimal numbers are case insensitive. In case zeros occur, they can be skipped.

Let's check out some IPv6 examples on the next slide.







IPv6 addresses examples:

2001:0db8:0020:130F:0000:0000:087C:140B

2001:0db8:0:160F::850C:140B







2.2.7.1 IPv6 header

0 0 1 2 3 4 5	15 6 7 0 1 2 3 4 5 6 7	16 0 1 2 3 4 5 6 7	0 12 3 4 5 6 7
Version (4)	Traffic Class (8)	Flow label (20)	
Payload length (16)		Next header (8)	Hop limit (8)
	Source A	address (128)	
	Destination	Address (128)	









2.2.7.2 IPv6 forms

IPv6 can be presented in following text representations:

Regular form: 1080:0:FF:0:8:800:200C:417A

Compressed form: FF01:0:0:0:0:0:0:0:43 becomes
 FF01::43 as a result of skipping zeros

• IPv4-compatible: 0:0:0:0:0:0:13.1.68.3 or ::13.1.68.3 after skipping zeros









2.2.7.3 IPv6 Reserved Addresses

IPv6 also has reserved addresses, which cannot be used like the reserved IPv4 ones.

For example:

- ::1/128 is a loopback address
- ::FFFF:0:0/96 are IPv4 mapped addresses

For more information, you can check <u>RFC3513</u>









2.2.7.4 IPv6 Structure

An **IPv6 address** can be split in half (64 bits each) into a **network part** and a **device part**.

Furthermore, the **first 64 bits** ends with a **dedicated 16-bits space** (one hex word) that can be used only for **specifying a subnet.**

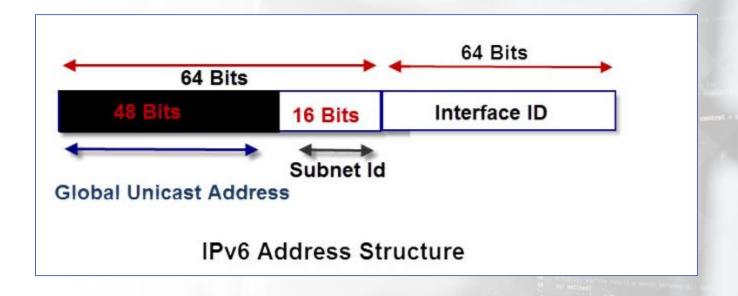








2.2.7.4 IPv6 Structure











2.2.7.5 IPv6 Scope

Address Types and Scope

IPv6 addresses have three types:

 Global Unicast Address – These addresses are global ones, and reside in global internet.

 Unique Local and Link Local — reside only in Internal Networks.









2.2.7.5 IPv6 Scope



Address Types and Scope

Global Unicast Address -- Scope
Internet- Routed on Internet
Unique Local -- Scope Internal
Network or VPN -Internally routable
but Not routed on Internet
Link Local - Scope network linkNot Routed internally or externally.









2.2.7.6 IPv6 Translation

IPv6 addresses can also be translated to binary.

One 4-digit hex word represents **16 binary digits**; we can see this demonstrated in the following way:

- Bin 000000000000000 = Hex 0000 (or just 0)
- Bin 11111111111111 = Hex FFFF
- Bin 1101010011011011 = Hex D4DB

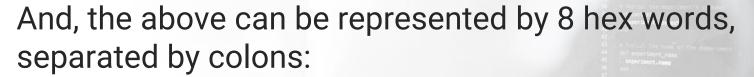






2.2.7.6 IPv6 Translation

Thus, 128-bit binary address looks like:



FFFF:FFFF:FFFF:FFFF:FFFF:FFFF









Like IPv4, an IPv6 address has a network portion and a device portion.

Unlike IPv4, an IPv6 address has a dedicated subnetting portion. On the next few slides, we'll show how the ranges are divided in IPv6.









Network Address Range

In IPv6, the first 48 bits are for Internet global addressing.









Subnetting Range

The 16 bits from the 49th to the 64th are for defining subnets.









Device (Interface) Range

The last 64 bits are for device (interface) ID's:









In **IPv6**, there are **prefixes** instead of subnets blocks. For example:

2001:1111:1234:1234::/64

In the above IPv6 address, the number after the slash (64) is the **number of bits that is used for a prefix**. Everything behind it can be used for **hosts** of the **subnet**.







As you may have noticed, /64 means that the first 64 bits are a prefix. And, as previously mentioned earlier, each 4-digit hex word is 16 bits, thus in following IPv6 address we can divide it as such:

Prefix Host

2001:1234:5678:1234

5678:ABCD:EF12:1234

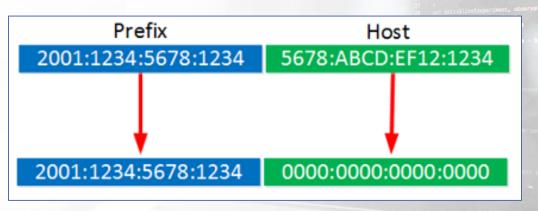








We confirmed that **2001:1234:5678:1234** is the prefix, but let's now focus on writing down a correctly formatted IPv6 address.











2001:1234:5678:1234:0000:0000:0000:0000 is a valid prefix, but it can be shortened by omitting zeros, into following form:

2001:1234:5678:1234::/64









You can practice more on this topic by using a subnet calculator.

Here is a calculator you can check out:

IPv6 Calculator











Routing







2.3 Routing

How does this support my pentesting career?

Understanding routing protocol attacks

Performing network traffic inspection

2.3 Routing

Addressing devices is just half of the work needed to reach a host. Your packets need to follow a valid **path** to reach it.

Routers are devices connected to different networks at the same time. They are able to forward IP datagrams from one network to another. The forwarding policy is based on routing protocols.









2.3 Routing

Routing protocols are used to determine the best path to reach a network. They behave like a postman who tries to use the shortest path possible to deliver a letter.

A router inspects the destination address of every incoming packet and then forwards it through one of its interfaces.









2.3.1 Routing Table

To choose the right forwarding interface, a router performs a lookup in the **routing table**, where it finds an IP-to-interface binding.

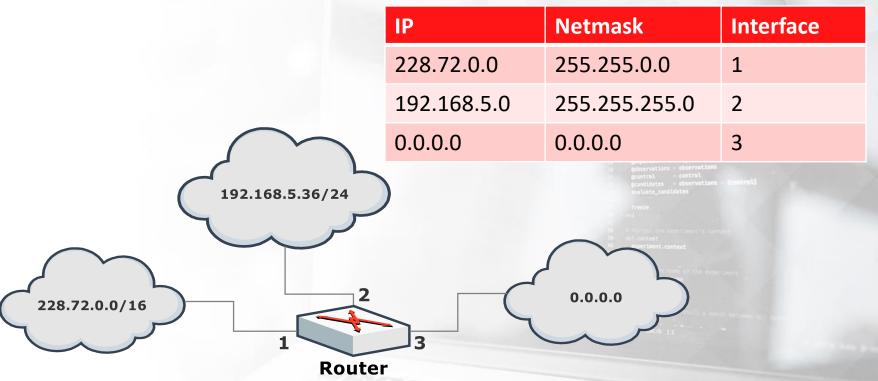
The table can also contain an entry with the **default address** (0.0.0.0). This entry is used when the router receives a packet whose destination is an *unknown network*.



















In this example, the routing table is made of three entries.

- Interface 1 is used to forward the packets to 228.72.0.0/16.
- Interface 2 is used to forward the packets to 192.168.5.0/24.
- Interface 3 is used as the default route for packets whose destination does not match any other entry in the table.

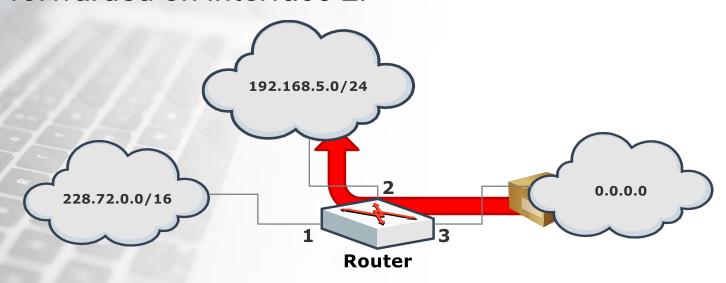








A packet arriving on interface 3 for 192.168.5.3 is forwarded on interface 2.

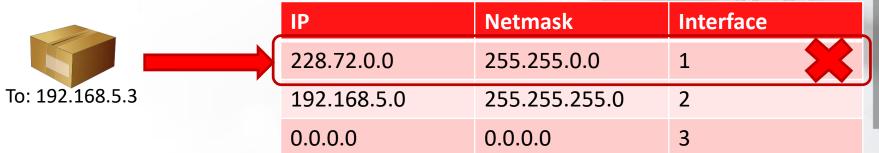








In fact, the first entry in the routing table does not match the destination network.











While the second does: 192.168.5.3 sits in the 192.168.5.0/24 network.

	IP	Netmask	Interface	
	228.72.0.0	255.255.0.0	1	
	192.168.5.0	255.255.255.0	2	√
To: 192.168.5.3	0.0.0.0	0.0.0.0	3	



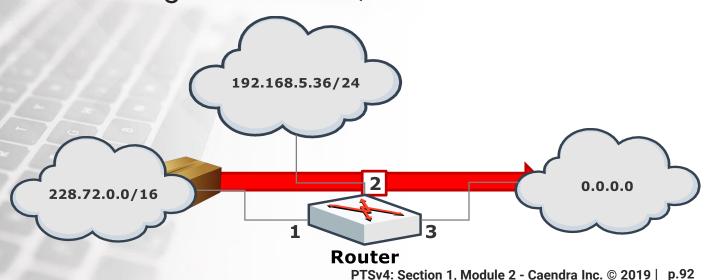






2.3.1.2 Default Route Example

A packet arriving on interface 1 for 72.13.37.2 is routed through interface 3, the default route.





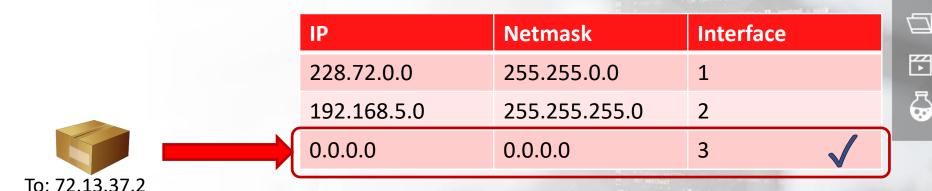






2.3.1.2 Default Route Example

There is no matching entry, so the router forwards the packet through interface 3.



PTSv4: Section 1, Module 2 - Caendra Inc. © 2019 | p.93

2.3.2 Routing Metrics

As in the real world, there could be more than a way to reach a destination.

So, during path discovery, routing protocols also assign a **metric to each link**.









2.3.2 Routing Metrics

This ensures that, if two paths have the same number of hops, the fastest route is selected.

The metric is selected according to the channel's estimated bandwidth and congestion.



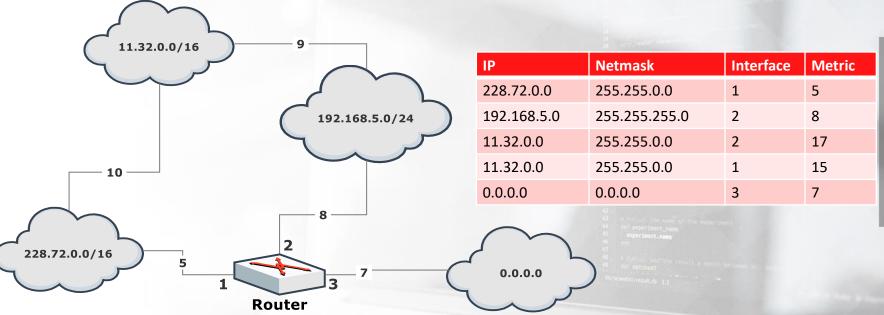






2.3.2.1 Routing Metrics Example

Let's look at how routing decisions are made according to metrics.







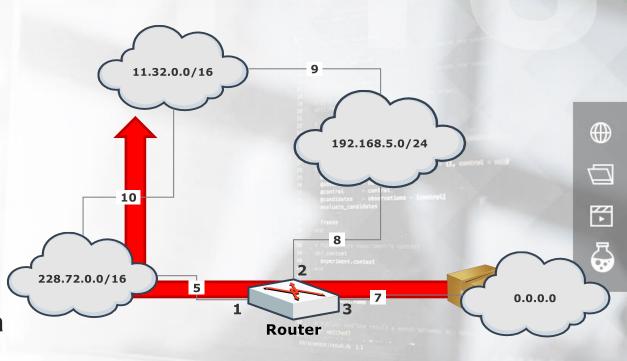




2.3.2.1 Routing Metrics Example

A packet arriving on interface 3 for 11.32.3.118 is routed through interface 1, as the metric for that route is 15.

Routing through interface 2 would have a metric of 17.



Routing tables are not only kept by routers; every host stores its own table.

