



Message Passing Interface (MPI)

K Hari Babu Department of Computer Science & Information Systems





Programming with MPI: Derived Data Types

Sending Data in a Single Message



- Sending a, b and n in one message will take lesser time compared to sending a, b and n individually
- MPI provides three basic approaches to consolidate data
 - Send in an array using count argument if data is of same type
 - Derived data types
 - MPI_Pack/MPI_Unpack are used to pack data into a contiguous buffer space
 - The pack/unpack routines are provided for compatibility with previous libraries.
 Derived data types is offered in later versions of MPI

```
09 \text{ if (my rank == 0) } 
10
            printf("Enter a, b, and n\n");
            scanf("%lf %lf %d", a,b,n);
12 for (dest = 1; dest < comm sz; dest++) {
            MPI Send(a, 1, MPI DOUBLE, dest, 0, MPI COMM WORLD);
13
            MPI Send(b, 1, MPI DOUBLE, dest, 0, MPI COMM WORLD);
14
15
            MPI Send(n, 1, MPI INT, dest, 0, MPI COMM WORLD);
16 }
17 } else { /* my rank != 0 */
            MPI Recv(a, 1, MPI DOUBLE, 0, 0, MPI COMM WORLD, MPI STATUS IGNORE);
18
            MPI Recv(b, 1, MPI DOUBLE, 0, 0, MPI COMM WORLD, MPI STATUS IGNORE);
20
            MPI Recv(n, 1, MPI INT, 0, 0, MPI COMM WORLD, MPI STATUS IGNORE);
22
24 }
```



- In MPI, a derived datatype can be used to represent any collection of data items in memory by storing both the <u>types</u> of the items and their <u>relative locations</u> in memory
- In our trapezoidal rule example, suppose that on process 0 the variables a, b, and n are stored in memory locations with the following addresses
 - Then the following derived datatype could represent these data items
 - {(MPI_DOUBLE,0),(MPI_DOUBLE,16),(MPI_INT,24)}
 - The first element of each pair corresponds to the type of the data, and the second element of each pair is the displacement of the data element from the beginning of the type

Variable	Address
а	24
b	40
n	48



```
int MPI_Type_create_struct(
  int count, //no of elements in the struct
  int array_of_blocklengths[],
  MPI_Aint array_of_displacements[],
  MPI_Datatype array_of_types[],
  MPI_Datatype *newtype
);
```

- We can use MPI_Type_create_struct to build a derived datatype that consists of individual elements that have different basic types
- The argument count is the number of elements in the datatype
 - for our example, it should be three
- Each of the next 3 array arguments should have count elements
- The first array, array of block lengths, allows for the possibility that the individual data items might be arrays or subarrays
 - for example, the first element were an array containing five elements, we would have array_of_blocklengths[0] = 5;
 - in our case array_of_blocklengths[3] = {1, 1, 1};



```
int MPI_Type_create_struct(
  int count, //no of elements in the struct
  int array_of_blocklengths[],
  MPI_Aint array_of_displacements[],
  MPI_Datatype array_of_types[],
  MPI_Datatype *newtype
);
```

array_of_blocklengths allows for the possibility that the individual data items might be arrays or subarrays.

- array_of_displacements, specifies the displacements, in bytes, from the start of the message
 - array of displacements[] = {0, 16, 24};
- We can use MPI Get address to calculate displacements

```
int MPI_Get_address(
void* location_p /* in */,
MPI_Aint* address_p /* out */);
/*returns address of the memory
location referenced by location_p*/
```

```
MPI_Aint a_addr, b_addr, n_addr;
MPI_Get_address(&a, &a_addr);
array_of_displacements[0] = 0;
MPI_Get_address(&b, &b_addr);
array_of_displacements[1] = b_addr - a_addr;
MPI_Get_address(&n, &n_addr);
array_of_displacements[2] = n_addr - a_addr;
```



```
int MPI_Type_create_struct(
  int count, //no of elements in the struct
  int array_of_blocklengths[],
  MPI_Aint array_of_displacements[],
  MPI_Datatype array_of_types[],
  MPI_Datatype *newtype
);
```

