## Lecture 17: More Counting

(If notation is unfamiliar, see Appendix!)

## 1 Inclusion/Exclusion

Recall the Sum Rule:  $|A \cup B| = |A| + |B|$  if A, B are disjoint. What if they're not?!

How many queens and/or hearts are in a standard deck of cards? 4Q + 13 = 17 cards? But then Q = 17 would be counted twice! Instead: 4 + 13 - 1 = 16.

In general,  $|A \cup B| = |A| + |B| - |A \cap B|$ . Easier to see using a Venn Diagram (draw picture).

Similar formula for 3 sets:

$$|A \cup B \cup C| = |A| + |B| + |C| - |A \cap B| - |A \cap C| - |B \cap C| + |A \cup B \cup C|.$$

Application: Let n = pqr, product of 3 different primes. How many numbers in  $\{1, 2, ..., n\}$  are relatively prime to n?

How many aren't? Let  $A_p$  be the set of numbers in  $\{1, 2, ..., n\}$  that are divisible by p, same for  $A_q$  and  $A_r$ . Our answer is  $n - |A_p \cup A_q \cup A_r|$ .

Inc exc:  $|A_p| = n/p$  and same for  $q, r, |A_p \cap A_q| = n/(pq)$  and same for other pairwise intersections, and  $|A_p \cap A_q \cap A_r| = n/(pqr) = 1$ . Formula gives

$$|A_p \cup A_q \cup A_r| = |A_p| + |A_q| + |A_r| - |A_p \cap A_q| - |A_p \cap A_r| - |A_q \cap A_r| + |A_p \cap A_q \cap A_r|$$

$$= n/p + n/q + n/r - n/(pq) - n/(pr) - n/(qr) + n/(pqr)$$

$$= qr + pr + qp - r - q - p + 1$$

$$= pqr - (pqr - qr - pr - qp + r + q + p - 1)$$

$$= n - (p - 1)(q - 1)(r - 1).$$

so answer is n minus that! Simplifies to (p-1)(q-1)(r-1).

Even more generally:

$$|A_1 \cup A_2 \cup \dots \cup A_n| = \sum_{i} |A_i|$$

$$- \sum_{i < j} |A_i \cap A_j|$$

$$+ \sum_{i < j < k} |A_i \cap A_j \cap A_k|$$

$$\dots$$

$$\pm |A_1 \cap A_2 \cap \dots \cap A_n|.$$

Add the initial sets, subtract the 2-way intersections, add the 3-way intersections, subtract 4-way intersections, etc., all the way to the n-way intersection.

Can also be expressed as follows:

**Theorem 1** (PIE). Let  $\mathcal{U} = \bigcup_{i \in [n]} A_i$  be a finite universe of discourse. Then

$$\sum_{I \subseteq [n]} (-1)^{|I|} \left| \bigcap_{i \in I} A_i \right| = 0,$$

where  $\bigcap \emptyset = \mathcal{U}$  by convention.

Proof. For  $x \in \mathcal{U}$ , let  $I_x \subseteq [n]$  be the set of indices i for which  $x \in A_i$ . For every  $I \subseteq I_x$ ,  $x \in \bigcap_{i \in I} A_i$ . This means x contributes +1 to the LHS for each such I of even size, and x contributes -1 to the LHS for each such I of odd size.  $I_x$  is non-empty by definition of  $\mathcal{U}$ , so contains some index i. Now symmetric difference with  $\{i\}$  gives a self-inverting bijection

between even-size and odd-size subsets  $I \subseteq I_x$ , so x contributes 0 in total to the LHS.

Example: If we take n=2 and expand out the LHS, this is saying that

Example: If we take n = 3 and expand out the LHS, this is saying that

#### 2 Pidgeyhole Principle

**Theorem 2** (Pidgeyhole Principle). If |A| > |B|, and  $f : A \to B$  is a total function, then f is not injective. In other words, there exist  $a_1, a_2 \in A$  such that  $a_1 \neq a_2$  and  $f(a_1) = f(a_2)$ .

More generally, any total relation  $R \subseteq A \times B$  is not injective: there must exist at least two distinct  $a_1, a_2 \in A$  that relate to the same  $b \in B$ .

Name comes from medieval times: cubbies for domestic pigeons to rest in. A is pigeons, B is pigeonholes. If more birds than cubbies, then some pigeons must share.