To check what they look like, you can use:

- ip route on Linux
- route print on Windows
- netstat -r on OSX









Checking the routing table on a Linux box:

</>

root@host:~# ip route

default via 192.168.51.1 dev eth0 proto static
192.168.51.0/24 dev wlan0 proto kernel scope link src 192.168.51.123









Checking the routing table on Microsoft Windows:

</> C:\Users\User>route print Interface List 11...08 00 27 bf ac c8Intel(R) PRO/1000 MT Desktop Adapter 1.....Software Loopback Interface 1 TPv4 Route Table Active Routes: Network Destination Netmask Gateway Interface Metric 0.0.0.0 0.0.0.0 10.0.2.2 10.0.2.15 10 10.0.2.0 255.255.255.0 10.0.2.15 On-link 266







Checking the routing table on Mac OSX:

		9								
</th <th colspan="10">User:~ user\$ netstat -r</th>	User:~ user\$ netstat -r									
	Routing tables									
	<pre>Internet:</pre>									
	Destination	Gateway	Flags	Refs	Use	Netif E	Expire			
Ę,	default	192.168.51.1	UGSc	13	0	en1				
	127	127.0.0.1	UCS	0	0	100				
	127.0.0.1	127.0.0.1	UH	1	16	100				
24	169.254	link#4	UCS	0	0	en1				
	192.168.51	link#4	UCS	4	0	en1				
	192.168.51.1	58:6d:8f:e5:e:d2	UHLWIir	14	24	en1	1200			
	192.168.51.109	2:f:b5:4b:76:cf	UHLWIi	0	0	en1	1148			

















2.4 Link Layer Devices and Protocols

How does this support my pentesting career?

- MAC spoofing
- Testing switches security
- Sniffing techniques
- Man in the middle attacks

2.4 Link Layer Devices and Protocols

Packet forwarding also happens in the **lowest layer** of the TCP/IP stack: the **link layer**.

While routers are aware of the best overall path to the destination, link layer devices and protocols deal only with the next hop.









2.4 Link Layer Devices and Protocols

In this section you will see:

How switches work

Network card's MAC addresses

The Address Resolution Protocol (ARP)









2.4.1 Link Layer Devices

Hubs and switches are network devices that forward frames (layer 2 packets) on a local network.

They work with link layer network addresses: MAC addresses.









2.4.2 Mac Addresses

IP addresses are the Layer 3 (Network layer) addressing scheme used to identify a host in a network, while **MAC** addresses uniquely identify a network card (Layer 2).

A MAC (Media Access Control) address is also known as the **physical address**.









2.4.2 Mac Addresses

MAC addresses are 48 bit (6 bytes) long and are expressed in hexadecimal form (HEX).

00:11:AA:22:EE:FF







2.4.2 Mac Addresses

To discover the MAC address of the network cards installed on your computer, you can use:

- ipconfig /all on Windows
- ifconfig on *nix operating systems, like MacOS
- ip addr on Linux







2.4.2 Mac Addresses

Every host on a network has both a MAC and an IP address.

Let us see how they are used together to send packets.

Remember: the lower layer serves the layer above.



7

Let's take a look at an example to see how MAC addresses are used.



MAC: 12:af:09:ee:bb:01

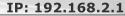


IP: 192.168.2.15

MAC: 0a:0c:11:33:5b:aa

IP: 10.32.1.1

MAC: d0:50:99:11:67:2



MAC: d0:50:99:11:67:30









Two different networks are connected together by a router:

- 10.32.1.0/24
- 192.168.2.0/24



IP: 10.32.1.4

MAC: 12:af:09:ee:bb:01

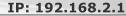


IP: 192.168.2.15

MAC: 0a:0c:11:33:5b:aa

IP: 10.32.1.1

MAC: d0:50:99:11:67:2f



MAC: d0:50:99:11:67:30









Every host on the network has both an IP and a MAC address. The router has two interfaces, each with its own addresses.



IP: 10.32.1.4

MAC: 12:af:09:ee:bb:01

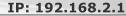


IP: 192.168.2.15

MAC: 0a:0c:11:33:5b:aa

IP: 10.32.1.1

MAC: d0:50:99:11:67:2



MAC: d0:50:99:11:67:30









If workstation A wants to send a packet to workstation B, which IP and MAC addresses will it use?



MAC: 12:af:09:ee:bb:01

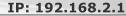


IP: 192.168.2.15

MAC: 0a:0c:11:33:5b:aa

IP: 10.32.1.1

MAC: d0:50:99:11:67:2f



MAC: d0:50:99:11:67:30









Workstation A will create a packet with:

- The destination IP address of workstation B in the IP header of the datagram.
- The destination MAC address of the router in the link layer header of the frame.
- The source IP address of workstation A
- The source MAC address of workstation A









The router will then take the packet and forward it to B's network, **rewriting the packet's MAC addresses**:

- The destination MAC address will be B's
- The source MAC address will be the router's

The router will not change the source and destination IP addresses.









When a device sends a packet:

 The destination MAC address is the MAC address of the next hop; this ensures that, locally, the network knows where to forward the packet.

 The destination IP address is the address of the destination host; this is global information and remains the same along the packet trip.









This method, in a way, recalls how you send a letter to a friend.

You need to know his or her home address (IP address) and the address of the nearest post office (MAC address) where you can drop the letter.









2.4.4 Broadcast MAC Address

There is also a special MAC address

FF:FF:FF:FF:FF

...which is the broadcast MAC address.

A frame (the name of the packets at Layer 2) with this address is delivered to all the hosts in the local network (within the same broadcast domain).









While routers work with IP addresses, switches work with MAC addresses. Switches also have multiple interfaces, so they need to keep a **forwarding table** that binds one or more MAC addresses to an interface.



MAC	Interface	TTL
00:11:22:33:44:55	1	30
AA:BB:CC:DD:EE:01	2	5
AA:CC:FF:0A:0C:12	2	5
11:22:33:1D:CC:0A	3	7









The forwarding table is called Content Addressable Memory (CAM) table. Many hosts can connect to a switch. Let's see how.



MAC	Interface	TTL
00:11:22:33:44:55	1	30
AA:BB:CC:DD:EE:01	2	5
AA:CC:FF:0A:0C:12	2	5
11:22:33:1D:CC:0A	3	7









The smallest switches you can encounter are home switches, usually integrated into a DSL home router. They usually have 4 ports.

Corporate switches may have up to 64 ports, and system administrators can connect multiple switches together to accommodate more hosts.







The main difference between one switch and another is the packet forwarding speed.

The speed of a switch varies from 10Mbps (megabits per second) to 10Gbps (gigabits per second). Nowadays, 1Gbps is the most common forwarding speed in commercial switches.





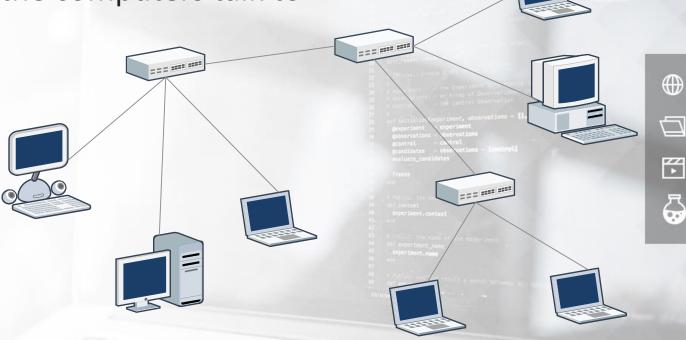


2.4.5.1 Multi-switch Network

In this diagram, all the machines are on the same network.

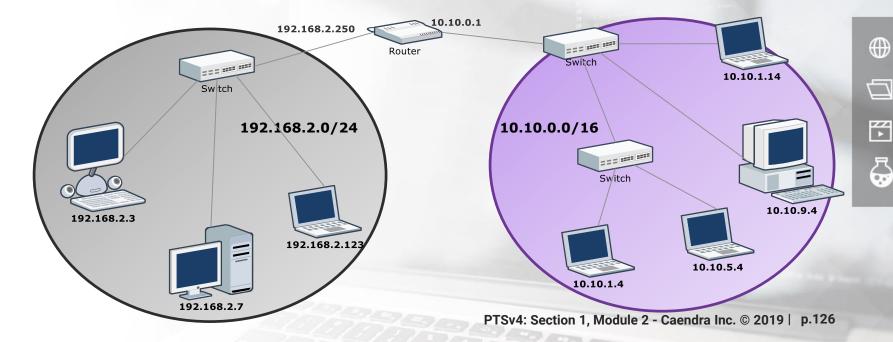
2.4.5.1 Multi-switch Network

Switches let all the computers talk to each other.



2.4.5.2 Segmentation

Switches, without VLANs, do not **segment** networks. Routers do.



2.4.5.2 Segmentation

Usually, every interface of a router is attached to a different subnet with a different network address.

Also, routers do not forward packets coming from one interface if they have a ff:ff:ff:ff:ff:ff broadcast MAC address (imagine if they did!).



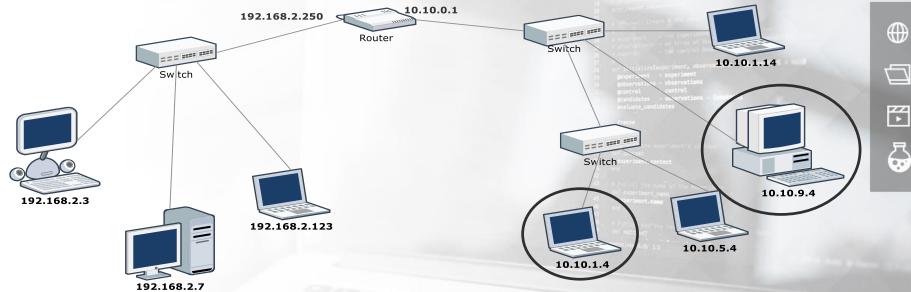






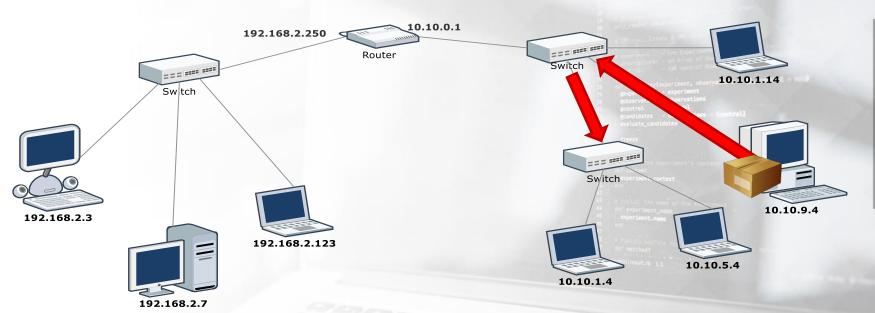
2.4.5.3 Multi-switch Example

What happens if 10.10.9.4 wants to send a packet to 10.10.1.4?



2.4.5.3 Multi-switch Example

The first switch receives the packet, performs a look-up in the CAM table and forwards it to the next switch.



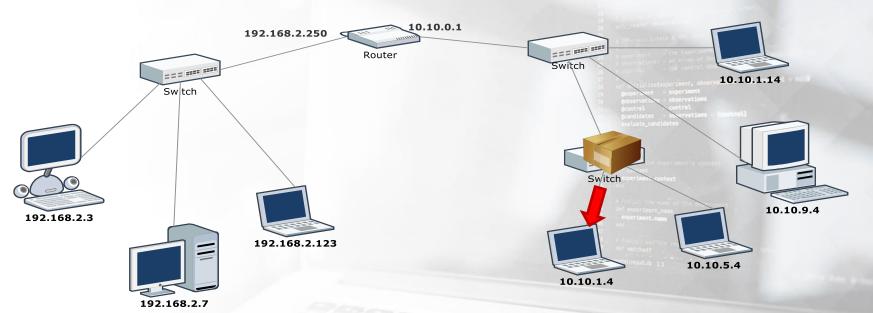






2.4.5.3 Multi-switch Example

The second switch forwards the packet to 10.10.1.4.

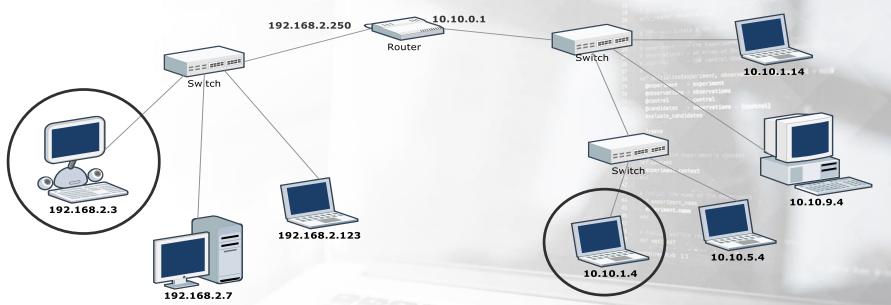








What happens if 10.10.1.4 wants to send a packet to 192.168.2.3?

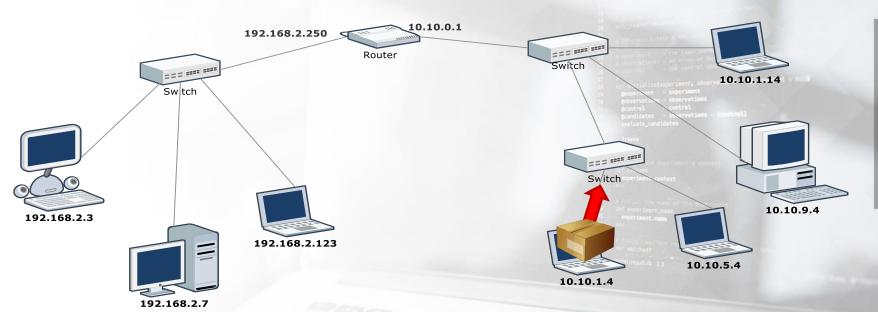








10.10.1.4 needs to send the packet to the router so that the first switch will forward the packet to the next one.

