

**5.1. Wedge sum and smash product** The disjoint union of two spaces is geometrically exactly what it sounds like: you just imagine the two spaces side by side.

### Definition 5.1: (Disjoint union)

Let  $X$  and  $Y$  be two topological space. The **disjoint union**,  $X \sqcup Y$ , is defined by

$$X \sqcup Y := \text{disjoint union of } X \text{ and } Y.$$

A subset  $U \subseteq X \sqcup Y$  is open if  $U \cap X$  and  $U \cap Y$  are open in  $X$  and  $Y$ , respectively.

### Exercise 5.2: (Disjoint union on any indexed set)

Let  $\{X_\alpha : \alpha \in \Lambda\}$  be a collection of topological space indexed by  $\Lambda$ . Let

$$X := \bigsqcup_{\alpha \in \Lambda} X_\alpha$$

denote the disjoint union of  $X_\alpha$ 's. For each  $\alpha \in \Lambda$ , we have the natural inclusion map

$$f_\alpha : X_\alpha \hookrightarrow X.$$

Show that the above defined topology is the strongest topology on  $X$  such that each  $f_\alpha$  is continuous.

### Definition 5.3: (Based or pointed space)

A **based space** or **pointed space** is a pair  $(X, x_0)$ , where  $X$  is a topological and  $x_0 \in X$  is a chosen **basepoint**.

If  $(X, x_0)$  and  $(Y, y_0)$  are based spaces, then a continuous map  $f : (X, x_0) \rightarrow (Y, y_0)$  is called a **based map** if  $f(x_0) = y_0$ .

### Definition 5.4: (Wedge sum)

Let  $(X, x_0)$  and  $(Y, y_0)$  be two based spaces. Then the **wedge sum**, denoted by  $X \vee Y$ , is defined as the quotient space of the disjoint union of  $X$  and  $Y$  by the identification  $x_0 \sim y_0$ .

$$X \vee Y := (X \sqcup Y) / \sim, \text{ where } x_0 \sim y_0.$$

Similarly, the wedge sum of collection of pointed spaces is defined as

$$\bigvee_{\alpha \in \Lambda} (X_\alpha, x_\alpha) := \left( \bigsqcup_{\alpha \in \Lambda} (X_\alpha, x_\alpha) \right) / \sim.$$

### Definition 5.5: (Smash product)

Let  $(X, x_0)$  and  $(Y, y_0)$  be two based spaces. The **smash product**, denoted by  $X \wedge Y$ , is defined as the quotient space of  $X \times Y$  with the identifications  $(x, y_0) \sim (x_0, y)$  for all  $x \in X$  and  $y \in Y$ .

Note that the wedge sum of  $X$  and  $Y$  be viewed as  $X \vee Y = X \times \{y_0\} \cup \{x_0\} \times Y$  and hence

$$X \wedge Y = \frac{X \times Y}{X \vee Y}.$$

## 5.2. Category theory

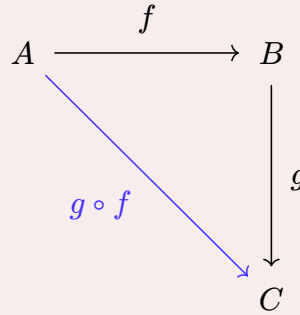
### Definition 5.6: (Category)

A **category**  $\mathcal{C}$  consists of following:

- i) A collection of **objects** denoted by  $\text{Ob}(\mathcal{C})$ .
- ii) For any two objects  $A, B \in \text{Ob}(\mathcal{C})$ , a set  $\text{Hom}_{\mathcal{C}}(A, B)$  called a **morphism**.
- iii) For any three objects  $A, B, C \in \text{Ob}(\mathcal{C})$ , we have a **composition**

$$\circ : \text{Hom}_{\mathcal{C}}(A, B) \times \text{Hom}_{\mathcal{C}}(B, C) \rightarrow \text{Hom}_{\mathcal{C}}(A, C),$$

$$(f, g) \mapsto g \circ f$$



satisfying the following properties.

- a) **(Associativity)** Let  $A, B, C, D \in \text{Ob}(\mathcal{C})$ . For morphisms  $f \in \text{Hom}_{\mathcal{C}}(A, B)$ ,  $g \in \text{Hom}_{\mathcal{C}}(B, C)$  and  $h \in \text{Hom}_{\mathcal{C}}(C, D)$ , we have

$$h \circ (g \circ f) = (h \circ g) \circ f.$$

- b) **(Identity)** For each object  $A$ , there is a morphism  $1_A \in \text{Hom}_{\mathcal{C}}(A, A)$  such that for any  $g \in \text{Hom}_{\mathcal{C}}(A, B)$  and any  $h \in \text{Hom}_{\mathcal{C}}(C, A)$ , we have

$$g \circ 1_A = g, \quad 1_A \circ h = h.$$

When the context is clear, we will write  $\text{Hom}(A, B)$  instead of  $\text{Hom}_{\mathcal{C}}(A, B)$ .

### Remark 5.7:

- i) The only restriction on  $\text{Hom}$  is that it be a set. In particular,  $\text{Hom}(A, B) = \emptyset$  is allowed, although the identity axiom shows that  $\text{Hom}(A, A) \neq \emptyset$  as it contains  $1_A$ .
- ii) Instead of writing  $f \in \text{Hom}(A, B)$ , we usually write  $f : A \rightarrow B$ .

### Example 5.8: (Examples of category)

Categories are everywhere in mathematics (and even in computer science). We will list some of them here.

- i)  $\mathcal{C} = \text{Sets}$ . Here

$\text{Ob}(\mathcal{C}) = \text{all sets},$

$\text{Hom}(A, B) = \{\text{all functions } f : A \rightarrow B\}$

$\circ = \text{composition is the usual function composition}$

ii)  $\mathcal{C} = \mathbf{Top}$ . Here

$\text{Ob}(\mathcal{C}) = \text{all topological spaces},$

$\text{Hom}(X, Y) = \{\text{all continuous functions } f : X \rightarrow Y\}$

$\circ = \text{composition is the usual function composition}$

iii)  $\mathcal{C} = \mathbf{Top}_*$ . Here

$\text{Ob}(\mathcal{C}) = \{(X, *_X) : X \text{ is a topological space and } *_X \in X\}$

$*_X \in X$  is called the basepoint,

$\text{Hom}((X, *_X), (Y, *_Y)) = \{\text{all continuous functions } f : (X, *_X) \rightarrow (Y, *_Y) \text{ with } f(*_X) = *_Y\}$

this is called basepoint preserving functions

$\circ = \text{composition is the usual function composition}$

iv)  $\mathcal{C} = \mathbf{Groups}$ : Category of groups and group homomorphisms.

v)  $\mathcal{C} = \mathbf{Ab}$ : Category of abelian groups and group homomorphisms.

#### Example 5.9: (Discrete Category)

Let  $X$  be a nonempty set. Consider a category  $\mathcal{C}$ ,

$\text{Ob}(\mathcal{C}) = \text{elements of } X$

$$\text{Hom}(x, y) = \begin{cases} \emptyset & \text{if } x \neq y \\ 1_x & \text{if } x = y. \end{cases}$$

This is called a *discrete category*.