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- Improves on $\mathrm{poly}(\log n)$ update time of earlier works such as [Andoni, Krauthgamer, Onak '10]

- **Theorem**: For $0 < p < 2$, can approximate $F_p(x)$ up to $1 \pm \varepsilon$ using **optimal** $O(\varepsilon^{-2} \log n)$ bits of space and $O(\log n)$ **update time**

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