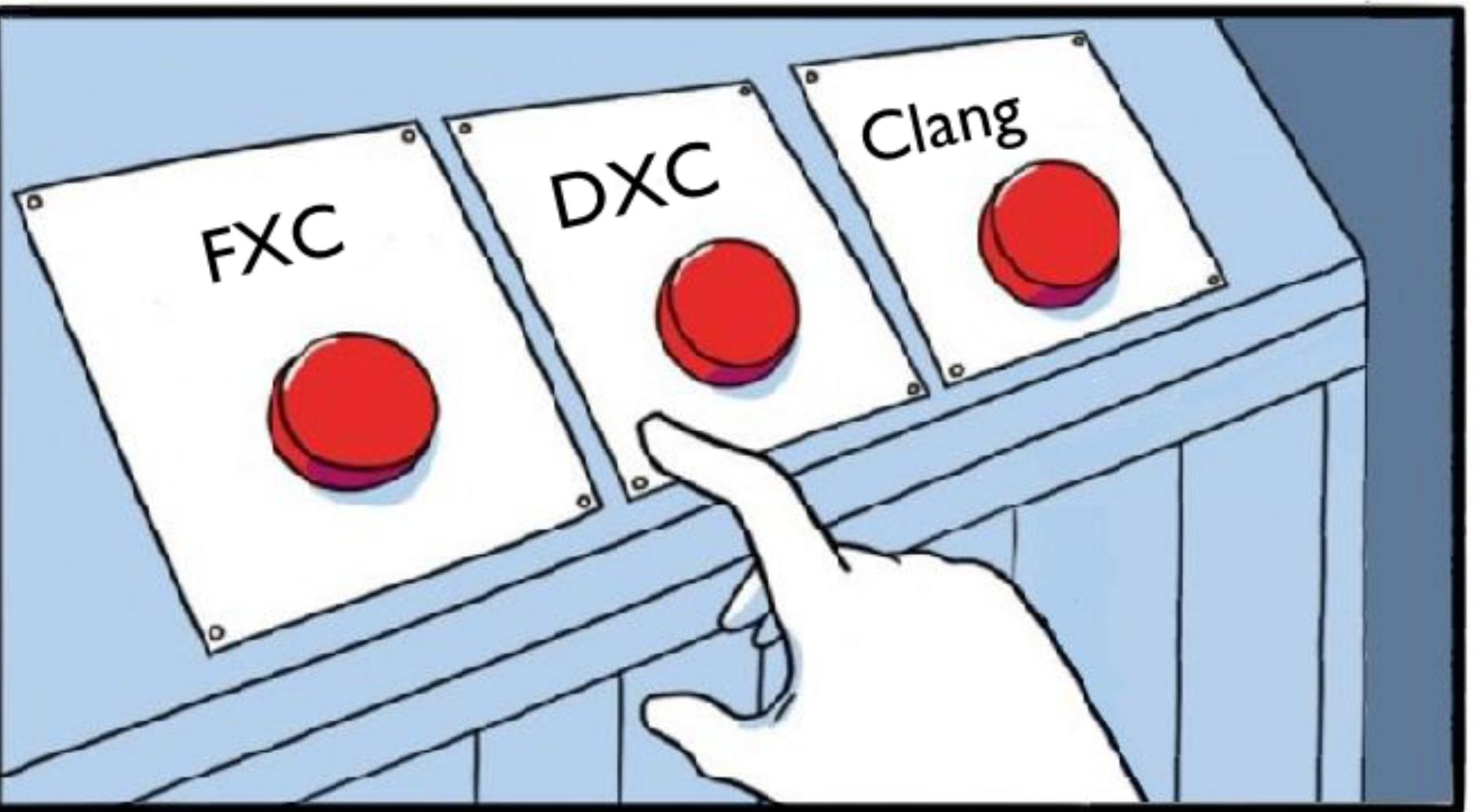
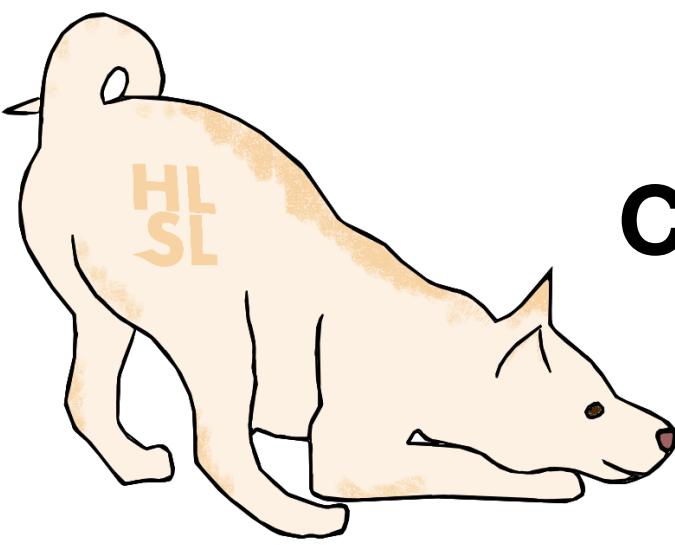


Two Compilers, One Language, No Specification

Defining a Language from
Disagreeing References

Chris Bieneman



JAKE-CLARK.TUMBLR

What is HLSL?

- GPU programming language introduced in 2002
- Evolved from NVIDIA's C for Graphics (Cg)
- Mostly used for computer graphics
 - Also used in GPGPU and ML applications



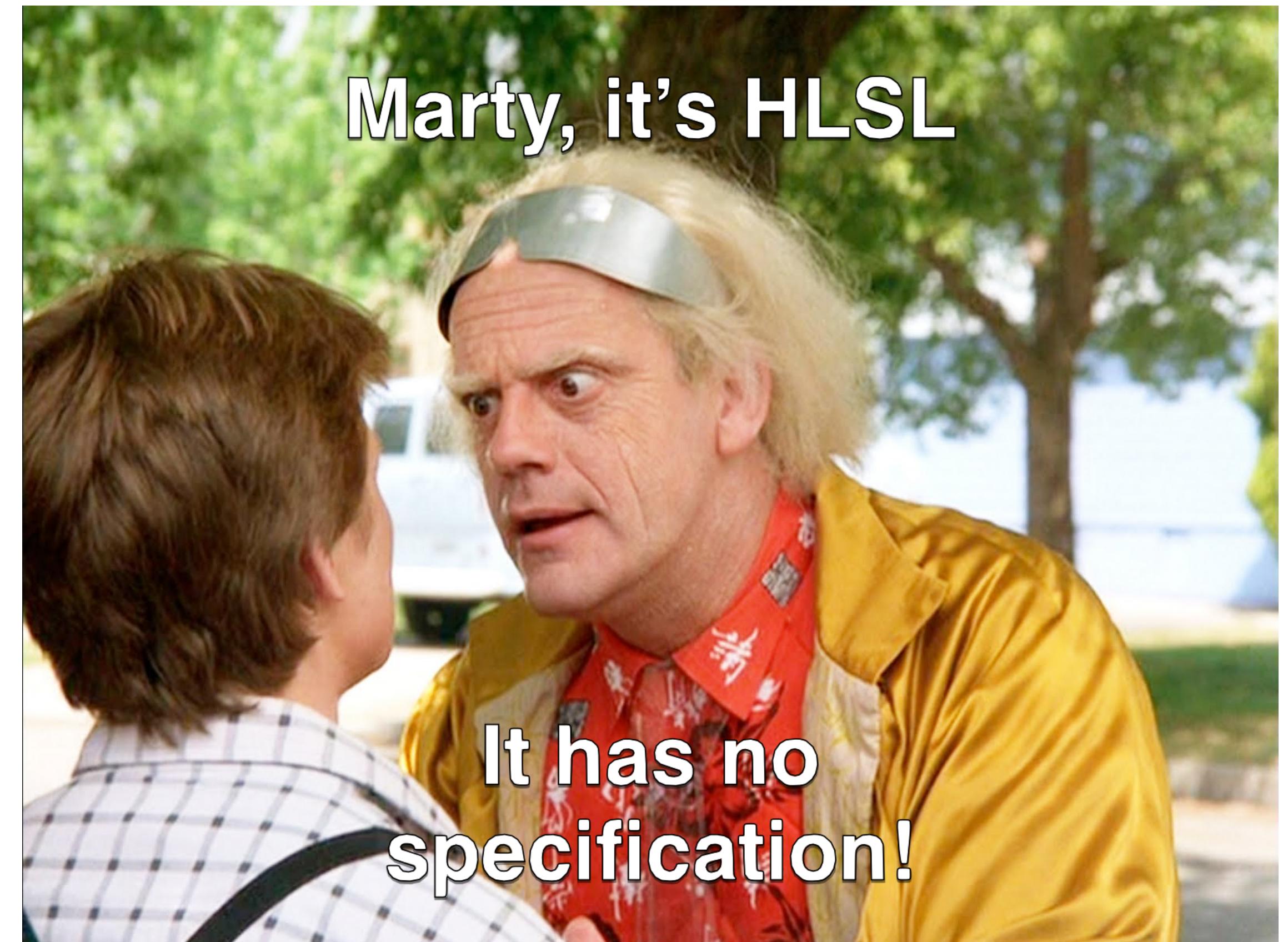
State of HLSL

- Reference Compiler #1 - FXC
 - Out of active support
 - DirectX 9-11 (and early 12)
 - HLSL versions 2002->2015
- Reference Compiler #2 - DXC
 - Fork of LLVM 3.7
 - DirectX 12 Ultimate
 - HLSL Versions 2016, 2017, 2018, 2021 & 202x
 - 202x & critical support only



The Problem

- We need a modern HLSL compiler
- We have a large base of existing HLSL code
- We have a lot of users that need to transition
- We don't want a bug-for-bug compatible replacement
 - GPU code portability is a common problem
- We need a written specification!



