Operating Systems

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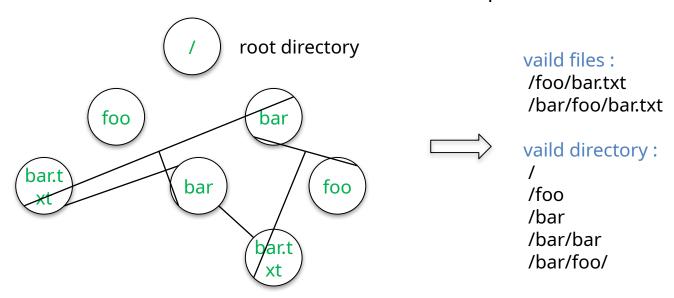




39. File and Directories

Concepts

- File
 - File is simply a linear array of bytes.
 - Each file has low-level name as 'inode number'
- Directory
 - A file
 - A list of <user-readable filename, low-level filename> pairs



An Example Directory Tree

Interface: Creating a file

Use open system call with O_CREAT flag.

```
int fd = open("foo", O_CREAT|O_WRONLY|O_TRUNC, S_IRUSR|S_IWUSR);
```

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- 0_CREAT : create file.
- 0_WRONLY: only write to that file while opened.
- 0_TRUNC: make the file size zero (remove any existing content).
- open system call returns file descriptor.
 - file descriptor is an integer, is used to access files.
 - Ex)read (file descriptor)

```
File descriptor table
struct proc {
...
struct file *ofile[NOFILE]; // Open files
...
};
```

Interface: Reading and Writing Files

An Example of reading and writing 'foo' file.

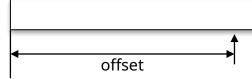
The result of strace to figure out cat is doing.

```
prompt> stace cat foo //strace to figure out what cat is doing ... open("foo", O_RDONLY|O_LARGEFILE) = 3 read(3, "hello\n", 4096) = 6 /*the number of bytes to read*/ write(1, "hello\n", 6) = 6 hello read(3, "hello\n", 4096) = 0 write(1, "hello\n", 6) = 0 ... prompt>
```

- open(): open file for reading with O_RDOLY and O_LARGEFILE flags.
 returns file descriptor 3 (0,1,2, is for standard input/output/error)
- read(): read bytes from the file.
- write(): write buffer to standard output.

Reading and Writing Files (Cont.)

OFFSET



- The position of the file where we start read and write.
- When a file is open, "an offset" is allocated.
- Updated after read/write
- How to read or write to a specific offset within a file?

```
off_t lseek(int fd, off_t offset /*location */, int whence);
```

- Third argument is how the seek is performed.
 - SEEK_SET: to offset bytes.
 - SEEK_CUR: to its current location plus offset bytes.
 - SEEK_END: to the size of the file plus offset bytes.

abstractions

```
struct file {
   int ref;
   char readable;
   char writable;
   struct inode *ip;
   uint off;
};
```

```
struct {
        struct spinlock lock;
        struct file file[NFILE];
} ftable;
```

System Calls	Return Code	Current Offset
<pre>fd = open("file", O_RDONLY);</pre>	3	0
read(fd, buffer, 100);	100	100
read(fd, buffer, 100);	100	200
read(fd, buffer, 100);	100	300
read(fd, buffer, 100);	0	300
close(fd);	0	_

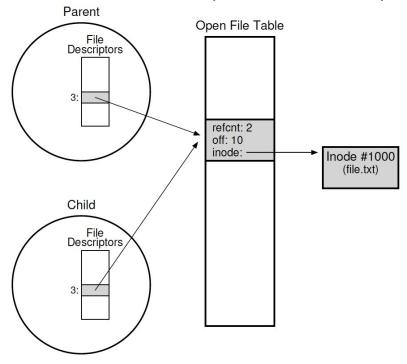
Sample traces

System Calls	Return Code	OFT[10] Current Offset	OFT[11] Current Offset
		Oliset	Oliset
fd1 = open("file", O_RDONLY);	3	0	_
<pre>fd2 = open("file", O_RDONLY);</pre>	4	0	0
read(fd1, buffer1, 100);	100	100	0
read(fd2, buffer2, 100);	100	100	100
close(fd1);	0	_	100
close(fd2);	0	_	_

System Calls	Return Code	Current Offset
<pre>fd = open("file", O_RDONLY);</pre>	3	0
<pre>lseek(fd, 200, SEEK_SET);</pre>	200	200
read(fd, buffer, 50);	50	250
close(fd);	0	_

fork() and dup()

Child process inherits the file descriptor table of the parent.



