

Memory Management – Linking, Swapping

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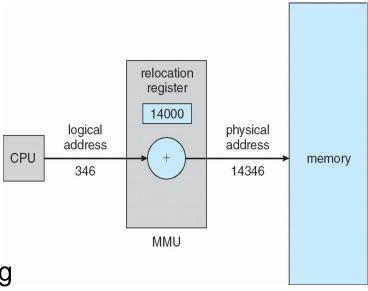


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Dynamic Allocation using Relocation Register

- Routine is not loaded until it is called
- Better memory-space utilization; unused routine is never loaded
- All routines kept on disk in relocatable load format
- Useful when large amounts of code are needed to handle infrequently occurring cases
- No special support from the operating system is required
 - Implemented through program design
 - OS can help by providing libraries to implement dynamic loading





Dynamic Linking

- Static linking system libraries and program code combined by the loader into the binary program image
- Dynamic linking –linking postponed until execution time
- Small piece of code, stub, used to locate the appropriate memory-resident library routine
- Stub replaces itself with the address of the routine, and executes the routine
- Operating system checks if routine is in processes' memory address
 - If not in address space, add to address space
- Dynamic linking is particularly useful for libraries
- System also known as shared libraries
- Consider applicability to patching system libraries
 - Versioning may be needed



Process Swapping

- A process can be swapped temporarily out of memory to a backing store, and then brought back into memory for continued execution
 - Total physical memory space of processes can exceed physical memory
- Backing store fast disk large enough to accommodate copies of all memory images for all users; must provide direct access to these memory images
- Roll out, roll in swapping variant used for prioritybased scheduling algorithms; lower-priority process is swapped out so higher-priority process can be loaded and executed
- Major part of swap time is transfer time; total transfer time is directly proportional to the amount of memory swapped
- System maintains a ready queue of ready-to-run processes which have memory images on disk



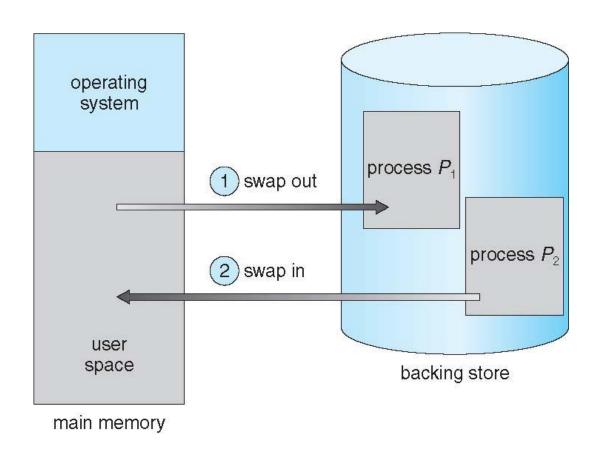
Process Swapping (Contd.)

- Does the swapped out process need to swap back in to same physical addresses?
- Depends on address binding method
 - Plus consider pending I/O to / from process memory space
- Modified versions of swapping are found on many systems (i.e., UNIX, Linux, and Windows)
 - Swapping normally disabled
 - Started if more than threshold amount of memory allocated
 - Disabled again once memory demand reduced below threshold



Schematic View of Swapping





Context Switch Time Including Swapping

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- If next processes to be put on CPU is not in memory, need to swap out a process and swap in target process
- Context switch time can then be very high
- 100MB process swapping to hard disk with transfer rate of 50MB/sec
 - Swap out time of 2000 ms
 - Plus swap in of same sized process
 - Total context switch swapping component time of 4000ms (4 seconds)
- Can reduce if reduce size of memory swapped by knowing how much memory really being used
 - System calls to inform OS of memory use via request_memory() and release_memory()

Context Switch Time Including Swapping (Contd.)

- Other constraints as well on swapping
 - Pending I/O can't swap out as I/O would occur to wrong process
 - Or always transfer I/O to kernel space, then to I/O device
 - Known as double buffering, adds overhead
- Standard swapping not used in modern operating systems
 - But modified version common
 - Swap only when free memory extremely low



Swapping on Mobile System

- Not typically supported
 - Flash memory based
 - Small amount of space
 - Limited number of write cycles
 - Poor throughput between flash memory and CPU on mobile platform
- Instead use other methods to free memory if low
 - iOS asks apps to voluntarily relinquish allocated memory
 - Read-only data thrown out and reloaded from flash if needed
 - Failure to free can result in termination
 - Android terminates apps if low free memory, but first writes application state to flash for fast restart
 - Both OSes support paging as discussed below





THANK YOU

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