Why do we need a spec?

```
static int Count;

void increment(inout int C) {
    C += 1;
    C += Count;
}

int Fn() {
    Count = 0;
    increment(Count);
    return Count;
}
```

Same Compiler, Different Results

DirectX stores 1, SPIR-V stores 2

The image shows two instances of the Compiler Explorer tool running the same HLSL code. The left instance is set to DXC (trunk) and -T cs_6_6, while the right instance is set to DXC (trunk) and -T cs_6_6 -spirv. Both instances show the same HLSL code, but the generated assembly output is different.

HLSL Source:

```
static int Count;
void increment(inout int C) {
    C += 1;
    C += Count;
}
int Fn() {
    Count = 0;
    increment(Count);
    return Count;
}
RWBuffer<int> Out;
[numthreads(1,1,1)]
void main() {
    Out[0] = Fn();
}
```

DXC (trunk) -T cs_6_6 Output:

```
%dx.types.Handle = type { i8* }
%dx.types.ResBind = type { i32, i32, i32, i8 }
%dx.types.ResourceProperties = type { i32, i32 }
%"class.RWBuffer<int>" = type { i32 }

define void @main() {
    %Out_UAV_buf = call %dx.types.Handle @_dx.op.createHandleFromBinding(i32 1, i64 0, metadata !42, %1 = call %dx.types.Handle @_dx.op.annotateHandle(i32 216, %dx.types.ResourceProperties)
    call void @_dx.op.bufferStore.i32(i32 69, %dx.types.Handle %1, ret void, !dbg !46
}

declare void @_llvm.dbg.value(metadata i32 1, i64 0, metadata !42, %int = OpTypeInt 32 1
declare void @_dx.op.bufferStore.i32(i32, %dx.types.Handle, i32, %uint = OpTypeInt 32 0
declare %dx.types.Handle @_dx.op.annotateHandle(i32, %dx.types.Handle, i32, %uint_0 = OpConstant %uint 0
declare %dx.types.Handle @_dx.op.createHandleFromBinding(i32, %dx.types.ResourceProperties, %ptr_UniformConstant_type_buffer_image = OpTypePointer Uniform
attributes #0 = { nounwind readonly }
attributes #1 = { nounwind }
```

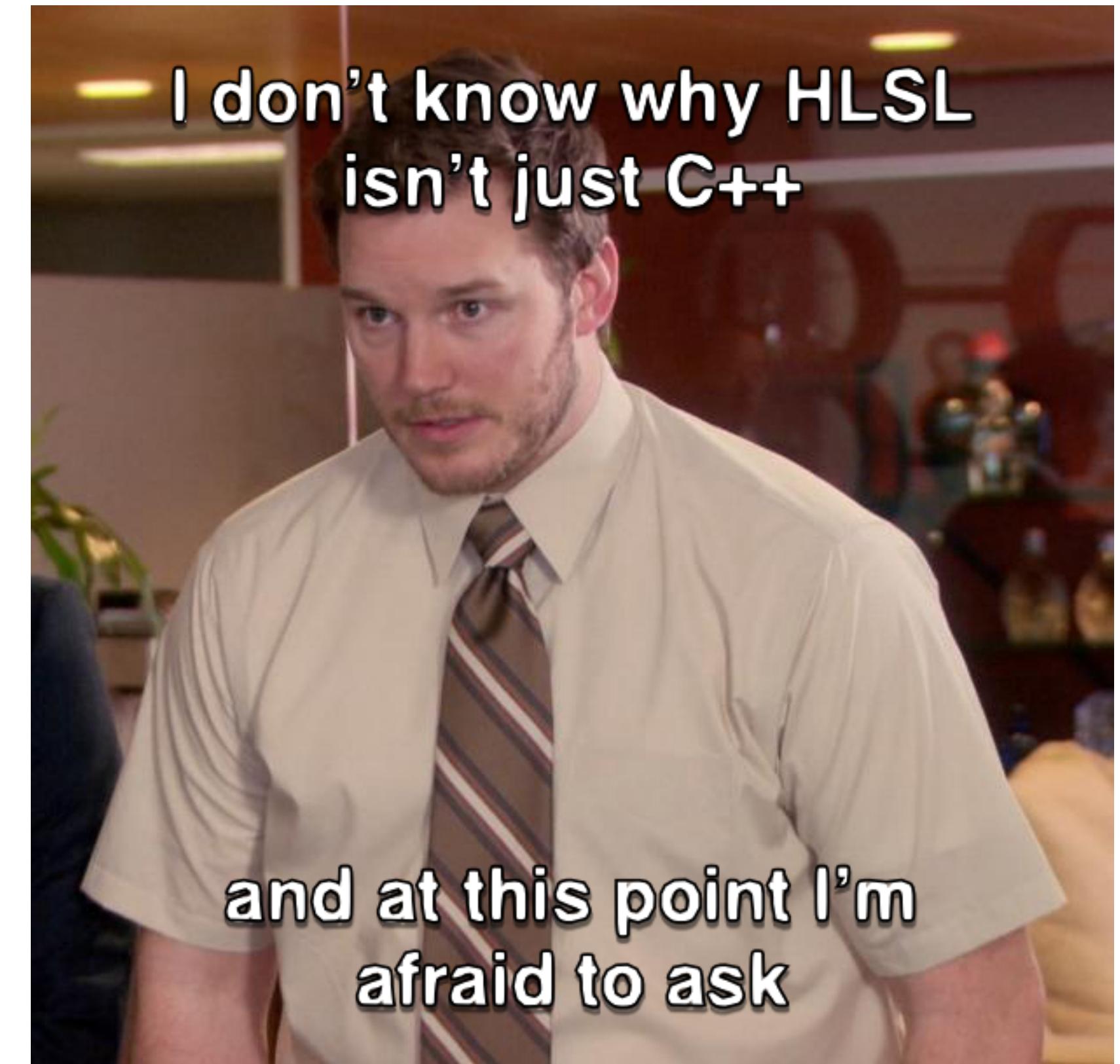
DXC (trunk) -T cs_6_6 -spirv Output:

```
OpCapability ImageBuffer
OpMemoryModel Logical GLSL450
OpEntryPoint GLCompute %main "main"
OpExecutionMode %main LocalSize 1 1 1
OpDecorate %Out DescriptorSet 0
OpDecorate %Out Binding 0
%int = OpTypeInt 32 1
%uint = OpTypeInt 32 0
%uint_0 = OpConstant %uint 0
%type_buffer_image = OpTypeImage %int Buffer 2 0 0 2 R32i
%_ptr_UniformConstant_type_buffer_image = OpTypePointer Uniform
%void = OpTypeVoid
%10 = OpTypeFunction %void
%Out = OpVariable %_ptr_UniformConstant_type_buffer_image
%int_2 = OpConstant %int 2
%main = OpFunction %void None %10
%12 = OpLabel
%13 = OpLoad %type_buffer_image %Out
OpImageWrite %13 %uint_0 %int_2 None
OpReturn
OpFunctionEnd
```

Why isn't HLSL just C or C++?

