# Our Results

- **Theorem**: Given $t$ and $r$, there is a streaming algorithm which can compute $\hat{x}[i]$ such that for $\alpha \leq 1$

- Obtained by derandomizing [Minton and Price '14]

- The algorithm uses $O(tr \log(n) + \log^2 n)$ bits of space

- $O(r \log n)$ update time

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# How to construct Sketching Matrices?