

---

# I. Curves, part 1

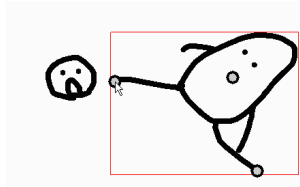
Summer Term 2025

Prof. Dr. Dr. Fellner  
Dr. Daniel Ströter

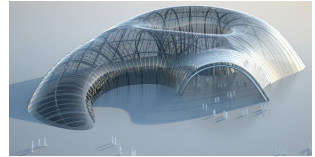
Why do we need curves in computer science/graphics?

## Why do we need curves in computer science/graphics?

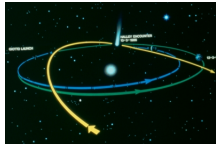
- Parametric design is art and architecture
- Describing trajectories
- Motion and smooth detail in movies
- Smooth shapes in product design/engineering
- ...



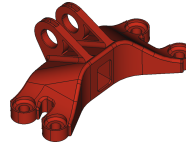
Igarashi (Eurographics '08)



Pottmann et al. "Freeform surfaces from single curved panels"



<https://sci.esa.int/web/giotto/-/36674-comet-encounter-diagram>



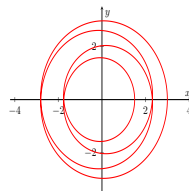
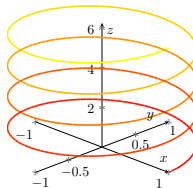
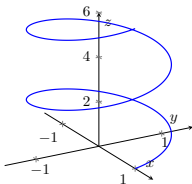
SimJEB Bracket 10

The function

$$\mathbf{X} : [a, b] \rightarrow \mathbb{R}^d, \mathbf{X}(t) = \begin{pmatrix} x_1(t) \\ \vdots \\ x_d(t) \end{pmatrix} \in \mathbb{R}^d, t \in [a, b]$$

is a **parametric** representation of the curve:

$$C_{\mathbf{X}} = \{\mathbf{X}(t) : t \in [a, b]\}.$$

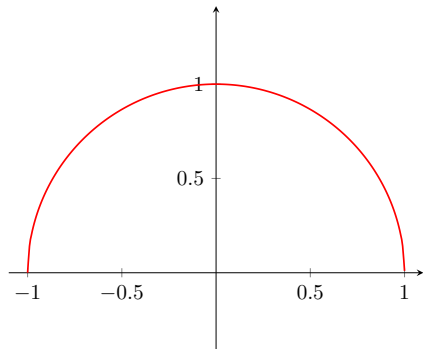


## Facts & Examples

The functions

$$\mathbf{x}(t) = \begin{pmatrix} \cos(\pi + t) \\ \sin(t) \end{pmatrix}, t \in [0, \pi]$$

$$\tilde{\mathbf{x}}(t) = \begin{pmatrix} t \\ \sqrt{1-t^2} \end{pmatrix}, t \in [-1, 1]$$



parameterize the same curve  $C_{\mathbf{x}} = C_{\tilde{\mathbf{x}}}$  (semicircle).

Likewise, the terms

$$\mathbf{x}(t) = \begin{pmatrix} t^2 \\ t^2 \end{pmatrix} = t^2 \begin{pmatrix} 1 \\ 1 \end{pmatrix}, t \in [0, 1]$$

$$\tilde{\mathbf{x}}(t) = \begin{pmatrix} t \\ t \end{pmatrix} = t \begin{pmatrix} 1 \\ 1 \end{pmatrix}, t \in [0, 1]$$

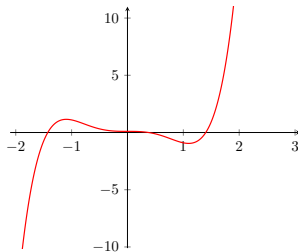
parameterize the same curves.

⇒ parametric representations of a curve are not unique!

For the real function,  $f : [a, b] \rightarrow \mathbb{R}$

$$\mathbf{x}_f(t) = \begin{pmatrix} t \\ f(t) \end{pmatrix} \in \mathbb{R}^2, t \in [a, b]$$

parameterizes a special curve  $C_{\mathbf{x}_f}$ :  
the graph of  $f$   
(also: a **functional** curve).



The curve  $\mathbf{X}: [a, b] \rightarrow \mathbb{R}^n$  is regular iff.  $\mathbf{X}'(t) \neq \mathbf{0}$  for all  $t \in [a, b]$ .

The curve  $\mathbf{X}: [a, b] \rightarrow \mathbb{R}^n$  is singular at  $\mathbf{X}(t)$  iff.  $\mathbf{X}'(t) = \mathbf{0}$ .

Neil's parabola ...

$$\mathbf{X}(t) = \begin{pmatrix} t^2 \\ t^3 \end{pmatrix}, t \in [-1, 1]$$

... is **non-regular**, because

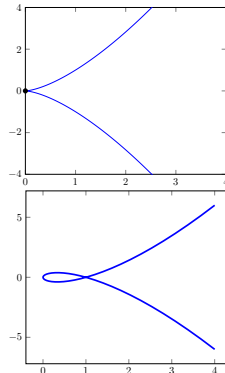
$$\mathbf{X}'(0) = \begin{pmatrix} 2 \cdot 0 \\ 3 \cdot 0^2 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \end{pmatrix}$$

For

$$\mathbf{X}(t) = \begin{pmatrix} t^2 \\ t^3 - t \end{pmatrix}, t \in [-2, 2]$$

applies  $\mathbf{X}(-1) = \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \mathbf{X}(1)$ .

$\mathbf{X}$  exhibits a **self-intersection** in  $t = -1$ .





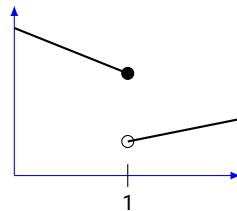
The function  $f(t)$  is **continuous** in  $t_0$  iff.