TL;DR: (1) History, (2) GPUs are weird

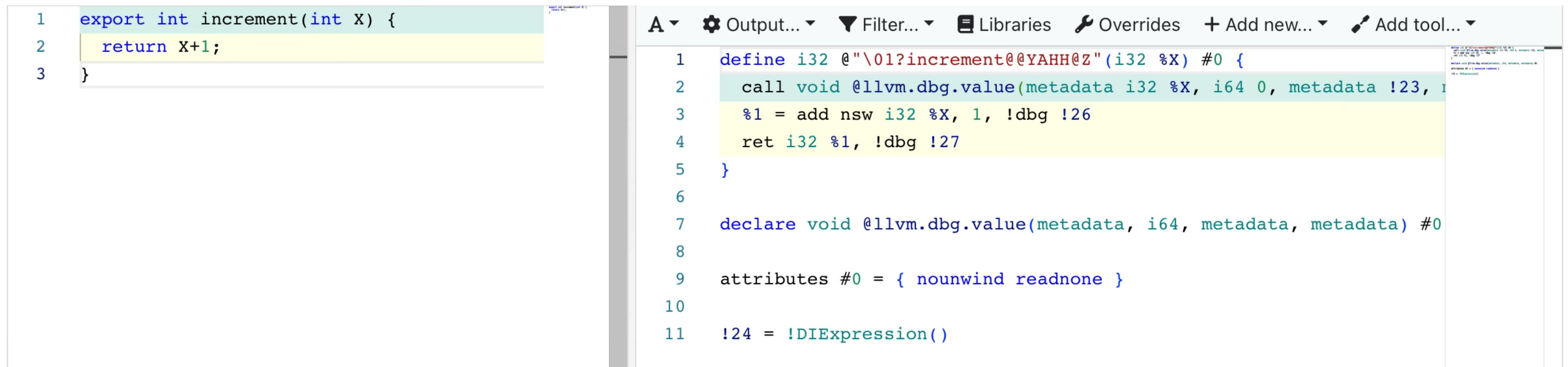
- Early GPUs were very limited
 - No direct memory access
 - No control flow
 - Limited (and strange) data types
- GPU architectures are wildly varied
- GPUs are fundamentally parallel



Similarity: Basic Functions

Reason: HLSL has always been C-like

- HLSL's C/C++ based syntax for functions looks very familiar
- Basic functions just kinda work how you expect
 - HLSL has argument-dependent lookup
- Until they don't...



The image shows a code editor with two panes. The left pane displays HLSL code:`1 export int increment(int x) {
2 return x+1;
3 }`

The right pane displays the generated LLVM Intermediate Representation (IR) code:`A - Output... Filter... Libraries Overrides + Add new... Add tool...
1 define i32 @"\01?increment@@YAH@Z"(i32 %x) #0 {
2 call void @llvm.dbg.value(metadata i32 %x, i64 0, metadata !23, !24)
3 %1 = add nsw i32 %x, 1, !dbg !26
4 ret i32 %1, !dbg !27
5 }
6
7 declare void @llvm.dbg.value(metadata, i64, metadata, metadata) #0
8
9 attributes #0 = { nounwind readnone }
10
11 !24 = !DIExpression()`

Difference: Implementation Inconsistency

Reason: Bugs

- DXC is sometimes inconsistent about overload resolution
- built-in functionality has unintended behavior differences

 **beanz** 2 hours ago
Given this HLSL code which overload should it resolve to?

```
bool fn(bool B) {  
    return WaveReadLaneFirst(B);  
}
```

bool(bool)
int(int)
float(float)
Error! (edited)



 **beanz** 2 hours ago
How about this one?

```
bool fn(bool B) {  
    return WaveActiveSum(B);  
}
```

bool(bool)
int(int)
float(float)
Error!



 **beanz** 2 hours ago
How about this one?

```
bool fn(bool B) {  
    return tan(B);  
}
```

bool(bool)
int(int)
float(float)
Error!



 **beanz** 1 hour ago
Bonus question: How should this behave?

```
bool f(int){}  
bool f(int64_t){}
```

```
bool fn(bool B) {  
    return f(B);  
}
```

bool(bool)
bool(int)
bool(int64_t)
Error!



 **beanz** 2 hours ago
One last one!

```
f(in)  
f(f)  
bool fn(bool B) {  
    return f(B);  
}
```

bool(bool)
int(int)
float(float)
Error!



forat(boat)

Solution: Implicit hlsl.h

Functions are functions

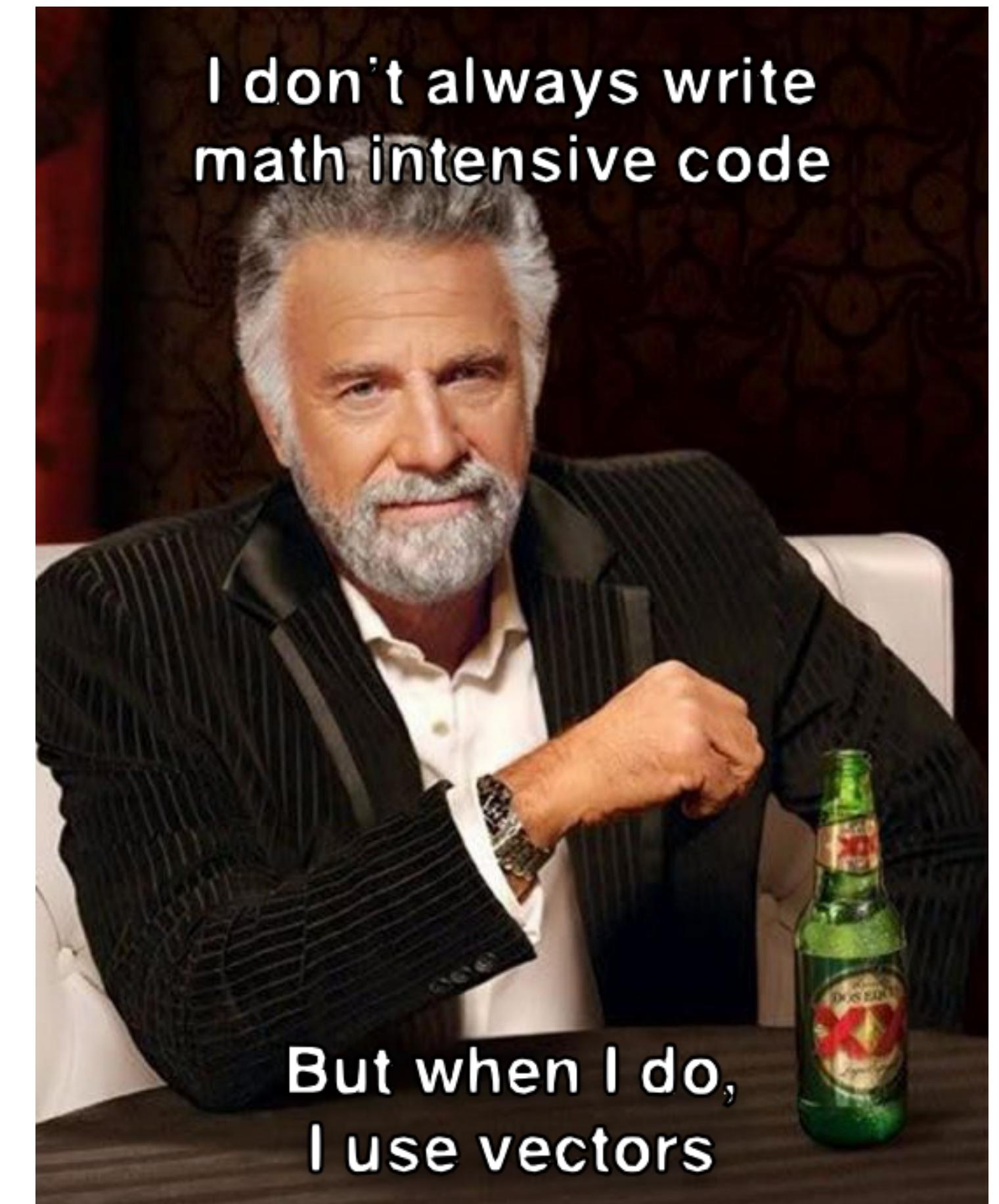
- Alias function declarations to builtins
- Go-to-definition works for library functions
- Consistent overload resolution behavior
- This will break existing code
 - Socializing this with users



Difference: Vectors all the way down!

Reason: colors, vertices, matrices...