Duplicating a file descriptor

```
int main(int argc, char *argv[]) {
   int fd = open("README", O_RDONLY);
   assert(fd >= 0);
   int fd2 = dup(fd);
   // now fd and fd2 can be used interchangeably
   return 0;
}
```

fsync()

- Persistency
 - write(): write data to the buffer. Later, save it to the storage.
 - some applications require more than eventual guarantee. Ex) DBMS
- fsync(): the writes are forced immediately to disk.

```
off_t fsync(int fd /*for the file referred to by the specified file*/)
```

An Example of fsync().

```
int fd = open("foo", O_CREAT | O_WRONLY | O_TRUNC);
int rc = write(fd, buffer, size);
rc = fsync(fd);
```

- If a file is created, it needs to be durably a part of the directory.
 - Above code requires fsync() to directory also.

Renaming Files

- rename(): rename a file to different name.
 - It implemented as an atomic call.
 - Ex) change from foo to bar

```
promt > mv foo bar
```

Saving a file in an editor

```
int fd = open("foo.txt.tmp", O_WRONLY|O_CREAT|O_TRUNC);
write(fd, buffer, size); // write out new version of file
fsync(fd);
close(fd);
rename("foo.txt.tmp", "foo.txt");
```

Getting Information About Files

- stat(): Show the File metadata
 - metadata is information about each file, ex: size, permission, ...
- stat structure is below:

```
struct stat {
1
        dev_t st_dev; /* ID of device containing file */
        ino t st ino; /* inode number */
        mode_t st_mode; /* protection */
4
        nlink_t st_nlink; /* number of hard links */
        uid_t st_uid; /* user ID of owner */
        gid_t st_gid; /* group ID of owner */
        dev_t st_rdev; /* device ID (if special file) */
        off_t st_size; /* total size, in bytes */
10
        blksize_t st_blksize; /* blocksize for filesystem I/O */
11
        blkcnt_t st_blocks; /* number of blocks allocated */
12
        time t st atime; /* time of last access */
13
        time_t st_mtime; /* time of last modification */
14
        time t st ctime; /* time of last status change */
15 };
```

Getting Information About Files (Cont.)

- An example of stat()
 - All information is in a inode

```
prompt> echo hello > file
prompt> stat file

File: 'file'
Size: 6 Blocks: 8 IO Block: 4096 regular file
Device: 811h/2065d Inode: 67158084 Links: 1
Access: (0640/-rw-r----) Uid: (30686/ root) Gid: (30686/ remzi)
Access: 2011-05-03 15:50:20.157594748 -0500
Modify: 2011-05-03 15:50:20.157594748 -0500
Change: 2011-05-03 15:50:20.157594748 -0500
```

Removing Files

- The result of strace to figure out what rm is doing.
 - rm is Linux command to remove a file
 - rm calls unlink() to remove a file.

```
1 prompt> strace rm foo
2 ...
3 unlink("foo")
4 ...
5 prompt>
```

Making Directories

- mkdir(): Make a directory
 - When a directory is created, it is empty.
 - Empty directory have two entries: . (itself), ..(parent)

```
prompt> strace mkdir foo
...
mkdir("foo", 0777) = 0
...
prompt>
```

```
1 prompt> ls -al
2 total 8
3 drwxr-x--- 2 roo root 6 Apr 30 16:17 ./
4 drwxr-x--- 26 root root 4096 Apr 30 16:17 ../
```