1.  $\lim_{t \rightarrow t_0} f(t)$  exists and
2.  $\lim_{t \rightarrow t_0} f(t) = f(t_0)$

In other words:  $f$  is continuous iff. for each  $\epsilon > 0$  there is a  $\delta > 0$  such that

$$|f(t) - f(t_0)| < \epsilon$$

in the local neighborhood of  $t_0$  with radius  $\delta$  (thus  $|t - t_0| < \delta$ ).



non-continuous function in  $t_0 = 1$

If  $f$  is continuous for all  $t \in [a, b]$ , then  $f$  is continuous ( $f \in C^0$ ) in  $[a, b]$ .

Special class of curves: **Polynomial curves**  $C_P$  determined by

$$\mathbf{P}(t) = \begin{pmatrix} p_1(t) \\ \vdots \\ p_d(t) \end{pmatrix} = \mathbf{a}_0 + t \mathbf{a}_1 + \dots + t^q \mathbf{a}_q \in \mathbb{R}^d, \quad t \in [a, b],$$

where  $\mathbf{a}_i \in \mathbb{R}^d$ ,  $i = 0, \dots, q$  (coefficient vectors).

$p_i$  are *polynomial functions*,  $q$  is the *degree* of  $\mathbf{P}$ .

Monomial representation (also referred to: Taylor representation in 0), because  $q + 1$  monomial

$$1, t, \dots, t^q$$

are used. Therefore:  $(q + 1) \cdot d$  degrees of freedom.

## Examples

- Neil's parabola

- 

$$\mathbf{P}(t) = \begin{pmatrix} t + 3t^2 \\ -2 - 3t^2 \\ 1 + t \end{pmatrix} = \begin{pmatrix} 0 \\ -2 \\ 1 \end{pmatrix} + \begin{pmatrix} 1 \\ 0 \\ 1 \end{pmatrix} \cdot t + \begin{pmatrix} 3 \\ -3 \\ 0 \end{pmatrix} \cdot t^2$$

- corresponding tangent vector (velocity)

$$\mathbf{P}'(t) = \begin{pmatrix} 1 + 6t \\ -6t \\ 1 \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \\ 1 \end{pmatrix} + \begin{pmatrix} 6 \\ -6 \\ 0 \end{pmatrix} \cdot t$$

- 2. derivative  $\mathbf{P}''(t) = \begin{pmatrix} 6 \\ -6 \\ 0 \end{pmatrix}$

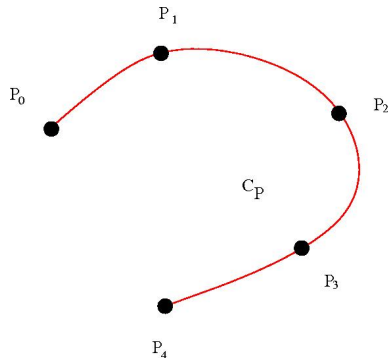
Beneficial properties of polynomial curves  $\mathbf{P} = \mathbf{P}(t) = \sum_{i=0}^q \mathbf{a}_i t^i$ :

- differentiable as often as required
- Relies on vector space: If  $\mathbf{P}_1, \mathbf{P}_2$  are polynomial curves and  $\alpha, \beta \in \mathbb{R}$ , then  $\alpha \cdot \mathbf{P}_1 + \beta \cdot \mathbf{P}_2$  is also a polynomial curve.
- Interpolation can always be applied
- Fast and exact evaluation
- ...

For an arbitrary choice of interpolation points  $a \leq t_0 < \dots < t_q \leq b$  and data points  $\mathbf{P}_0, \dots, \mathbf{P}_q \in \mathbb{R}^d$  there is **a unique** polynomial curve  $\mathbf{P}$  of degree  $q$  such that

$$\mathbf{P}(t_i) = \mathbf{P}_i, \quad i = 0, \dots, q.$$

In the example:  $q = 4$



# Choice of Interpolation Points: Typical Parametrization

Typically only the data points  $\mathbf{P}_0, \dots, \mathbf{P}_q \in \mathbb{R}^d$  are given.  
Frequent choice of parametrization:

**equidistant**  $t_i = a + i \cdot \frac{b-a}{q}, i = 0, \dots, q.$

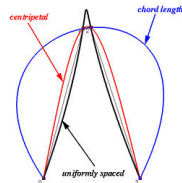
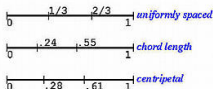
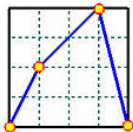
Other choices:

**cordal**  $t_{i+1} - t_i = \|\mathbf{P}_{i+1} - \mathbf{P}_i\|_2, i = 0, \dots, q - 1$

**centripetal**  $t_{i+1} - t_i = \sqrt{\|\mathbf{P}_{i+1} - \mathbf{P}_i\|_2}, i = 0, \dots, q - 1$

**Tschebyscheff**  $t_i = \cos\left(\frac{2(q-i)+1}{q+1} \frac{\pi}{2}\right), i = 0, \dots, q.$

( $\|\cdot\|_2$  is Euclidean norm)



C.K. Shene, <http://www.cs.mtu.edu/~shene/COURSES/cs3621/NOTES/INT-APP/PARA-centripetal.html>

## Interpolation in monomial base

Determine polynomial  $\mathbf{P}(t) = \sum_{i=0}^q \mathbf{a}_i t^i$  such that  $\mathbf{P}(t_i) = \mathbf{P}_i$ , for  $i = 0, \dots, q$   
leads to **Vandermonde** matrix  $\mathbf{V}$

$$\begin{pmatrix} 1 & t_0 & t_0^2 & \dots & t_0^q \\ 1 & t_1 & t_1^2 & \dots & t_1^q \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & t_q & t_q^2 & \dots & t_q^q \end{pmatrix} \begin{pmatrix} \mathbf{a}_0 \\ \mathbf{a}_1 \\ \vdots \\ \mathbf{a}_q \end{pmatrix} = \begin{pmatrix} \mathbf{P}_0 \\ \mathbf{P}_1 \\ \vdots \\ \mathbf{P}_q \end{pmatrix}, \text{ i.e. } \mathbf{V}\mathbf{a} = \mathbf{P}$$

Coefficient vector  $\mathbf{a}$  computed as  $\mathbf{a} = \mathbf{V}^{-1}\mathbf{P}$

- It works out, because determinant  $\det(\mathbf{V}) \neq 0$  for  $t_0 < t_1 < t_2 < \dots$
- but:  $\mathbf{V}$  can be ill-conditioned  $\rightarrow$  numerically unstable
- expensive, because inverse of  $\mathbf{V}$  needs to be computed (Cost:  $\mathcal{O}(q^3)$ ).

