**Normalization**

Here are the most commonly used normal forms:

* First normal form(1NF)
* Second normal form(2NF)
* Third normal form(3NF)
* Boyce & Codd normal form (BCNF)

**First normal form (1NF)**

As per the rule of first normal form, an attribute (column) of a table cannot hold multiple values. It should hold only atomic values.

**Example**: Suppose a company wants to store the names and contact details of its employees. It creates a table that looks like this:

|  |  |  |  |
| --- | --- | --- | --- |
| emp\_id | emp\_name | emp\_address | emp\_mobile |
| 101 | Herschel | New Delhi | 8912312390 |
| 102 | Jon | Kanpur | 8812121212  9900012222 |
| 103 | Ron | Chennai | 7778881212 |
| 104 | Lester | Bangalore | 9990000123  8123450987 |

Two employees (Jon & Lester) are having two mobile numbers so the company stored them in the same field as you can see in the table above.

This table is **not in 1NF**as the rule says “each attribute of a table must have atomic (single) values”, the emp\_mobile values for employees Jon & Lester violates that rule.

To make the table complies with 1NF we should have the data like this:

|  |  |  |  |
| --- | --- | --- | --- |
| emp\_id | emp\_name | emp\_address | emp\_mobile |
| 101 | Herschel | New Delhi | 8912312390 |
| 102 | Jon | Kanpur | 8812121212 |
| 102 | Jon | Kanpur | 9900012222 |
| 103 | Ron | Chennai | 7778881212 |
| 104 | Lester | Bangalore | 9990000123 |
| 104 | Lester | Bangalore | 8123450987 |

**Second normal form (2NF)**

A table is said to be in 2NF if both the following conditions hold:

* Table is in 1NF (First normal form)
* No non-prime attribute is dependent on the proper subset of any candidate key of table.

An attribute that is not part of any candidate key is known as non-prime attribute.

**Example**: Suppose a school wants to store the data of teachers and the subjects they teach. They create a table that looks like this: Since a teacher can teach more than one subjects, the table can have multiple rows for a same teacher.

|  |  |  |
| --- | --- | --- |
| teacher\_id | subject | teacher\_age |
| 111 | Maths | 38 |
| 111 | Physics | 38 |
| 222 | Biology | 38 |
| 333 | Physics | 40 |
| 333 | Chemistry | 40 |

**Candidate Keys**: {teacher\_id, subject}  
**Non prime attribute**: teacher\_age

The table is in 1 NF because each attribute has atomic values. However, it is not in 2NF because non prime attribute teacher\_age is dependent on teacher\_id alone which is a proper subset of candidate key. This violates the rule for 2NF as the rule says “**no** non-prime attribute is dependent on the proper subset of any candidate key of the table”.

To make the table complies with 2NF we can break it in two tables like this:  
**teacher\_details table:**

|  |  |
| --- | --- |
| teacher\_id | teacher\_age |
| 111 | 38 |
| 222 | 38 |
| 333 | 40 |

**teacher\_subject table:**

|  |  |
| --- | --- |
| teacher\_id | subject |
| 111 | Maths |
| 111 | Physics |
| 222 | Biology |
| 333 | Physics |
| 333 | Chemistry |

Now the tables comply with Second normal form (2NF).

**Third Normal form (3NF)**

A table design is said to be in 3NF if both the following conditions hold:

* Table must be in 2NF
* [Transitive functional dependency](https://beginnersbook.com/2015/04/transitive-dependency-in-dbms/) of non-prime attribute on any super key should be removed.

An attribute that is not part of any [candidate key](https://beginnersbook.com/2015/04/candidate-key-in-dbms/) is known as non-prime attribute.

In other words 3NF can be explained like this: A table is in 3NF if it is in 2NF and for each functional dependency X-> Y at least one of the following conditions hold:

* X is a [super key](https://beginnersbook.com/2015/04/super-key-in-dbms/) of table
* Y is a prime attribute of table

An attribute that is a part of one of the candidate keys is known as prime attribute.