- The array_of_types should store the MPI datatypes of the elements
 - array_of_types[3] = {MPI_DOUBLE, MPI_DOUBLE, MPI_INT}
- With these initializations, we can build the new datatype with the call

```
MPI_Datatype input_mpi_t;
MPI_Type_create_struct(3, array_of_blocklengths,
array_of_displacements, array_of_types,
&input_mpi_t);
```

 Before we can use input_mpi_t in a communication function, we must first commit it with a call to

int MPI_Type_commit(MPI_Datatype* new_mpi_t p /*in/out*/);

Using Derived Data Type



- Now, in order to use new_mpi_t, we make the following call to MPI_Bcast on each process:
- MPI_Bcast(&a, 1, input_mpi_t, 0, comm);

Example



```
91
        * Function:
92
                        Build mpi type
93
                        Build a derived datatype so that the three
        * Purpose:
94
                        input values can be sent in a single message.
        * Input args: a p: pointer to left endpoint
96
                        b p: pointer to right endpoint
97
                        n p: pointer to number of trapezoids
98
        * Output args: input mpi t p: the new MPI datatype
99
       */
100
      □void Build mpi type (
101
             double*
                                          /* in */,
                            ар
                                          /* in */,
102
             double*
                          bр
                                          /* in */.
103
             int*
                          n p
             MPI Datatype* input mpi t_p /* out */) {
104
105
          int array of blocklengths[3] = {1, 1, 1};
106
107
          MPI Datatype array of types[3] = {MPI DOUBLE, MPI DOUBLE, MPI INT};
          MPI Aint a addr, b addr, n_addr;
108
109
          MPI Aint array of displacements[3] = {0};
110
111
          MPI Get address(a p, &a addr);
112
          MPI Get address(b p, &b addr);
113
          MPI Get address(n p, &n addr);
114
          array of displacements[1] = b addr-a addr;
115
          array of displacements[2] = n addr-a addr;
116
          MPI Type create struct(3, array of blocklengths,
117
                array of displacements, array of types,
118
                input mpi t p);
119
          MPI Type commit(input mpi t p);
120
      } /* Build mpi type */
```

Example



```
132
     ─void Get input(
133
                 my rank /* in */,
             int
134
                comm sz /* in */,
             int
            double* a_p /* out */,
135
            double* b_p /* out */,
136
            int* np
                            /* out */) {
137
138
          MPI Datatype input mpi t;
139
          Build_mpi_type(a_p, b_p, n_p, &input_mpi_t);
140
141
          if (my rank == 0) {
142
143
            printf("Enter a, b, and n\n");
144
             scanf("%lf %lf %d", a_p, b_p, n_p);
145
146
          MPI_Bcast(a_p, 1, input_mpi_t, 0, MPI_COMM_WORLD);
147
148
          MPI Type free (&input mpi t);
149
      | /* Get input */
```

 We're usually not interested in the time taken from the start of program execution to the end of program execution

```
78
         MPI Barrier (comm);
         start = MPI Wtime();
79
         Mat vect mult(local A, local x, local y, local m, n, local n, comm);
         finish = MPI Wtime();
81
         loc elapsed = finish-start;
82
83
         MPI Reduce (&loc elapsed, &elapsed, 1, MPI DOUBLE, MPI MAX, 0, comm);
84
     # ifdef DEBUG
85
         Print vector ("y", local y, m, local m, my rank, comm);
86
      -# endif
87
88
          if (my rank == 0)
89
             printf("Elapsed time = %e\n", elapsed);
90
```

• Time taken by parallel algorithm is the time of the slowest process. We can't get exactly this time because we can't insure that all the processes start at the same instant. However, we can come reasonably close. The MPI collective communication function MPI Barrier insures that no process will return from calling it until every process in the communicator has started calling it.