### Lagrange Representation

$$\mathbf{P}(t) = \sum_{i=0}^q \ell_i(t) \mathbf{P}_i, \quad t \in [a, b]$$

relies on Lagrange base  $\ell_i$ ,  $i = 0, \dots, q$ , (degree  $q$ ):

$$\ell_i(t) = \prod_{j=0, j \neq i}^q \frac{t - t_j}{t_i - t_j} = \frac{(t - t_0) \dots (t - t_{i-1})(t - t_{i+1}) \dots (t - t_q)}{(t_i - t_0) \dots (t_i - t_{i-1})(t_i - t_{i+1}) \dots (t_i - t_q)}, \quad t \in [a, b]$$

The Lagrange polynomials form the partition of unity:

$$\sum_i^q \ell_i(t) = 1$$



The construction of  $\ell_i$  implies:  $\ell_i(t_j) = \delta_{ij} := \begin{cases} 1, & i = j \\ 0, & i \neq j \end{cases}$ .

Thus the polynomial

$$\mathbf{P}(t) = \sum_{i=0}^q \ell_i(t) \mathbf{P}_i, \quad t \in [a, b]$$

is the unique solution to  $\mathbf{P}(t_j) = \mathbf{P}_j, j = 0, \dots, q$ .

- easy to calculate
- no matrix inversion
- relatively expensive evaluation of  $\ell_i$

**Example:** Your homework!

## Newton Representation

$$\mathbf{P}(t) = \sum_{i=0}^q \omega_i(t) \Delta(t_0, \dots, t_i), \quad t \in [a, b]$$

relies on Newton base  $\omega_i$ ,  $i = 0, \dots, q$ , (degree  $i$ ):

$$\omega_0(t) = 1, \quad \omega_i(t) = (t - t_0) \dots (t - t_{i-1}), \quad t \in [a, b]$$

**Newton's idea:** For  $j = 1, \dots, q$ , determine  $\mathbf{P}^{[j]}$  as a solution to  $\mathbf{P}^{[j]}(t_i) = \mathbf{P}_i$ ,  $i = 0, \dots, j$  by (re-)usage of  $\mathbf{P}^{[j-1]}$ .

$$\text{q=0} \quad \mathbf{P}^{[0]}(t) = \mathbf{P}_0$$

$$\text{q=1} \quad \mathbf{P}^{[1]}(t) = \mathbf{P}_0 + \overbrace{(t - t_0)}^{\omega_1(t)} \frac{\mathbf{P}_1 - \mathbf{P}_0}{t_1 - t_0}$$

$$q=0 \quad \mathbf{P}^{[0]}(t) = \mathbf{P}_0$$

$$q=1 \quad \mathbf{P}^{[1]}(t) = \mathbf{P}_0 + \overbrace{(t - t_0)}^{\omega_1(t)} \frac{\mathbf{P}_1 - \mathbf{P}_0}{t_1 - t_0}$$

$$q=2 \quad \text{Approach } \mathbf{P}^{[2]}(t) = \mathbf{P}^{[1]}(t) + \overbrace{(t - t_0)(t - t_1)}^{\omega_2(t)} \mathbf{B}$$

Straightforwardly, it holds that:  $\mathbf{P}^{[2]}(t_i) = \mathbf{P}^{[1]}(t_i)$  for  $i = 0, 1$ .  
Remaining condition  $\mathbf{P}^{[2]}(t_2) = \mathbf{P}_2$  leads to

$$\begin{aligned} \mathbf{B} &= \left( \left( \frac{\mathbf{P}_2 - \mathbf{P}_0}{t_2 - t_0} \right) - \left( \frac{\mathbf{P}_1 - \mathbf{P}_0}{t_1 - t_0} \right) \right) / (t_2 - t_1) \\ &= \left( \left( \frac{\mathbf{P}_2 - \mathbf{P}_1}{t_2 - t_1} \right) - \left( \frac{\mathbf{P}_1 - \mathbf{P}_0}{t_1 - t_0} \right) \right) / (t_2 - t_0) \end{aligned}$$

For general  $q \geq 0$

$k$ -th difference regarding  $a \leq t_j \dots t_{j+k} \leq b$  is recursively defined by  $\Delta(t_i) := \mathbf{P}_i$ ,  $i = j, \dots, j+k$ , and

$$\Delta(t_j, \dots, t_{j+k}) := \frac{\Delta(t_{j+1}, \dots, t_{j+k}) - \Delta(t_j, \dots, t_{j+k-1})}{t_{j+k} - t_j}.$$

Computation as triangle scheme – Cost:  $\mathcal{O}(q^2)$

Formula of Newton interpolation:

$$\begin{aligned} \mathbf{P}(t) &= \Delta(t_0) + \omega_1(t) \Delta(t_0, t_1) + \dots + \omega_q(t) \Delta(t_0, \dots, t_q) \\ &= \sum_{i=0}^q \omega_i(t) \Delta(t_0, \dots, t_i), \quad t \in [a, b]. \end{aligned}$$

$$\begin{array}{cccc|cccc}
& & & t_0 & \Delta(t_0) & & & \\
& & & t_1 - t_0 & & \Delta(t_0, t_1) & & \\
& & t_2 - t_0 & t_1 & \Delta(t_1) & & \Delta(t_0, t_1, t_2) & \\
& & \vdots & t_2 - t_1 & & \Delta(t_1, t_2) & \vdots & \\
& \ddots & \vdots & \vdots & t_2 & \Delta(t_2) & \vdots & \vdots & \ddots \\
t_q - t_0 & & \vdots & \vdots & \vdots & \vdots & \vdots & \Delta(t_0, \dots, t_q) \\
& \ddots & \vdots & \vdots & \vdots & \vdots & \vdots & \ddots \\
& & t_q - t_{q-2} & \vdots & t_{q-1} & \Delta(t_{q-1}) & & \Delta(t_{q-2}, t_{q-1}, t_q) \\
& & t_q - t_{q-1} & t_q & \Delta(t_q) & \Delta(t_{q-1}, t_q) & & 
\end{array}$$

