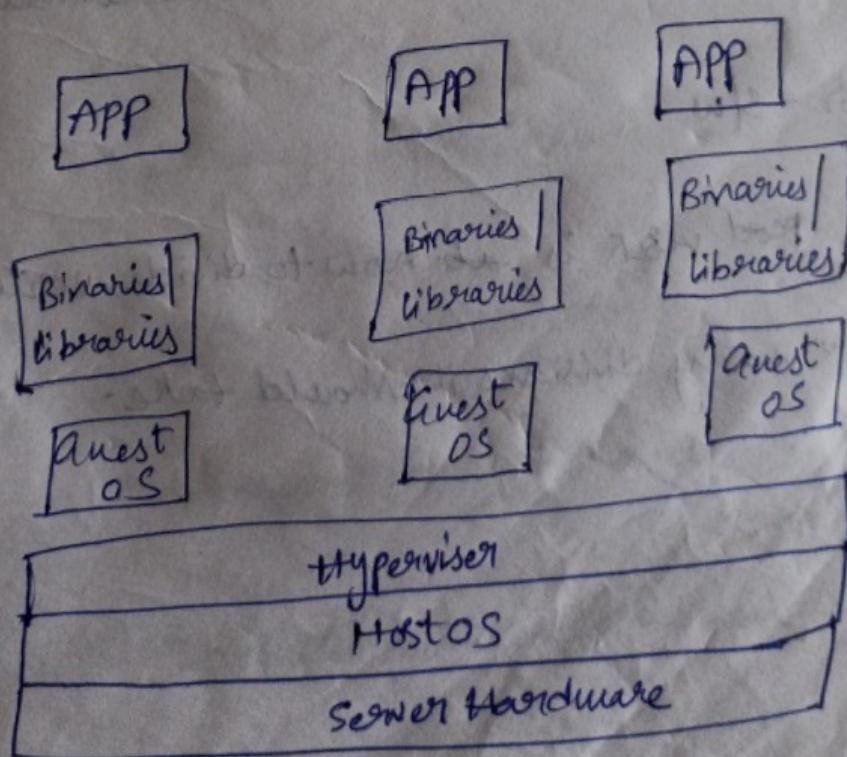


UNIT-11.

virtualization: It is a technique for creating a virtual platform of storage devices and the servers/OS.

- It helps the user make use of multiple machines sharing one single physical instance of any resource across the network of other users respectively using their machines.
- Moreover, virtualization technologies provide a virtual environment for not only executing applications but also for storage, memory, and networking.

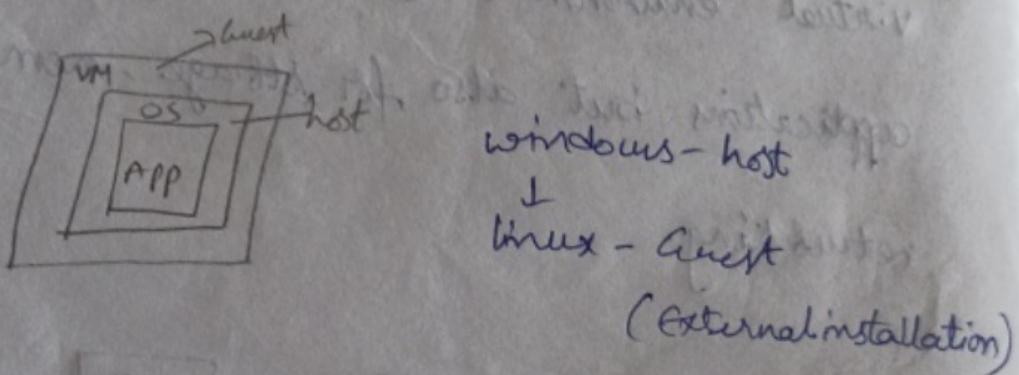


Virtualization.

* we can implement any services in cloud computing.
through virtualisation.

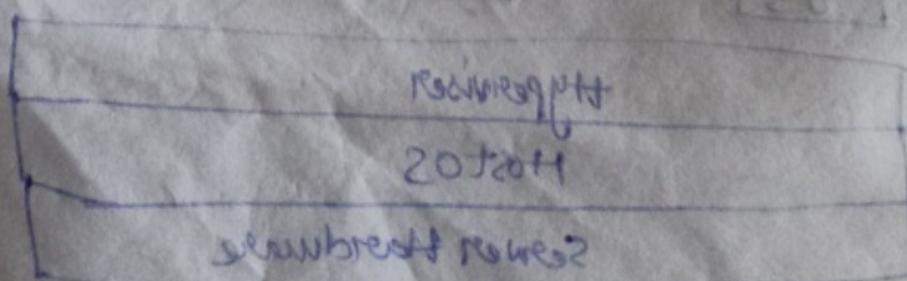
Ex I had a laptop, wanted to install Linux in
MacOS and to use both SW at a time
Now that laptop should be divided virtually
not physically.