Odd-even Transposition Sort



 Odd-even transposition sort is a sorting algorithm that's similar to bubble sort, but that has more opportunities for parallelism

```
137

──void Odd even sort(
              int a[] /* in/out */,
138
139
              int n
                        /* in
                                   */) {
140
           int phase, i, temp;
141
           for (phase = 0; phase < n; phase++)</pre>
142
143
              if (phase % 2 == 0) { /* Even phase */
144
                 for (i = 1; i < n; i += 2)
                    if (a[i-1] > a[i]) {
145
                       temp = a[i];
146
147
                       a[i] = a[i-1];
148
                        a[i-1] = temp;
149
150
              } else { /* Odd phase */
151
                 for (i = 1; i < n-1; i += 2)
                    if (a[i] > a[i+1]) {
152
153
                       temp = a[i];
154
                       a[i] = a[i+1];
155
                        a[i+1] = temp;
156
157
158
             Odd even sort */
159
```

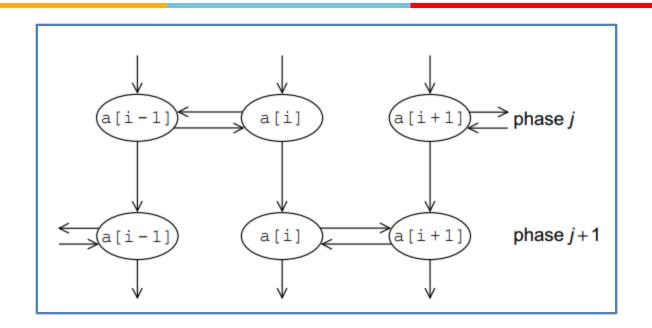
Suppose a = $\{9, 7, 8, 6\}$.

Subscript in Array

Phase 0 1 2 3

0 9 \leftrightarrow 7 8 \leftrightarrow 6
7 9 6 8
1 7 9 \leftrightarrow 6 8
2 7 \leftrightarrow 6 9 \leftrightarrow 8
2 7 \leftrightarrow 6 9 \leftrightarrow 8
3 6 7 \leftrightarrow 8 9
6 7 8 9





When n = p: depending on the phase, process i can send its current value, a[i], either to process i – 1 or process i + 1. At the same time, it should receive the value stored on process i – 1 or process i + 1, respectively, and then decide which of the two values it should store as a[i] for the next phase.

It's unlikely that n = p

How should this be modified when each process is storing

n/p > 1 elements?

At any time, each process should have exactly n/p elements.

```
Sort local keys;
 2
      for (phase = 0; phase < comm sz; phase++) { } 
 3
           partner = Compute partner (phase, my rank);
 4
           if (I'm not idle) {
 5
                Send my keys to partner;
                Receive keys from partner;
                if (my rank < partner)</pre>
 8
                    Keep smaller keys;
 9
                else
10
                    Keep larger keys;
11
12
```

	Process				
Time	0	1	2	3	
Start After Local Sort After Phase 0 After Phase 1 After Phase 2 After Phase 3	15, 11, 9, 16 9, 11, 15, 16 3, 7, 8, 9 3, 7, 8, 9 1, 2, 3, 4 1, 2, 3, 4	3, 14, 8, 7 3, 7, 8, 14 11, 14, 15, 16 1, 2, 4, 5 5, 7, 8, 9 5, 6, 7, 8	4, 6, 12, 10 4, 6, 10, 12 1, 2, 4, 5 11, 14, 15, 16 6, 10, 11, 12 9, 10, 11, 12	5, 2, 13, 1 1, 2, 5, 13 6, 10, 12, 13 6, 10, 12, 13 13, 14, 15, 16 13, 14, 15, 16	

innovate achieve lead

Parallel odd-even transposition sort

how do we compute the partner rank?

```
if (phase % 2 == 0) /* Even phase */
          if (my rank % 2 != 0) /* Odd rank */
              partner = my rank - 1;
          else /* Even rank */
              partner = my rank + 1
      else /* Odd phase */
 6
 7
          if (my rank % 2 != 0) /* Odd rank */
              partner = my rank + 1;
 8
          else /* Even rank */
 9
10
              partner = my rank - 1;
      if (partner == -1 | | partner == comm sz)
11
12
      partner = MPI PROC NULL;
```

MPI_PROC_NULL is a constant defined by MPI. When it's used as the source or destination rank in a point-to-point communication, no communication will take place and the call to the communication will simply return.