- First class vector types are a requirement
 - Vector types need to work with built-in operators
 - Defined conversion sequences
 - Constrain conversions of scalar types... because they're also implicit vectors



SIMT Execution

Single Instruction Multiple Threads

	Arr	Sz	Min	Executing?
0	{5, 0, 0, 0}	1		yes
1	{5, 6, 2, 3}	4		yes
2	{9, 6, 0, 0}	2		yes
3	{2, 6, 1, 0}	3		yes



```
export int reduce_min(int Arr[4], int Sz) {
    int Min = Arr[0];
    for (int I = 1; I < Sz; ++I)
        Min = min(Min, Arr[I]);
    return Min;
}
```

SIMT Execution

Single Instruction Multiple Threads

	Arr	Sz	Min	Executing?
0	{5, 0, 0, 0}	1	5	yes
1	{5, 6, 2, 3}	4	5	yes
2	{9, 6, 0, 0}	2	9	yes
3	{2, 6, 1, 0}	3	2	yes



```
export int reduce_min(int Arr[4], int Sz) {
    int Min = Arr[0];
    for (int I = 1; I < Sz; ++I)
        Min = min(Min, Arr[I]);
    return Min;
}
```

SIMT Execution

Single Instruction Multiple Threads

	Arr	Sz	Min	Executing?
0	{5, 0, 0, 0}	1	5	no
1	{5, 6, 2, 3}	4	5	yes
2	{9, 6, 0, 0}	2	9	yes
3	{2, 6, 1, 0}	3	2	yes

I
1



```
export int reduce_min(int Arr[4], int Sz) {
    int Min = Arr[0];
    for (int I = 1; I < Sz; ++I)
        Min = min(Min, Arr[I]);
    return Min;
}
```

SIMT Execution

Single Instruction Multiple Threads

	Arr	Sz	Min	Executing?
0	{5, 0, 0, 0}	1	5	no
1	{5, 6, 2, 3}	4	5	yes
2	{9, 6, 0, 0}	2	6	yes
3	{2, 6, 1, 0}	3	2	yes

I
1



```
export int reduce_min(int Arr[4], int Sz) {
    int Min = Arr[0];
    for (int I = 1; I < Sz; ++I)
        Min = min(Min, Arr[I]);
    return Min;
}
```

SIMT Execution

Single Instruction Multiple Threads

	Arr	Sz	Min	Executing?
0	{5, 0, 0, 0}	1	5	no
1	{5, 6, 2, 3}	4	5	yes
2	{9, 6, 0, 0}	2	6	no
3	{2, 6, 1, 0}	3	2	yes

1
2



```
export int reduce_min(int Arr[4], int Sz) {
    int Min = Arr[0];
    for (int I = 1; I < Sz; ++I)
        Min = min(Min, Arr[I]);
    return Min;
}
```

SIMT Execution

Single Instruction Multiple Threads

	Arr	Sz	Min	Executing?
0	{5, 0, 0, 0}	1	5	no
1	{5, 6, 2, 3}	4	2	yes
2	{9, 6, 0, 0}	2	6	no
3	{2, 6, 1, 0}	3	1	yes

1
2



```
export int reduce_min(int Arr[4], int Sz) {
    int Min = Arr[0];
    for (int I = 1; I < Sz; ++I)
        Min = min(Min, Arr[I]);
    return Min;
}
```

SIMT Execution

Single Instruction Multiple Threads

	Arr	Sz	Min	Executing?
0	{5, 0, 0, 0}	1	5	no
1	{5, 6, 2, 3}	4	2	yes
2	{9, 6, 0, 0}	2	6	no
3	{2, 6, 1, 0}	3	1	no

I
3



```
export int reduce_min(int Arr[4], int Sz) {
    int Min = Arr[0];
    for (int I = 1; I < Sz; ++I)
        Min = min(Min, Arr[I]);
    return Min;
}
```

SIMT Execution

Single Instruction Multiple Threads

	Arr	Sz	Min	Executing?
0	{5, 0, 0, 0}	1	5	no
1	{5, 6, 2, 3}	4	2	yes
2	{9, 6, 0, 0}	2	6	no
3	{2, 6, 1, 0}	3	1	no

I
3



```
export int reduce_min(int Arr[4], int Sz) {
    int Min = Arr[0];
    for (int I = 1; I < Sz; ++I)
        Min = min(Min, Arr[I]);
    return Min;
}
```

SIMT Execution

Single Instruction Multiple Threads

	Arr	Sz	Min	Executing?
0	{5, 0, 0, 0}	1	5	no
1	{5, 6, 2, 3}	4	2	no
2	{9, 6, 0, 0}	2	6	no
3	{2, 6, 1, 0}	3	1	no

1
4



```
export int reduce_min(int Arr[4], int Sz) {
    int Min = Arr[0];
    for (int I = 1; I < Sz; ++I)
        Min = min(Min, Arr[I]);
    return Min;
}
```

SIMT Execution

Single Instruction Multiple Threads

	Arr	Sz	Min	Executing?
0	{5, 0, 0, 0}	1	5	no
1	{5, 6, 2, 3}	4	2	no
2	{9, 6, 0, 0}	2	6	no
3	{2, 6, 1, 0}	3	1	no

I
4

```
export int reduce_min(int Arr[4], int Sz) {  
    int Min = Arr[0];  
    for (int I = 1; I < Sz; ++I)  
        Min = min(Min, Arr[I]);  
    return Min;  
}
```

Problem: Hidden Costs

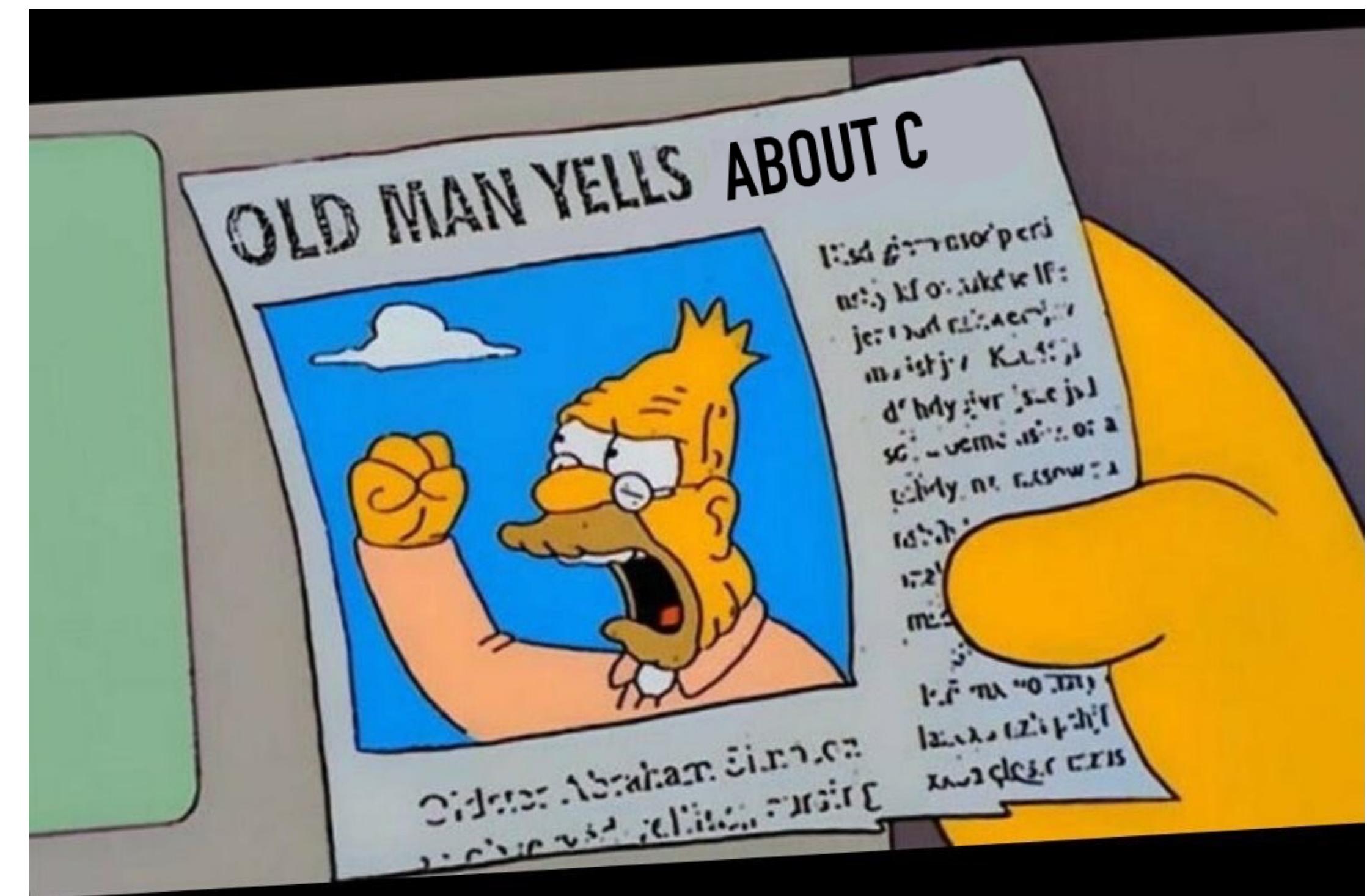
- Each operation happens n times
- Each variable is actually n variables
- Dynamic control flow is always the worst of threads



Usual Arithmetic Conversions

Most hated of C features, worse on vectors

- C implicitly promotes to word-sized values
 - This is extremely expensive on a GPU!
- HLSL also has vectors of arithmetic types
 - We need vector element and dimension conversions!



