pata base creation &

pata: collection of facts which provides information about something

Database: correction of data, specification,
Types of data: Structure, constraint

Traditional Recent

DBMS: Set of applications that allows users to construct & manage datab-

\* General purpose software that simplifie s. Creating. Building, modifying & sharing data bases across users & applications. Metadata: Telling data about data. Data base catalogue: Tabular form. Miniworld: Real-time world. Concurrency: Accessing one's data. why DBMS:

\* Traditional file systems were used before DBMS.

\* DBMS got vid off many limitations faced by traditional file systems.

\* Familiar with physical specifics.

No concurrency.

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	Basis	File system	DBMS
1	Structure	Software that manages & organise the file in storage	software for managing database
	Data	Redundant data	
	Redundance	is present	podata sandatod
	Backup	NO Backup &	provides backup
	Recovery	Recovery	E Recovery if
I	Query.	No efficient	Efficient query
	processing		processing
-	anage datal	ming + obstance	t anson -
	consistency	Less data consistency	* General purp
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	pplications	o a around ador	Metadata: Tellin
		que: Fabular for	Data base cotator
+	2-2-2-3-11		Minisportds Real
	security	ssing one's data	concurrency: Acce
			why DBMS:
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	ricy zacar.	conditional of the	DETOTE DEME
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	Data	detidnal file	
	Independe	the physical is	No concurre
-	user		
	Access		
-			Mark Bridge Strange St

Meaning
sharing
Data abstraction
Integrity constraints

Data: known facts & have an implicit meaning.

Mini-world: Some part of real world stored in a database.

DBMs: software that facilitate operat--ions on Sata base.

Features of DBMs:

Normalization

User-defined limitations/rules

Security

Data Backup

Data organisation

Goals of Database-design:

minimize the redundancy, lossless join,
Dependency preservation.

Draw-Backs of DBMS:

\* Cost, complexity, through understanding of SQL is required.

\* simple, well defined, Not expected to change are characteristics iff DBMs is not required.

\* Access to data by multiple users is not required.

Database users: > End user. > DB Administrator > DB designer \* use & control the database content, who design, develop & maintain (Actors) \* who design & develop Three sevel Architecture of Database: (Blue print) Profesternal schema Conceptual conceptual level -> External level schema) Li le le seured in a datobas (Internal physical lever End users schema) Notiliant toll properties memad \* Data Base End users Normalization DE TULES Regartsh. External level concéptual level ppro pto a appropriate of patellastales physical level) and simina Dependency preservational Information Data Base 1 to a son a word physical level: Data is actually stored in database. singples well defineds not exa refignale are characteristics ns not required. Access to data by multiple users is not required.

Architecture in DBMs: Tier-1: Tier-2: Tier-3 C2 Application) > DB ODBC, JDBC, C3 Cayer poor performance (more users), query processing, management of transactions. of the leaf the the State State Prepare Grant o APB a nos stories Actes (Reconstructing data) concurrency: Discretions of using simulta-Truncate parapage - by securing Application layer: Partially processed data Remances the relation in (sand o (ugers level) evel database ischema Conceptual schema hbs (10) stormer (Physical) schema) al what data is DB DB Low level language conceptual Physical view schema Tabucar Rawfolders user interaction pata Independence: change in the data shouldn't be affected by the program. Data is seperated from the program. This would save time & cost requires as it doesn't affects the data at other levels of the data base.

Hence, change in the data at a level won't affect the program execution (or) application of program.

Data base Architecture (A)

\* Differences between procedural & Non-procedure

Data définition [anguage: (DDL) [schema-structure] statements that will be used to implement the database schema. set of statements in DOL to implement database scheme:

-> create (construct a table) -> Acter ( Reconstructing data) -> Drop (Delete a relation) (or) whole > Truncate ( Deletes by securing Rename relation structure) ( Renames the relation in the data base)

