

**Memory Management – Page Table Structure** 

**Dr Rahul Nagpal**Computer Science



# Memory Management – Page Table Structure

**Dr. Rahul Nagpal**Computer Science

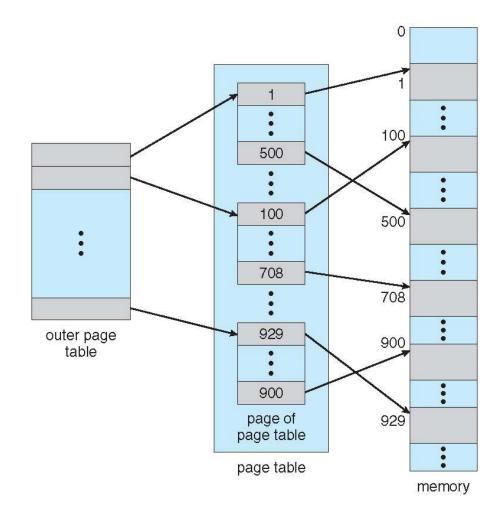
### Structure of Page Table

PES UNIVERSITY ONLINE

- Memory structures for paging can get huge using straight-forward methods
  - Consider a 32-bit logical address space as on modern computers
  - Page size of 4 KB (2<sup>12</sup>)
  - Page table would have 1 million entries (2<sup>32</sup> / 2<sup>12</sup>)
  - If each entry is 4 bytes -> 4 MB of physical address space / memory for page table alone
    - That amount of memory used to cost a lot
    - Don't want to allocate that contiguously in main memory
- Hierarchical Paging
- Hashed Page Tables
- Inverted Page Tables

### Hierarchical Page Table – Two Level Page table

- Break up the logical address space into multiple page tables
- A simple technique is a two-level page table
- We then page the page table



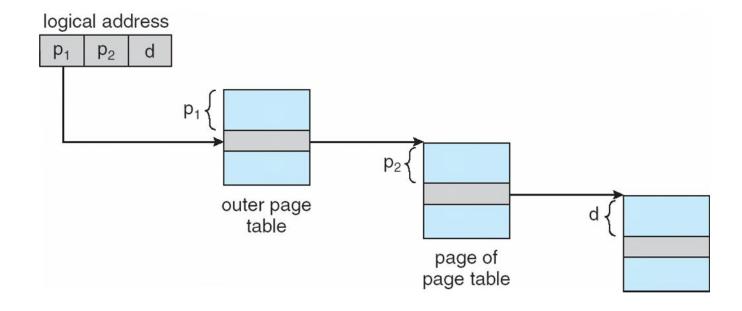


### Two Level Paging Example

- A logical address (on 32-bit machine with 1K page size) is divided into:
  - a page number consisting of 22 bits
  - a page offset consisting of 10 bits
- Since the page table is paged, the page number is further divided into:
  - a 12-bit page number
  - a 10-bit page offset
- Known as forward-mapped page table
- where  $p_1$  is an index into the outer page table, and  $p_2$  is the displacement within the page of the inner page table



### **Address Translation Scheme**





### 64 Bit Logical Address Space

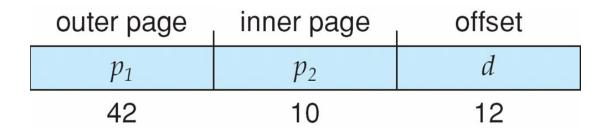
- Even two-level paging scheme not sufficient
- If page size is 4 KB (2<sup>12</sup>)
  - Then page table has 2<sup>52</sup> entries
  - If two level scheme, inner page tables could be 2<sup>10</sup> 4-byte entries
  - Address would look like

outer page	inner page	page offset	
$p_1$	$p_2$	d	
42	10	12	

- Outer page table has 2<sup>42</sup> entries or 2<sup>44</sup> bytes
- One solution is to add a 2<sup>nd</sup> outer page table
- But in the following example the 2<sup>nd</sup> outer page table is still 2<sup>34</sup> bytes in size
  - And possibly 4 memory access to get to one physical memory location



## Three Level Paging Scheme



2nd outer page	outer page	inner page	offset
$p_1$	$p_2$	$p_3$	d
32	10	10	12





# **THANK YOU**

**Dr Rahul Nagpal** 

**Computer Science** 

rahulnagpal@pes.edu