

## Matching Pairs Game on Pt-51

1. [20 points] In this project, you will be writing a game program for a memory matching game. It is similar to the Concentration card game [https://en.wikipedia.org/wiki/Concentration\\_\(card\\_game\)](https://en.wikipedia.org/wiki/Concentration_(card_game)). The player will type on a keyboard connected to the Pt-51 board via UART.

- When the game starts, the LCD displays the following message on its two lines, where XX is either blank or a positive integer.

```
*****  0
*****  L:XX
```

- The 0 on the first line indicates the number of moves the player has made in a game. It should increment after every move.
- The L:XX on the second line will be used to show the lowest number of moves used to finish the game so far. In the beginning, XX will be blank. After a player finishes the first game, XX will have the number of moves used by the player to finish the game. After two or more games have been finished by the player, XX has the minimum number of moves used by the player to complete a game.
- The 16 \* characters hide two 1s, two 2s, two 3s, two 4s, two 5s, two 6s, two 7s, two 8s.
- The goal of the player is to uncover all the numbers while matching them with already uncovered numbers, using as few moves as possible. At the end of a game, the LCD might look like the following:

```
15274623  26
64185783  L:
```

- A player uncovers a covered position by pressing an alphabet key. The mapping from LCD positions to the alphabet key is shown below.

```
qwertyui
asdfghjk
```

Note that the keys are taken from the first two rows of a QWERTY keyboard.

- If the player presses a key corresponding to an already uncovered position, this key press should be ignored. The number of moves is not incremented.
- If the player presses a key corresponding to a covered position, the number at the position is revealed on the LCD. The number of moves is incremented on the first line of the LCD.
  - If the newly uncovered number is the first one to be uncovered, then it remains uncovered.
  - If the newly uncovered number matches an existing uncovered number, then the newly uncovered number remains uncovered.
  - If the newly uncovered number does not match an existing uncovered number, then the newly uncovered number is replaced by \* after a 3 second delay.

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- When all the positions have been uncovered, the current game ends and a new game begins. The number of moves is used to update the **XX** on the second line.
  - To make the game interesting, the locations of the hidden numbers need to change from game to game. We associate a unique integer from 0 to 15 with each of the 16 positions as shown below.

0	1	2	3	4	5	6	7
8	9	10	11	12	13	14	15

- We can use a linear feedback shift register (LFSR) to generate pseudorandomness. Consider a LFSR with 4 registers (each containing a bit). We can use a single nibble to store the current state of the LFSR. If the current state of the 8 bits is  $(b_3, b_2, b_1, b_0)$ , the next state is

$$(b_3 \oplus b_0, b_3, b_2, b_1),$$

where  $\oplus$  represents bitwise XOR. With this next-state rule, the LFSR will cycle through all the non-zero 4-bit states with a period of 15. This is an example of a maximal-length LFSR [https://en.wikipedia.org/wiki/Linear-feedback\\_shift\\_register](https://en.wikipedia.org/wiki/Linear-feedback_shift_register).

- Initialize the LFSR state  $(b_3, b_2, b_1, b_0)$  to a non-zero value (this is done only once when the program loads for the first time).
  - Let  $B_0$  be the decimal integer corresponding to  $b_3b_2b_1b_0$ .
  - Place the first 1 at the location  $B_0$ . For example, if  $b_3b_2b_1b_0 = 0101$  then place the 1 at location 5 (the sixth location on the first line of the LCD). If  $b_3b_2b_1b_0 = 1010$  then place the 1 at location 12 (the fifth location on the second line of the LCD).
  - Calculate the next state of the LFSR. Let the next state be  $b'_3b'_2b'_1b'_0$ .
  - Let  $B_1$  be the decimal integer corresponding to  $b'_3b'_2b'_1b'_0$ .
  - Place the second 1 at the location  $B_0 + B_1 \bmod 16$ . As  $B_1$  is an LFSR state, it is non-zero. Hence the second 1 will be placed in a location which is different from the first 1.
  - Calculate the next state of the LFSR. Let  $B_2$  be the decimal integer corresponding to this state.
  - Place the first 2 at the location  $B_0 + B_2 \bmod 16$ .
  - Calculate the next state of the LFSR. Let  $B_3$  be the decimal integer corresponding to this state.
  - Place the second 2 at the location  $B_0 + B_3 \bmod 16$ .
  - And so on, until the second 8 is placed.
  - The characters are placed at  $B_0$  and  $B_0 + B_i \bmod 16$  for  $i = 1, 2, \dots, 15$ . As the  $B_i$ 's are distinct integers in the set  $\{1, 2, \dots, 15\}$ , the locations of the characters will be distinct.
- Preserve the last state of the LFSR for the next game. This state will be used as the  $B_0$  in the next game.