Example: strictly more than 26 people in the room, so two of us must have names that start with the same first letter.

Example: n differently colored pairs of socks; how many single socks do I need to pick before I'm guaranteed to have a matching pair? Pidgeyhole Principle says n+1 is enough. And can't be less, because I might accidentally pick one from each pair. So n+1 is the exact answer.

Example: there exist two non-bald Bostonians ( $\approx 650,000$ ) with the same number of hairs on their head ( $\leq 200,000$ ). Assuming less than 2/3 of Boston is bald, we have more Bostonians than possible hair counts, so must be true by Pidgeyhole.

Note: **nonconstructive!** We know they exist, but we don't know who! Pidgeyhole Principle says there must be a collision, but doesn't give an easy way to actually find the people.

Example: no lossless compression scheme that strictly shortens all n-bit strings. There are  $2^n$  bitstrings with length n, but only  $2^{n-1} + 2^{n-2} + \cdots + 2^1 + 2^0 = 2^n - 1$  shorter strings (including the empty string). Any total function from bigger set to smaller set must have collisions, so it's not lossless.

**Theorem 3** (Generalized Pidgeyhole Principle). If  $|A| > k \cdot |B|$ , then every total relation/function from A to B must have at least k+1 elements in A that map to the same element in B.

With either version, proofs are often short but can sometimes require cleverness! Gotta pick your pigeons, holes, and/or the map between them carefully; not all choices are useful.

One more example: on an  $8 \times 8$  chessboard, we fill 33 of the 64 cells with Rooks. Show we can find 5 of them that don't attack each other, i.e., lie in 5 distinct rows and 5 distinct columns.

Clever idea: Choose 8 pigeonholes, where each one is a subset of 8 cells like this:

2	3	4	5	6	7	8	1
3	4	5	6	7	8	1	2
4	5	6	7	8	1	2	3
5	6	7	8	1	2	3	4
6	7	8	1	2	3	4	5
7	8	1	2	3	4	5	6
8	1	2	3	4	5	6	7
1	2	3	4	5	6	7	8

Label each rook with the number on its cell. There are 8 labels and 33 rooks, so at least 5 of them must have the same label. These 5 are in different rows and columns; done!

## 3 Combinatorial Proofs / Double Counting

$$\sum_{k=0}^{n} \binom{n}{k} = 2^{n}.$$

*Proof.* How many subsets of  $\{1, 2, ..., n\}$  are there?

Idea: For each  $0 \le k \le n$  there are  $\binom{n}{k}$  subsets of size k, so together these should add to the total number of subsets,  $2^n$ .

More precisely: Let S be the set of all subsets of  $\{1, 2, ..., n\}$ . We will count |S| in two different ways.

First,  $|S| = 2^n$  since each element is either in or out.

Let's instead work by cases, depending on the size of the subsets. Let  $S_k$  be the set of subsets of size k, and note that  $S_0, S_1, \ldots, S_n$  form a **partition** of S: every member of S belongs to exactly one of the  $S_k$ . By the sum rule, this means  $|S| = \sum_{k=0}^n |S_k|$ . But  $|S_k| = \binom{n}{k}$ , so we get  $2^n = |S| = \sum \binom{n}{k}$ , as claimed.

This is a special case of a useful theorem. We've been using binomial coefficients a lot already; here's their eponym:

**Theorem 4** (Binomial Theorem). For any x, y, and for  $n \in \mathbb{N}$ :

$$(x+y)^n = \sum_{k=0}^n \binom{n}{k} x^k y^{n-k},$$

where  $0^0 = 1$  by convention.

*Proof.* Expand  $(x+y)^n$  to get a sum of  $2^n$  terms of the form  $a_1a_2\cdots a_n$ , where  $a_i$  are either x or y. By commutativity, we can group like terms of the form  $x^ky^{n-k}$ . There are  $\binom{n}{k}$  such terms; this is the number of ways we can choose k of the n indices i for which  $a_i = x$ .  $\square$ 

Example: If n = 3, we have  $(x + y)^3 = y^3 + 3xy^2 + 3x^2y + x^3$ .

The coefficient of  $(x^0)y^3$  is  $\binom{3}{0} = 1$ .

