

- 6.4. You are given a string of n characters $s[1 \dots n]$, which you believe to be a corrupted text document in which all punctuation has vanished (so that it looks something like “itwasthebestoftimes...”). You wish to reconstruct the document using a dictionary, which is available in the form of a Boolean function $\text{dict}(\cdot)$: for any string w ,

$$\text{dict}(w) = \begin{cases} \text{true} & \text{if } w \text{ is a valid word} \\ \text{false} & \text{otherwise} \end{cases}$$

- Give a dynamic programming algorithm that determines whether the string $s[\cdot]$ can be reconstituted as a sequence of valid words. The running time should be at most $O(n^2)$, assuming calls to dict take unit time.
- In the event that the string is valid, make your algorithm output the corresponding sequence of words.

Scratch

$S = \text{itwasthebestofqnlimes}$

Problem Formulation

$T[i] =$ Whether the sequence $T[1 \dots i]$ is a valid string

Recurrence

$T[i] =$ We will loop back from $i-1$ to 1 using iterator j and check whether at any j , $T[j]$ is true and $\text{dict}(T[j+1 \dots i])$ is also true. Or if $\text{dict}(T[1 \dots i])$ is true. If yes, we will set $T[i]$ to true, else false.

$T[i] = \text{any}((\text{for } j \text{ in } 1 \text{ to } i-1, \text{ if } T[j] \text{ is true and } \text{dict}(T[j+1 \dots i]) \text{ is true}) \text{ or } \text{dict}(T[1 \dots i]) \text{ is true})$

Base case

For all i in 1 to n , $T[i] = \text{false}$

Sample Run (rhymes with Temple Run)

$S = \text{itwasthebestoftimes}$

$T = \text{TTF F TFFT FFFTFTFF FT T}$

$S = \text{itwasthebestofqnlimes}$

$T = \text{TTF F TFFT FFFTFT FFFFF F F}$

Return Value

$T[n]$ indicates whether the whole string is valid or not.

For resolving the valid sequence if the whole sequence is known to be valid, we need to iterate over T

from 1 to n, detect the word ending indices by detecting transition from T to F, insert a space after the word ending index and output the resultant string.

PseudoCode

```
for i = 1-> n      ----- (1)
    T[i] = False
for i = 1-> n      ----- (2)
    if (dict(S[1....i]))
        T[i] = True
    Else
        for j = i-1-> 1      ----- (3)
            if (T[j] == True) and (dict(T[j+1....i]) == True)
                T[i] = True
If T[n] == True
    k = 1
    for i = 1->n-1      ----- (4)
        Sentence[k] = S[k]
        if (T[i] == True) and (T[i+1] == False)
            k = k+1
            Sentence[k] = ' '
        k = k+1
    Sentence[k] = S[k]
    Return T[n], Sentence
Else
    Return T[n], "Invalid String"
```

Runtime Complexity

Loop 1 takes $O(n)$ time.

The loops 2 and 3 are nested, hence will take $O(n^2)$ time in worst case.

The Loop 4 will also take $O(n)$.

All other statements take $O(1)$ time.

Hence, Total time complexity is $O(n^2)$.