**Example:** determine polynomial  $p: [-2, 2] \rightarrow \mathbb{R}$  of degree 4 such that

$$p(i) = |i|, \quad i = -2, \dots, 2.$$

Computation of shared differences (triangle scheme):

				-2	2															
				1	-1	1	-1	0												
				2	1	0	1	-1	1/3											
				3	2	1	0	1	-1	1/3	-1/6									
				4	3	2	1	0	1	0	-1/3	1/6								
				2	1	0	1	1	0	1	-1/3	1/6	-1/6							
				1	0	1	1	1	0	1	-1/3	1/6	-1/6	1/6						
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6					
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0	1	1	0	1	-1/3	1/6	-1/6	1/6	-1/6	1/6				
				2	1	0														

$$\begin{array}{cccc|cccc}
 & & & -2 & 2 & & & \\
 & & 1 & & & -1 & & \\
 & 2 & & -1 & 1 & & 0 & \\
 3 & & 1 & & & -1 & & 1/3 \\
 4 & 2 & & 0 & 0 & & 1 & & -1/6 \\
 & 3 & 1 & & & 1 & & -1/3 \\
 & 2 & & 1 & 1 & & 0 & \\
 & 1 & & & & 1 & & \\
 & & & 2 & 2 & & & 
 \end{array}$$
$$\begin{aligned} p(t) &= 2 + 1(t-2) + 0(t-2)(t-1) + (-1/3)(t-2)(t-1)t + (-1/6)(t-2)(t-1)t(t+1) \\ &= \boxed{0} + \boxed{(-1)}t + \boxed{0}t(t+1) + \boxed{1/3}t(t+1)(t+2) + \boxed{(-1/6)}t(t+1)(t+2)(t-1) \end{aligned}$$

-  **GRIS**

Naive, direct evaluation:  $\mathcal{O}(q^2)$

Monomial representation

$$\mathbf{P}(\xi) = \sum_{i=0}^q \underbrace{\xi \dots \xi}_{\xi^i} \mathbf{a}_i \quad (i-1) \text{ multiplications}$$

Lagrange representation

$$\mathbf{P}(\xi) = \sum_{i=0}^q \underbrace{\frac{(\xi - t_0) \dots (\xi - t_{i-1})(\xi - t_{i+1}) \dots (\xi - t_q)}{(t_i - t_0) \dots (t_i - t_{i-1})(t_i - t_{i+1}) \dots (t_i - t_q)}}_{\ell_i(\xi)} \mathbf{P}(t_i) \quad (2q-1) \text{ multiplications}$$

Newton representation

$$\mathbf{P}(\xi) = \sum_{i=0}^q \underbrace{(\xi - t_0) \dots (\xi - t_{i-1})}_{\omega_i(\xi)} \mathbf{b}_i \quad (i-1) \text{ multiplications}$$



Suppose  $\mathbf{P}$  is provided in Newton representation

$$\begin{aligned}\mathbf{P}(t) &= \mathbf{b}_0 + (t - t_0) \mathbf{b}_1 + (t - t_0)(t - t_1) \mathbf{b}_2 + \dots \\ &\quad + (t - t_0) \dots (t - t_{q-2}) \mathbf{b}_{q-1} + (t - t_0) \dots (t - t_{q-1}) \mathbf{b}_q \\ &= \mathbf{b}_0 + (t - t_0) (\mathbf{b}_1 + (t - t_1) (\mathbf{b}_2 + \dots \\ &\quad + (t - t_{q-2}) (\mathbf{b}_{q-1} + (t - t_{q-1}) \mathbf{b}_q) \dots)) .\end{aligned}$$

## Algorithm: modified Horner scheme

Compute  $\mathbf{P}(\xi)$ : Set  $\mathbf{b}_i^{[0]} = \mathbf{b}_i$ ,  $i = 0, \dots, q$ ,  $\mathbf{b}_q^{[1]} := \mathbf{b}_q^{[0]}$ ,  
and compute for each  $i = q, \dots, 1$  successively

$$\mathbf{b}_{i-1}^{[1]} = \mathbf{b}_{i-1}^{[0]} + (\xi - t_{i-1}) \mathbf{b}_i^{[1]}.$$

Compute  $\mathbf{P}(\xi)$ : Set  $\mathbf{b}_i^{[0]} = \mathbf{b}_i$ ,  $i = 0, \dots, q$ ,  $\mathbf{b}_q^{[1]} := \mathbf{b}_q^{[0]}$ ,  
and compute for each  $i = q, \dots, 1$  successively

$$\mathbf{b}_{i-1}^{[1]} = \mathbf{b}_{i-1}^{[0]} + (\xi - t_{i-1}) \mathbf{b}_i^{[1]}.$$