Reading Directories

- readdir()
 - Directory is a file, but with a specific structure.
 - When reading a directory, we use specific system call other than read().
 - A sample code to read directory entries.

```
int main(int argc, char *argv[]) {
1
     DIR *dp = opendir("."); /* open current directory */
2
     assert(dp != NULL);
3
     struct dirent *d;
4
     while ((d = readdir(dp)) != NULL) { /* read one directory entry
5
   * /
      printf("%d %s\n", (int) d \rightarrow d ino, d \rightarrow d name);
6
7
     }
     closedir(dp); /*close current directory */
8
     return 0;
9
10
    }
```

Reading Directories

Structure of the directory entry

```
struct dirent {
   char d_name[256]; /* filename */
   ino_t d_ino; /* inode number */
   off_t d_off; /* offset to the next dirent */
   unsigned short d_reclen; /* length of this record */
   unsigned char d_type; /* type of file */
};
```

Deleting Directories

- rmdir(): Delete a directory.
 - rmdir()requires directory be empty before it deleted.
 - If you call rmdir() to a non-empty directory, it will fail.

Hard Links

- link(): Link old file and a new file.
 - Create hard link named file2.

```
prompt> echo hello > file
prompt> cat file
hello
prompt> ln file file2 /* create a hard link, link file to file2 */
prompt> cat file2
hello
```

- The result of link()
 - Two files have same inode number, but two human name(file, file2)

```
prompt> ls -i file file2
67158084 file /* inode value is 67158084 */
67158084 file2 /* inode value is 67158084 */
prompt>
```

Hard Links (Cont.)

- How to create hard link file?
 - Step1. Make an inode, track all information about the file.
 - Step2. Link a human-readable name to file.
 - Step3. Put link file into a current directory.
- After creating a hard link to file, old and new files have no difference.
- Thus, to remove a file, we call unlink().

unlink Hard Links

- What unlink() is doing?
 - Check reference count within the inode number.
 - Remove link between human-readable name and inode number.
 - Decrease reference count.
 - When only it reaches zero, It delete a file (free the inode and related blocks)

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unlink Hard Links (Cont.)

The result of unlink()

```
/* create file*/
prompt> echo hello > file
prompt> stat file
... Inode: 67158084 Links: 1 ... /* Link count is 1 */
prompt> ln file file2
                                /* hard link file2 */
prompt> stat file
... Inode: 67158084 Links: 2 ... /* Link count is 2 */
prompt> stat file2
... Inode: 67158084 Links: 2 ... /* Link count is 2 */
prompt> ln file2 file3
                                /* hard link file3 */
prompt> stat file
... Inode: 67158084 Links: 3 ... /* Link count is 3 */
                                   /* remove file */
prompt> rm file
prompt> stat file2
... Inode: 67158084 Links: 2 ... /* Link count is 2 */
prompt> rm file2
                                  /* remove file2 */
prompt> stat file3
... Inode: 67158084 Links: 1 ... /* Link count is 1 */
prompt> rm file3
```

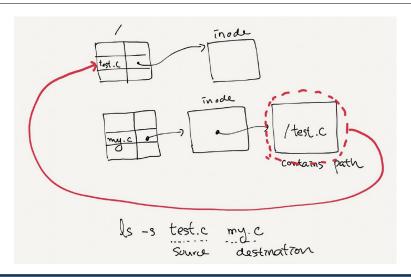
Symbolic Links

Symbolic link

- Special file that contains path to the source directory.
- Hard Link cannot create to a directory.
- Hard Link cannot create to a file to other partition.

An example of symbolic link

```
prompt> echo hello > file
prompt> ln -s file file2 /* option -s : create a symbolic link, */
prompt> cat file2
hello
```



Symbolic Links (Cont.)

Symbolic link is different file type.

Symbolic link is subject to the dangling reference.

```
prompt> echo hello > file
prompt> ln -s file file2
prompt> cat file2
hello
prompt> rm file
prompt> cat file2
cat: file2: No such file or directory
```

Summary

- Create file
- read/write/lseek
- mkdir/readdir
- fsync
- hardlink/softlink