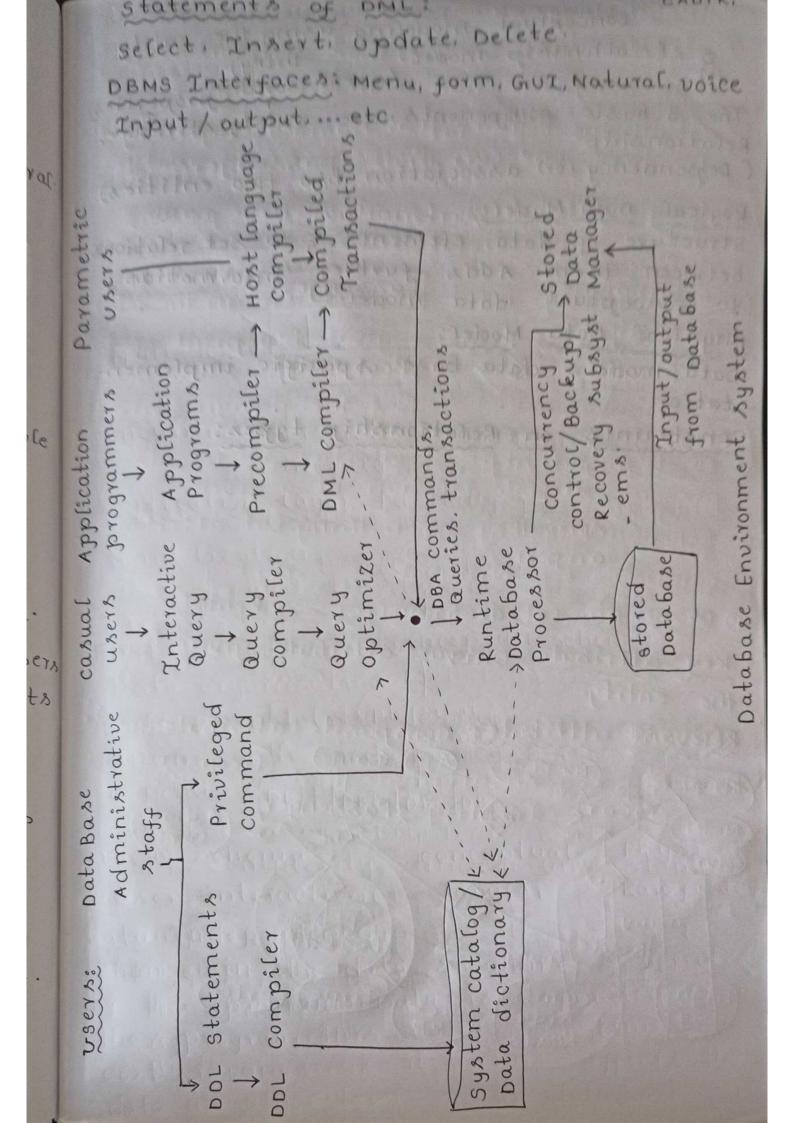
DDL follows three level database schema Data Manipulation Language (DML): Accows use to access & alter data in a data base. Let users to retrieve, insert, remove, (or) edit data in a database.

(i) Procedural DMLS:

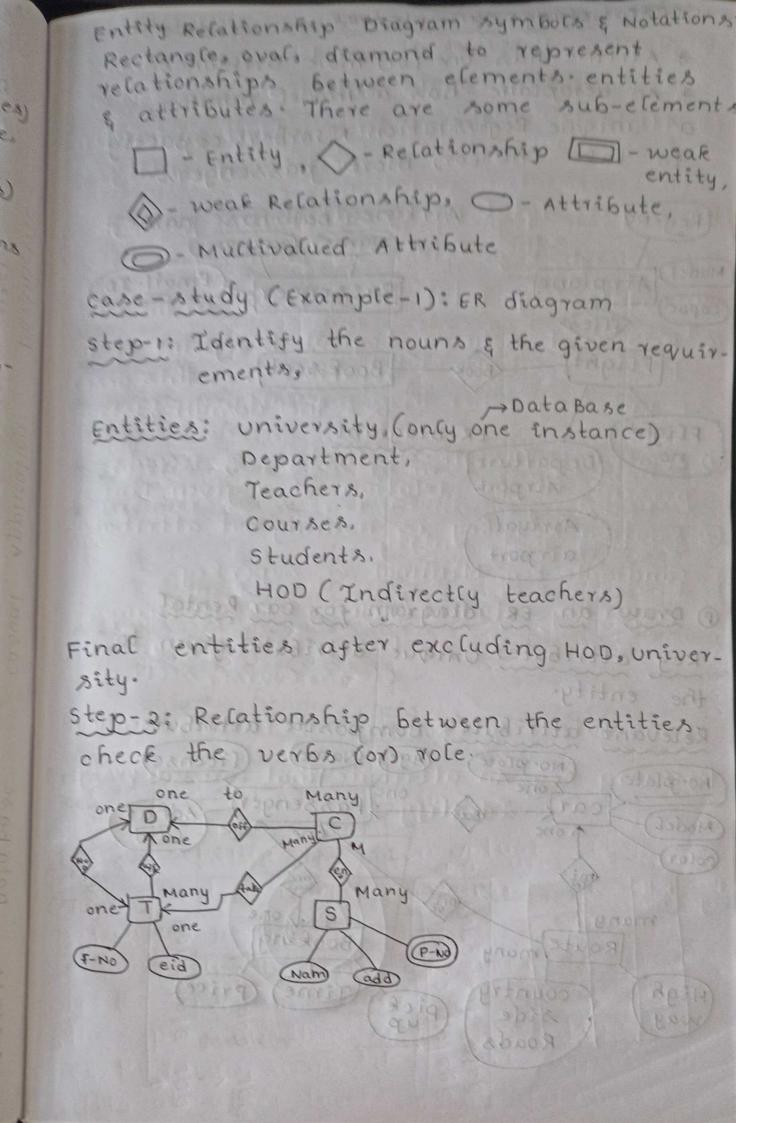
(i) Procedural DMLS: Low level languages, defines what data is needed, how to obtain that data called as one-at-a-time DML's. (ii) Non-procedural DMLS:

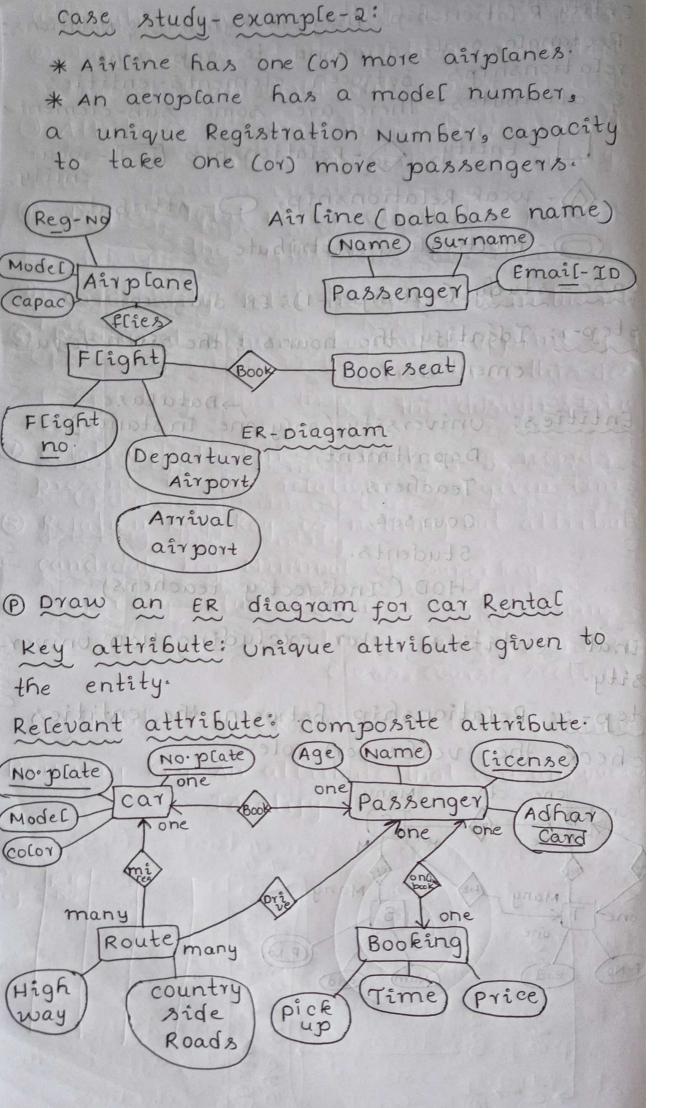
Precisely define the data requirement without specifying the way to access it Set-a-time DMLs. Hod Branchord Set