Scheme:

$$\begin{array}{ccccccc}
 \mathbf{b}_q = \mathbf{b}_q^{[0]} & \mathbf{b}_{q-1} = \mathbf{b}_{q-1}^{[0]} & \mathbf{b}_{q-2} = \mathbf{b}_{q-2}^{[0]} & \dots & \mathbf{b}_2 = \mathbf{b}_2^{[0]} & \mathbf{b}_1 = \mathbf{b}_1^{[0]} & \mathbf{b}_0 = \mathbf{b}_0^{[0]} \\
 & +(\xi - t_{q-1}) \mathbf{b}_q^{[1]} & +(\xi - t_{q-2}) \mathbf{b}_{q-1}^{[1]} & \dots & +(\xi - t_2) \mathbf{b}_3^{[1]} & +(\xi - t_1) \mathbf{b}_2^{[1]} & +(\xi - t_0) \mathbf{b}_1^{[1]} \\
 \hline
 \mathbf{b}_q^{[1]} & \mathbf{b}_{q-1}^{[1]} & \mathbf{b}_{q-2}^{[1]} & \dots & \mathbf{b}_2^{[1]} & \mathbf{b}_1^{[1]} & \mathbf{b}_0^{[1]} = \mathbf{P}(\xi)
 \end{array}$$

Cost:  $\mathcal{O}(q)$

# Complete modified Horner scheme

Iterate

$\mathbf{b}_q = \mathbf{b}_q^{[0]}$	$\mathbf{b}_{q-1} = \mathbf{b}_{q-1}^{[0]}$	$\mathbf{b}_{q-2} = \mathbf{b}_{q-2}^{[0]}$	$\dots$	$\mathbf{b}_2 = \mathbf{b}_2^{[0]}$	$\mathbf{b}_1 = \mathbf{b}_1^{[0]}$	$\mathbf{b}_0 = \mathbf{b}_0^{[0]}$
$\mathbf{b}_q^{[1]}$	$\mathbf{b}_{q-1}^{[1]}$	$\dots$	$\dots$	$\mathbf{b}_2^{[1]}$	$\mathbf{b}_1^{[1]}$	$\mathbf{b}_0^{[1]} = \mathbf{P}(\xi)$
$\mathbf{b}_q^{[2]}$	$\mathbf{b}_{q-1}^{[2]}$	$\dots$	$\dots$	$\mathbf{b}_2^{[2]}$	$\mathbf{b}_1^{[2]} = \mathbf{P}'(\xi)$	
$\mathbf{b}_q^{[3]}$	$\mathbf{b}_{q-1}^{[3]}$	$\dots$	$\dots$	$\mathbf{b}_2^{[3]} = \frac{\mathbf{P}''(\xi)}{2}$		
$\vdots$						
$\mathbf{b}_q^{[q]}$	$\mathbf{b}_{q-1}^{[q]} = \frac{\mathbf{P}^{(q-1)}(\xi)}{(q-1)!}$					
$\mathbf{b}_q^{[q+1]} = \frac{\mathbf{P}^{(q)}(\xi)}{q!}$						

Determine  $\mathbf{P}(\xi), \mathbf{P}'(\xi), \dots, \mathbf{P}^{(q)}(\xi)$  with total expenditure of  $\mathcal{O}(q^2)$

# Complete modified Horner scheme Example

Suppose polynomial  $p: [-1, 2] \rightarrow \mathbb{R}$  of degree 3 is given by  $p(t) = 1 + \frac{1}{2} (t+1)t + \frac{5}{6} (t+1)t(t-1)$   
 $t_0 = -1, t_1 = 0, t_2 = 1, t_3 = 2$ , Evaluate at  $\xi = 0$

$b_i^{[0]}, i = 3, 2, 1, 0 :$	$\frac{5}{6}$	$\frac{\frac{1}{2}}{(0-1)} \frac{5}{6}$	$\frac{0}{(0-0) (-\frac{1}{3})}$	$\frac{1}{(0+1) \cdot 0}$
$b_i^{[1]}, i = 3, 2, 1 :$	$\frac{5}{6}$	$\frac{-\frac{1}{3}}{(0-0) - \frac{5}{6}}$	$\frac{0}{(0+1) (-\frac{1}{3})}$	$1 = \frac{p(0)}{0!}$
$b_i^{[2]}, i = 3, 2 :$	$\frac{5}{6}$	$\frac{-\frac{1}{3}}{(0+1) \frac{5}{6}}$	$-\frac{1}{3} = \frac{p'(0)}{1!}$	
$b_3^{[3]} :$	$\frac{5}{6}$	$\frac{1}{2} = \frac{p''(0)}{2!}$		
$b_3^{[4]} :$	$\frac{5}{6} = \frac{p'''(0)}{3!}$			

Thus,  $p(t) = 1 - \frac{1}{3} t + \frac{1}{2} t^2 + \frac{5}{6} t^3$

Monomial , (Taylor) , Lagrange , Newton representation of  $P$

$$\mathbf{a}_i = \frac{P^{(i)}(0)}{i!} , (\mathbf{a}_i = \frac{P^{(i)}(\xi)}{i!}) , P(t_i) , \mathbf{b}_i = \Delta(t_0, \dots, t_i)$$

Shortcomings (computer graphics):

Monomial, (Taylor), Lagrange, Newton representation of  $P$

$$a_i = \frac{P^{(i)}(0)}{i!}, (a_i = \frac{P^{(i)}(\xi)}{i!}), P(t_i), b_i = \Delta(t_0, \dots, t_i)$$

**Shortcomings (computer graphics):**

- difficult geometric interpretation ● ●
- not well-aligned with parameter interval  $[a, b]$ , asymmetrical ● ● ●
- (occasionally) unstable ● ● ●
- sub-optimal evaluation ●

Monomial, (Taylor), Lagrange, Newton representation of  $\mathbf{P}$

$$\mathbf{a}_i = \frac{\mathbf{P}^{(i)}(0)}{i!}, (\mathbf{a}_i = \frac{\mathbf{P}^{(i)}(\xi)}{i!}), \mathbf{P}(t_i), \mathbf{b}_i = \Delta(t_0, \dots, t_i)$$

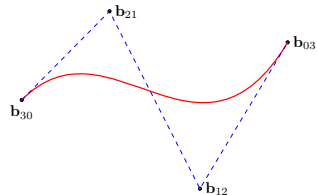
**Shortcomings (computer graphics):**

- difficult geometric interpretation ● ●
- not well-aligned with parameter interval  $[a, b]$ , asymmetrical ● ● ●
- (occasionally) unstable ● ● ●
- sub-optimal evaluation ●
- ... ● ● ●