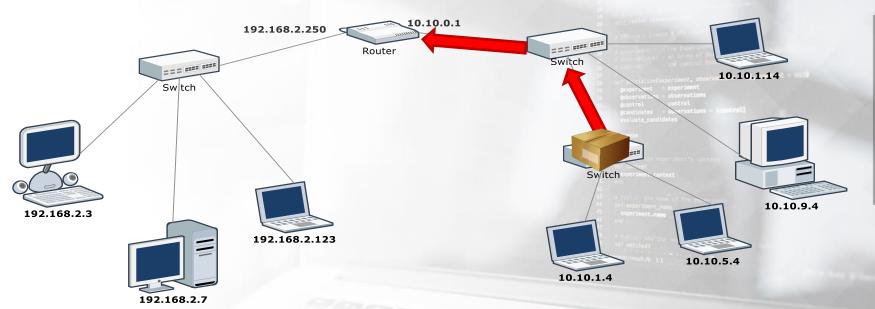








The packet then arrives at the router that, after a look up in the routing table, forwards it to the 192.168.2.0/24 network.



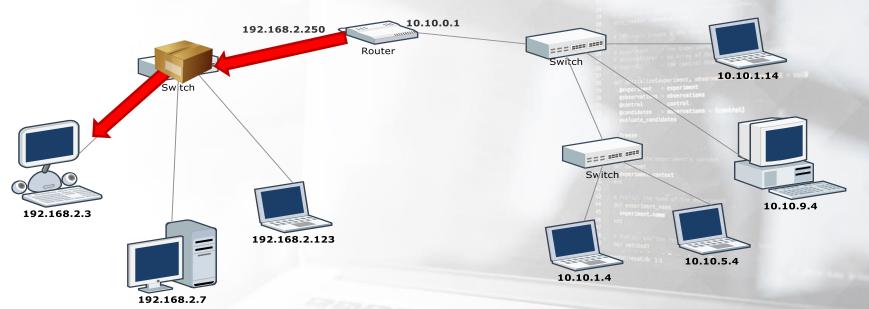








Finally, the packet is delivered.









A forwarding table binds MAC addresses to interfaces.

In the following slides you will see:

- The structure of the table
- The way a switch constructs the MAC address interface binding
- How forwarding works









A typical forwarding table contains:

- The MAC address
- The interfaces the switch can use to deliver packets to a specific MAC address
- A time to live (TTL)

MAC	Interface	TTL
00:11:22:33:44:55	1	30
AA:BB:CC:DD:EE:01	2	5
AA:CC:FF:0A:0C:12	2	5
11:22:33:1D:CC:0A	3	7









The forwarding table, or Content Addressable Memory table (CAM table), is stored in the device's RAM and is constantly refreshed with new information.

MAC	Interface	TTL
00:11:22:33:44:55	1	30
AA:BB:CC:DD:EE:01	2	5
AA:CC:FF:0A:0C:12	2	5
11:22:33:1D:CC:0A	3	7









Looking at the table you can tell that:

- A single host is attached to Interface 1 and 3 respectively
- Two hosts are attached to interface 2 (probably via another switch).

MAC	Interface	TTL
00:11:22:33:44:55	1	30
AA:BB:CC:DD:EE:01	2	5
AA:CC:FF:0A:0C:12	2	5
11:22:33:1D:CC:0A	3	7









There might be multiple hosts on the same interface and interfaces without any host attached.

In our example interface, 4 has no hosts attached.

MAC	Interface	TTL
00:11:22:33:44:55	1	30
AA:BB:CC:DD:EE:01	2	5
AA:CC:FF:0A:0C:12	2	5
11:22:33:1D:CC:0A	3	7









The TTL determines how long an entry will stay in the table. This is important because the **CAM table has a finite size**.

So, as soon as an entry expires it is removed from the table.

MAC	Interface	TTL
00:11:22:33:44:55	1	30
AA:BB:CC:DD:EE:01	2	5
AA:CC:FF:0A:0C:12	2	5
11:22:33:1D:CC:0A	3	7









2.4.5.6 CAM Table Population

Switches learn new MAC addresses dynamically; they inspect the header of every packet they receive, thus identifying new hosts.

While routers use complex routing protocols to update their routing rules, switches just use the source MAC address of the packets they process to decide which interface to use when forwarding a packet.









2.4.5.6 CAM Table Population

The source MAC address is compared to the CAM table:

- If the MAC address is not in the table, the switch will add a new MAC-Interface binding to the table
- If the MAC-Interface binding is already in the table, its TTL gets updated
- If the MAC is in the table but bound to another interface the switch updates the table







2.4.5.7 Forwarding

To forward a packet:

- 1 The switch reads the destination MAC address of the frame.
- 2 It performs a look-up in the CAM table.
- 3 It forwards the packet to the corresponding interface.
- 4 If there is no entry with that MAC address, the switch will forward the frame to all its interfaces.









2.4.6 ARP

When a host wants to send a packet to another host, it needs to know the IP and the MAC address of the destination in order to build a proper packet.

You wouldn't be able to send your friend a letter if you don't know his/her address, right? What happens if the source host knows the IP address, but not the MAC address of the destination host?









This situation occurs in many circumstances, for example at every power up.

 A PC in an office knows a bunch of IP addresses, like the fileserver, the printers, and the webserver, but not their corresponding MAC addresses.

The host needs to know the MAC addresses of the other network nodes, and it can learn them by using the **Address Resolution Protocol (ARP)**.









With ARP a host can build the correct IP Address – MAC address binding.

This is one of the most fundamental protocols any modern network uses, so make sure to fully understand it.









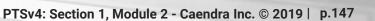
When a host (A) wants to send traffic to another (B), and it only knows the IP address of B:

- 1. A builds an **ARP request** containing the IP address of *B* and FF:FF:FF:FF:FF as destination MAC address. This is fundamental because the switches will forward the packet to every host.
- 2. Every host on the network will receive the request.
- 3. B replies with an ARP reply, telling A its MAC address.

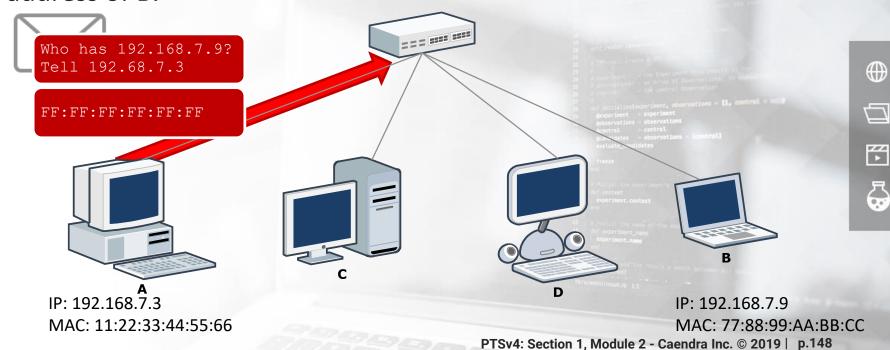




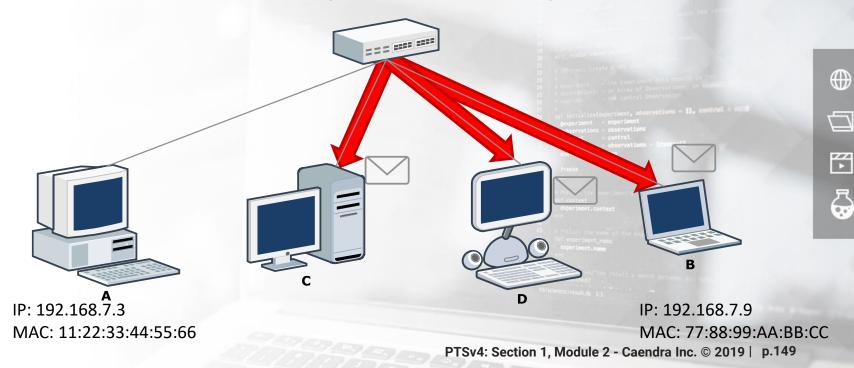




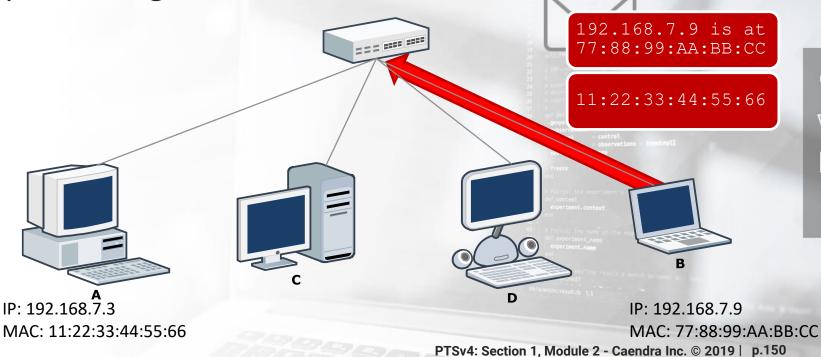
'A' sends a packet to the broadcast MAC address, asking for the MAC address of B.



The switch forwards the packet to all its ports.

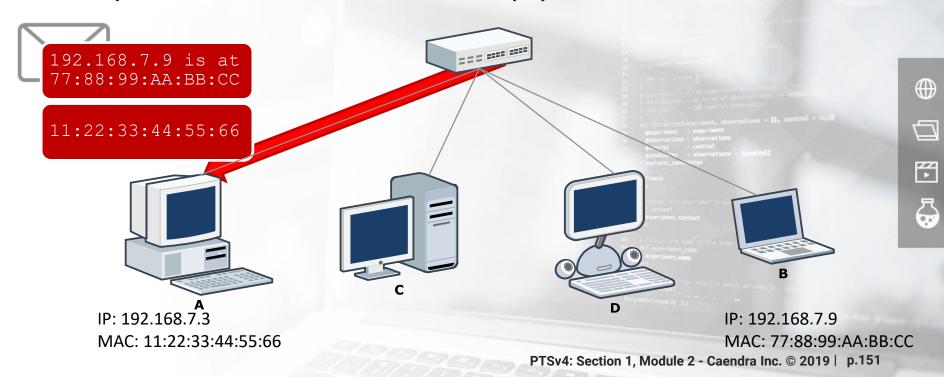


3 B replies telling A its MAC address.



 \Box

Finally, the switch forwards the reply to A.



'A' will save the IP – MAC binding in its ARP cache. Further traffic to 'B' will not need a new ARP resolution protocol round.

ARP cache entries have a TTL too, as the size of the device RAM is finite. A host discards an entry at the power off or when the entry's TTL expires.







2.4.6.1 Checking the ARP Cache

You can check the ARP cache of your host by typing:

- arp -a on Windows.
- arp on *nix operating systems
- ip neighbour on Linux

```
$ ip neighbour

192.168.17.202 dev eth0 lladdr d0:d4:12:e1:ef:5a STALE

192.168.17.1 dev eth0 lladdr 00:50:7f:78:fc:40 STALE

192.168.17.99 dev eth0 lladdr 00:d0:4b:92:2d:89 STALE

192.168.17.14 dev eth0 lladdr 60:a4:4c:a8:be:5b STALE

192.168.17.18 dev eth0 lladdr 20:cf:30:c7:ad:ae STALE

192.168.17.30 dev eth0 lladdr 20:cf:30:ea:22:13 STALE

192.168.17.66 dev eth0 lladdr a4:ee:57:e8:2e:0b STALE

192.168.17.254 dev eth0 lladdr c8:4c:75:a4:79:a6 REACHABLE

192.168.17.12 dev eth0 lladdr 60:a4:4c:a8:bd:la STALE

192.168.17.19 dev eth0 lladdr 54:04:a6:a0:6e:ad STALE

192.168.17.24 dev eth0 lladdr bc:5f:f4:ef:63:51 STALE
```







2.4.7 Hubs

Hubs were used in computer networks before switches. They have the same **purpose** but not the same **functionality**.

Hubs are simple repeaters that do not perform any kind of header check and simply forward packets by just repeating electric signals. They receive electric signals on a port and repeat the same signals on all the other ports.









2.4.7 Hubs

This means that every node on a hub-based network receives the same electric signals, thus the same packets.

Nowadays, hubs are very rare as they have mostly been replaced by switches.



















- TCP Session Attacks
- Advanced DoS attacks
- Network scanning

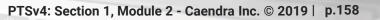
In this section, you will see how the **transport layer** works, and how the application layer uses its services to identify server and client processes.

The **Transmission Control Protocol** (TCP) and the **User Datagram Protocol** (UDP) are the most common transport protocols used on the Internet.









Before checking out the different services that a transport layer protocol can offer to the application layer, let's consider something important about networks.

Computer networks can be **unreliable**. This means that some **packets can be lost** during their trip from source to destination. A packet can be lost because of network congestion, temporary loss of connection and other technical issues.









When designing a transport layer protocol, the designer must choose how to deal with these limitations. For example, TCP:

- Guarantees packet delivery. Because of that, an application that needs a guaranteed delivery will use TCP as the transport protocol.
- Is also **connection oriented**. It must establish a connection before transferring data.

Keep in mind these facts during your study!









TCP is the most used transport protocol on the Internet. The vast majority of applications use it, and the IP protocol suite is often called **TCP/IP**.

Email clients, web browsers and FTP clients are some common applications using TCP.









On the other hand, UDP is much more simple than TCP:

- It does not guarantee packet delivery.
- It is connectionless.









UDP is faster than TCP, as it provides a **better throughput** (number of packets per second); in fact, it is used by **multimedia applications** that can tolerate packet loss but are throughput intensive.

For example, UDP is used for VoIP and video streaming: applications where you can tolerate a little glitch in the audio or video.









Here we can see a comparison table between TCP and UDP.

ТСР	UDP
Lower throughput	Better throughput
Connection-oriented	Connectionless
Guarantees delivery	Does not guarantee packet delivery









2.5.1 Ports

Applications and their processes use TCP and UDP to send and receive data over the network. When an IP datagram reaches a host, how can the transport layer **know what the destination process is**?

We'll now introduce ports.









2.5.1 Ports

Ports are used to identify a single network **process** on a machine. If you want to unequivocally identify a process on a network, you need to know the **IP**: **Port**> **pair**.

