KONGU ENGINEERING COLLEGE, PERUNDURAI, ERODE- 638 060

EVEN SEMESTER 2022-23

CONTINUOUS ASSESSMENT TEST I – March 2023

Regulation 2020

Key

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Consider the fo	llowing relation	sume tha	t the schei	ma of A U	B is the s	same as t	that of A. (GATE)	[CO1,K3
Table	A		Table l	В			ele C	
	Age	Id	Name	Age	Id	Phone	Area	
15 Shreya		25 98	Rohit	24 40 20	10 99	2200 2100	02 01	
Id Name 12 Arun 15 Shreya 99 Rohit 25 Hari 98 Rohit	Age 60 24 11 40 20 given relation	al alge	ole bra expr	ession wil	ne Ar 02 02	2 2 2		
F	Result of AUB Id Name 12 Arun 15 Shreya 99 Rohit 25 Hari 98 Rohit The result of g	15 Shreya 24 99 Rohit 11	15 Shreya 24 99 Rohit 11 Result of AUB will be following tab 10 Name Age 12 Arun 60 15 Shreya 24 15 Shreya 24 16 Hari 40 17 Hari 40 18 Rohit 20 19 Rohit 20 10 Name Age 11 Arun 60 12 Arun 60 13 Shreya 24 14 Arun 60 15 Shreya 24 15 Hari 40 16 Rohit 20 17 Rohit 20 18 Rohit 20 18 Rohit 20	15 Shreya 24 25 Hari 98 Rohit 99 Rohit 99 Rohit 11 98 Rohit 99 Rohit 12 Arun 60 15 Shreya 24 99 Rohit 11 25 Hari 40 98 Rohit 20 The result of given relational algebra expr	15 Shreya 24 25 Hari 40 98 Rohit 20 99 Rohit 11	15 Shreya 24 25 Hari 40 99	15 Shreya 24 99 Rohit 11 25 Hari 40 99 2100	15 Shreya 24 99 Rohit 11 25 Hari 40 99 2100 01