Solution: Constrain and Extend UAC

- Remove “promotable” sub-word types
- Add vector element and dimension conversions
- Aligns with C language, but considers HLSL’s characteristics

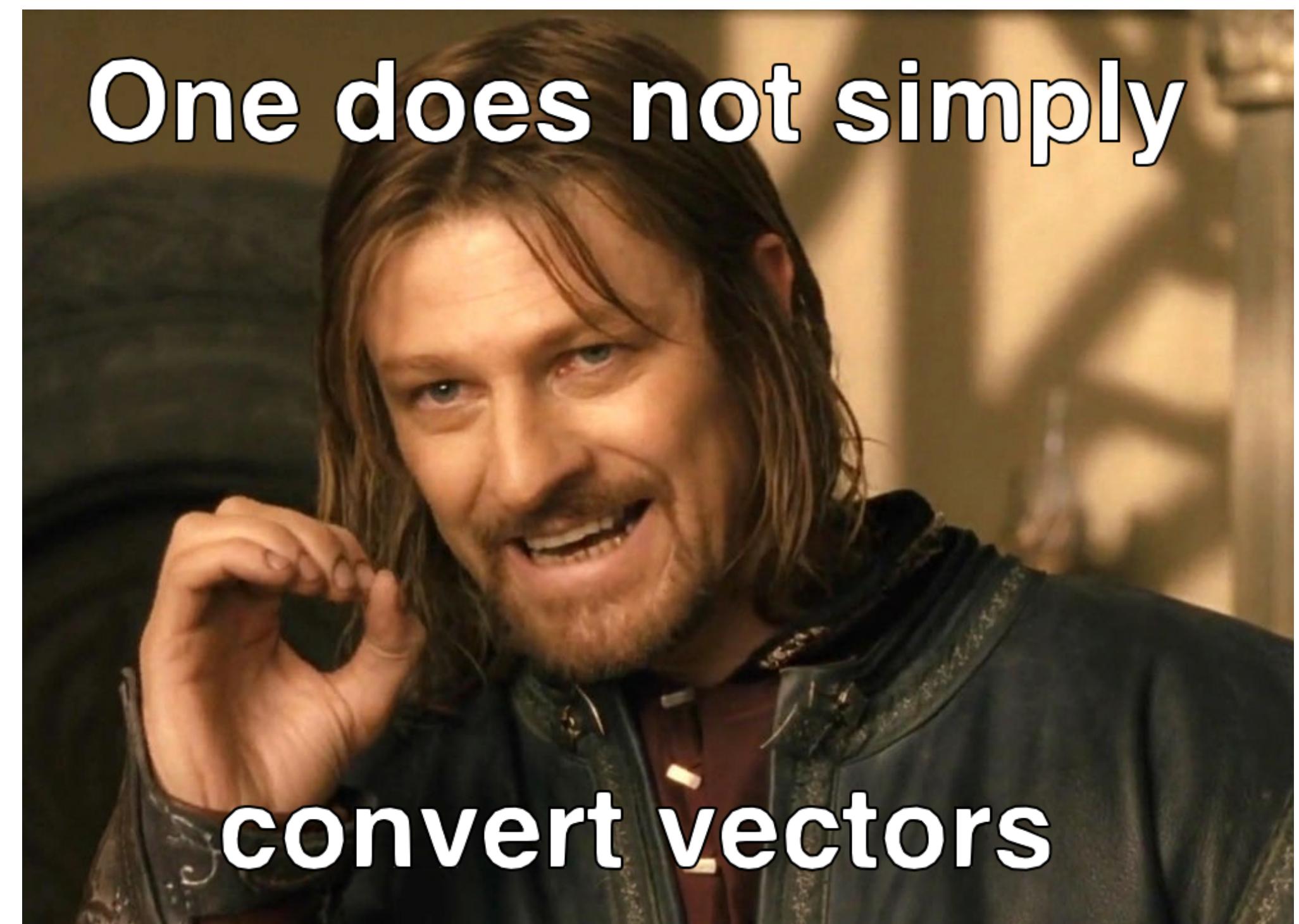
- If either operand is of scoped enumeration type no conversion is performed, and the expression is ill-formed if the types do not match.
- If either operand is a `vector<T,X>`, vector truncation or scalar extension is performed with the following rules:
 - If both vectors are of the same length, no dimension conversion is required.
 - If one operand is a vector and the other operand is a scalar, the scalar is extended to a vector via a Splat conversion (4.9) to match the length of the vector.
 - Otherwise, if both operands are vectors of different lengths, the vector of longer length is truncated to match the length of the shorter vector (4.10).



Difference: Implicit Dimension Conversion

Reason: Lots of vector types!

- C++ Implicit conversion sequences have 3 components
 - Lvalue Transformation
 - Promotion/Conversion
 - Qualification Adjustment
- HLSL adds 4th component for dimension adjustment
 - Extends or truncates dimensionality for scalar, vector or matrix types



New Problem: Moar Conversions!

- Immediately complicates overload scoring...
 - 3 ranks turns to 9

Conversion	Category	Rank	Subclause
No conversions required	Identity		
Lvalue-to-rvalue conversion	Lvalue Transformation	Exact Match	4.1
Array-to-pointer conversion			4.2
Function-to-pointer conversion			4.3
Qualification conversions	Qualification Adjustment	Promotion	4.4
Integral promotions			4.5
Floating point promotion			4.6
Integral conversions	Conversion	Conversion	4.7
Floating point conversions			4.8
Floating-integral conversions			4.9
Pointer conversions			4.10
Pointer to member conversions			4.11
Boolean conversions			4.12

Conversion	Category	Rank	Reference
No conversion	Identity	Exact Match	4.1
Lvalue-to-rvalue			4.2
Array-to-pointer			4.12
Qualification	Lvalue Transformation	Promotion	4.12
Scalar splat (without conversion)			4.9
Integral promotion			4.5 & 4.13.1
Floating point promotion	Qualification Adjustment	Scalar Extension	4.6 & 4.13.2
Component-wise promotion			4.11
Scalar splat promotion			4.9
Integral conversion	Scalar Extension	Conversion	4.5
Floating point conversion			4.6
Floating-integral conversion			4.7
Boolean conversion	Scalar Extension Promotion	Conversion	4.8
Component-wise conversion			4.11
Scalar splat conversion			4.9
Vector truncation (without conversion)	Dimensionality Reduction	Conversion Extension	4.10
Vector truncation promotion			4.10
Vector truncation conversion			4.10

Difference: Best Match Resolution

Reason: Lineage from Cg

- In C++ all these cases would be ambiguous!
- DXC uses a scoring system
 - Disambiguates cases where “worse” conversions are present
 - Produces unintuitive results
 - Scores have a limit

```
1 int fn2(float, int) { return 1; }
2 int fn2(int, float) { return 2; }
3
4 int fn3(float, int, float) { return 10; }
5 int fn3(int, float, int) { return 20; }
6
7 export void call() {
8     float f;
9     fn2(f,f);    // ambiguous!
10    fn3(f,f,f); // float, int, float
11
12    int i;
13    fn2(i,i);    // ambiguous!
14    fn3(i,i,i); // int, float, int
15 }
```

Solution: Define ICS Rules

- Try and align with C++ as much as possible
 - Use C++'s better/worse/indistinguishable
- New behavior does not match DXC
 - DXC used a number-based scoring system
- New behavior has more ambiguous cases, but it also handles more cases (operator overloads & const-ness)
- Behavior changes that break compiling are preferred!



beanz Aug 21st at 8:21 PM

Today's chapter of HLSL of Horror:

HLSL can't resolve overloads of functions unless they are exact matches when the functions have more than 128 parameters.

<https://godbolt.org/z/EW9GKzfdE>

godbolt.org

[Compiler Explorer - HLSL](#)

```
int fn(int, int, int, int, int, int, int, int, int,  
      int, int, int, int, int, int, int, int, int,  
      int, int, int, int, int, int, int, int, int,  
      int, int, int, int, int, int, int, int, int,  
      int, int, int, int, int, int, int, int, int,  
      int, int, int, int, int, int, int, int, int,  
      int, int, int, int, int, int, int, int, int,
```

[Show more](#)