As an example, you can compare the port to the recipient's name on a letter; the street address (IP) identifies the building, while the person name identifies the final recipient of the letter.

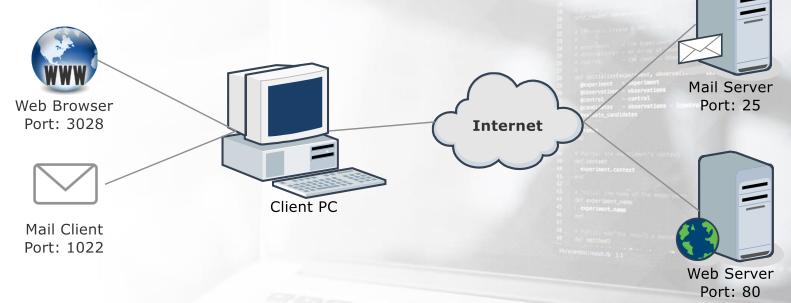








In this image, you can see how every client application on *Client PC* uses a different port.

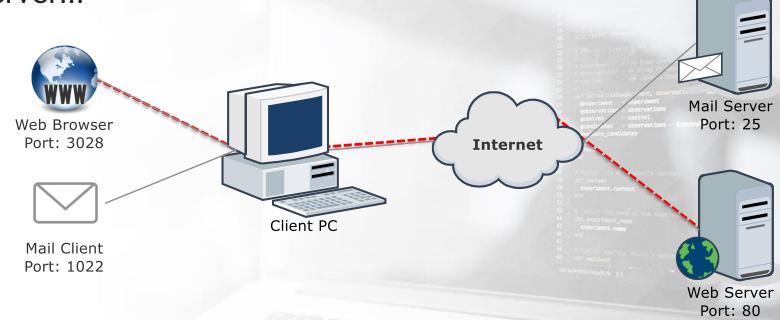








The browser uses local port 3028 to connect to the web server...

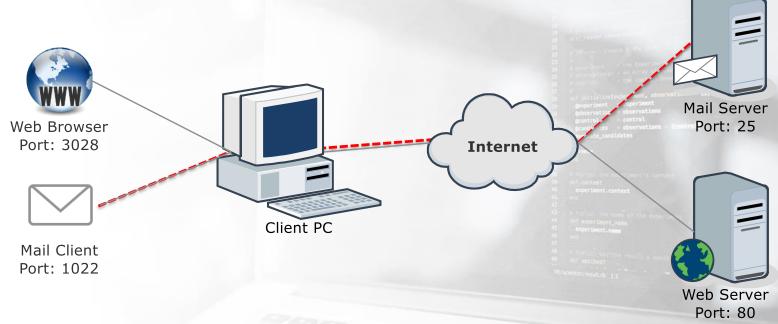








... while the mail client uses local port 1022.



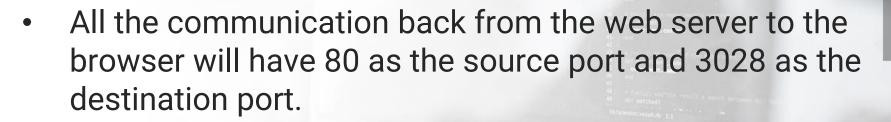






In the previous example:

 All the communication from the web browser to the web server will have 3028 as the source port and 80 as the destination port.











Similarly, for the mail client and server:

 All the communication from the mail client to the server will have 1022 as the source port and 25 as the destination port.

 All the communication back from the mail server to the mail client will have 25 as the source port and 1022 as the destination port.



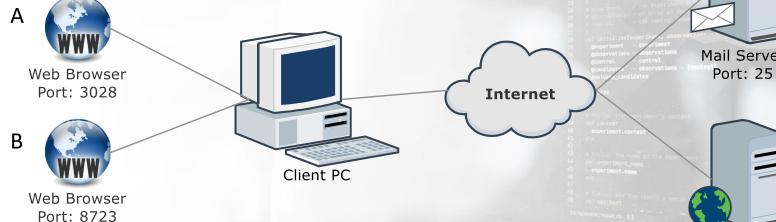


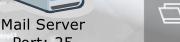




Furthermore, you may also have multiple instances of the same application running at the same time. Every process will





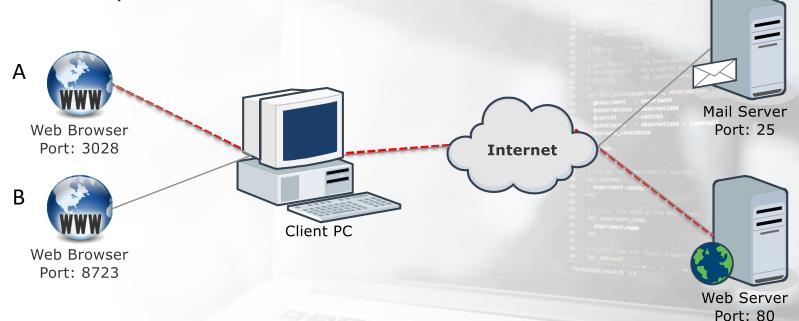




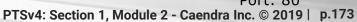


Port: 80 PTSv4: Section 1, Module 2 - Caendra Inc. © 2019 | p.172

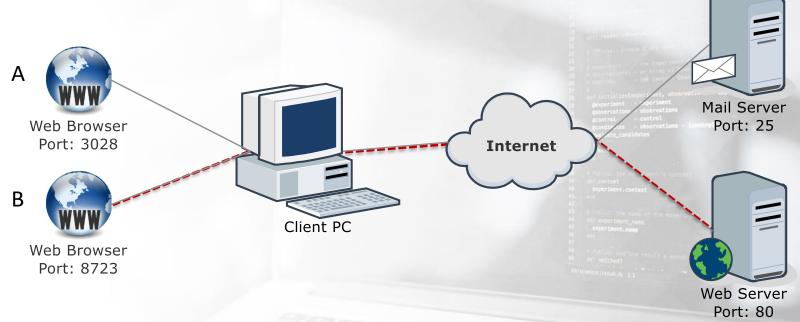
In this example, 'A' communicates with the web server using 3028 as the source port...







... while 'B' uses port 8723.





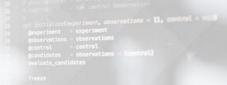




2.5.1 Ports

To correctly address a process on a network, you have to refer to the <IP>:<Port> pair. For example:

- 192.168.5.3:80
- 10.11.12.1:443
- 172.16.8.9:22





But, how can you know the right port for a common service?

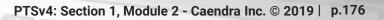
Ports in the ranging from **0-1023**, the first 1024 that is, are called **well-known ports** and are used by servers for the most common services.

For example, when a web browser connects to a server via HTTPS, the user does not have to manually specify 443 as the destination port.









Each common protocol has a well-known port in the 0-1023 range. Common server processes, or daemons, use well-known ports most of the time.

Ports are assigned by IANA and are referenced in this document.









You do not need to know all the service port assignments, but you should at least remember the most common, such as:

- SMTP **(25)**
- SSH (22)
- POP3 **(110)**
- IMAP (143)
- HTTP (80)
- HTTPS (443)
- NETBIOS (137, 138, 139)

- SFTP (115)
- Telnet (23)
- FTP (21)
- RDP (3389)
- MySQL (3306)
- MS SQL Server (1433)







As briefly introduced before, a **daemon** is a program that runs a service. System administrators can change the daemon configuration, **changing the port** the service listens to for connection. They do that to make services recognition a little bit harder for hackers.

For example, you could find an FTP daemon listening on port 4982 instead of 21 or SSH listening on port 8821.









2.5.3 TCP and UDP headers

Let's now see how ports are used by applications.

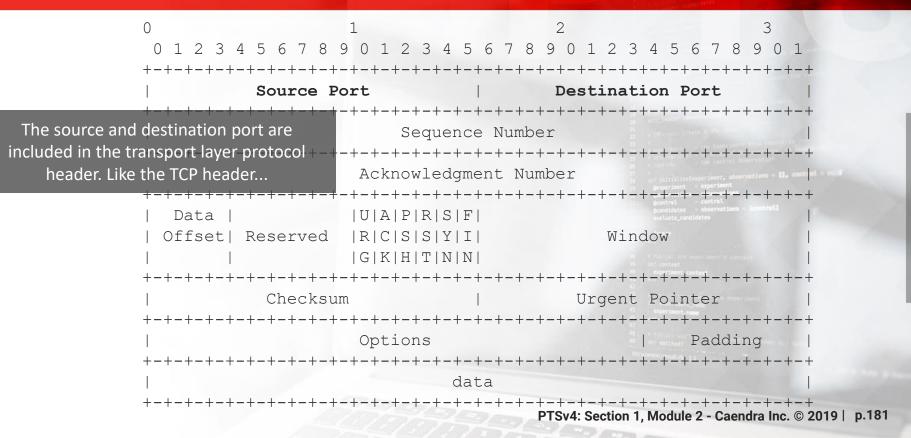
How can server and client applications know which port to use? They use two fields in the TCP or UDP header: the **source** and **destination** ports.







2.5.3.1 TCP Header

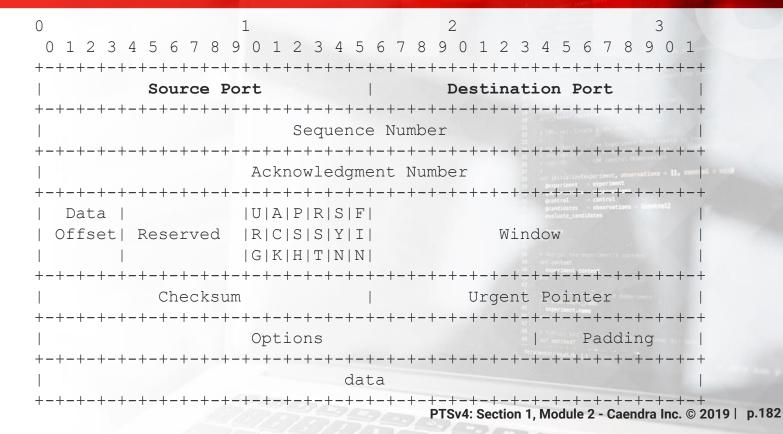








2.5.3.1 TCP Header



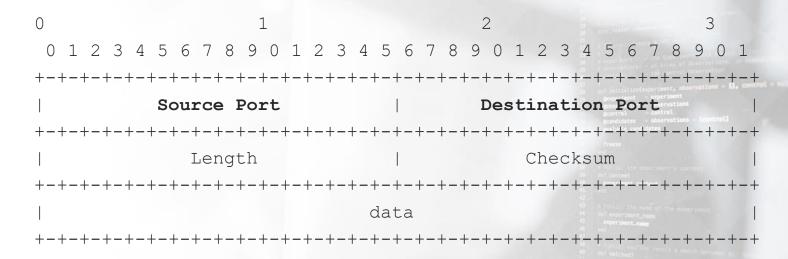






2.5.3.2 UDP Header

...or the UDP header.

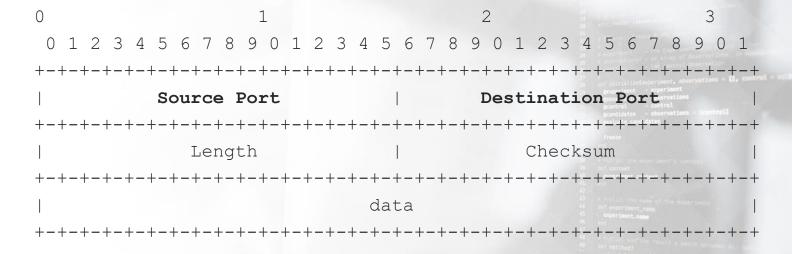








2.5.3.2 UDP Header











2.5.4 Netstat Command

To check the listening ports and the current (TCP) connections on a host you can use:

- netstat -ano on Windows
- netstat -tunp on Linux
- netstat -p tcp -p udp together with lsof -n -i4TCP -i4UDP on MacOS

Use these commands to show information about the processes listening on the machine and processes connecting to remote servers.







2.5.4 Netstat Command

Another great tool for Windows is <u>TCPView</u> from Sysinternals.

TCPView shows:

- Process name
- PID
- Protocol
- Local and remote addresses

- Local and remote ports
- State of the connection (if applicable)









We have seen that TCP is **connection oriented**. Now, let's look at how TCP connections work, as well as highlight the most important factors involved, from the penetration tester's point of view, in a 3-way handshake.

To establish a connection between two hosts running TCP, they must perform three steps: the **three-way handshake**. They can then start the actual data transmission.









The header fields involved in the handshake are:

- Sequence number
- Acknowledgement numbers
- SYN and ACK flags

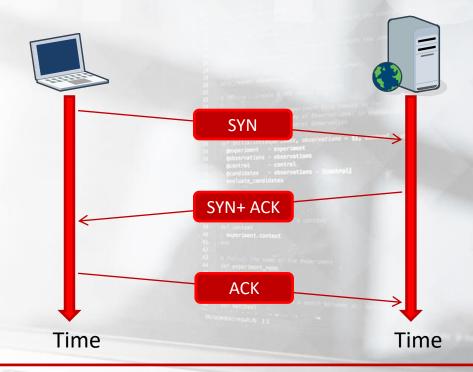
```
Source Port
Sequence Number
Acknowledgment Number
Data
         |U|A|P|R|S|F|
Offset| Reserved
         |R|C|S|S|Y|I
         |G|K|H|T|N|N|
```







The steps in the handshake are used to synchronize the sequence and acknowledgment numbers between the server and the client.

