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PES University, Bangalore (Established under Karnataka Act No. 16 of 2013)

UE17CS302

END SEMESTER ASSESSMENT (ESA) - December 2019

Introduction to Operating Systems

Time: 3 Hrs Answer All Questions Max Marks: 100

. a)	Explain Pi	ocess State transi	ition with a diagram.	5	
b)	queue are a) What	e pointers to the P would be the effe	ound-Robin scheduling algorithm in which the entries in the ready Process Control Blocks. ect of putting two pointers to the same process in the ready queue? r advantage and a disadvantage of this scheme?	·4	
c)					
d)	Suppose the following processes arrive for execution as shown below.				
	Process Arrival Time Burst Time				
	P1	0	6		
	P2	1	2		
	Р3	2	5		
	0.0000000000000000000000000000000000000		ng Time and Average Turnaround Time based on non-preemptive m. You must show the formulas and all steps leading to the result.		
_	1				
a)	to the se synchroni	cond operation. I zation.	ant to each execute their first operation before any thread proceeds Below is the code that each thread runs in order to achieve this variable initialized to zero and mutex is a mutex binary semaphore,	4	
a)	to the se synchroni count is a initialized	cond operation. I zation. n integer shared to to one, that prote	Below is the code that each thread runs in order to achieve this	4	
a)	count is a initialized step1Don You are to made to	cond operation. I zation. n integer shared to one, that prote e is a semaphore	Below is the code that each thread runs in order to achieve this variable initialized to zero and mutex is a mutex binary semaphore, ects this shared variable initialized to zero. It is wrong and does not work correctly. Suggest the changes to be snippet below to fix it by using only semaphores and no other	4	
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				SRN	T				
c)	Consider the	following cla	ssical synchron	nization problem called the barbershop problem A					
()	c) Consider the following classical synchronization problem called the barbershop problem. barbershop consists of a barber in a room with N chairs. If a customer enters the barbershop consists of a barber in a room with N chairs.								
	and all chairs are occupied, then the customer leaves the shop. If the barber is busy, but chairs are available, then the customer sits in one of the free chairs and awaits his turn. The barber								
				ustomer after he finishes one haircut. If there are no					
				bes to sleep. If the barber is asleep when a customer					
	arrives, the customer wakes up the barber to give him a haircut. A waiting customer vacates his chair after his haircut completes.								
	chair after his haircut completes.								
	Write the pseudocode for the customer and barber threads with suitable synchronization.								
	You must use only semaphores to solve this problem. Use the standard notation of in								
	up/down or v	emaphore variable.							
	The following	variables (3	semaphores a	nd a count) are provided to you for your solution. You					
				y additional variables if required.					
	I The second of Street and Street		CONTROL CONTROL DE MANAGEMENT DE LA CONTROL	arber = 0; int waiting count = 0;					
d)			- 31	litions to cause deadlock in a system. Explain each of	. 8				
",		in a sentence		7	(6+				
	ii. State	any two pos	sible remedie	es to the deadlock problem in the Dining-Philosophers					
	probl	em.		***					
l. a)	Consider the	following na	ge reference s	tring.	5				
. -,		rono ming pag	se reference s		"				
	122121	E 6 2 1 2	, 3, 7, 6, 3, 2, 1	2 2 6					
	1, 2, 3, 4, 2, 1	, 3, 0, 2, 1, 2,	, 3, 7, 0, 3, 2, 1	1, 2, 3, 6					
	How many page faults would occur for Optimal replacement algorithm assuming four frames								
	result.	i the frames	are initially e	mpty. Your answer must show all steps leading to the					
b)									
-,	(A) 15-75 (C)			is the effective access time (in ns) if the TLB hit ratio is	5				
	90% and there is no page-fault?								
c)	i. Consi	der a logical	address space	of eight pages of 1024 words each, mapped onto a	5				
"	physical memory of 32 frames.								
	How many bits are in the logical address?								
	How many bits are in the physical address?								
	ii. Consider the following segment table:								
	Segme	nt Base	Length	1	8				
	Jegine	219							
4	0	215							
	0) 14						
	1	2300							
	-	2300 90	100						
	1 2	2300	100 7 580						
	1 2 3 4 What	2300 90 1327 1952 are the phys) tuples?	100 7 580 2 96	for the following logical addresses given as (segment,	**				
	1 2 3 4 What offset a) 0, 4	2300 90 1327 1952 are the phys) tuples?	100 7 580 2 96	for the following logical addresses given as (segment,	**				
	1 2 3 4 What	2300 90 1327 1952 are the phys) tuples? 30	100 7 580 2 96	for the following logical addresses given as (segment,	**				

		SRN						
	d)	d) Assume the following: a 32-bit virtual address space, with a 1KB page size and a linea table with each page table entry size = 4 bytes.						
		i. How many bits are in the offset portion of the virtual address?						
		ii. How many bits are in the Virtual Page Number portion of the virtual address?						
		iii. How many entries are in the table?						
		iv. What is the total size of the table?						
		v. In a live system, if the page table size = 10 MB and there are 100 processes how much memory would be occupied by the page tables?						
4.	a)	i. Consider a file system with 12 direct pointers, 1 indirect pointer and 1 double-indirect	5					
		pointer in the i-node. Assume that disk blocks are 8 KB size and each pointer to a disk block requires 4 bytes. What is the largest possible file that can be supported with this design?	(3 + 2)					
		ii. If the same file system supports an additional triple-indirect pointer, what is the largest file that can be supported?						
		Note : Give the answer as an expression and calculate the final numeric value for both the questions						
	b)	What is the purpose of System Hardening? Explain any 3 best practices for System Hardening.	5					
	c)	Consider a system that supports the strategies of contiguous, linked, and indexed allocation. What criteria should be used in deciding which strategy is best suited for a particular file?	5					
	d)	i. What is the basic design principle behind FFS? How does FFS determine where to put file data blocks?	5 (3 + 2)					
		ii. What is the difference between inode bitmap and inode table?						
5.	a)	Explain with a diagram how Direct Memory Access (DMA) technique improves system performance.	5					
	b)	Explain any three techniques used by operating systems to improve disk drive performance.	5					
	c)	Explain the three main delays in getting data from disk.	6					
	d)	Suppose we have a disk with 200 tracks (numbered from 0 to 199) and the head is initially at track 100. There is a queue of disk access requests for tracks 27, 129, 110, 186, 147, 41, 10, 64 and 120.	4					
		If Shortest-Seek Time First (SSTF) is being used for scheduling the disk access, show all the requests serviced and calculate the average number of tracks visited per request.						