	Process				
Time	0	1	2	3	
Start After Local Sort After Phase 0 After Phase 1 After Phase 2 After Phase 3	15, 11, 9, 16 9, 11, 15, 16 3, 7, 8, 9 3, 7, 8, 9 1, 2, 3, 4 1, 2, 3, 4	3, 14, 8, 7 3, 7, 8, 14 11, 14, 15, 16 1, 2, 4, 5 5, 7, 8, 9 5, 6, 7, 8	4, 6, 12, 10 4, 6, 10, 12 1, 2, 4, 5 11, 14, 15, 16 6, 10, 11, 12 9, 10, 11, 12	5, 2, 13, 1 1, 2, 5, 13 6, 10, 12, 13 6, 10, 12, 13 13, 14, 15, 16 13, 14, 15, 16	



- how do we exchange elements?
 - If a process is not idle, we might try to implement the communication with a call to MPI_Send and a call to MPI_Recv

```
MPI_Send(my_keys, n/comm_sz, MPI_INT, partner, 0, comm);
MPI_Recv(temp_keys, n/comm_sz, MPI_INT, partner, 0, comm, MPI_STATUS_IGNORE);
```

- This might result in deadlock.
- MPI standard allows MPI_Send to behave in two different ways:
 - it can simply copy the message into an MPI-managed buffer and return,
 - or it can block until the matching call to MPI Recv starts
 - Furthermore, many implementations of MPI set a threshold at which the system switches from buffering to blocking. That is, messages that are relatively small will be buffered by MPI Send, but for larger messages, it will block.
- If the MPI_Send executed by each process blocks, no process will be able to start executing a call to MPI_Recv leading to deadlock



- A program that relies on MPI-provided buffering is said to be unsafe
 - Such a program may run without problems for various sets of input, but it may hang or crash with other sets. If we use MPI_Send and MPI_Recv in this way, our program will be unsafe, and it's likely that for small values of n the program will run without problems, while for larger values of n, it's likely that it will hang or crash
 - o we can use an alternative to MPI Send i.e. MPI Ssend
 - MPI_Ssend is guaranteed to block until the matching receive starts

ort





Parallel odd-even transposition sort

- How can we modify the communication in the parallel oddeven sort program so that it is safe?
 - Communication must be restructured. The most common cause of an unsafe program is multiple processes simultaneously first sending to each other and then receiving
 - need to restructure the communications so that some of the processes receive before sending

```
if (my rank % 2 == 0) {
    MPI_Send(msg, size, MPI_INT, (my_rank+1) % comm_sz, 0, comm);

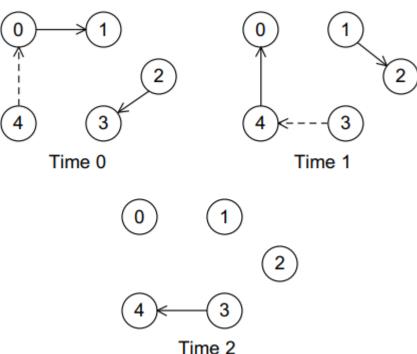
MPI_Recv(new_msg, size, MPI_INT, (my_rank+comm_sz-1) % comm_sz,
    0, comm, MPI_MPI_STATUS_IGNORE IGNORE).

else {
    MPI_Recv(new_msg, size, MPI_INT, (my_rank+comm_sz-1) % comm sz,
    0, comm, MPI_STATUS_IGNORE).

MPI_Send(msg, size, MPI_INT, (my_rank+1) % comm_sz, 0, comm);
}
```

o this will work if comm_sz is even. But when it is odd?