This complete process is known as virtualisation.



* hypervisor - SW

This SW need work is ab. how to divide virtually
and how many divisions it should take.



• mitigation

Implementation levels of virtualisation:-

There are 5 levels of virtualisation. They are:-

- ① ISA (Instruction set Architecture level).
- ② Hardware abstraction layer
- ③ Operating system level
- ④ Library support level
- ⑤ Application level

ISA: It is used to run many legacy codes.

These codes runs on any VM using ISA.

user application level

library support level

operating system level

Hardware Abstraction level

Instruction set Architecture level

JVM

LxRun

virtual environment

VMware

BIRD

① ISA:- This level is an interface b/w S/W & H/W.

using this instruction S/W communicate with H/W.

→ when virtualisation is carried at this level,
we create an emulator.

↓ Level instant mapping

→ It receives all instructions from VM &
interprets them , then it maps those instruction
to the instruction of host machine.

→ This technique is simple but every instruction
from VM should be interpreted before mapping.
disadvantages are ;-

So time consuming is more

Performance is low .

(Q) HAL:

- In ISA level, performances goes down, due to interpretation of every instruction.
- To overcome that we map virtual resources with physical resources.
- When accessing of ^{HW} come to virtual resources, they forwarded directly to physical resources.
- Here every instruction is not interpreted, but we check whether instruction is privileged or non-privileged.
- Some memory instruction.
- If instruction is non-privileged, normal execution is done.
- If it is privileged, then ctrl is passed to VMM which handles them accordingly.
- ↓
Virtual machine monitor.

③ Operating system level:-

The layer is present b/w operating system and the applications is known as operating system layer.

* At this level, virtualization can be done

using isolated containers which act as real servers and are created on only one physical server.

④ Library support level:-

Virtualization at library level can be done simply by managing the APIs associated with the applications and the system.

Ex :- wine tool developed for allowing windows application over UNIX system.

⑤ Application level:-

At this level user application is virtualized as a virtual machine. It is also called as process level virtualization because OS considers each application as a process.

* In this, an application is made as virtual

machine and attached to the OS, as a layer which manages different virtual machines.

Ex:- Sand boxing, application isolation.

VIRTUALIZATION STRUCTURES/TOOLS AND MECHANISMS:

- * Before virtualisation, the operating system manages the hardware.
- * After virtualisation, a virtualization layer is inserted between the hardware and the operating system.
In such case, the virtualization layer is responsible for converting portions of the real hardware into virtual hardware.

Physical hardware
(separate app) → no of logical ones.
Simultaneously

- * Different operating systems such as Linux and windows can run on the same physical machine simultaneously.

- * Depending on the position of the virtualization layer, there are several classes of VM architectures.

- Hypervisor architectures
- para-virtualization architecture
- host-based virtualisation architecture.

* The hypervisor is known as VMM (Virtual machine monitor).

* They both perform the same virtualization operations.

Hypervisor:- (by using this we are able to create virtual h/w)
(multiple top-level multiplex called guest)

* A hypervisor is a h/w virtualization technique allowing multiple operating systems, called guests to run on a host machine.

Two types of hypervisor. They are:-

→ Bare metal hypervisor.

→ Hosted hypervisor.

Type-1:- Bare metal hypervisor:-

* Sits on the bare metal computer h/w like CPU, memory etc.

(partition top h/w into h/w)

* All guest operating systems are a layer of ~~are~~ above the hypervisor.

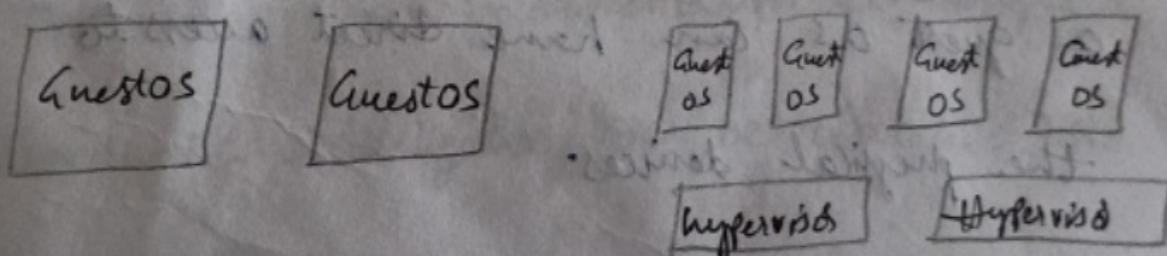
Ex:- CP/CMS (Control Program) (cambridge Monitor system) developed by IBM.

Type - 1 :- Hosted hypervisor :-

* In this one separate extra layer is installed known as Host OS.

