Robust Subspou Clusterins

Problem formulation: there ove L subspace, {S,,..., S} ER" and have dimensions di,..., de {L, di, Si} are unknown.

given $Y \in \mathbb{R}^n$, |Y| = N, Y = Y, $UY_2 U - ... UY_L$ and Y_L is collection of N_L vectors from Selor close to)

Croal: Obtain S,, --, Se Rom Y.

Egyn: y = x +3, [there are N such vertors]

such that = = S_ka_k for some k, 3 indefendant "wise"

Preliminary accomplishs: @ Enzll= = 0211×112, 0<0+<1

© nar de € Co N logN

Challenges: @ how "far" are cubepous Se?

6 How big ore # sample Ne?

Def : [Principal cos (θ^i) = mox mox $\frac{u^iv}{|uv|^2||vv||} = \frac{v_i^iv_i^2}{|uv|^2||vv||} = \frac{v_i^iv_i^2}{|vv|^2||vv||} = \frac{v_i^iv_i^2}{|vv|^2} = \frac{v_i^iv_i^2}$

· Cirqular volus of UTM

Def": Normolized affinity: aff (c, s') = \(\sum_{iz1}^{\text{Normolized}} \text{ Lind,d'}\)

- · off = 0 > mand early brob > Ut = V (Chan-wise)
- . off ≈1 > hard prob, > V & V (chor-wise)

Def": Compline Deneity: Pe of subspace Se in defined as for Ps = Ns. It Ps <1, no hope of recovering CIC

NORMALIZED MODEL: Y = X + Z, Y, Z FIR N

accome that $||x||_2 = 1 \Rightarrow ||y_1||^2 \approx \sqrt{||y_2||^2} ||x_1||_2^2$ (in expendation)

Zii ~ N(0,07/n)

SOLUTION: The Shorse Substan Chestering Scheme!

- A. Compute affinity matrix encoding similarities 4w sauply, and tonctive weighted graph W
- B. Construct cheters vein checked chistoring
- C. Apply PCA to each chucker

INTURION (for 2=0 case): Exposes x; as a share lin-word
of x;, j \(\frac{1}{2} \). (must hold since x; would be from come six as x;)

- So, solve win IBII, Sit X; = XB, B; =0.
 BER
- stack B's as columns of matrix B, W= 181+181*

For noisy case, $x_i = xB \iff y_i - 3_i = (Y-Z)B$ $= y_i = YB + (z_i - z_B) \implies \text{noise! (per hurbation)}$ Use LASCO

Use LASSO

win = 1 y; - YBII + AIBI, S.t B; = 0

BERN Z Y Y; - YBII + AIBI,

Performance Methics: Deff: False disovery, (i,i) obeying Bij #0 is take disovery it yi and yi don't belong to same ss

Defi: True disovery: (i,i) obegin Bii +0 is T.D it yi, y; E came Cs.

- It there are no F.D's, then subspace detection property holds.

- then B is permutation similar to block-diagonal matrix.

Data - dependent regularization: (consider mise hu one)

nininizer of \mathcal{D} - $\hat{\mathcal{B}}$ and \mathcal{U} $\hat{\mathcal{B}}_{eq}$ be solution with equality combainty. ie solving \mathcal{D} with $\lambda \to 0^+$. Then

$$\frac{1}{2} \| x - x \hat{\beta} \|_{1}^{2} \leq K (\hat{\beta}, \lambda) \leq K (\hat{\beta}_{eq}, \lambda) = \lambda \| \hat{\beta}_{eq} \|_{1}$$

@ If II PII. E [0.5d, 0.8d), |xc-x\beta|12 > C since |\beta|5 0.8d.