20

11

20

Rohit

Rohit

98

1.6

9.9

99

2200

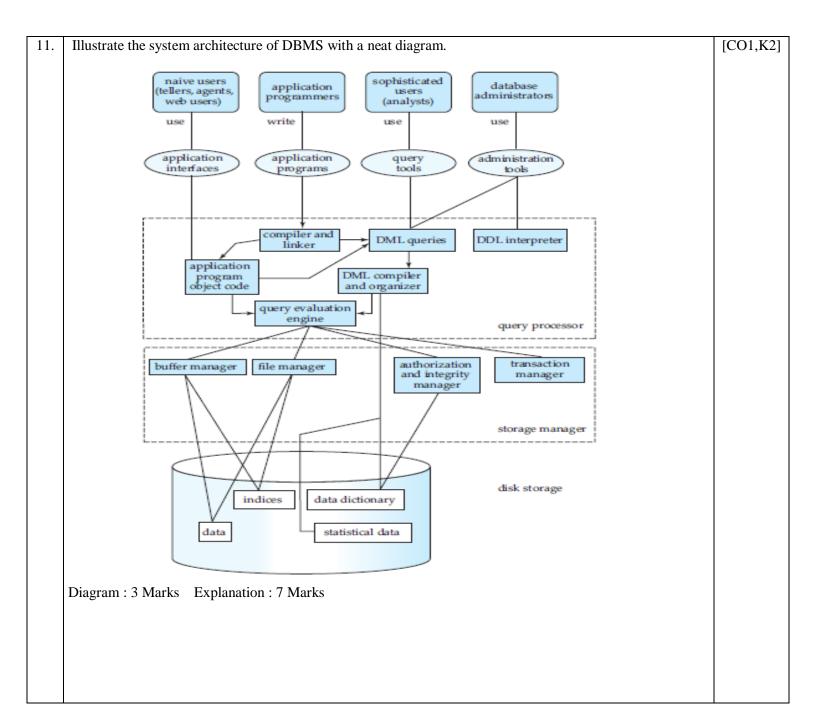
2100

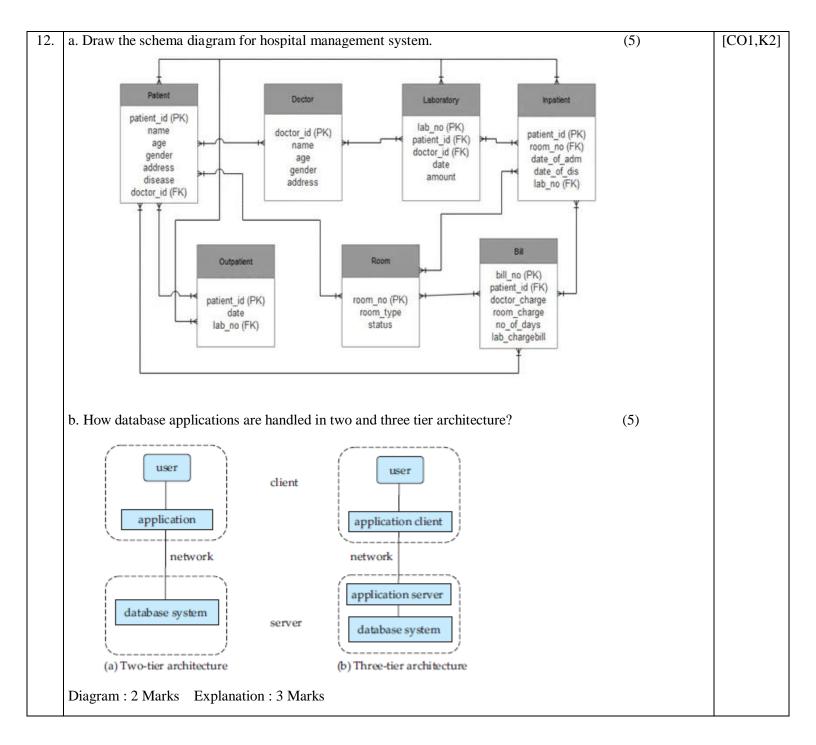
2100

82

61

Where P; (GATE) π _{a1, a2,, an} (σ _P (r ₁ ×r ₂ ×r ₃ ×r _m)) 8. Consider the relations r ₁ (P, Q, R) and r ₂ (R, S, T) with primary keys P and R respectively. The relation r ₁ contains 3000 tuples and r ₂ contains 2500 tuples. Find the maximum size of the join r ₁ ⋈ r ₂ . (GATE) 2500	
Select SID from supplier where rating > 10; 7. Convert into relational algebra for the following SQL query. Select distinct a₁, a₂,, aₙ From r₁, r₂,, rՠ Where P; (GATE) π a1, a2,, a⋒ (σp (r1×r2×r3×rm)) 8. Consider the relations r₁(P, Q, R) and r₂(R, S, T) with primary keys P and R respectively. The relation r₁ contains 3000 tuples and r₂ contains 2500 tuples. Find the maximum size of the join r₁⋈ r₂. (GATE) 2500 9. Consider the following relational schema: (GATE) COURSES (cno, cname) STUDENTS (rollno, sname, age, year) REGISTERED FOR (cno, rollno) The underlined attributes indicate the primary keys for the relations. The 'year' attribute for the STUDENTS relation indicates the year in which the student is currently studying (First year, Second year	
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COURSES (cno, cname) STUDENTS (rollno, sname, age, year) REGISTERED FOR (cno, rollno) The underlined attributes indicate the primary keys for the relations. The 'year' attribute for the STUDENTS relation indicates the year in which the student is currently studying (First year, Second year	2,K3]
SELECT year, min(age) FROM STUDENTS GROUP BY year;	2,K3]
10. A table T1 in a relational database has the following rows and columns: (GATE) Rollno marks 10 2 20 3 30 4 Null The following sequence of SQL statements was successfully executed on table T1. Update T1 set marks = marks + 5 Select avg(marks) from T1 What is the output of the select statement? Rollno marks 1 15 2 25 3 35 4 Null	2,K3]
$PART - B (3 \times 10 = 30 \text{ Marks})$	





13.	Consider the following relational database	[CO1,K3]					
13.		[CO1,K3]					
	emp (empname, street, city)						
	works(empname, compname, salary)						
	company(compname, city)						
	manager (empname, managername)						
	Give an expression in the relational algebra to express each of the following queries:						
	a. Find the names of all employees who work for First Bank Corporation.						
	Πempname (σcompname = "First Bank Corporation" (works))						
	b. Find the names and cities of residence of all employees who work for First Bank Corporation.						
	Πempname, city (emp \bowtie (σcompname = "First Bank Corporation" (works)))						
	c. Find the names, street address, and cities of residence of all employees who work for First Bank Corporation and earn more than \$10,000.						
	Hempname , street, city ($\sigma_{(compname)} = \text{``First Bank Corporation''} \land salary > 10000) works employee)$						
	d. Find the names of all employees in this database who live in the same city as the company for which they work.						
	Петрпате (employee works company)						
	e. Find the names of all employees who live in the same city and on the same street as do their managers.						
	$T \leftarrow \Pi_{\text{street,city}} \left(\sigma_{emp.empname=manager.managername}(employee^{\bowtie} manager) \right)$						
14.	Consider the following relational database	[CO2,K3]					
	emp (empname, street, city)						
	works(empname, compname, salary)						
	company(compname, city)						
	manager (empname, managername)						
	Give an expression in SQL for each of the following queries						
	i)Modify the database so that Ram now lives in Delhi.						
	Update emp set city='Delhi" where empname='Ram';						
	Opuate emp set city— Deim' where emphame— Ram;						
	ii) Find the names and cities of residence of all employees who work for First Bank Corporation. Select empname, city from emp natural join works where compname=' First Bank Corporation';						
	iii) Find all employees in the database who earn more than every employee of small bank corporation. select empname from works where salary > all (select salary from works where compname = 'Small Bank Corporation');						
	iv) Find all employees who earn more than the average salary of all employees of their company. select empname from works T where salary > (select avg (salary) from works S where T.compname = S.compname);						
	v) Remove the company relation from the database. drop company;						
	•						