(Host OS) $\xrightarrow{\text{hardware}}$ Hypervisor $\xrightarrow{\text{hardware}}$ Physical hardware

* Host OS is unaware of virtualization.



Hypervisor

Host OS

Hardware

Hardware

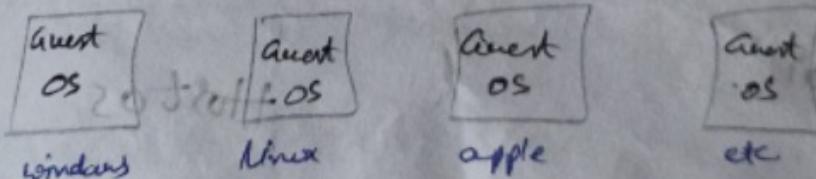
Type 1 (Bare-metal)

Type 2 (Hosted)

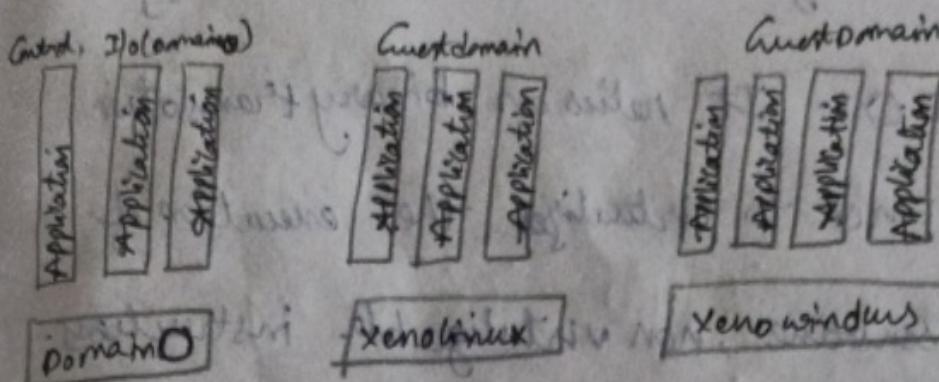
Example for Bare metal hypervisor is XEN architecture.

XEN ARCHITECTURE:

- * Xen is an open source hypervisor program developed by Cambridge University.
- * It makes possible to run many instances of an operating system or indeed different operating systems parallel on a single machine (the host).
- * It just provides a mechanism by which guest OS can have direct access to the physical devices.



Architecture of XEN:-



Xen (Hypervisor)

hardware devices

BINARY TRANSLATION WITH FULL VIRTUALIZATION,-

* Depending on implementation technologies,

hardware virtualization can be classified into

two categories

They are:

→ full virtualization

→ host-based virtualization.

- * Full virtualization does not need to modify the host OS. It relies on binary translation to trap and to virtualize the execution of certain sensitive, non-virtualizable instructions.
- * The guest OSes and their applications consist of non-critical and critical instructions.
- * In a host-based system, both a host OS and a guest OS are used. A virtualization software layer is built between the host OS and guest OS.

Paravirtualization :-

It is a type of virtualization where software instructions from the guest OS running inside a virtual machine can use "hypervcalls" that communicate directly with the hypervisor.

Virtualization of CPU:-

H/w virtualization:-

H/w virtualization refers to the creation of virtual versions of computers and operating systems.

- developed by intel and AMD for their server platforms
- And designed to improve the performance of the processor.

→ This term H/w virtualization is known as n/w-assisted virtualization.

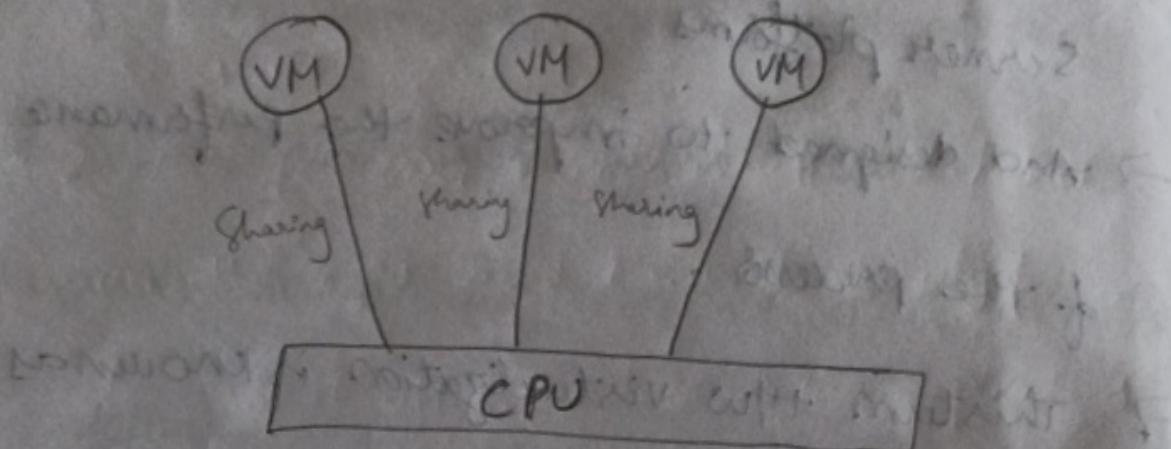
Advantages:-

- H/w Virtualization has many advantages because controlling virtual machines is much easier than controlling a physical machine.
- lower cost.
- decrease the quantity of rack space.
- reduce no. of servers.