**Example**: Suppose a company wants to store the complete address of each employee, they create a table named employee\_details that looks like this:

|  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- |
| emp\_id | emp\_name | emp\_zip | emp\_state | emp\_city | emp\_district |
| 1001 | John | 282005 | UP | Agra | Dayal Bagh |
| 1002 | Ajeet | 222008 | TN | Chennai | M-City |
| 1006 | Lora | 282007 | TN | Chennai | Urrapakkam |
| 1101 | Lilly | 292008 | UK | Pauri | Bhagwan |
| 1201 | Steve | 222999 | MP | Gwalior | Ratan |

**Super keys**: {emp\_id}, {emp\_id, emp\_name}, {emp\_id, emp\_name, emp\_zip}…so on  
**Candidate Keys**: {emp\_id}  
**Non-prime attributes**: all attributes except emp\_id are non-prime as they are not part of any candidate keys.

Here, emp\_state, emp\_city & emp\_district dependent on emp\_zip. And, emp\_zip is dependent on emp\_id that makes non-prime attributes (emp\_state, emp\_city & emp\_district) transitively dependent on super key (emp\_id). This violates the rule of 3NF.

To make this table complies with 3NF we have to break the table into two tables to remove the transitive dependency:

**employee table:**

|  |  |  |
| --- | --- | --- |
| emp\_id | emp\_name | emp\_zip |
| 1001 | John | 282005 |
| 1002 | Ajeet | 222008 |
| 1006 | Lora | 282007 |
| 1101 | Lilly | 292008 |
| 1201 | Steve | 222999 |

**employee\_zip table:**

|  |  |  |  |
| --- | --- | --- | --- |
| emp\_zip | emp\_state | emp\_city | emp\_district |
| 282005 | UP | Agra | Dayal Bagh |
| 222008 | TN | Chennai | M-City |
| 282007 | TN | Chennai | Urrapakkam |
| 292008 | UK | Pauri | Bhagwan |
| 222999 | MP | Gwalior | Ratan |

**Boyce Codd normal form (BCNF)**

It is an advance version of 3NF that’s why it is also referred as 3.5NF. BCNF is stricter than 3NF. A table complies with BCNF if it is in 3NF and for every [functional dependency](https://beginnersbook.com/2015/04/functional-dependency-in-dbms/) X->Y, X should be the super key of the table.

**Example**: Suppose there is a company wherein employees work in **more than one department**. They store the data like this:

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| emp\_id | emp\_nationality | emp\_dept | dept\_type | dept\_no\_of\_emp |
| 1001 | Austrian | Production and planning | D001 | 200 |
| 1001 | Austrian | stores | D001 | 250 |
| 1002 | American | design and technical support | D134 | 100 |
| 1002 | American | Purchasing department | D134 | 600 |

**Functional dependencies in the table above**:  
emp\_id -> emp\_nationality  
emp\_dept -> {dept\_type, dept\_no\_of\_emp}

**Candidate key**: {emp\_id, emp\_dept}

The table is not in BCNF as neither emp\_id nor emp\_dept alone are keys.

To make the table comply with BCNF we can break the table in three tables like this:  
**emp\_nationality table:**

|  |  |
| --- | --- |
| emp\_id | emp\_nationality |
| 1001 | Austrian |
| 1002 | American |

**emp\_dept table:**

|  |  |  |
| --- | --- | --- |
| emp\_dept | dept\_type | dept\_no\_of\_emp |
| Production and planning | D001 | 200 |
| stores | D001 | 250 |
| design and technical support | D134 | 100 |
| Purchasing department | D134 | 600 |

**emp\_dept\_mapping table:**

|  |  |
| --- | --- |
| emp\_id | emp\_dept |
| 1001 | Production and planning |
| 1001 | stores |
| 1002 | design and technical support |
| 1002 | Purchasing department |

**Functional dependencies**:  
emp\_id -> emp\_nationality  
emp\_dept -> {dept\_type, dept\_no\_of\_emp}

**Candidate keys**:  
For first table: emp\_id  
For second table: emp\_dept  
For third table: {emp\_id, emp\_dept}

This is now in BCNF as in both the functional dependencies left side part is a key.