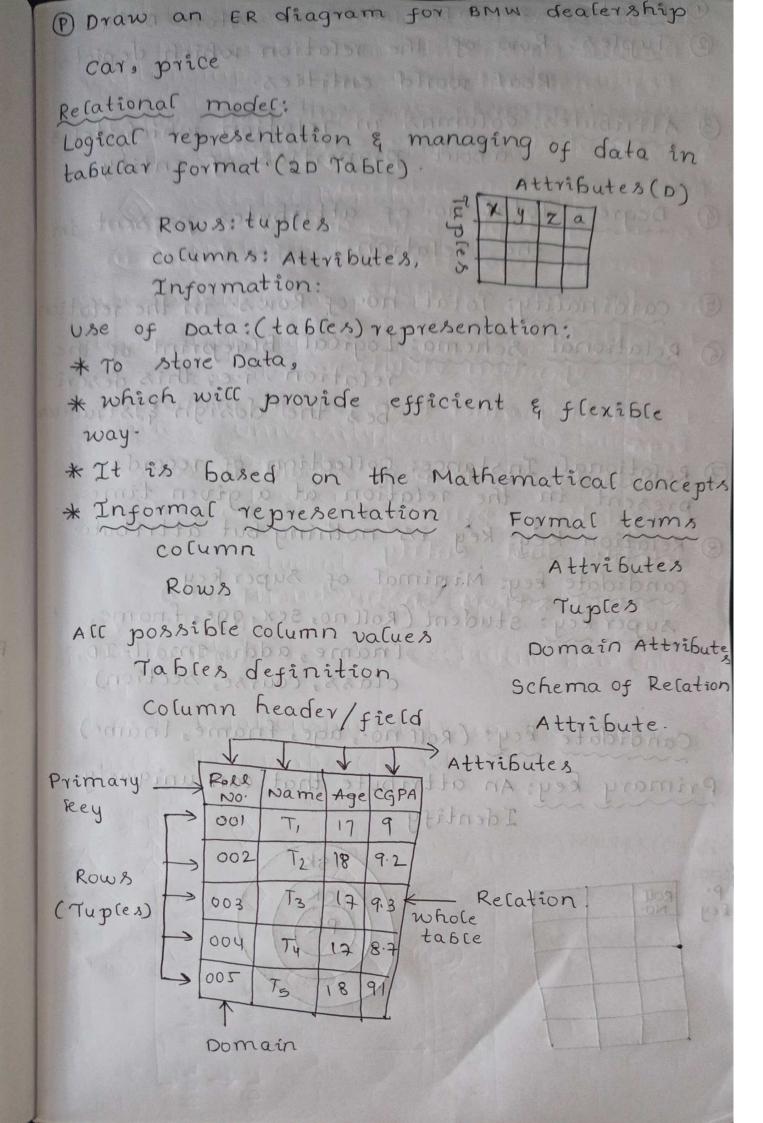
cost requires as it doesn't affects the data at other levels of the data base.



Conceptual modeling using the entity Relationship model thing) (A real world (properties The 3 basic components are Entity, Attribute, Relationship. ( Dependency (or) association b/w two entities) Logical Data Model: structure of data elements & to set relations between them. Adds further information to conceptual data model. Physical Data Model: Data model's data base specific implemen-- tation. Constraints on Relationship types:







1 Relation: 20 Table used to store Data! 2 Tuples: Row of the relation which depicts Real world entities. 3 Attributes: columns in relation which tell, about properties of entity. Attailected (1991) Degree: Total no. of attributes present in the relation .... 5 cardinality: Total no of Rows in the relation © Relational schema: Logical blueprint of the relation, i.e; this descri-- bes the design & struct 7 Relational Instance: Collection of records present in the relation at a given time. B Relation key: key is nothing but attribute. Candidate key: Minimal of Super key super key: student (Roll no, sex, age, fname, Iname, addr, email-ID, Schema of Relation class, course, section). Candidate key: (Roll no., age, fname, lname). Primary key: An attribute that can uniquely Identify SK ROLL CK key NO. PK DOMING UR

2. key constraint: (uniqueness)

RN F-name Age

8 Entitly integrity constrtitle (NXI volume

002 4. Reservation untegrida custoff X

3. Entity Integrity: constraint

4. Referential Integrity constraint:

RN	FN	Age

RNO add Eid babbb on Branch of Colored Code (3)	RN	add	Eld	) nd				
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opdate operation violations

CREATE DATABASE soudent\_DB; DROP DATABASE Student-DB;

Insert operation: The insert operation provides a list of attribute values for new tuple t that is to be inserted into a relation R. SELECT \* FROM INSTRUMOST: SIGNIFIES

INSERTABOUT

Pelete operation:

Decetes specific row that fits the criteria specified in Delete query.

Syntax:

DELETE FROM (table-name) WHERE (condi tion:

Update (or) Modify operations:

update (or) modify to change the values of one (or) more attributes.

> UPDATE table-name SET column 1 = value1, column 2 = value 2,... WHERE condition;

Insert operation violations: 1. Domain constraint 2. key constraint 3. Entity integrity constraint (Nucl values) 4. Referential integrity constraint Delete operation violation: occurs if tuple being deleted is referenced by foreign keys from the other tuples in the database solutions 1. restrict, 2. cascade. 3. Nucl Update operation violations: 1. updating other than foreign, primary DROP DATABASE SEUdent-DB; Eeys. 2. Data type & Domain should be correct 3. Updating primary key is as good as deceting that entity and of the \* SELECT \* FROM Instructor; (Prints) all values & tuples) noithropo stolag prietre specific now that fits the criteria specified in Octobe query sites DELETE FROM (table-name) WHERE (condi speake (or) Modify operations: 111 update (or) modify to change the rolues of one Con more attributes. UPDATE tobbe nome SET columns = values, column o = volue e... MHERE CONDITIONS