CPU virtualization:-

Central processing unit virtualization:

- It allows a single CPU to be divided into multiple virtual CPUs for use by multiple VMs.



- All VMs acts as physical machines and distribute their hosting resources like having various virtual processors.
- sharing of physical resources takes place to each VM when all hosting services get request.
- Finally VMs get a share of single CPU allocated to them, being a single processor

acting as dual - processor.

Types of CPU virtualization:-

- ① s/w Based CPU virtualization.
- ② H/w - Assisted CPU virtualization.
- ③ virtualization and processor-specific behaviour.
- ④ performance implications of CPU virtualization.

① s/w Based CPU virtualization:-

- Application code gets executed on the processor and the privilege code gets translated first, and that translated code gets executed directly on the processor.
- The code that gets translated is very large in size and also slow at the same time on execution.

② H/w-Assisted CPU virtualization!

- Here Guest user uses a different version of code and mode of execution known as a guest mode.

→ Here there is no requirement for translation while using for h/w assistance.

→ System calls runs faster than expected.

③ virtualization & processor - specific behaviour

→ The VM still helps in detecting the processor models on which the system runs.

→ The processor model is different based on the CPU and the wide variety of features it offers, whereas the applications that produce the output generally utilize such features.

④ Performance implications of CPU virtualization

→ It adds the amount of overhead based on the workloads and virtualization used.

→ Any application depends mainly on CPU power waiting for the instructions to

get Executed first .

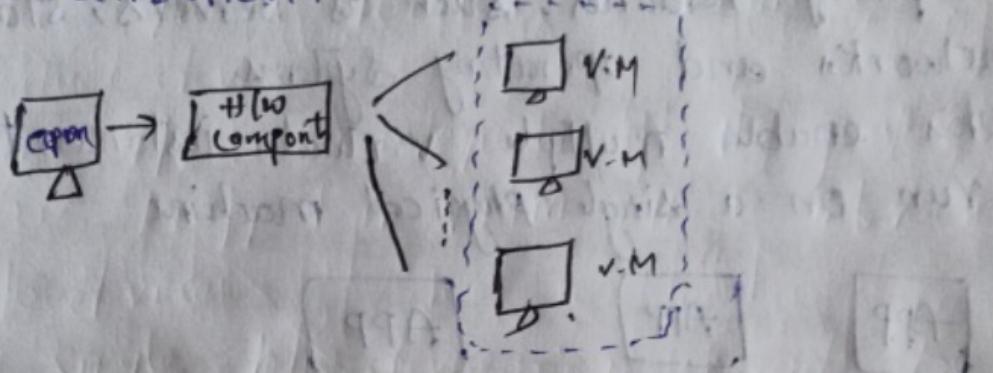
→ Such applications require the use of CPU virtualization that gets the command of executions that are needed to be executed first . *

→ This overhead takes the overall processing time and results in an overall degradation in performance and CPU virtualisation execution .

VIRTUALIZATION OF MEMORY

Virtualization:

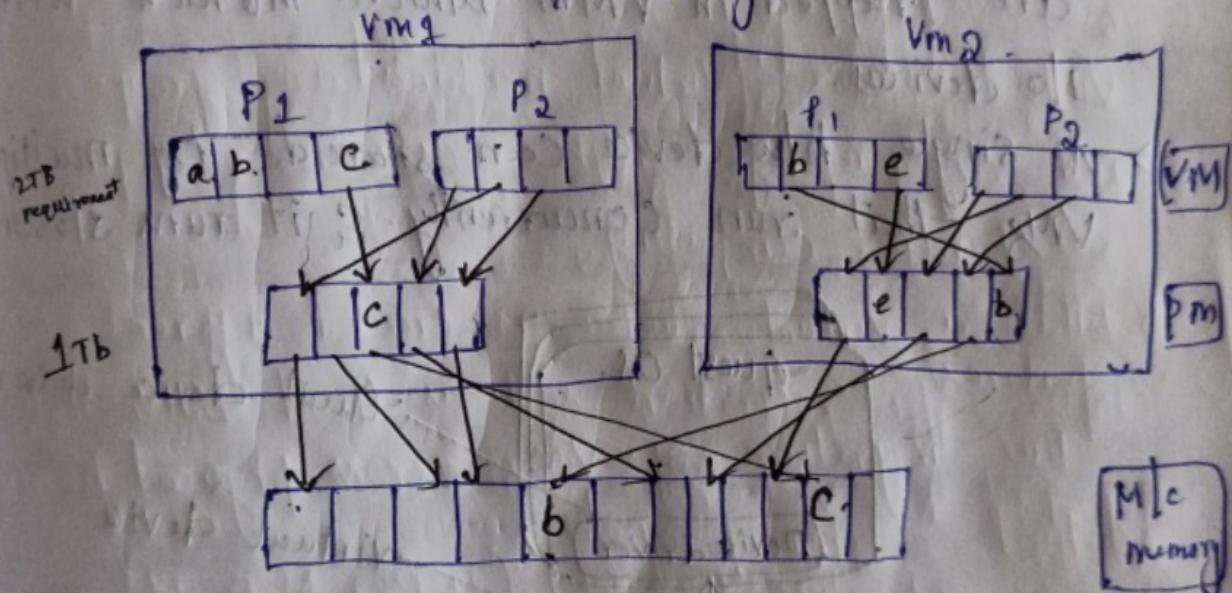
- The process ^{which} allows a computer to share its h/w components to multiple virtual machines (Virtual Environment). Is known as Virtualization.



Virtualization of Memory:

- Virtual memory virtualization is similar to VM support provided by modern OS.
- In a traditional execution the OS maintains mappings of VM to MM using page tables, which is a one-stage mapping from VM to MM.
- All modern x86 CPUs include a MMU and a Translation Lookaside buffer (TLB) to optimize VM performance.
- However virtual memory virtualization involves sharing the physical system memory in RAM and dynamically allocating it to the PM of VMs.

- That means a two-stage mapping process should be maintained by the guest OS and VMM respectively - 1) VM to PM 2) PM to MM.
- Furthermore MMU virtualization should be supported, which is transparent to the guest OS.
- The guest OS continues to control the mapping of VT to the PMA of VM_X. But the guest OS cannot directly access the actual memory.
- The VMM is responsible for mapping guest PM to actual M/c memory.



I/O VIRTUALIZATION:

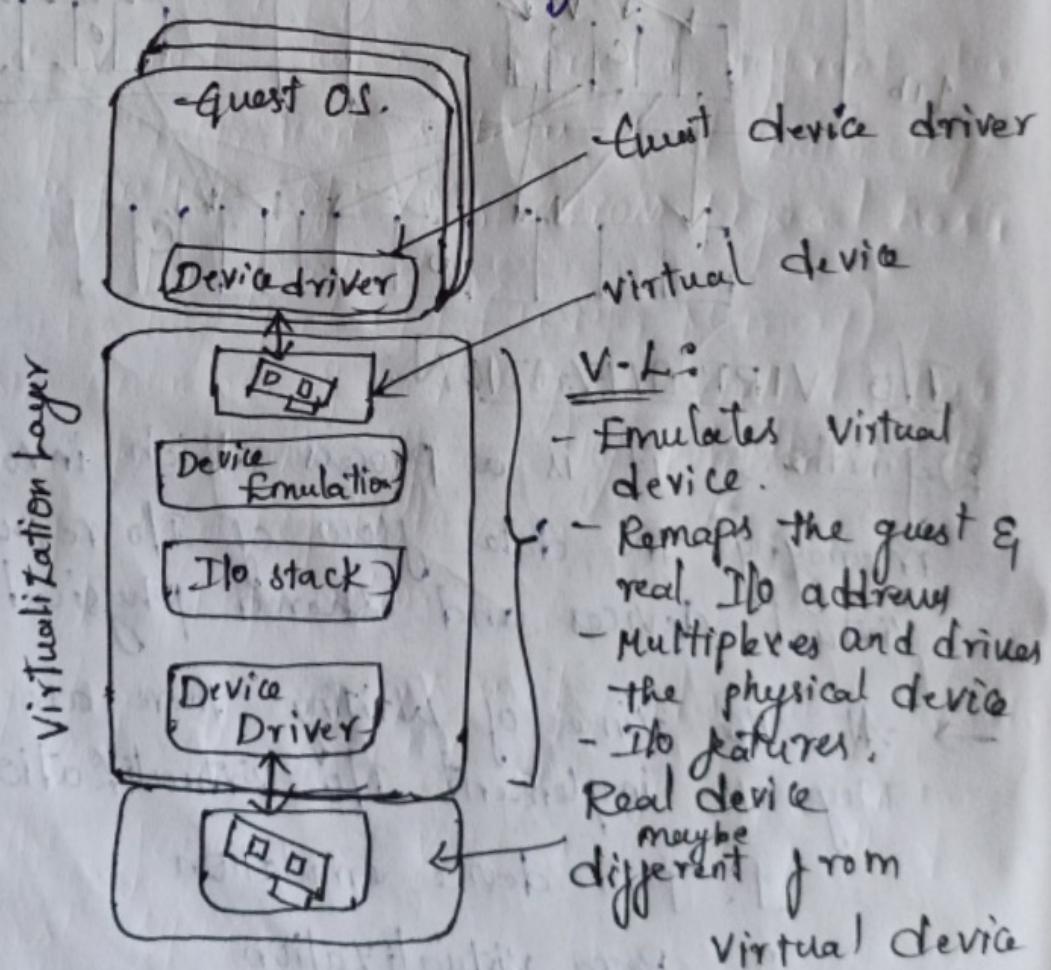
Definition: It is a process which involves managing the data flow of I/O requests b/w virtual devices and shared physical h/w.