1 1 0



zeux Aug 21st at 8:27 PM

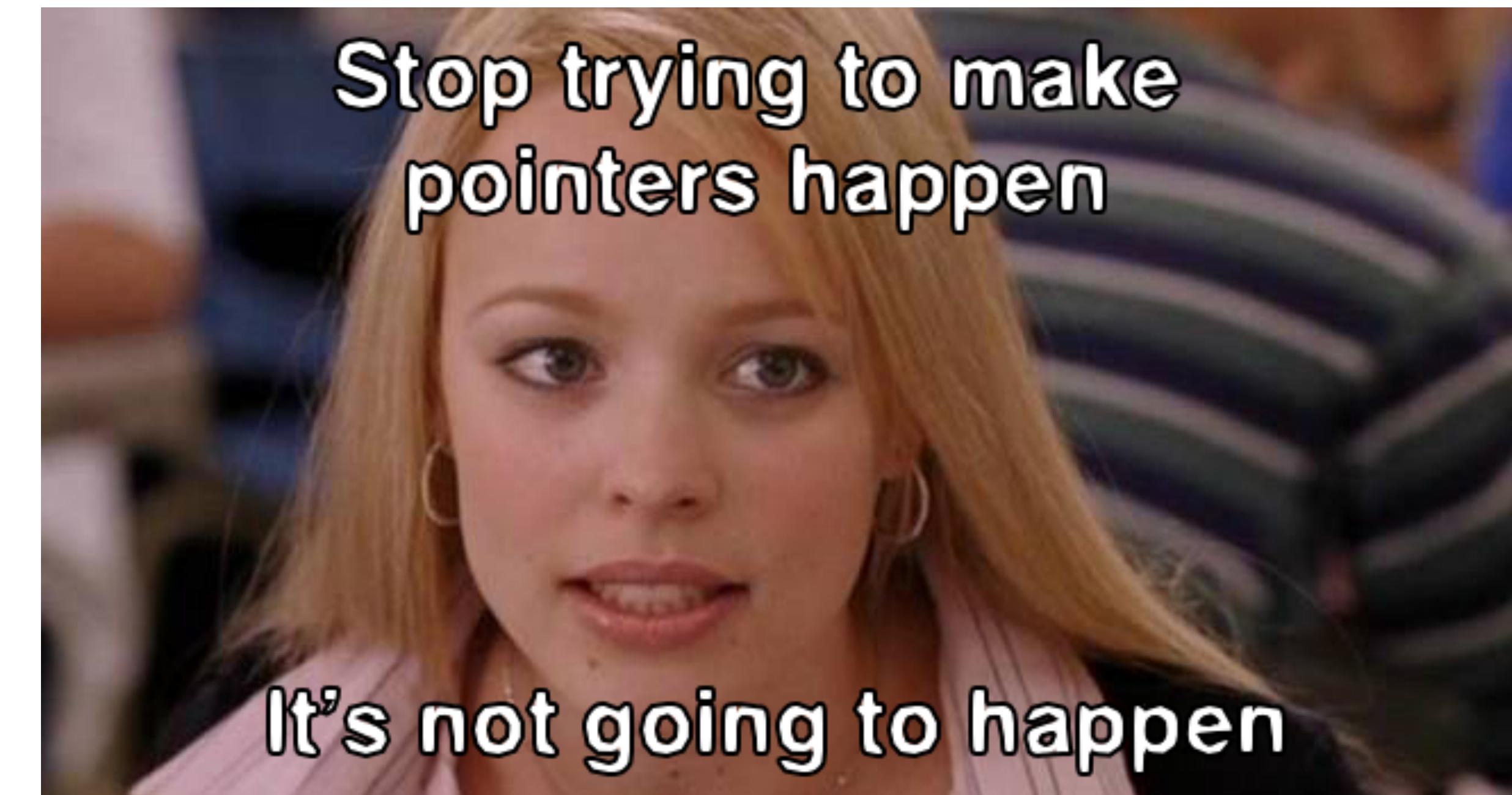
hot take, this is more defensible than anything else you shared here thus far!

2 3 0

Difference: No pointers!

Reason: part history, part complexity

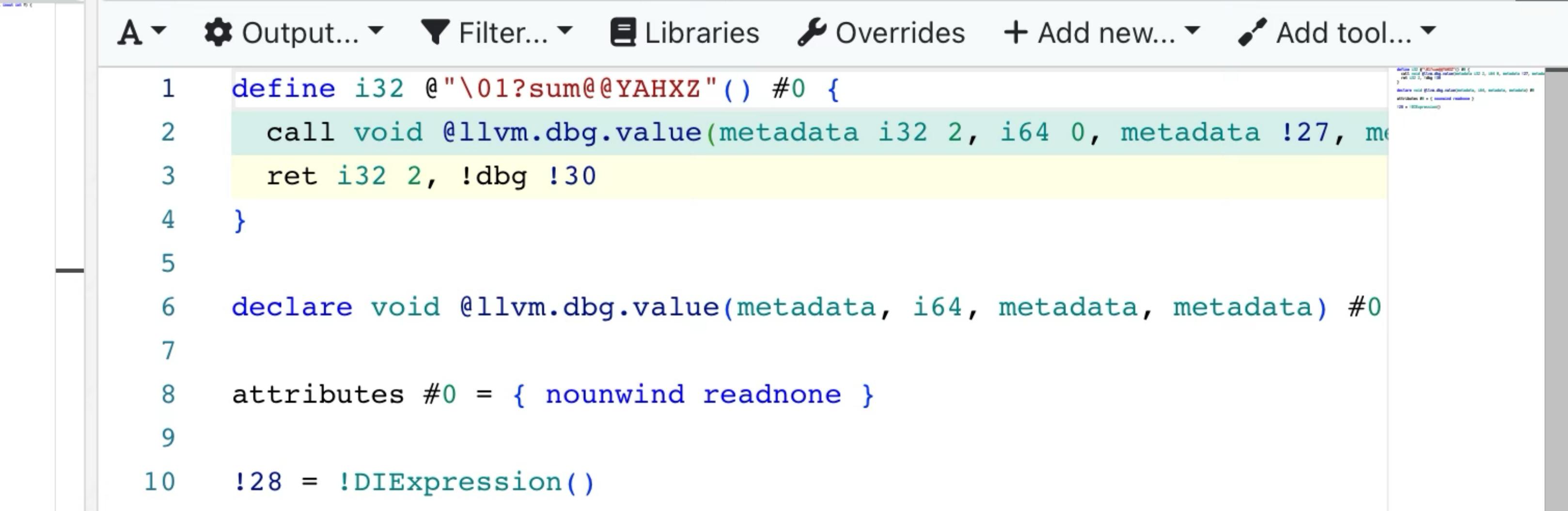
- Early GPUs didn't really have direct memory access
- Modern GPU memory has fragmented memory spaces
- Pointers and unified memory are becoming more common



Problem: Multiple Function Outputs

- HLSL's inout and out parameter keywords are still passed by value
- This normalizes addresses for parameters to thread-local memory

```
1 void fn(inout int X, inout int Y) {  
2     Y = 2;  
3     X = 1;  
4 }  
5  
6 export int sum() {  
7     int Val = 0;  
8     fn(Val, Val);  
9     return Val;  
10 }
```



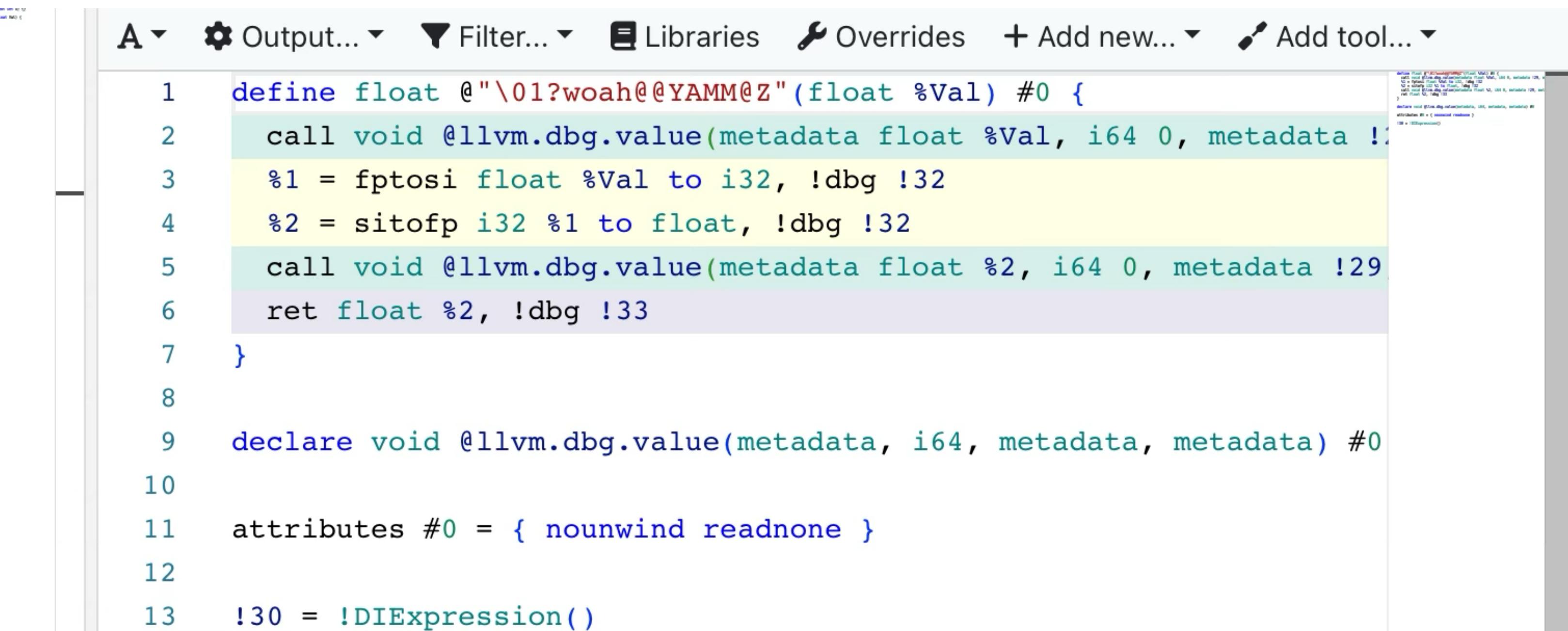
The screenshot shows a debugger interface with two panes. The left pane displays the original HLSL code. The right pane shows the generated LLVM assembly code. The assembly code includes a call to `@llvm.dbg.value` for each parameter and a declaration of the same function.

```
A▼ Output...▼ Filter...▼ Libraries Overrides + Add new...▼ Add tool...▼  
1 define i32 @"\01?sum@YAHZXZ"() #0 {  
2     call void @llvm.dbg.value(metadata i32 2, i64 0, metadata !27, me  
3     ret i32 2, !dbg !30  
4 }  
5  
6 declare void @llvm.dbg.value(metadata, i64, metadata, metadata) #0  
7  
8 attributes #0 = { nounwind readnone }  
9  
10 !28 = !DIExpression()
```

