Operating Systems Laboratory (CS39002) -- Assignment 5

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Data Structures USed:

1. Room

Members:

->int room_id

unique id for each room

->int guest_id

unique id for each guest who is in the room

->int guest_count

to keep track of number of guest who have already used the room so after 2 guest we need to activate cleaner to clean the room

-> pair<int,int> stay_time

stay_time.first stored the stay duration for the first guest who has entered the room after

it has been cleaned or during the start of the program and stay_time.second stores the stay duration for the second guest who has entered the room after the first guest has left the room.

Priority Heuristic:

We are randomly alloting unique priority to each guest between 1 to Y(Number of guests).

Global Variable:

->stay_count

To indicate the number of times all rooms used by the guests , when stay_count become equal to 2*N i.e when each of the rooms has been used twice then the cleaner threads will be act ivated by unblocking the clean_start semaphore by setting it to 1 such that the cleaner threads start cleaning

Semaphores Used:

N - number of rooms

Used Semaphores:

->bin_room_sems[N] :

To lock the rooms so that only one quest can access the room at a time

->sig_room_sems[N] :

Used for blocking the room when used by some guest and free it when guest leaves the room or a higher priority guest has to be allocated the room

->cleaning :

To block guest threads and facilitate the functioning of cleaner threads till all rooms are cleaned

->stay_count_sem :

Binary semaphore to lock stay_count varible while incrementing by one guest thread

->clean start :

To indicate that cleaner threads has started cleaning the room and no guest can enter the room

->clean_end : To indicate that cleaner threads has finished cleaning the room and guest can enter the room

Semaphores Application

1.All bin_room_sems are initially initialized to 1 , after a room 'i' is alloted to a guest

the guest count is increased and the wait by $bin_room_sem[i]$ becomes unblocking and $bin_room_sem[i]$ becomes 0 so now no other guest could access this room till the room id = -1 (empty)

- 2.sig_room_sems all are intially intialized to 0 so when each guest thread will be using th e room till the sig_room_sem[i] for a room 'i' is et to 1 either when a more priority guest comes and also the guest count for that room is leass than 2 or when the guest currently in the room leaves the room after some given sleep time.
- 3.stay_count_sem is used to lock the stay_count for increment , now when each room has been used 2 times each that the stay_count has become 2*N , so now using a if conditon in this critical section locked by semaphore stay_count_sem , we make the binary semaphore clean_st art , clean_end set in order to activate the cleaner threads. Also within the stay_count_s em Semaphores when the above if (stay_count == 2*N) conditon is satisfied then set the cle aning counting semaphore to N
- 4.cleaning so this counting semaphore has been incremented to N by the guest threads after the room has been used twice each so now sem_wait(&cleaning) will gets unblocked for N time s in order to facilitate cleaning for N rooms, now if it gets unblocked then if clean_star t is set to 1 then we do a sem_post on sig_room_sems[i] for each room 'i' such that in case a room is currently occupied by a guest who is in sleep then he has to leave the room when the cleaning starts. Also we make the clean_start to 0 so that no other guest could enter the room till the cleaning is done. Now we do a sem_wait on clean_end so that the cleaner t hread could wait till the cleaning is done for all the N rooms. After the cleaning is done we make the clean_end to 0 so that the guest threads could enter the room again and we make the cleaning to 0 so that the cleaner threads could wait for the cleaning to be done again for N rooms.
- 5. After all rooms have been cleaned we will just check if all_clean is true and here the b locking call by clean_end binary semaphore will become unblocking as clean_send was set so now we will just make all the bin_room_sems[i] for each room[i] to be set such that the rooms could be occupied again now.