- At the time of writing there are three ways to implement I/O virtualization:

1. Full device emulation
2. Para virtualization
3. Direct I/O

1. FULL DEVICE EMULATION

- It emulates the well-known, real-world devices
- All the functions of a device such as device enumeration, identification, interrupts & DMA are replicated in s/w.
- This s/w is located in VMM acts as virtual device.
- The I/O access requests of the guest OS are trapped in VMM which interacts with I/O devices.
- A single h/w device can be shared by multiple VMs that run concurrently & it runs slower



2. PARA VIRTUALIZATION:

- It consists of frontend driver and backend driver
- The frontend driver is running in Domain U and backend driver is running in Domain O.
- * Domain O: The guest OS which has control ability, is called Domain O.
- * Remaining are called Domain U.
- The frontend driver manages the I/O requests of the guest OSes.
- Backend driver is responsible for managing the real I/O devices and multiplexing the I/O requests data of guest OSes, different VM.
- Achieves better device performance than full device emulation.

3. DIRECT I/O:

- It allows virtual machines to access devices directly.
- It can achieve close-to-native performance without high CPU costs.

VIRTUAL CLUSTERS AND RESOURCE MANAGEMENT

Physical versus virtual cluster:

cluster: cluster is a group of servers, computers and other resources that act like a single system.

Physical cluster: It is a collection of servers (physical machines) interconnected by a physical n/w such as LAN.

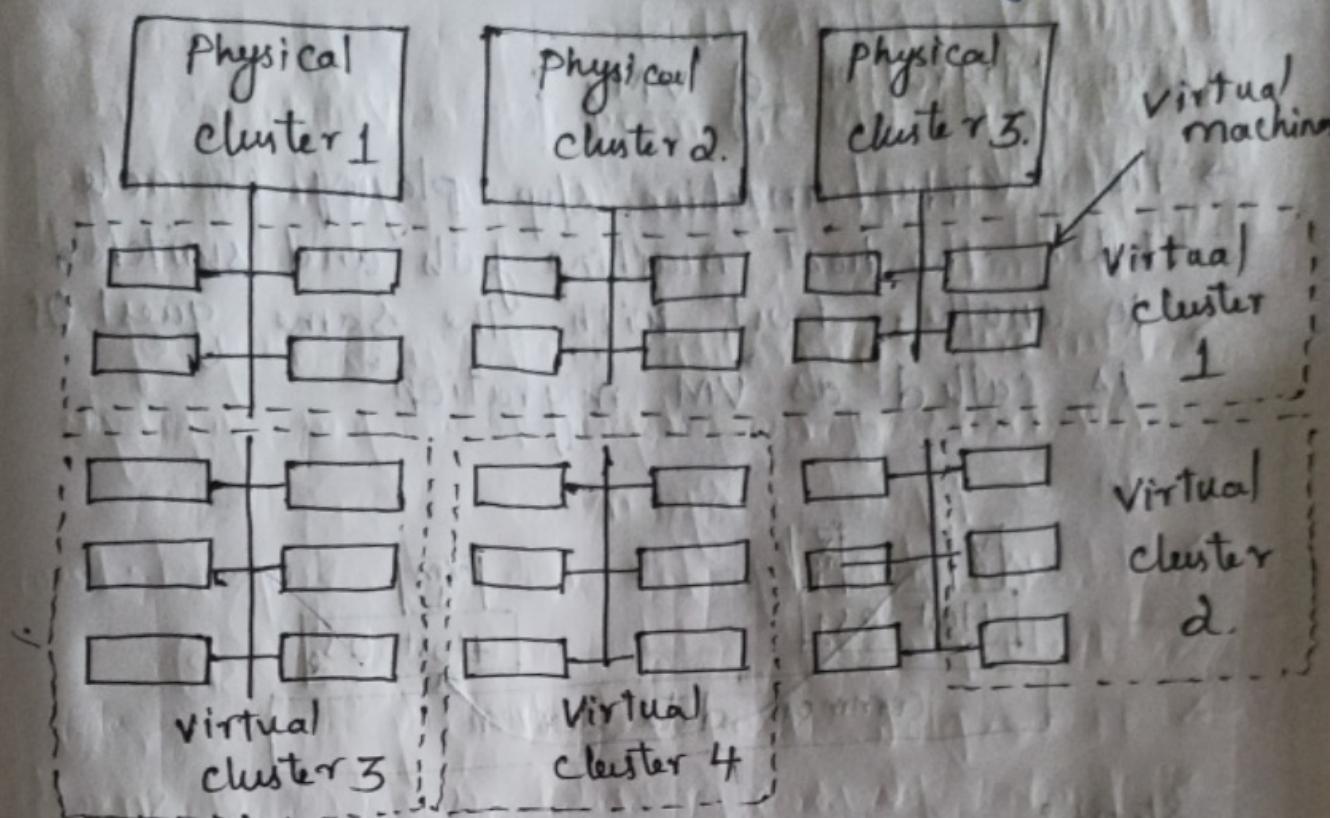
Virtual cluster: Virtual machines are grouped and configured for high performance computing (HPC). Parallel computing.