Complication: Outputs Can Cast?!?!

- HLSL's output parameters can also have cast

```
1 void silly_trunc(inout int X) {}  
2  
3 export float woah(float Val) {  
4     silly_trunc(Val);  
5     return Val;  
6 }
```



The screenshot shows a debugger interface with assembly code. The assembly code is as follows:

```
1 define float @"\01?woah@@YAMM@Z"(float %Val) #0 {  
2     call void @llvm.dbg.value(metadata float %Val, i64 0, metadata !28)  
3     %1 = fptosi float %Val to i32, !dbg !32  
4     %2 = sitofp i32 %1 to float, !dbg !32  
5     call void @llvm.dbg.value(metadata float %2, i64 0, metadata !29)  
6     ret float %2, !dbg !33  
7 }  
8  
9 declare void @llvm.dbg.value(metadata, i64, metadata, metadata) #0  
10  
11 attributes #0 = { nounwind readnone }  
12  
13 !30 = !DIExpression()
```

Casting Expiring Values

- cx-values are a spec-level construct only
- Vocabulary for casted temporary values
- x-value with a bound l-value
- may be initialized from an l-value or it may be uninitialized
- on expiration it “writes back” to the bound l-value
- it can optionally cast on both initialization and write back



Solution: HLSLOutArgExpr

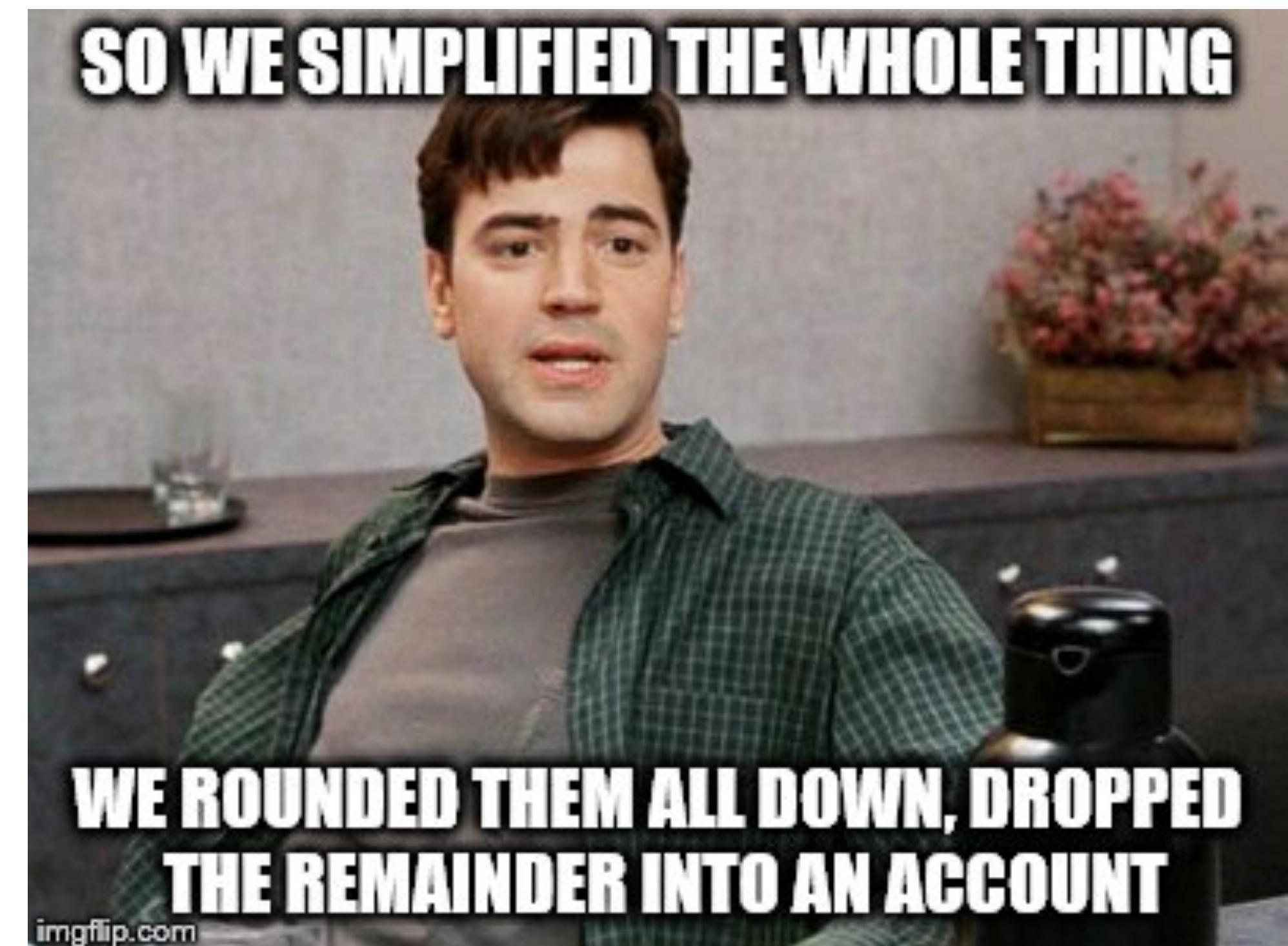
- New AST node to represent the cx-value
- Use Parameter ABI type information added for Swift
- Leverage call argument write back support added for Objective-C
- All casts are represented in the AST!

```
CallExpr 'void'  
|-ImplicitCastExpr 'void (*)(int &)' <FunctionToPointerDecay>  
| `DeclRefExpr 'void (int &)' lvalue Function 'Init' 'void (int &)'  
`-HLSLOutArgExpr <col:10> 'int' lvalue inout  
  |-OpaqueValueExpr 0xSOURCE <col:10> 'int' lvalue  
  | `DeclRefExpr <col:10> 'int' lvalue Var 'V' 'int'  
  |-OpaqueValueExpr 0xTEMPORARY <col:10> 'int' lvalue  
  | `ImplicitCastExpr <col:10> 'int' <LValueToRValue>  
  | `OpaqueValueExpr 0xSOURCE <col:10> 'int' lvalue  
  |   `DeclRefExpr <col:10> 'int' lvalue Var 'V' 'int'  
  `-BinaryOperator <col:10> 'int' lvalue '='  
    |-OpaqueValueExpr 0xSOURCE <col:10> 'int' lvalue  
    | `DeclRefExpr <col:10> 'int' lvalue Var 'V' 'int'  
    `-ImplicitCastExpr <col:10> 'int' <LValueToRValue>  
      `-OpaqueValueExpr 0xTEMPORARY <col:10> 'int' lvalue  
        `ImplicitCastExpr <col:10> 'int' <LValueToRValue>  
          `-OpaqueValueExpr 0xSOURCE <col:10> 'int' lvalue  
            `DeclRefExpr <col:10> 'int' lvalue Var 'V' 'int'
```

Difference: literal types

Reason: Low-precision math is fast

- Not unique to HLSL!
 - WGSL calls these abstract numeric types
- Allow compile-time expressions to use higher precision
- AST-level literals are treated as 64-bit for constant evaluation
- Implicitly truncated to a target type... usually



Problem: Unresolved Types

- DXC's literal handling is extremely complicated!
 - Implementation had strange bugs
 - C++ isn't designed for multi-expression type inference

```
RWStructuredBuffer<float> Output;  
  
[numthreads(1,1,1)]  
void CSMain(uint x : SV_DispatchThreadID) {  
    Output[x] = x ? (x ? 1.0 : 2.0) : 2.0 * (x ? 1.0 : 2.0);  
}
```

```
RWByteAddressBuffer buffer : register(u0, space0);  
  
[numthreads(1, 1, 1)]  
void main() {  
    buffer.Store4(0, 1.xxxx);  
}
```