Consequently: Bernstein-Bézier representation of polynomial curve  $\mathbf{P}$

$$\mathbf{P}(t) = \sum_{i+j=q} B_{ij}(t) \mathbf{b}_{ij}, \quad t \in [a, b]$$

$\mathbf{b}_{ij} \in \mathbb{R}^d$  **Bernstein coefficients:** control points



a cubic Bézier curve

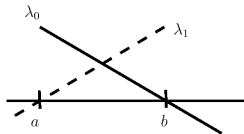
$\mathbf{b}_{q0}, \mathbf{b}_{q-1\,1}, \dots, \mathbf{b}_{0q}$  form a path: control polygon

$B_{ij}$ ,  $i + j = q$ : **Bernstein polynomials**

$$B_{ij}(t) = \frac{q!}{i!j!} \lambda_0^i(t) \lambda_1^j(t)$$

$\lambda_0, \lambda_1$ : **barycentric coordinates**

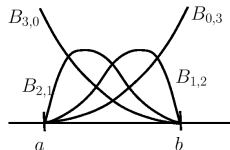




$$\lambda_0(t) = \frac{b-t}{b-a}, \quad \lambda_1(t) = \frac{t-a}{b-a}, \quad t \in \mathbb{R}$$

## Mathematical properties:

- $t \in [a, b]$  if and only if  $\lambda_0(t) \geq 0$  and  $\lambda_1(t) \geq 0$
- $\lambda_0(t) + \lambda_1(t) = 1$  and  $a \lambda_0(t) + b \lambda_1(t) = t$
- For all polynomial curves  $\mathbf{P}$  of degree 1 (linear, straight lines):  $\lambda_0(t) \mathbf{P}(a) + \lambda_1(t) \mathbf{P}(b) = \mathbf{P}(t)$



$$B_{ij}(t) = \frac{q!}{i!j!} \lambda_0^i(t) \lambda_1^j(t) = \frac{q!}{i!j!} \left(\frac{b-t}{b-a}\right)^i \left(\frac{t-a}{b-a}\right)^j$$

$q + 1$  distinct curves ( $i + j = q$ ) – basis of polynomials (degree  $q$ )

## Mathematical properties:

- $B_{ij}(t) > 0$ ,  $t \in (a, b)$  (positivity)
- $B_{ij}(t) = 0$  if and only if ( $t = a$  and  $i \neq q$ ) or ( $t = b$  and  $i \neq 0$ ) (roots)
- $B_{ij}(t) = B_{ji}(a + b - t)$  (symmetry)

Important for geometry of B-form ( . . . in a few minutes)

$$\sum_{i+j=q} B_{ij}(t) = 1 \text{ and } \sum_{i+j=q} \left( \frac{i}{q} a + \frac{j}{q} b \right) B_{ij}(t) = t$$

For all polynomial curves **P** of degree 1 (linear, straight lines):  $\sum_{i+j=q} \mathbf{P} \left( \frac{i}{q} a + \frac{j}{q} b \right) B_{ij}(t) = \mathbf{P}(t)$

**Recursive definition:**  $B_{ij}(t) = \lambda_0(t) \cdot B_{i-1,j}(t) + \lambda_1(t) \cdot B_{i,j-1}(t)$

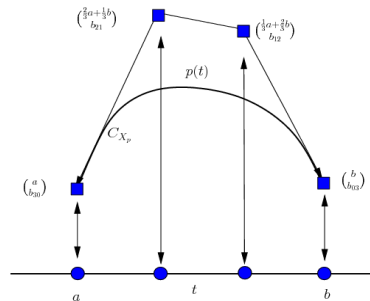
**Derivative:**  $B'_{ij}(t) = \frac{q}{b-a} ( B_{i,j-1}(t) - B_{i-1,j}(t) )$

**Maxima:** (in the *domain points*)  $\max\{B_{ij}(t) : t \in [a, b]\} = B_{ij}(\frac{i}{q} a + \frac{j}{q} b)$

$$p(t) = \sum_{i+j=q} b_{ij} B_{ij}(t), \quad t \in [a, b], \quad b_{ij} \in \mathbb{R}$$

Graph  $C_{\mathbf{x}_p}$  of  $p$  is given by

$$\begin{aligned} \mathbf{x}_p(t) &= \begin{pmatrix} t \\ p(t) \end{pmatrix} \\ &= \sum_{i+j=q} B_{ij}(t) \begin{pmatrix} \frac{i}{q} a + \frac{j}{q} b \\ b_{ij} \end{pmatrix} \end{aligned}$$



Note: domain points  $\xi_{ij} = \frac{i}{q} a + \frac{j}{q} b$ ,  $i + j = q$ , are uniquely associated with  $b_{ij}$ .

$$\mathbf{P}(t) = \sum_{i+j=q} B_{ij}(t) \mathbf{b}_{ij}, \quad t \in [a, b]$$

$\mathbf{b}_{ij} \in \mathbb{R}^d$  “control” the geometry of  $\mathbf{P}$ .

## End point interpolation

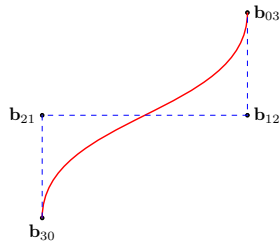
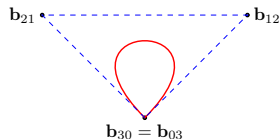
- $\mathbf{P}(a) = \mathbf{b}_{q0}$  and  $\mathbf{P}(b) = \mathbf{b}_{0q}$ .

## Affine invariance

- is a consequence of the partition of unity  $\sum B_{ij} = 1$ .