B
$$||\hat{B}_{ev}||_{1} = O(\sqrt{a})$$
 $[dx \frac{1}{\sqrt{a}}]$ (it also has no false disoveries!)
 $||\hat{B}_{ev}||_{1} = O(\sqrt{a})$ and if $||\hat{A}_{ev}||_{2} = O(\sqrt{a})$

THEORETICAL RESULTS:

1A Assurption: . offinity. Se obeys affinity condition it

max off (S_R, S_L) $\in \frac{K_0}{\log N}$ K_0 - fixed combate $R: R \neq L$

• Se obegs campling andition if $P_a \ge p^+$ p^- - fixed constant

B Main Results:

Theorem (No talse discoveries): Assume that the subspace attached to it colored beyon A-c, S.C above and or in a small numerical contact. In Algorithm 2, take $Z=2\sigma$ and $f(t) \neq \frac{\sigma}{2t}$. Then, with probability at half $1-2e^{in}-6e^{ix_2d(i)}-e^{-iN(i)A(i)}-\frac{23}{N^2}$ there is no talse discovery in the ith cot of B.

Theorem (Many true disoveries): Under assumptions of theorem 3.1, with f(t) also obegins $f(t) < \frac{1}{2}$ for some also, with prob. at least $1-2e^{-Y_1 n}-6e^{-y_2 d(i)}-e^{-(Pri) z d(i)}-\frac{23}{N^2}$ there are at least $\frac{G_1}{G_2}$ free disoveries in the ith column.

Proofs:

root of theorem 3.1:

- 118411,

Lemma 8.2: Let val(cfep1) be optimal valve of Algo. 2 (data driven regularization] with Zz 20. Assume P.>p*. Then, work

10 Jan (val (step) < 2 Ja,

· upper bod hold w.b > 1-exin -e-12d,

· lover bnd holds wip > 1-e-43d1-10/N2

Veily lem. 8.2, \(= f(1811) > \(\frac{120}{348711} \) \(\frac{5}{810}, \)

now, need to show that thou are no fake disoveries. First, let us use years you fully and so step 1 is of the form

if there are no talse disoveries, above is equivalent to

nin = 1/y-Y"B" | + > | B" | . W org nin be B" |

now, it cuffices to show Bu obey, I Y'D' (y-y'DBU) Im < > , +1+1

Lemma 8.6: Fix A ERdEN and TC St,..., N3. Suppose than in a soll x4 to

win 1 11y - Ax112 + >11x11, Sit x7c=0

obegins 1 ATC (y-Ax*) 110 (x) then any optimal x to min \(\frac{1}{2}\) \left(+ \lambda 11x\).

now, $[y^{(a)}](y-y^{(i)}\hat{g}^{(i)})^{\frac{1}{2}} = x^{(a)}(y_{1}-y_{1}^{w}\hat{g}^{(i)}) + x^{w}(y_{1}-y_{1}^{w}\hat{g}^{(i)})$ + Z(1) (y, - Y, B(1)) + Z(1) (y, - Y, B(1))

Lemma 8.1: Let A ERAXN, be analyis with column complete uniformly of roudon from the unit ofhere of Rd', we will wave veter complete v. AR from unit ofhere of Rdz and independent of A. ZER' was be deterministic notion. We have up 1-2-2N, To

1 A Si wil a & Thosa log h VI Ils

Term 1: || x(a) (y - y) &) || , A = x(a) + ERdexM

W= (y, - y1) Bas ER , wb 1- 4 N2

and $11x^{(a)^T}(y_{11}-y_{11}^{(0)})_{a}^{(b)})_{a}^{(b)} \in \lambda$ by att $(c_1, c_2) \stackrel{\Delta}{=} \lambda I$,

from lum 8.5, 11y- Yus Brills & CASa,

and Arm in 8.5, 1xw (y2-Y") B")110 & in 5 = I2

Term 3: 11 zw (y, - Y, Bull o C C > o (d, log N) = > I3

Term 4: || ≥w (y1 - y1) Bin) || € 62 \(\frac{\int_n}{n} \\ \text{\term 4} \) | = \(\frac{\int_n}{n} \)

2) we need (I, +Iz) >+ Ia+ Ia <> 2) it suffice to have

>> h 0

and veing $d_1 \leqslant \frac{n}{2n}$ coughtes proof

lemma 8:4: The projection of the recidual vector rzy-Y"B" onto either 5, 895, has visitorm orientation.