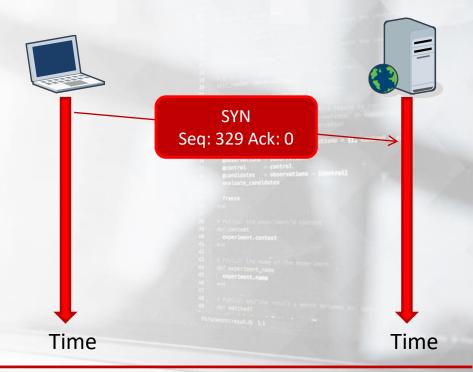








During the first step, the client sends a TCP packet to the server with the SYN flag enabled and a random sequence number.



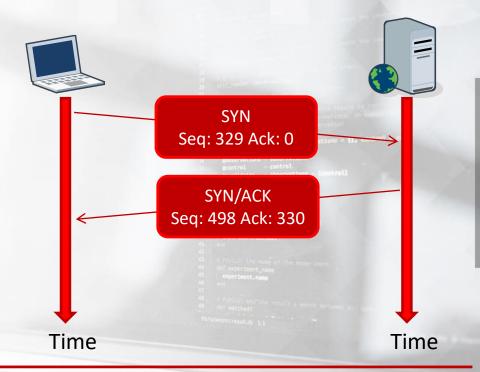








In the second step, the server replies by sending a packet with both the SYN and ACK flag set and another random sequence number.

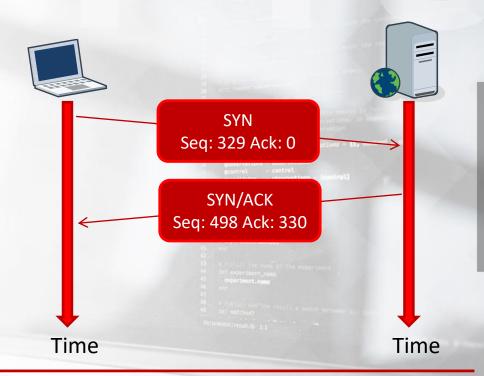








The ACK number is always a simple increment of the SYN number sent by the client.

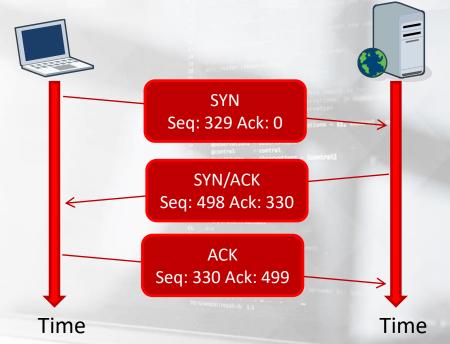


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7

Finally, the client completes the synchronization by sending an ACK packet.

Note that the client behaves just like the server when sending ACK packets.









2.5.6 References

For additional information, please check out the links below:

- IP Layer Network Administration with Linux.
- TCP/IP Tutorial and Technical Overview.
- Packet Analysis Reference Guide v3.0.









^{*}To access the above attachments, go to the course in your members area and click the resources drop-down in the appropriate module line.

Firewalls and Network Defense









2.6 Firewalls and Network Defense

There are many different appliances on the market that a system administrator can use to protect the network.

These devices use different techniques and work on different layers to perform access control, attack detection and prevention.









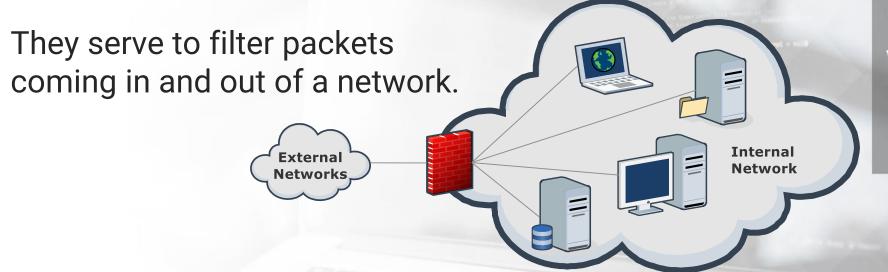
2.6 Firewalls and Network Defense

How does this support my pentesting career?

- Evading firewalls
- Advanced stealth scanning
- Filtering evasion

2.6.1 Firewalls

Firewalls are specialized software modules running on a computer or a dedicated network device.









2.6.1 Firewalls

Firewalls help system administrators and desktop users to control the access to network resources and services.

A firewall can work on different layers of the OSI model, thus providing different features and protections.









2.6.1 Firewalls

It is imperative that you understand how firewalls work and what kind of threats they prevent.

Many people believe that firewalls and antiviruses are all they need to stay secure, in the following slides you will see why this idea is wrong.









The most basic feature of a firewall is packet filtering.

With packet filtering, an administrator can create rules which will filter packets according to certain characteristics like:



- Destination IP address
- Protocol

- Source port
- Destination Port

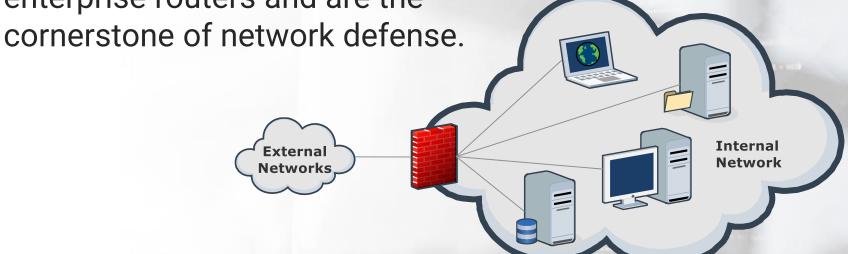








Packet filters run on home DSL routers as well as high-end enterprise routers and are the









Packet filters inspect the header of every packet to choose how to treat the packet. The more common actions are:

- Allow: allow the packet to pass
- Drop: drops the packet without any diagnostic message to the packet source host
- Deny: do not let the packet pass, but notify the source host

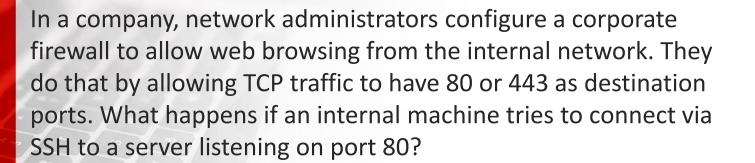








Inspecting the header of a packet does not give you any information on the actual packet content.



The traffic will pass even if it is not web traffic!









Let's now look at a typical example where packet inspection does not stop a hacker from exploiting a server.

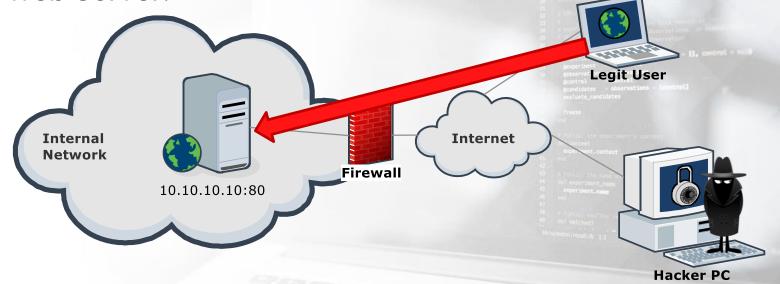








A company hosts a web server. The firewall will allow all incoming traffic from the Internet and direct it to port 80 of the web server.

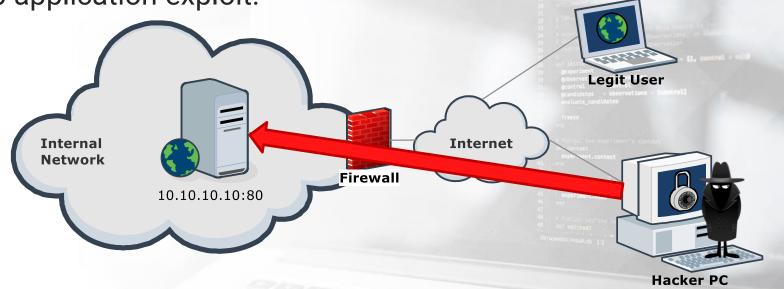








In the same way, application exploits will go through, because the firewall cannot see the difference between web browsing and a web application exploit.

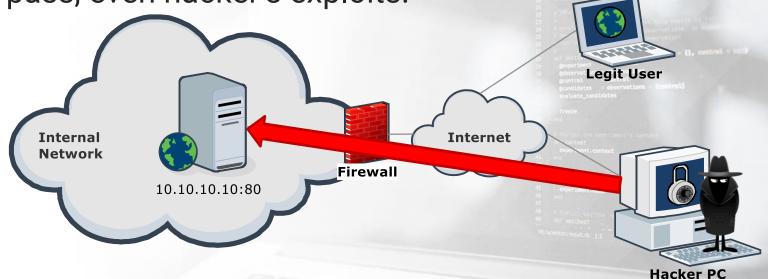








The firewall can only filter traffic by using IP addresses, ports, and protocols. **Any** kind of application layer traffic will pass, even hacker's exploits.









An application layer exploit could be an XSS, a buffer overflow, a SQL injection and much, much more.

Packet filtering is not enough to stop layer 7 attacks.

Let's take a look at another scenario.









A hacker manages to exploit a workstation in a company network and wants to install a Trojan horse to remotely manage that machine.

The Trojan, by default, can be configured to **accept** connections on port 123 TCP or UDP or to **connect back** to the hacker's machine on port 123 TCP or UDP.







A typical firewall configuration rule is to let HTTP traffic (TCP.dst = 80) pass, so internal workstations can browse the Internet. The remaining traffic is dropped.

What happens to the Trojan horse?

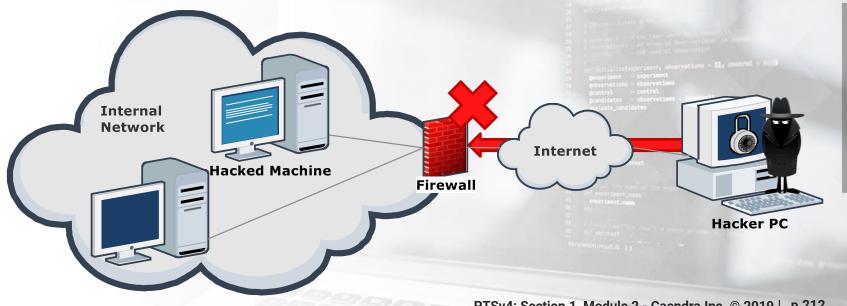








The hacker cannot connect to the infected machine, as there are no rules to allow traffic coming to port 123.

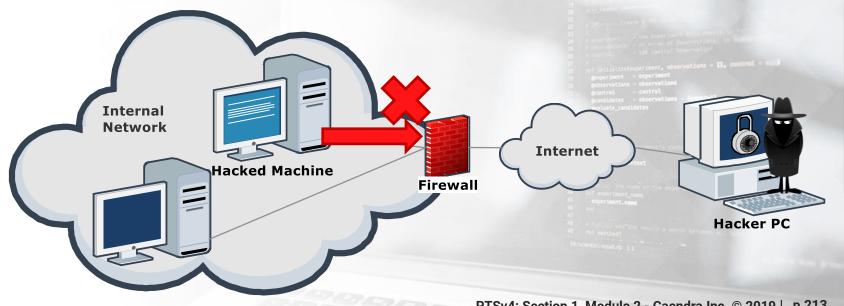








Similarly, the firewall will deny the traffic from the infected machine to the hacker's PC.









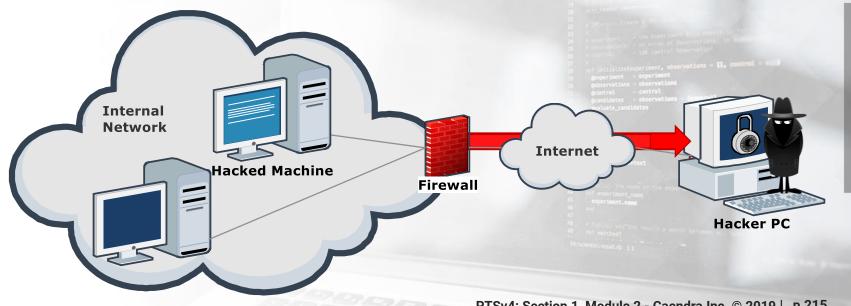
What happens if the hacker, who is smart and knows all the well-known ports by heart, configures the Trojan to **connect back** to his or her machine on port 80?







The firewall will allow the connection!









2.6.3 Application Layer Firewalls

Application level firewalls work by checking all the OSI 7 layers.

They provide a more comprehensive protection because they inspect the actual content of a packet, not just its headers. For example:

- Drop any peer-to-peer application packet.
- Prevent users from visiting a site.



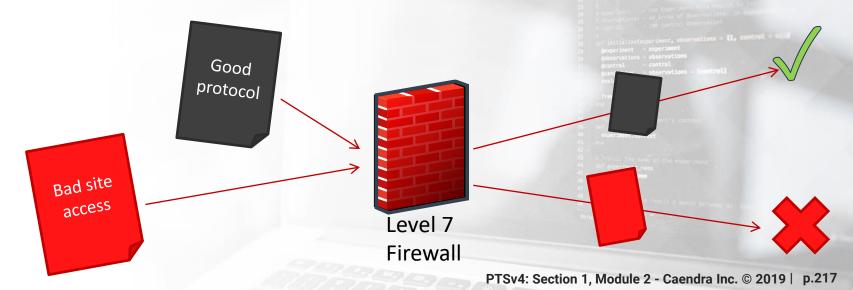






2.6.3 Application Layer Firewalls

Level 7 firewalls are indeed able to understand most of the application layer protocols in use nowadays. Organizations use them not only to protect their network from hackers but also to filter unwanted traffic.