- When a virtual cluster is created, different cluster features can be used such as load balancing, live migration of VM's across physical hosts.

Virtual cluster properties:

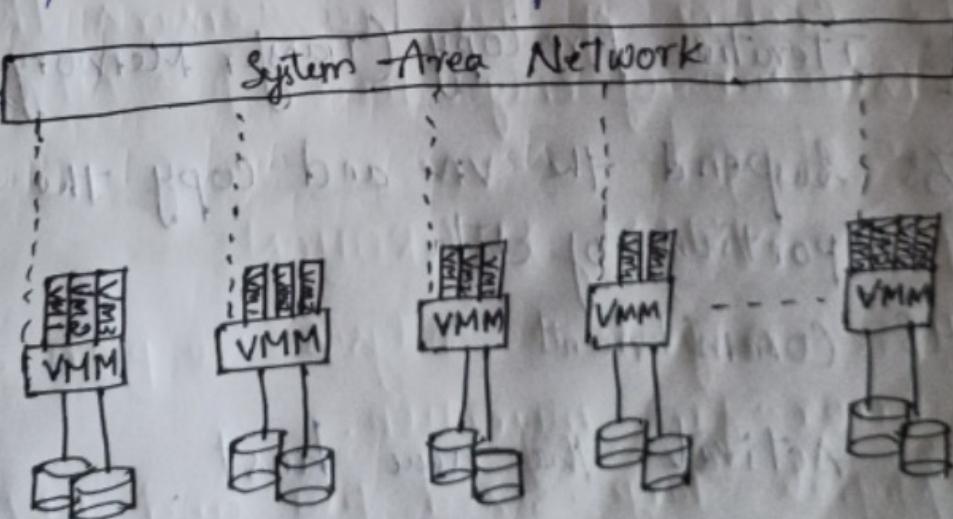
- The size (number of nodes) of a virtual cluster can grow or shrink dynamically, similar to the way an overlay n/w varies in size in a Peer-to-peer (P2P) n/w.
- The failure of any physical nodes may disable some VM's installed on the failing nodes. But the failure of VM's will not pull down the host system.

→ VMs can be replicated in multiple servers for the purpose of promoting distributed parallelism, fault tolerance, and disaster recovery.



SYSTEM AREA NETWORKS (SAN)

→ There are high-performance, connection oriented n/w's that link computer clusters.

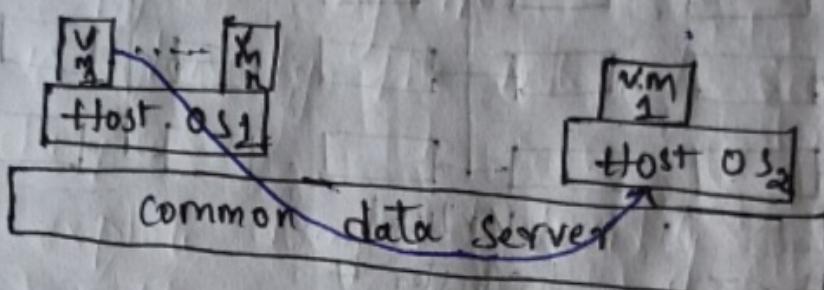


virtual cluster nodes for different application
like A, B, C, D

Live VM migration steps and Performance effects

Live Migration of VMs :

→ The virtual machine stop playing its role when the host machine fails and switched to another host with the same guest OS is called as "VM migration".