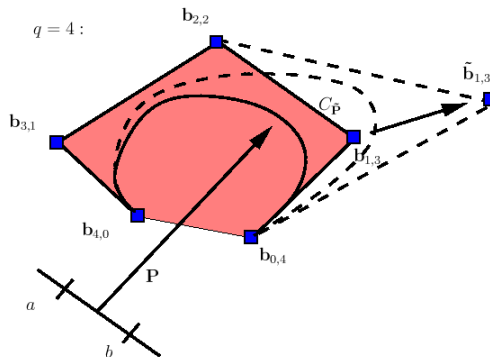
## Convex hull property

- due to  $B_{ij} \geq 0$  for  $t \in [a, b]$ , the curve is either inside or intersects the control polygon



Examples of cubic Bézier curves

Relocation of a single control point  
changes the geometry of the curve primarily at the local parts near the control point



quartic Bézier curve

## Algorithm: Determine $\ell$ -th derivative of $\mathbf{P}$

Set  $\widehat{\mathbf{b}}_{ij}^{[0]} = \mathbf{b}_{ij}$ ,  $i + j = q$  and compute for  $\ell = 1, \dots, q$  successively

$$\widehat{\mathbf{b}}_{ij}^{[\ell]} = \widehat{\mathbf{b}}_{i+1,j}^{[\ell-1]} - \widehat{\mathbf{b}}_{i+1,j}^{[\ell-1]}, \quad i + j = q - \ell.$$

Applies for  $\ell \in \{0, \dots, q\}$ :  $\mathbf{P}^{(\ell)} = \frac{q!}{(q-\ell)!} \frac{1}{(b-a)^\ell} \sum_{i+j=q-\ell} B_{ij} \widehat{\mathbf{b}}_{ij}^{[\ell]}$

In particular, it holds (derivative at  $\mathbf{P}(a)$ ):

$$\mathbf{P}^{(\ell)}(a) = \frac{q!}{(q-\ell)!} \frac{1}{(b-a)^\ell} \widehat{\mathbf{b}}_{q-\ell,0}^{[\ell]} = \frac{q!}{(q-\ell)!} \frac{1}{(b-a)^\ell} \sum_{k=0}^{\ell} (-1)^{\ell-k} \binom{\ell}{k} \mathbf{b}_{q-k,k}$$

Example ( $q = 3$ ):

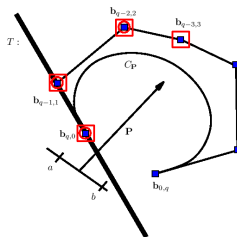
$$\mathbf{P}(t) = B_{30}(t) \mathbf{b}_{30} + B_{21}(t) \mathbf{b}_{21} + B_{12}(t) \mathbf{b}_{12} + B_{03}(t) \mathbf{b}_{03}$$

$$\mathbf{P}'(t) = \frac{3}{b-a} ( B_{20}(t) (\mathbf{b}_{21} - \mathbf{b}_{30}) + B_{11}(t) (\mathbf{b}_{12} - \mathbf{b}_{21}) + B_{02}(t) (\mathbf{b}_{03} - \mathbf{b}_{12}) )$$

$$\mathbf{P}''(t) = \frac{6}{(b-a)^2} ( B_{10}(t) (\mathbf{b}_{12} - 2 \mathbf{b}_{21} + \mathbf{b}_{30}) + B_{01}(t) (\mathbf{b}_{03} - 2 \mathbf{b}_{12} + \mathbf{b}_{21}) )$$

$$\mathbf{P}'''(t) = \frac{6}{(b-a)^3} ( \mathbf{b}_{03} - 3 \mathbf{b}_{12} + 3 \mathbf{b}_{21} - \mathbf{b}_{30} )$$





Control points  $\mathbf{b}_{q,0}, \mathbf{b}_{q-1,1}, \dots, \mathbf{b}_{q-\ell,\ell}$  determine  $\mathbf{P}^{(\ell)}(a)$ .

Derivatives  $\mathbf{P}(a), \mathbf{P}'(a), \dots, \mathbf{P}^{(\ell)}(a)$  determine  $\mathbf{b}_{q-\ell,\ell}$  (**Hermite!**)

Tangents in end points of the curve (e.g.,  $a$ ):

$$T = \{\mathbf{z} \in \mathbb{R}^d : \mathbf{z} = \underbrace{\mathbf{P}(a)}_{\mathbf{b}_{q,0}} + \tau \underbrace{\mathbf{P}'(a)}_{\frac{q}{b-a} (\mathbf{b}_{q-1,1} - \mathbf{b}_{q,0})}, \tau \in \mathbb{R}\}$$

We seek to compute:  $\mathbf{P}(\xi)$  for a certain  $\xi \in \mathbb{R}$

## de Casteljau algorithm

Set  $\mathbf{b}_{ij}^{[0]} = \mathbf{b}_{ij}$ ,  $i + j = q$ , and compute for  $\ell = 1, \dots, q$ , successively

$$\mathbf{b}_{ij}^{[\ell]} = \lambda_0(\xi) \mathbf{b}_{i+1,j}^{[\ell-1]} + \lambda_1(\xi) \mathbf{b}_{i,j+1}^{[\ell-1]}, \quad i + j = q - \ell.$$

$$\mathbf{P}(\xi) = \mathbf{b}_{00}^{[q]}$$

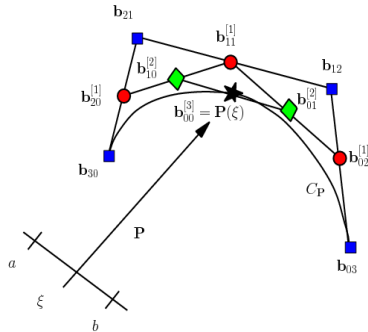
Idea of proof: (Inductive contention, Recursion of B-Polynomials, index-shift)

$$\begin{aligned} \mathbf{P}(\xi) &= \sum_{i+j=q} B_{ij}(\xi) \mathbf{b}_{ij} = \sum_{i+j=q} (\lambda_0(\xi) B_{i-1,j}(\xi) + \lambda_1(\xi) B_{i,j-1}(\xi)) \mathbf{b}_{ij} \\ &= \sum_{i+j=q-1} B_{ij}(\xi) (\lambda_0(\xi) \mathbf{b}_{i+1,j}^{[0]} + \lambda_1(\xi) \mathbf{b}_{i,j+1}^{[0]}) = \sum_{i+j=q-1} B_{ij}(\xi) \mathbf{b}_{ij}^{[1]}. \end{aligned}$$

For the polynomial curve of degree  $q - 1$ , one can obtain  $\tilde{\mathbf{P}}(\xi) = \mathbf{b}_{00}^{[q]}$  using de Casteljau and it holds that  $\mathbf{P}(\xi) = \tilde{\mathbf{P}}(\xi)$ .

