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- **Theorem**: For $p > 2$, there is an algorithm using **optimal** $\tilde{O}(n^{1-2/p})$ bits of space and an update time of $O(1)$ to approximate $F_p(x)$ up to constant factors

- Improves on poly$(\log n)$ update time of earlier works such as [Andoni, Krauthgamer, Onak '10]

- **Theorem**: For $0 < p < 2$, can approximate $F_p(x)$ up to $1 \pm \varepsilon$ using **optimal** $O(\varepsilon^{-2} \log n)$ bits of space and $O(\log n)$ update time

- Valid only for $\varepsilon < 1/n^c$

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