Steps :

Step 0: pre-migration

Step -1: Reservation

Step -2: Iterative Pre-copy (Transfer Memory)

Step -3: Suspend the VM and copy the last portion of data

Step -4: Commitment

Step -5: Activate the new host

VM running normally
on Host A

Stage 0: Pre-Migration

Active VM on host A

Alternate physical host may be
selected for migration. Block devices
mirrored and free resources maintained.

Stage 1: Reservation

Initialize a container on the
target host

Overhead due to
copying

Stage 2: Iterative Pre-Copy.
Enable shadow paging
copy dirty pages in
successive "rounds".

Downtime
(VM out of service)

Stage 3: Stop and copy.
Suspend VM on host A
Generate ARP to redirect traffic
Synchronize all remaining to Host B
VM state to Host B

ARP (Address Resolution
Protocol)

Stage 4: Commitment

VM state on host A is released

VM running normally
on Host B

Stage 5: Activation

VM starts on Host B

Connects to local devices

Resumes normal operation.

Live migration process of a VM from one
host to another.

MIGRATION OF MEMORY, FILES AND NETWORK RESOURCES

MEMORY MIGRATION:

- Moving the memory of a VM from one physical host to another is called memory migration.
- Memory migration can be in a range of hundreds of megabytes to a few gigabytes in a typical system.
- It is done by the Internet Suspend-Resume Technique. (ISR)
- It exploits temporal locality
- Each file is represented as a tree of small sub-files
- A copy of this tree exists in both the suspended and resumed VM.
- The caching of this tree ensures the transmission of only those files which have been changed.

FILE SYSTEM MIGRATION:

- Pre-migration analysis to plan an error-free migration.
- Migrate large-sized files and folders without any restriction
- Migrate specific data from file server

- to destination using filters
- CloudMigrator can be used to migrate file system files to Google Drive, Microsoft OneDrive or SharePoint.
- This technique significantly reduces the amount of actual physical data that has to be moved.

NETWORK MIGRATION:

- Network migration involves moving data and programs from one network to another as an upgrade or add-on to a network system.
- The process of migration makes it possible to set up migrated files on a new network or to blend two independent networks together.
- The need for n/w migration may result from security issues, corporate restructuring, increased storage needs and many other issues.

DYNAMIC DEPLOYMENT OF VIRTUAL CLUSTERS.

- LDAP (Light Weight Directory Access Protocol)
- It is a set of open protocols used to access and modify centrally stored information over a network.

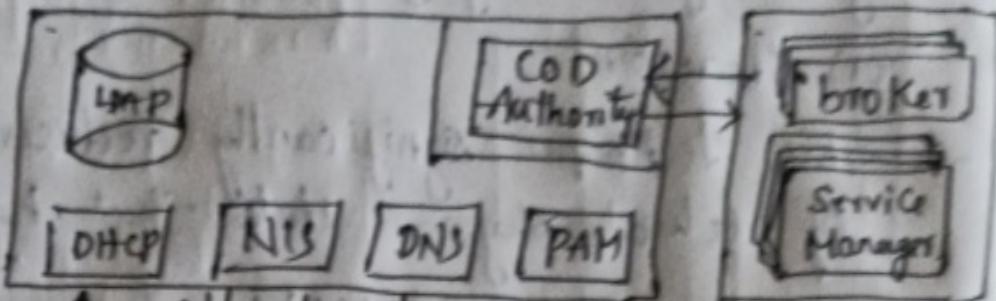
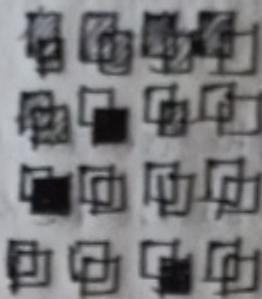
DHCP (Dynamic Host Configuration Protocol)

- It is a protocol that provides quick, automatic, and central management for the distribution of IP addresses within a n/w.

physical
cluster

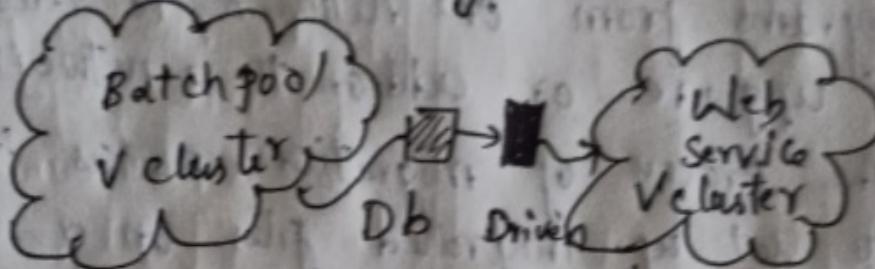
Name/boot/accm servers
backed by LDAP db.

SHARP



N/w boot
Automatic configuration
Resource negotiation.

Dynamic
virtual
clusters



network init
(select OSes)

Allocate new
physical servers
or guest virtual
machines (xen).