SSN College of Engineering

Department of Information Technology

UIT2201 — Programming and Data Structures

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Exercise — 04

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I. AIM:

The purpose of this exercise is to design and analyze algorithms and perform empirical analysis of algorithms as well.

- 1. Let p(x) be a polynomial of degree n.
- (a) Implement a simple $O(n^2)$ -time algorithm using Python for computing p(x), for a given value of x.
- (b) Implement a $O(n \log n)$ algorithm for computing p(x), based upon a more efficient calculation of x.
- (c) Now, consider rewriting p(x) as $p(x) = a0 + x (a1 + x (a2 + x (a3 + \cdots + x (an-1 + xan) \cdot \cdots)))$ which is known as the Horner's method. Write a Python function to compute p(x) using this method. Analyze the time complexity of your code and express the same in asymptotic notation.
- (d) Perform empirical analysis of run time of all the three versions: Execute the functions for different values of n (degree of the polynomial) and tabulate the results (note that each entry should be an average over several runs, say m). Use randomly generated values of a0, a1, , an-1 for each value of n. Perform ratio analysis with well known complexity classes to confirm the growth rate of running times of all the three versions.

II. CODE:

```
# -*- coding: utf-8 -*-

"""

This module provides a series of functions that calculate the value of a polynomial, given the polynomial coefficient terms and the value of x in 3 different time complexities, namely O(n^2), O(nlogn) and O(n). This is a part of the exercises given under the course UIT2201 (Programming and Data Structures).

In this source code I have executed my own logic. The code follows good coding practices.
```

```
Your comments and suggestions are welcome.
Created on Wed Apr 29 2023
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def complexity_n2(coefficients,x):
    The given function takes in the coefficients of \ensuremath{\mathsf{a}}
    polynomial as well as the value of \boldsymbol{x} to calculate
    the value of polynomial at a given x as well as
    number of operations taken.
    The input is not modified in any way and there
    are no side effects.
    args:
        coefficients: the coefficients of the polynomial
        x: the value of x to be substituted
    Returns:
        A tuple of polynomial value and number of
        operations performed.
    degree = len(coefficients) - 1
    total_sum = 0
    count = 0
    count += 2
    for coeff in coefficients:
        count += 1
        prod = 1
        for i in range(degree):
           count += 1
           prod *= x
        total_sum += prod*coeff
        degree -= 1 count += 2
    return (total_sum, count)
def power(x, y):
    The given function calculates the value of \boldsymbol{x} raised
    to the power y in time O(\log n).
    The input is not modified in any way and there are no
    side effects.
    args:
        x: the base
        y: the power
    Returns:
        Value of x raised to the power y.
```

```
temp = \overline{power(x, \underline{int}(y / 2))}
        return temp * temp
        return x * temp * temp
def complexity_nlogn(coefficients,x):
    The given function takes in the coefficients of a
    polynomial as well as the value of \boldsymbol{x} to calculate
    the value of polynomial at a given \boldsymbol{x} as well as
    number of operations taken. This implementation
    calculates value in O(nlogn) time.
    The input is not modified in any way and there
       coefficients: the coefficients of the polynomial
        x: the value of x to be substituted
        The polynomial value at given value \boldsymbol{x}.
    degree = len(coefficients) - 1
    total_sum = 0
    for coeff in coefficients:
        prod = power(x,degree)
        total_sum += prod*coeff
        degree -= 1
    return total_sum
def horner_method(coefficients,x):
    The given function takes in the coefficients of a
    polynomial as well as the value of \boldsymbol{x} to calculate
    the value of polynomial at a given \boldsymbol{x} as well as
    number of operations taken. This is the implementation
    of horner method which calculates value in O(n) time.
    The input is not modified in any way and there
    are no side effects.
    args:
        coefficients: the coefficients of the polynomial
        x: the value of x to be substituted
    Returns:
        The polynomial value at given value x.
```

```
global homer_ct
   result = coefficients[0]
   homer_ct += 1
   for i in range(1, len(coefficients)):
      homer_ct += 1
      result = result*x + coefficients[i]
   return result
f __name__ == '__main__':
  homer_ct = 0
   coeff = [x for x in range(1,11)]
   print(f"Coefficients of the polynomial being evaluated is: {coeff}")
   value, count = complexity_n2(coeff,2)
   \label{eq:print}  \text{print("Value of polynomial calculated by algorithm with 0(n^2) time complexity: ",value)} 
   print("Number of comparisons is :",count)
   print()
  print("Value of polynomial calculated by lgorithm with O(nlogn) time complexity: ",complexity_nlogn(coeff,2))
   print("Number of comparisons is :",ct)
  print("Value of polynomial calculated by horner method: ",horner_method(coeff,2))
   print("Number of comparisons is :",homer_ct)
   print()
```

III. OUTPUT:

Coefficients of the polynomial being evaluated is: [1, 2, 3, 4, 5, 6, 7, 8, 9, 10]

Value of polynomial calculated by algorithm with $O(n^2)$ time complexity: 2036 Number of comparisons is : 77

Value of polynomial calculated by Igorithm with O(nlogn) time complexity: 2036 Number of comparisons is: 45

Value of polynomial calculated by horner method: 2036 Number of comparisons is: 10

IV. EMPIRICAL ANALYSIS:

	ICAL ANI	124515:		
I re	gorithm (O(n2):		
h	P(n)	PinIn	P(n) /n2	F(n) In log n
10	77	7.7	0.77	2.3179
100	5252	52.52	0.5252	7.91
1000	502507	5025.02	0.502502	50.42
10000	5002 5002	5002.5002	0.5003	376.468
20000	200050002	10000.5	0.500 13	700.06
	con see 7	hat The	dosest fit	is f(n)/h²
we d		rity of a	closest fit is a lyonithm is a	
we d	re comple	rity of a	elgorithm is c	
We din	lgorithm	city of a	elgorithm is c	O(h²).
We din II . I	Ngorithm P(n)	0 (n lo	lgorithm is c gn): P(n) /n²	P(n) /nlogh
We din II . I	emple UgariThm P(n) 45	0 (n lo	elgorithm is 0 $g(n)$: $f(n)/h^2$ 0.45	P(n) /nlogh 1.3546
We di. Zim II . I	P(n) 45	O(n lo) P(n)/n 4.5 7.73	lgorithm is a gn): P(n)/h² 0.45 0.0773	P(n) /nlog h 1.3546 1.1635

III Algorithm Horner's method:

h	F(h)	finila	f(n)/n2	f(n)/hlogi
10	10	,	0-1	0.3010
100	100	,	0.01	0.1505
1000	1000	,	0.001	0.1003
10000	10000	1	0.0001	0.07525
20000	20000	1	0.00005	0.06999