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8/40

Figure6-3 drawing: -10

1. Question 2:

$$F(n) = \Omega(g(n)) \text{ implies } g(n) = O(f(n))$$

Answer: True

By definition

$$f(n) = \Omega(g(n)) \text{ which implies } O \leq c \cdot g(n) \leq f(n)$$

$$g(n) = O(f(n)) \text{ which implies } O \leq g(n) \leq c \cdot f(n)$$

$$\text{Let us assume that } f(n) = 100n^2, g(n) = n^2$$

$$f(n) \geq c \cdot g(n)$$

$$100 \cdot n^2 \geq c \cdot n^2$$

Consider the constant $c = 50$

$$100n^2 \geq 50n^2$$

$$2 \geq 1$$

 $g(n) = O(f(n))$ which is equal to

$$c \cdot f(n) \geq g(n)$$

$$c \cdot 100 \cdot n^2 \geq n^2$$

$$50 \cdot 100 \cdot n^2 \geq n^2$$

$$5000 \geq 1$$

Based on the above equations $F(n) = \Omega(g(n))$ implies $g(n) = O(f(n))$ is true.**-8****let c2 = 1/c1**

2. Question 3:

$$T(n) = 2T(n/2) + n^4$$

$$T(n) = 2T(n/2) + n^4 - \text{step 1}$$

Let us assume that $n = n/2$

$$T(n/2) = 2T(n/4) + (n/2)^4$$

$$T(n/2) = 2T(n/4) + n^4/16 - \text{step 2}$$

Substitute 2 in 1

$$T(n) = 2[2T(n/4) + n^4/2^4] + n^4$$

$$T(n) = 2^2 T(n/2^2) + 1/2^3 \cdot n^4 + n^4 - \text{step 3}$$

Substitute $n = n/2$ in step 2

$$T(n/4) = 2T(n/8) + (n/2)^4 \cdot 1/16$$

$$= 2T(n/8) + n^4/2^8 - \text{step 4}$$

Substitute 4 in 3

$$T(n) = 2^2[2T(n/8) + n^4/2^8] + (n^4/2^3) + n^4$$

$$= 2^3 \cdot T(n/2^3) + n^4/2^6 + n^4/2^3 + n^4$$

$$= 2^3(T(n/2^3)) + n^4(1/8^2 + 1/8 + 1)$$

After K substitutions

$$= 2^k T(n/2^k) + ((2^k)^4) \cdot \sum_{k=1}^{k-1} (1/8)$$

$$\text{Let } n = 2^k, k = \log n$$

Use Master method: -7

$$\begin{aligned}
T(n) &= nT(1) + n^4, (1/8)^k \\
&= n + n^4 \cdot 1^{\log n} / 8^{\log n} \\
&= n + n^4 \cdot (0.125)^{\log n}
\end{aligned}$$

As n increases $(0.125)^{\log n}$ decreases and it can be ignored.

$$\text{So, } T(n) = \theta(n^4)$$

3. Question 4:

$$T(n) = 16T(n/4) + n^2$$

$$T(n) = 16T(n/4) + n^2 - \text{step 1}$$

Substitute $n = n/4$ in step 1

$$T(n) = 16T(n/16) + (n^2)/(4^2) - \text{step 2}$$

Substitute step 2 in step 1

$$T(n) = 16(16T(n/16) + n^2/16) + n^2$$

$$= (16^2)T(n/16) + n^2 + n^2 - \text{step 3}$$

substitute $n = n/4$ in step 2

$$T(n/16) = 16T(n/64) + n^2/256 - \text{step 4}$$

Substitute step 4 in step 3

$$T(n) = 16^2(16T(n/64) + n^2/16^2) + n^2 + n^2$$

$$= 16^3T(n/64) + 3n^2$$

$$= 4^6 \cdot T(n/4^3) + 3n^2$$

After k substitutions

$$= 4^{2k}T(n/4^k) + kn^2$$

$$= n^2 \cdot T(1) + n^2 \log_4 n$$

$$4^k = n$$

$$k = \log_4 n$$

By ignoring the lower order terms and constants

$$T(n) = \theta(n^2 \log n)$$

Use Master method: -7