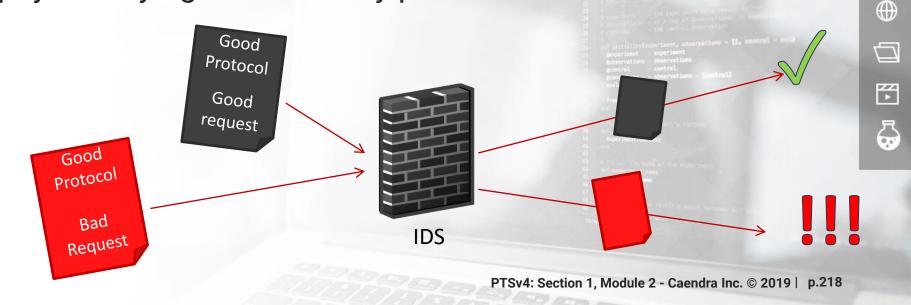








There is not just traffic detection, but intrusion detection! **Intrusion Detection Systems** (IDS) inspect the application payload trying to detect any potential attack.



An IDS is specialized software used for **detecting ongoing intrusions**. It checks for attack vectors like ping sweeps, port scans, SQL injections, buffer overflows and so on.

IDS can also **identify traffic** generated by a virus or a worm. Pretty much every kind of network threat can be detected by a well-configured IDS.







An IDS, like an antivirus, detects risky traffic by means of signatures. The vendor provides frequent signature updates as soon as new attack vectors are found in the wild. Without the right signatures an IDS cannot detect and report an intrusion; the IDS cannot detect something if it does not already know.

There are also **false positives**. They happen when legit traffic is flagged as malicious.









Detection is performed by a multitude of **sensors**, software components that inspect network traffic.

Sensors passively intercept intrusions and communicate them to the **IDS manager**, software in charge of maintaining policies and which provides a management console to the system administrator.









IDSs do not substitute firewalls.

They support firewalls by providing a further layer of security protecting the network from mainstream and well-known attack vectors.







IDSs fall into two main categories:

Network Intrusion
Detection
Systems (NIDS)

Host Intrusion
Detection
Systems (HIDS)









2.6.4.1 NIDS

Network intrusion detection systems inspect network traffic by means of sensors which are usually placed on a router or in a network with a high intrusion risk, like a DMZ.







2.6.4.2 HIDS

On the other hand, host IDS sensors monitor application logs, file-system changes and changes to the operating system configuration.

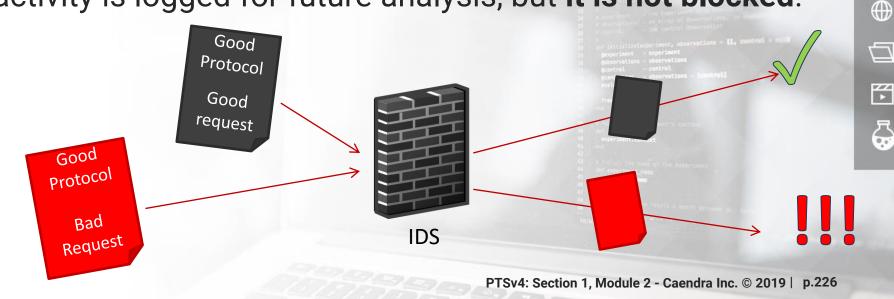






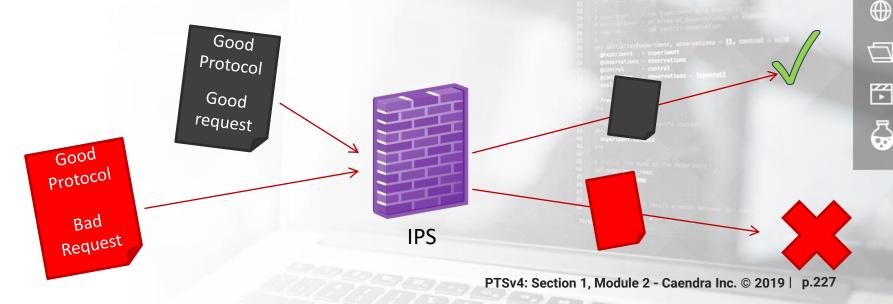
2.6.5 IPS

IDSs, unlike firewalls, can detect suspicious activities and report them to the network administrator. Suspicious activity is logged for future analysis, but **it is not blocked**.



2.6.5 IPS

Intrusion Prevention Systems (IPS) can drop malicious requests when the threat has a risk classification above a pre-defined threshold.



During penetration testing activities, you might want to identify if a firewall-like mechanism is used in the environment.

If you suspect presence of a firewall, you might want to check for anomalies in TCP Three-Way Handshake that was introduced previously in this module.

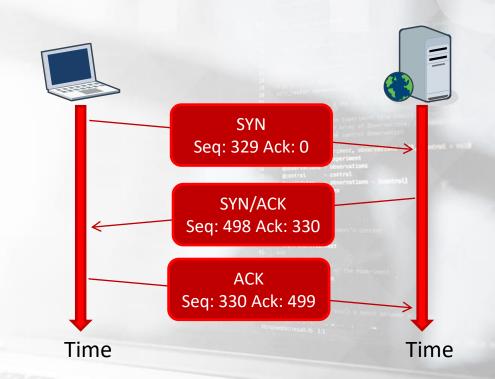


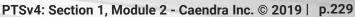






To recall, proper
Three-way handshake
looks like:





 \Box

X

When a firewall is in place, the following behavior may be spotted:

TCP SYN are sent, but there no TCP SYN/ACK replies

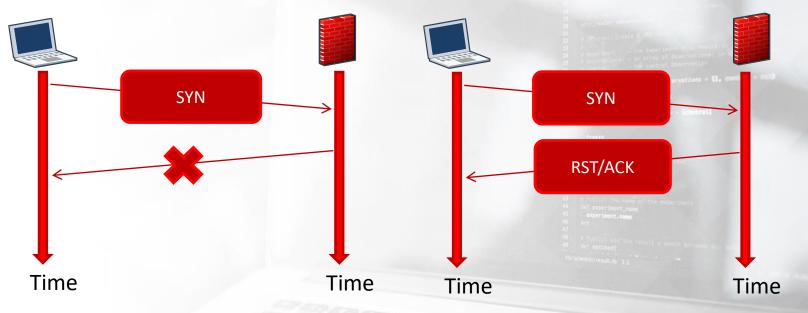
 TCP SYN packets are sent but a TCP RST/ACK reply is received



















Note, that this type of observation does not determine whether the detected obstacle is a **firewall**, an **IDS**, or **any other** device; this just helps you to identify **environmental constraints**.









Firewalls not only filter packets but can also be used to implement **Network Address Translation** or **NAT**.

Note: Your home router is most probably running NAT protocol to connect all your home devices to the internet without having to have a public IP assigned for each of them.

What is NAT and why is it needed?









As you know, every machine on the Internet must have a unique IP address. This does not mean that every device that can **access** the internet must have a unique **public** IP address.

Network Address Translation (NAT) and IP masquerading are two techniques used to provide access to a network from another network.

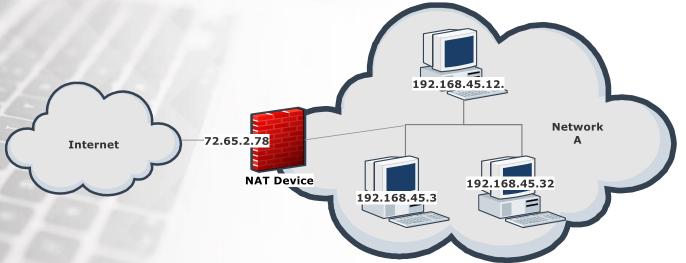








Network A can be a private network using a NAT device to access the Internet. A machine on the internet cannot directly access a machine in Network A.



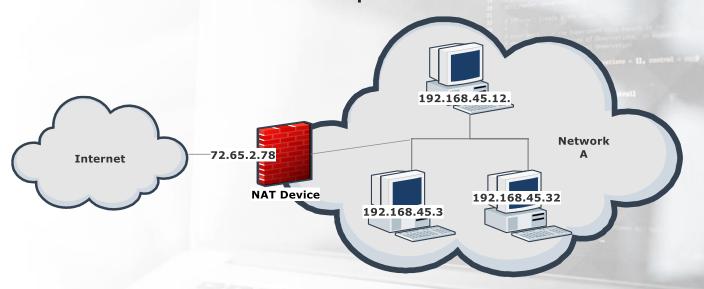








But, a machine in *Network A* can access the Internet, if the NAT device allows the traffic to pass.









Every machine inside *Network A* will use the NAT device as **its default gateway**, thus routing its Internet traffic through it. The NAT device then rewrites the **source IP address** of every packet setting it to 72.65.2.78 (in our example), thus masquerading the original client's IP address.

A machine on the Internet will never know the original client's IP address.









2.6.8 Hera Lab - Find the Secret Server

It is now practice time!

In the following lab, you will learn how security through obscurity cannot stop a skilled attacker. Try to solve the challenge by yourself; if you get stuck, you can check the solutions in the member's area.







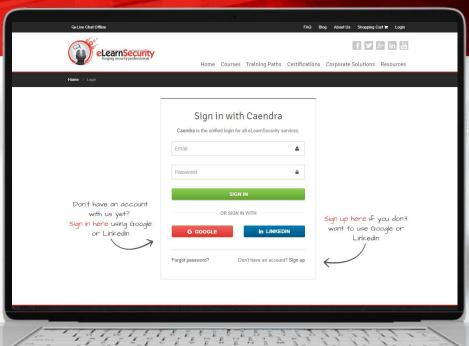


2.6.8 Hera Lab – Find the Secret Server

Find the Secret Server

In this lab, you will:

- Understand routing information.
- Explore different web servers.
- Connect to an apparently unreachable server!



*Labs are only available in Full or Elite Editions of the course. To upgrade, click <u>HERE</u>. To access, go to the course in your members area and click the labs drop-down in the appropriate module line or to the virtual labs tabs on the left navigation.

2.6.9 Resources

- <u>Netfilter</u> the Linux kernel packet filtering framework.
- Book: Firewall Fundamentals the essential guide to understanding and using firewalls to protect personal computers and your network.
- OSSEC IDS an Open Source Host-based Intrusion Detection System.
- IDS FAQ answers to simple questions related to detecting intruders who attack systems through the network.

















2.7 DNS



- SSL/TLS certificates validation relies on DNS
- Mounting spoofing attacks
- Performing information gathering

1800000 Q 1 4) 零 图 0

2.7 DNS

The **Domain Name System**, or **DNS**, is the only application layer protocol you will see in this module.

The DNS primarily converts human-readable names, like www.elearnsecurity.com, to IP addresses and is a fundamental support protocol for the Internet and computer networks in general. It is widely recognized that the entire internet security is relying upon DNS.









2.7 DNS

For this course, you will need to know how the DNS service provides name resolution because every common operation on the Internet such as opening a web site, sending an email, and sharing a document involves the use of a DNS to resolve resource names to IP addresses (and vice versa).









A DNS name such as www.elearnsecurity.com or members.elearnsecurity.com can be broken down into the following parts:

- Top level domain (TLD)
- Domain part
- Subdomain part (if applicable)
- Host part



















These parts form a hierarchy:

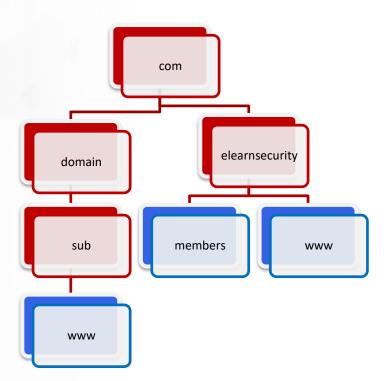
Top Level Domain **Domain Name** Subdomain Name **Host Name**







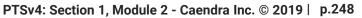
So, we can rewrite the previous names as:



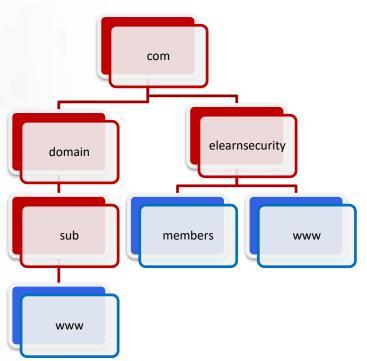








Where the blue squares are the hosts and the red ones are the top/sub/domain names.











Name resolution is performed by **resolvers**, servers that contact the top level domain (TLD) DNS servers and follow the hierarchy of the DNS name to resolve the name of a host.

Resolvers are DNS servers provided by your ISP or publicly available like OpenDNS or Google DNS.









2.7.2 DNS Names Resolution

To convert a DNS name into an IP address, the operating system must contact a **resolver** server to perform the DNS resolution.

The resolver breaks down the DNS name in its parts and uses them to convert a DNS name into an IP address.









2.7.2.1 DNS Resolution Algorithm

- 1 Firstly, the resolver contacts one of the **root name servers**; these servers contain information about the top level domains.
- Then, it asks the TLD name server what's the name server that can give information (authoritative name server) about the **domain** the resolver is looking for.
- If there are one or more **subdomains**, step 2 is performed again on the authoritative DNS server for every subdomain
- 4 Finally, the resolver asks for the name resolution of the **host** part.









A computer needs to open a web page on www.example.com.

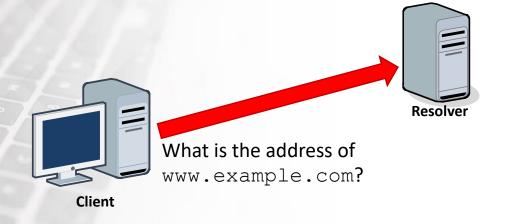








To do that, it contacts the resolver configured by the local administrator (e.g., OpenDNS).