 For 5 processes, dashed arrows shows waiting communications though expected to be complete in the same phase.



 MPI provides an alternative to scheduling the communications ourselves— MPI_Sendrecv:

```
lint MPI Sendrecv(
          void* send_buf_p /* in */,
 2
          int send buf size /* in */,
          MPI Datatype send buf type /* in */,
          int dest /* in */,
          int send tag /* in */,
          void* recv buf p /* out */,
          int recv buf size /* in */,
          MPI Datatype recv buf type /* in */,
          int source /* in */,
10
          int recv tag /* in */,
11
          MPI Comm communicator /* in */,
          MPI Status* status p /* in */
13
```

- This function carries out a blocking send and a receive in a single call. The dest and the source can be the same or different.
- MPI implementation schedules the communications so that the program won't hang or crash.

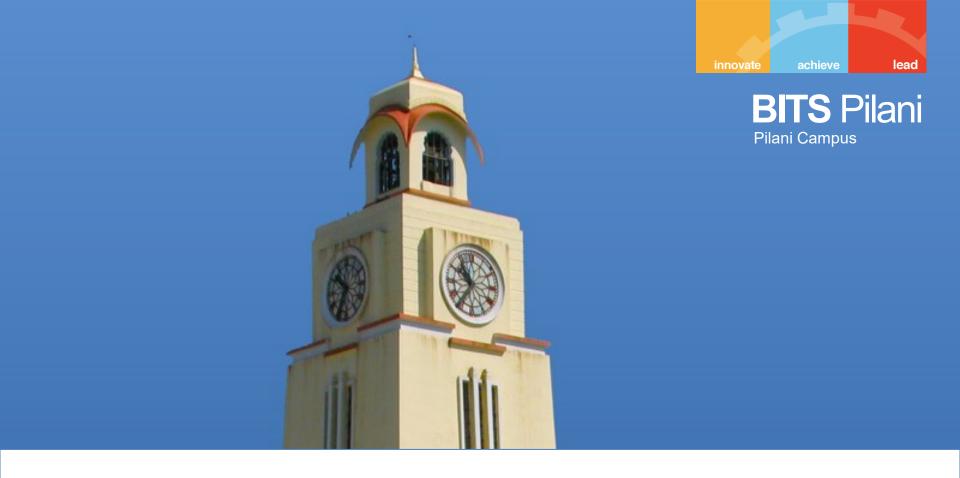
lead

```
One iteration of Odd-even transposition sort
306
          Purpose:
307
        * In args:
                        local n, phase, my rank, p, comm
        * In/out args: local A
308
309
         * Scratch:
                        temp B, temp C
       */
310
      woid Odd even iter(int local A[], int temp B[], int temp C[],
311
312
               int local n, int phase, int even partner, int odd partner,
313
               int my rank, int p, MPI Comm comm) {
          MPI Status status;
314
315
          if (phase % 2 == 0) {
316
317
              if (even partner >= 0) {
318
                 MPI Sendrecv(local A, local n, MPI INT, even partner, 0,
319
                    temp B, local n, MPI INT, even partner, 0, comm,
320
                    &status);
                 if (my rank % 2 != 0)
321
322
                   Merge high (local A, temp B, temp C, local n);
323
                 else
324
                   Merge low(local A, temp B, temp C, local n);
325
          } else { /* odd phase */
326
327
              if (odd partner >= 0) {
328
                 MPI Sendrecv(local A, local n, MPI INT, odd partner, 0,
329
                    temp B, local n, MPI INT, odd partner, 0, comm,
330
                    &status);
331
                 if (my rank % 2 != 0)
332
                   Merge low(local A, temp B, temp C, local n);
333
                 else
                   Merge high (local_A, temp_B, temp_C, local_n);
334
335
336
```

References



- https://computing.llnl.gov/tutorials/mpi/
- "An introduction to Parallel Programming", Peter S. Pacheco (Chapter 3)



Thank You