The coefficient of  $xy^2$  is  $\binom{3}{1} = 3$ .

The coefficient of  $x^2y$  is  $\binom{3}{2} = 3$ .

The coefficient of  $x^3(y^0)$  is  $\binom{3}{3} = 1$ .

Similarly:

**Theorem 5** (Multinomial Theorem). For any  $x_1, x_2, \ldots, x_m$ , and for  $n \in \mathbb{N}$ :

$$\left(\sum_{i=1}^{m} x_i\right)^n = \sum_{k_1 + k_2 + \dots + k_m = n} \binom{n}{k_1, k_2, \dots, k_m} \prod_{i=1}^{m} x_i^{k_i}.$$

As before,  $0^0 = 1$  by convention, and the summation is over non-negative integers  $k_i$  that sum to n.  $\binom{n}{k_1, k_2, \ldots, k_m}$  is the multinomial coefficient defined by

$$\binom{n}{k_1, k_2, \dots, k_m} = \frac{n!}{k_1! k_2! \cdots k_m!}.$$

Another useful identity:

$$\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k}$$

Could prove this by algebra using factorials, but no intuition! Here's a combinatorial proof:

*Proof.* How many subset of  $\{1, 2, ..., n\}$  have size k? Let S be this set of size-k subsets. Then  $|S| = \binom{n}{k}$ . But let's count these in a different way. Every size-k subset of  $\{1, 2, ..., n\}$  either includes n or doesn't. Let A be the set of ones that include n, and B the set of ones that don't, so S is the disjoint union of A and B.

Note:  $|B| = \binom{n-1}{k}$ , because we're not allowed to use n. Also,  $|A| = \binom{n-1}{k-1}$ , because we must pick n, and then k-1 other numbers from  $\{1, \ldots, n-1\}$ . So |S| = |A| + |B|, which is exactly the identity above.

Note: this fact shows that if we put the numbers  $\binom{n}{k}$  in a big triangle, each is the sum of the 2 above it. This is Pascal's Triangle. Above we showed that the sum of row k is  $2^k$ , and in homework you'll show that the sums of diagonals gives Fibonacci Numbers!

Fun exercise: Find a combinatorial proof that

$$\sum_{k=0}^{n} \binom{n}{k}^2 = \binom{2n}{n}$$

#### **Appendix: Notation**

Just as  $\sum_{i=1}^{n} x_i := x_1 + x_2 + \ldots + x_n$ , we can similarly define *n*-way products, unions, and intersections:

$$\prod_{i=1}^{n} x_i := x_1 \times x_2 \times \dots \times x_n$$

$$\bigcup_{i=1}^{n} S_i := S_1 \cup S_2 \cup \dots \cup S_n$$

$$\bigcap_{i=1}^{n} S_i := S_1 \cap S_2 \cap \dots \cap S_n$$

There are also different notations for indexing:

$$[n] := \mathbb{N} \cap (0, n] = \{1, 2, \dots, n\}$$

$$\sum_{i \in [n]} x_i \text{ is another way to write } \sum_{i=1}^n x_i.$$

More generally, for a finite indexing set S,  $\sum_{x \in S} f(x)$  means  $\sum_{i=1}^{|S|} f(\phi(i))$ , where  $\phi : [|S|] \to S$  is any bijection.

A common abuse of notation is to write  $\sum_{P(x)} f(x)$  instead of  $\sum_{x \in \{y: P(y)\}} f(x)$ .

One can even omit the index entirely:  $\sum S$  means  $\sum_{x \in S} x$ .

All of these notations extend naturally to  $\prod$ ,  $\bigcup$ ,  $\bigcap$ .

 $\sum_{i=1}^{m} \emptyset = 0$  by convention, because 0 is the *identity* of +. Basically, it should be true that  $\sum_{i=1}^{m} x_i + \sum_{i=m+1}^{n} x_i = \sum_{i=1}^{n} x_i$ . Taking m = 0 tells us  $\sum_{i=1}^{0} x_i = \sum_{i=1}^{n} \emptyset = 0$ .

Similarly, for  $\prod$  we have  $\prod \emptyset = 1$ , and  $\bigcup \emptyset = \emptyset$ , and  $\bigcap \emptyset = \mathcal{U}$ , where  $\mathcal{U}$  is the universe of discourse, or "everything".

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