Solution: HLSL 202x

Bridge to Clang

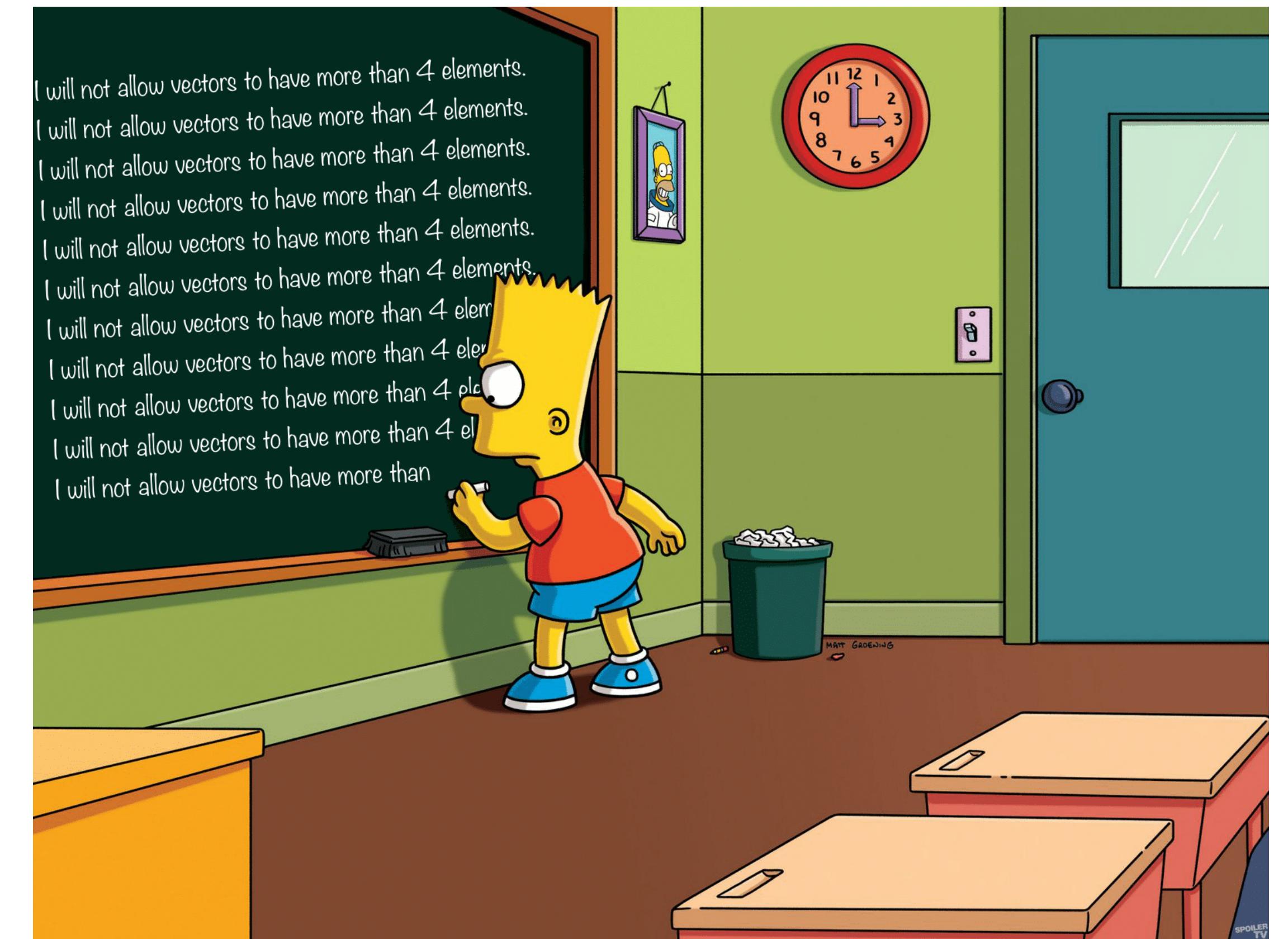
- Without a specification we can't make Clang match
- When possible we're implementing breaking changes in DXC
- Fix bugs in DXC and not introduce complexity or debt to Clang



Problem: Constraints on Templates

Reason: History

- HLSL for years has had template-like syntax for types
- Some built-in types have constraints on valid template arguments
 - vectors must have arithmetic or boolean type
 - vector size must be ≤ 4
 - Resource types have rules for different types



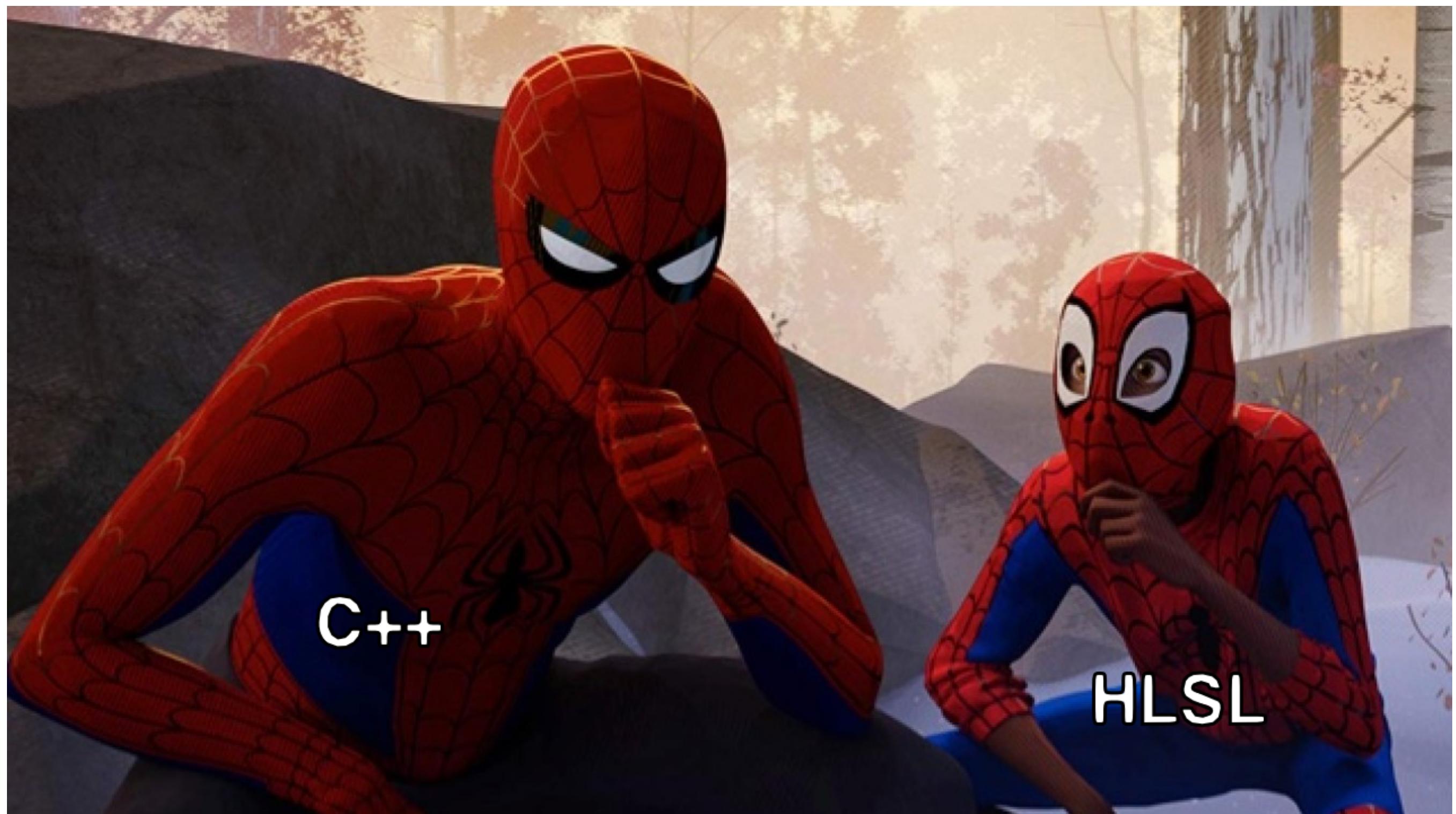
Solution: A Concept of a Plan

- Built-in types are injected into the AST
- Inject a concept declaration & concept specialization
- Avoids needing to add custom checking for types in template instantiation
- Concepts are super useful for GPUs!



Path Forward

- Diligence in writing a spec and docs
- Creative about how we achieve source compatibility
- Considerate about the future of the language
- Open and engaging with our users



Questions?