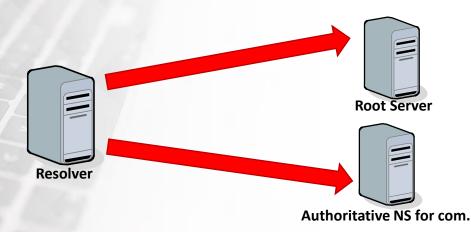






The resolver contacts a **root server** and asks about the authoritative name server(s) for the com. domain.

Then the resolver contacts that authoritative name server and asks what is an authoritative name server for the example.com. domain.











The resolver then asks what the address of www is.

Finally, the resolver sends the IP address back to the client.









2.7.3 Resolvers and Root Servers

How can a resolver know how to contact a **root name** server?

IP addresses of the root servers are **hardcoded in the configuration** of the resolver. System administrators keep the list updated, otherwise, the resolver would not be able to contact a root server!









2.7.4 Reverse DNS Resolution

The domain name system can also perform the inverse operation; it can convert an **IP address to a DNS name**.

Keep in mind that this is not always the case; the administrator of a domain must have enabled and configured this feature for the domain to make it work.









2.7.4 Reverse DNS Resolution

Many tools use the reverse DNS if it's available.

The Linux ping utility performs a reverse DNS query after receiving every response from the target.

```
$ ping www.yahoo.com
PING fd-fp3.wgl.b.yahoo.com (46.228.47.115) 56(84) bytes of data.
64 bytes from ir1.fp.vip.ir2.yahoo.com (46.228.47.115): icmp_req=1 ttl=49 time=125 ms

--- fd-fp3.wgl.b.yahoo.com ping statistics ---
2 packets transmitted, 1 received, 50% packet loss, time 1001ms
rtt min/avg/max/mdev = 125.706/125.706/0.000 ms
```







2.7.5 More about the DNS

In this section, you have seen how the DNS is used to perform name resolution. The domain name system is not just that, it is used to identify what the mail servers for a domain are, to know what is the right server for a specific role and much more.

Please refer to RFC1034 and RFC1035 to dig more into DNS!

NOTE: the DNS is also very important to the security of the whole internet because breaking DNS security means breaking SSL and TLS. This, however, is beyond the scope of this training course.











Wireshark







The best way to deeply understand the topics of this module is to see the actual protocols in action. You can do that by using a sniffer tool.

This section will enhance your Networking and Wireshark skills.



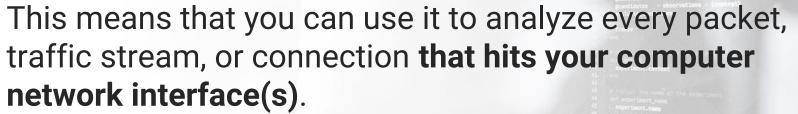








As you know, Wireshark is a network sniffer and protocol analyzer.











Knowing this tool is extremely important to understand how networking works.

Wireshark is widely used by network administrators, networking protocol researchers, and hackers.









Wireshark can capture all the traffic **seen** by the network card of the computer running it.

To understand what traffic a network card sees, you have to know that most network cards, also known as Network Interface Cards (NIC), can work in **promiscuous or monitor mode**.









2.8.1 NIC Promiscuous Mode

During normal operations, a network card **discards** any packet addressed to another NIC. In **promiscuous mode**, a network card will accept and process **any** packet it receives.

For example, in a hub-based network, a NIC will receive traffic addressed to other machines. The NIC usually drops these packets but accepts them while in promiscuous mode.









2.8.1 NIC Promiscuous Mode

With the introduction of switched networks, sniffing other machines Ethernet traffic got harder. You have to perform an attack such as ARP poisoning or MAC flooding in order to do that.

WiFi medium (the air), instead, is broadcast by nature, so it's possible to still detect traffic destined to a different host. In this chapter, we will concentrate on Ethernet traffic only.









Wireshark is free software that can run on practically all modern operating systems. You can download it from https://www.wireshark.org/ or https://www.wireshark.org/download/ for older versions.

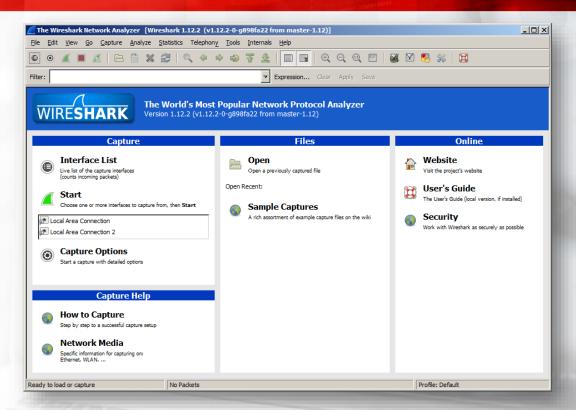
In the following slides, we are going to see how to configure it to use its main features. The version covered is 1.12.2.







Here we see Wireshark's main window.









Clicking on Interface List opens a window with a list of your network cards (wired, wireless, VPNs, virtual interfaces, etc.).









Clicking on Capture Options opens...



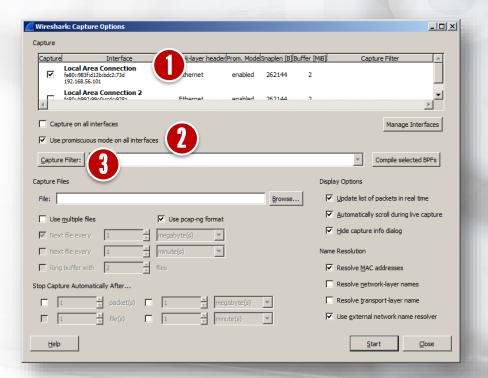






... the capture options window, which has a huge impact on your capture session as you can configure:

- Which interfaces to use during the capture
- NIC promiscuous mode
- Capture filtering



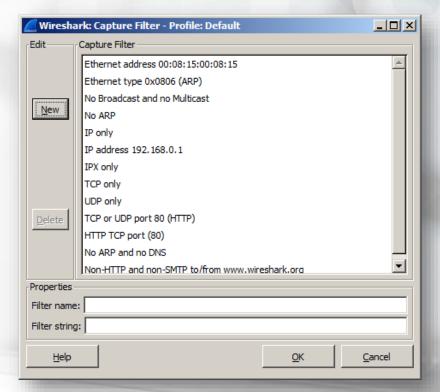






Capture filters will make
Wireshark discard packets
that do not match the filter.
These filters impact how
many packets your computer
must process and how big
the capture file will be.

This is very useful to limit captured traffic in high traffic networks.

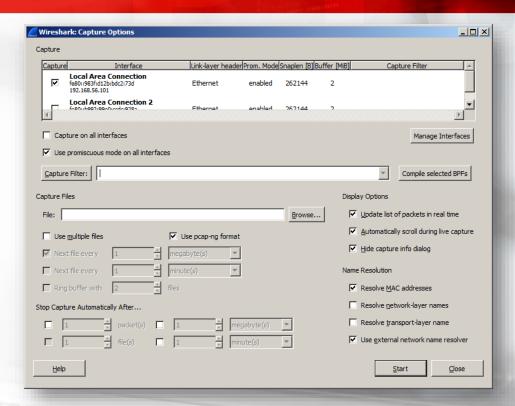








During your first captures, leave the filter blank and just click **start**.









To perform these very same operations, you can just select the capture interface and then click on **capture options** or **start** from the main window.



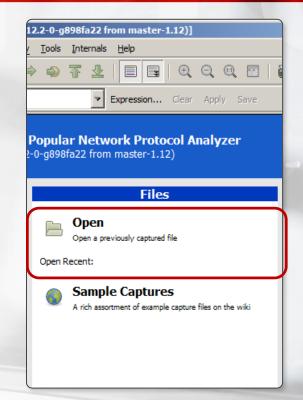






PCAP files store an entire capture (from a previous capture session).

If you already have a PCAP file, you can open it using this button.

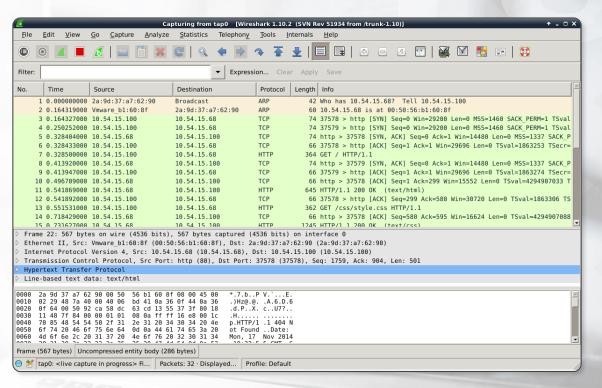








In either case, doing a live capture or opening a previous one, you will see this interface.



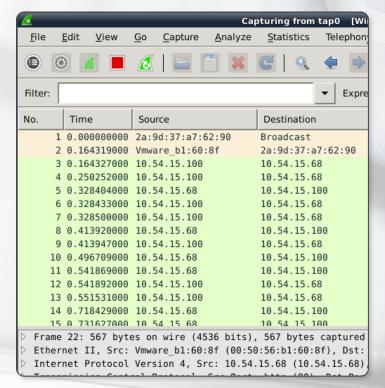






The first two columns of the upper pane contain:

- The **number** of the captured packet
- The arrival time of the packet in seconds.
 The arrival time is relative to the start of the capture.



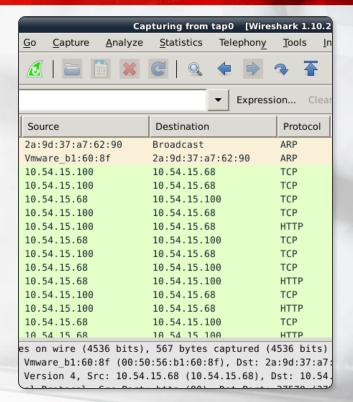






You can then see the **source**, **destination** and **protocol** columns.

Note how the source and destination address vary according to the protocol.





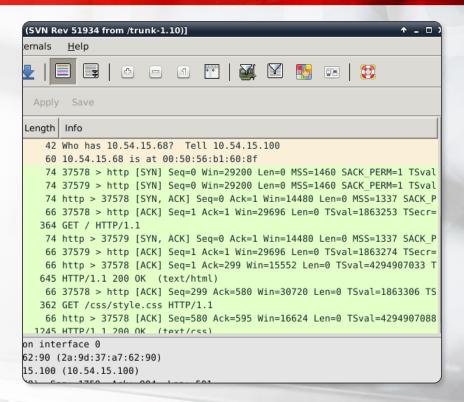






In the last two columns, you can find the **size of the packet** and some related **information**.

The *info* column is protocol specific.



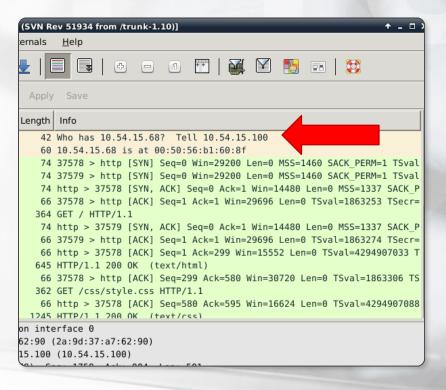








For example, the first two packets are ARP requests and replies.



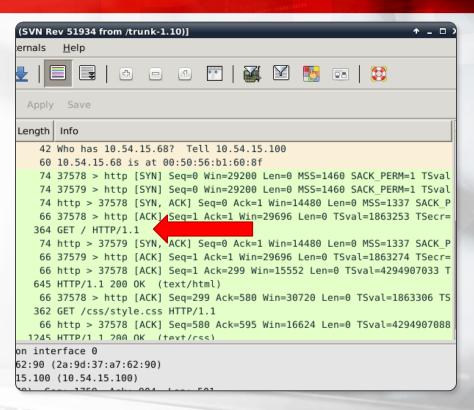








This packet is an HTTP request.









The center pane gives you access to all the protocol layers used by a packet.

This actually allows you to read the entire packet layer by layer!

```
Frame 22: 567 bytes on wire (4536 bits), 567 bytes captured (4536 bits) on interface 0
```







Ethernet II, Src: Vmware_b1:60:8f (00:50:56:b1:60:8f), Dst: 2a:9d:37:a7:62:90 (2a:9d:37:a7:62:90)

> Internet Protocol Version 4, Src: 10.54.15.68 (10.54.15.68), Dst: 10.54.15.100 (10.54.15.100)

Transmission Control Protocol, Src Port: http (80), Dst Port: 37578 (37578), Seq: 1759, Ack: 904, Len: 501

Hypertext Transfer Protocol

You can drill down to get any information you want about a packet.

For example, this packet has the *ACK* TCP flag on.

```
Frame 6: 66 bytes on wire (528 bits), 66 bytes captured (528 bits) on interface 0
Ethernet II, Src: 2a:9d:37:a7:62:90 (2a:9d:37:a7:62:90), Dst: Vmware b1:60:8f (00:50:56:b1:60:8f)
Internet Protocol Version 4, Src: 10.54.15.100 (10.54.15.100), Dst: 10.54.15.68 (10.54.15.68)
Transmission Control Protocol, Src Port: 37578 (37578), Dst Port: http (80), Seq: 1, Ack: 1, Len: 0
  Source port: 37578 (37578)
  Destination port: http (80)
  [Stream index: 0]
  Seguence number: 1 (relative seguence number)
  Acknowledgment number: 1
                              (relative ack number)
  Header length: 32 bytes

¬ Flags: 0x010 (ACK)

    000. .... = Reserved: Not set
     ...0 .... = Nonce: Not set
     .... 0... = Congestion Window Reduced (CWR): Not set
     .... .0.. .... = ECN-Echo: Not set
     .... ..0. .... = Urgent: Not set
     .... ...1 .... = Acknowledgment: Set
     .... .... 0... = Push: Not set
     .... .... .0.. = Reset: Not set
     .... .... ..0. = Syn: Not set
     .... Not set
  Window size value: 29
  [Calculated window size: 29696]
  [Window size scaling factor: 1024]

    ○ Checksum: 0x30b9 [validation disabled]

Doptions: (12 bytes), No-Operation (NOP), No-Operation (NOP), Timestamps

    [SEQ/ACK analysis]
```









In the bottom pane, you can see the actual packet payload. In this example we see an HTTP GET request.

```
00 50 56 b1 60 8f 2a 9d
                              37 a7 62 90 08 00 45 00
                                                         .PV.`.*. 7.b...E.
      01 5e 5d b7 40 00 40 06 a8 cf 0a 36 0f 64 0a 36
                                                         .^].@.@. ...6.d.6
     0f 44 92 ca 00 50 13 55 33 b8 58 dc 5c ef 80 18
                                                         .D...P.U 3.X.\...
      00 1d 43 de 00 00 01 01 08 0a 00 1c 6e 55 ff ff
                                                         ..C.... ....nU...
0040
     14 6f 47 45 54 20 2f 20 48 54 54 50 2f 31 2e 31
                                                         .oGET / HTTP/1.1
      0d 0a 48 6f 73 74 3a 20 31 30 2e 35 34 2e 31 35
                                                         ..Host: 10.54.15
                                                         .68..Use r-Agent:
      2e 36 38 0d 0a 55 73 65 72 2d 41 67 65 6e 74 3a
0070
     20 4d 6f 7a 69 6c 6c 61 2f 35 2e 30 20 28 58 31
                                                         Mozilla /5.0 (X1
0080
     31 3b 20 4c 69 6e 75 78
                              20 78 38 36 5f 36 34 3b
                                                         1; Linux x86 64;
                                                         rv:31.0 ) Gecko/
      20 72 76 3a 33 31 2e 30 29 20 47 65 63 6b 6f 2f
      32 30 31 30 30 31 30 31 20 46 69 72 65 66 6f 78
00a0
                                                         20100101 Firefox
     2f 33 31 2e 30 20 49 63 65 77 65 61 73 65 6c 2f
                                                        /31.0 Ic eweasel/
      33 31 2e 32 2e 30 0d 0a
                              41 63 63 65 70 74 3a 20
                                                         31.2.0.. Accept:
00d0
      74 65 78 74 2f 68 74 6d 6c 2c 61 70 70 6c 69 63
                                                        text/htm l,applic
00e0
      61 74 69 6f 6e 2f 78 68 74 6d 6c 2b 78 6d 6c 2c
                                                        ation/xh tml+xml,
      61 70 70 6c 69 63 61 74 69 6f 6e 2f 78 6d 6c 3b
                                                         applicat ion/xml;
0100
      71 3d 30 2e 39 2c 2a 2f
                              2a 3b 71 3d 30 2e 38 0d
                                                        q=0.9,*/*;q=0.8.
      0a 41 63 63 65 70 74 2d 4c 61 6e 67 75 61 67 65
                                                         .Accept- Language
      3a 20 65 6e 2d 55 53 2c 65 6e 3b 71 3d 30 2e 35
                                                         : en-US, en;q=0.5
      0d 0a 41 63 63 65 70 74 2d 45 6e 63 6f 64 69 6e
                                                         ..Accept -Encodin
      67 3a 20 67 7a 69 70 2c 20 64 65 66 6c 61 74 65
                                                         q: gzip, deflate
      0d 0a 43 6f 6e 6e 65 63 74 69 6f 6e 3a 20 6b 65
                                                         ..Connec tion: ke
     65 70 2d 61 6c 69 76 65 0d 0a 0d 0a
                                                         ep-alive ....
```









2.8.4 Filtering

A traffic capture can be overwhelming, even on a network with just a couple of dozens of nodes.

Wireshark can **filter** traffic at **capture** or at **display** time. Each method has its own pros and cons.



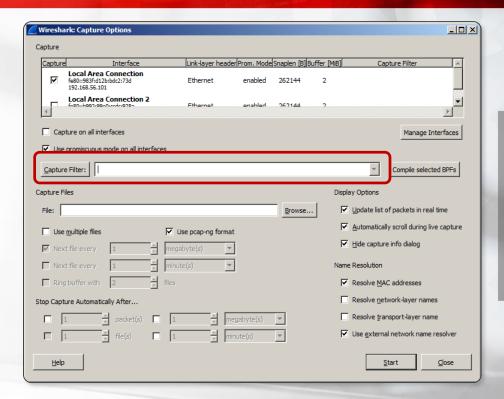






2.8.4.1 Capture Filters

You can set capture filters **before** starting the capture so that Wireshark will capture only packets matching the filters.









2.8.4.1 Capture Filters

Here are some basic capture filters.

Syntax	Description
ip	Only packets using IP as layer 3 protocol.
not ip	The opposite of the previous syntax.
tcp port 80	Packets where the source or destination TCP port is 80.
net 192.168.54.0/24	Packets from and to the specified network.
src port 1234	The source port must be 1234; the transport protocol does not matter.
src net 192.168.1.0/24	The source IP address must be in the specified network.
host 192.168.45.65	All the packets from or to the specified host.
host www.examplehost.com	All the packets from or to the specified hostname.









2.8.4.1 Capture Filters

Capture filters will downsize the amount of traffic gathered.

The final capture will be **smaller**, and it will contain **only** the **needed traffic**.









However, capture filters might not catch interesting traffic! Display filters instead allow you to inspect and apply very granular filters to every field of the captured packets. Wireshark then displays only the packets matching the filters.

You can always remove or fine tune a display filter, something you can't do with the capture filter (you would have to re-start the capture from scratch).

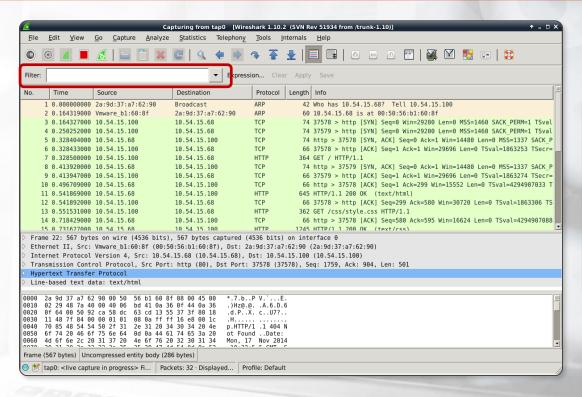








You can use the filter textbox to apply a display filter.

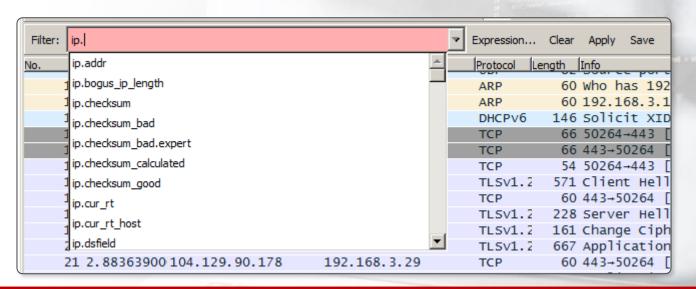








You can start typing a filter and Wireshark will give you valid protocol fields.

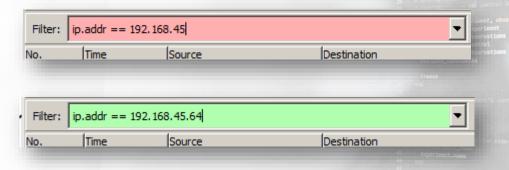








The background of the text-box will turn red if the filter is invalid or green when the filter is valid.











A display filter is made by:

Syntax	Description
<protocolname></protocolname>	Displays any packet using that protocol.
<protocolname>[.field]</protocolname>	Displays any packet with the specified field present.
<protocolname>[.field] [operand value]</protocolname>	Displays any packet whose protocol field matches the operand and value.
<protocolname>[.field] AND <protocolname>[.field] [operand value]</protocolname></protocolname>	You can combine multiple expressions by using logical operators.









Below is an example.

Syntax	Description
ip	Displays IP packets.
ip.addr	Displays IP packets with a populated source or destination address.
ip.addr == 192.168.12.13	Displays IP packets with 192.168.12.13 as source or destination address.
ip.addr == 192.168.12.13 or arp	The above or ARP packets.









You can find Wireshark display filter reference here.

For any other information, please refer to the Wireshark User's Guide.









2.8.5 Video - Using Wireshark

Using Wireshark

In this video, you will see how to configure and use Wireshark, how capture and display filters work and this powerful tool's main features.



*Videos are only available in Full or Elite Editions of the course. To upgrade, click <u>HERE</u>. To access, go to the course in your members area and click the resources drop-down in the appropriate module line.

2.8.6 Video – Full Stack Analysis with Wireshark

Full Stack Analysis with Wireshark

This video uses the traffic captured in the previous video to perform traffic analysis, which will improve your understanding of low-level interaction between routing, IP, ARP, TCP, application-level protocols and more!



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2.8.7 Sample Traffic Captures

If you want to practice these topics a little more, you can record some traffic from your computer or you can download a capture from Wireshark website.









2.8.8 Lab - Data Exfiltration

Apply your fresh knowledge about networking to the Data Exfiltration lab.

You can test it against a environment secured by a firewall, while on the other hand, learn how to bypass this kind of obstacles.







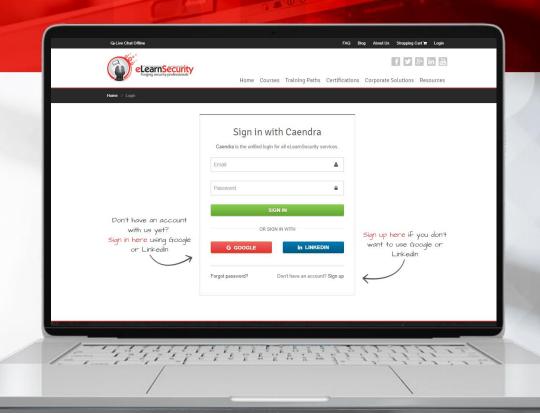


2.8.8 Lab – Data Exfiltration

Data Exfiltration

The environment is secured by a firewall.

Can you bypass it?



*Labs are only available in Full or Elite Editions of the course. To upgrade, click <u>HERE</u>. To access, go to the course in your members area and click the labs drop-down in the appropriate module line or to the virtual labs tabs on the left navigation.

PTSv4: Section 1, Module 2 - Caendra Inc. © 2019 | p.301

















Packet Analysis Reference Guide v3.0

You can download it from the **Resources** drop-down menu of the Networking module (Section Preliminary Skills – Prerequisites)

RFC 5735 - Special Use IPv4 Addresses

http://tools.ietf.org/html/rfc5735

Classful Subnet Calculator

http://www.subnet-calculator.com/

Service Name and Transport Protocol Port Number Registry

http://www.iana.org/assignments/service-names-port-numbers/service-names-port-numbers.xhtml











ISO/OSI Model

http://support.microsoft.com/kb/103884

Variable Length Subnet Table For IPv4

http://tools.ietf.org/html/rfc1878

CIDR Subnet Calculator

http://www.subnet-calculator.com/cidr.php

TCPView

http://technet.microsoft.com/en-us/sysinternals/bb897437

















TACO IPv6 Router - a Case Study in Protocol Processor Design

https://www.researchgate.net/publication/31596630_TACO_IPv6_Router_-a_Case_Study_in_Protocol_Processor_Design

Internet Protocol Version 6 (IPv6) Addressing Architecture

https://tools.ietf.org/html/rfc3513

<u>IPv6 Explained for Beginners</u>

http://www.steves-internet-guide.com/ipv6-guide/

How to find IPv6 Prefix

https://networklessons.com/ipv6/how-to-find-ipv6-prefix/

















IP Layer Network Administration with Linux

You can download it from the **Resources** drop-down menu of the Networking module (Section Preliminary Skills – Prerequisites)

The TCP/IP Guide: A Comprehensive, Illustrated Internet Protocols Reference

http://www.amazon.com/TCP-Guide-Comprehensive-Illustrated-Protocols/dp/159327047X/ref=sr_1_3?s=books&ie=UTF8&qid=1297880998&sr=1-3

Firewall Fundamentals

http://www.amazon.com/Firewall-Fundamentals-Wes-Noonan/dp/1587052210/ref=cm_cr_pr_sims_t

Wireshark Display Filter Reference

https://www.wireshark.org/docs/dfref/

















TCP/IP Tutorial and Technical Overview

You can download it from the **Resources** drop-down menu of the Networking module (Section Preliminary Skills – Prerequisites)

IPTables tutorial

https://www.frozentux.net/iptables-tutorial/iptables-tutorial.html

Wireshark

https://www.wireshark.org/

Wireshark User's Guide

https://www.wireshark.org/docs/wsug_html_chunked/



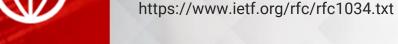














http://wiki.wireshark.org/SampleCaptures



https://www.ietf.org/rfc/rfc1035.txt

IPv6 Calculator

https://www.ultratools.com/tools/ipv6CIDRToRange

















Netfilter

http://www.netfilter.org/

Ossec IDS

http://www.ossec.net/

FAQ: Network Intrusion Detection Systems

http://www.linuxsecurity.com/resource_files/intrusion_detection/network-intrusion-detection.html











Videos



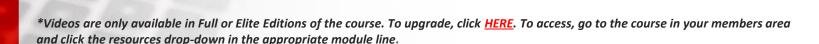
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Labs

Find the Secret Server

Use your TCP/IP routing knowledge to access a secret web server.









Data Exfiltration

The environment is secured by a firewall. Can you bypass it?