Example:  $q = 3$

$$\mathbf{P} \equiv B_{30} \mathbf{b}_{30} + B_{21} \mathbf{b}_{21} + B_{12} \mathbf{b}_{12} + B_{03} \mathbf{b}_{03}$$



$$\begin{array}{ccccccc}
 & & \mathbf{b}_{q0}^{[0]} & & & & \\
 & & & > & \mathbf{b}_{q-1\ 0}^{[1]} & & \\
 & \mathbf{b}_{q-1\ 1}^{[0]} & & > & \mathbf{b}_{q-2\ 0}^{[2]} & & \\
 & & & > & \mathbf{b}_{q-2\ 1}^{[1]} & & \dots \\
 & \mathbf{b}_{q-2\ 2}^{[0]} & & & \dots & & \mathbf{b}_{10}^{[q-1]} \\
 & & & & & & > \mathbf{b}_{00}^{[q]} = \mathbf{P}(\xi) \\
 & \vdots & & \vdots & & & \mathbf{b}_{01}^{[q-1]} \\
 & \mathbf{b}_{1\ q-1}^{[0]} & & > & \mathbf{b}_{0\ q-2}^{[2]} & & \\
 & & > & \mathbf{b}_{0\ q-1}^{[1]} & & & \\
 & \mathbf{b}_{0\ q}^{[0]} & & & & & 
 \end{array}$$

- Cost is  $\mathcal{O}(q^2)$ .
- For  $\xi \in [a, b]$ : convex combinations, therefore numerically stable.
- In addition, it provides the derivatives  $\mathbf{P}'(\xi)$ ,  $\mathbf{P}''(\xi)$ ,  $\dots$ ,  $\mathbf{P}^{(q)}(\xi)$ .
  - ▣  $i$ -th derivative (in general):  $\frac{1}{(b-a)^i} \cdot \frac{q!}{(q-i)!} \sum_{k=0}^i \binom{i}{k} (-1)^{i-k} \cdot b_{i-k, k}^{[q-i]}$
  - ▣ ,e.g.,  $\mathbf{P}'(\xi) = \frac{q}{b-a} \cdot (b_{01}^{[q-1]} - b_{10}^{[q-1]})$ ,  $\mathbf{P}''(\xi) = \frac{q(q-1)}{(b-a)^2} \cdot (b_{02}^{[q-2]} - 2b_{11}^{[q-2]} + b_{20}^{[q-2]})$ .

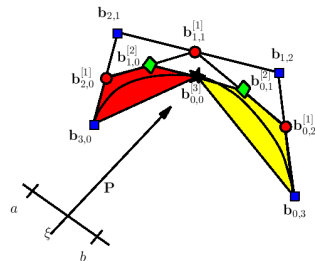
de Casteljau performs two changes of B-Bases !

$\mathbf{b}_{q,0}^{[0]}, \mathbf{b}_{q-1,0}^{[1]}, \dots, \mathbf{b}_{1,0}^{[q-1]}, \mathbf{b}_{0,0}^{[q]}$  are control points of  $\mathbf{P}$ 's B-Basis regarding  $[a, \xi]$ .

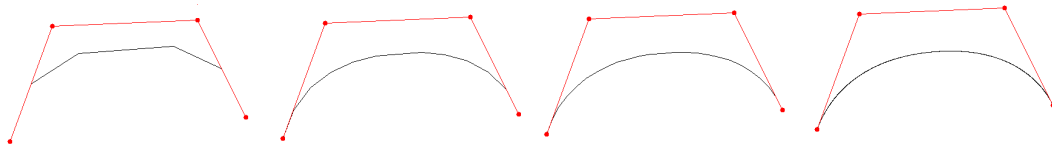
$\mathbf{b}_{0,0}^{[q]}, \mathbf{b}_{0,1}^{[q-1]}, \dots, \mathbf{b}_{0,q-1}^{[1]}, \mathbf{b}_{0,q}^{[0]}$  are control points of  $\mathbf{P}$ 's B-Basis regarding  $[\xi, b]$

**Example** ( $q = 3$ ):

$$\begin{aligned} \mathbf{P} &= B_{30} \mathbf{b}_{30} + B_{21} \mathbf{b}_{21} + B_{12} \mathbf{b}_{12} + B_{03} \mathbf{b}_{03} \\ &= \tilde{B}_{30} \mathbf{b}_{30} + \tilde{B}_{21} \mathbf{b}_{20}^{[1]} + \tilde{B}_{12} \mathbf{b}_{10}^{[2]} + \tilde{B}_{03} \mathbf{b}_{00}^{[3]} \\ &= \tilde{B}_{30}^* \mathbf{b}_{00}^{[3]} + \tilde{B}_{21}^* \mathbf{b}_{01}^{[2]} + \tilde{B}_{12}^* \mathbf{b}_{02}^{[1]} + \tilde{B}_{03}^* \mathbf{b}_{03} \end{aligned}$$



Display of polynomial curve in B-Form (3 – 4 recursion steps) through Subdivision using de Casteljau:



v.l.n.r.: 1,2,3,4 recursion steps; control polygon converge towards the curve

Lane & Riesenfeld 1980: quadratic convergence of control polygon  $\mathcal{K}(\mathbf{P})$  towards the curve  $\mathbf{P}$

Another application: **Bézier-Clipping**; approximate intersection point computation of polynomial curves through intersecting the (refined) control polygons.

- G. Farin: *Curves and Surfaces for CAGD*, [https://books.google.com/books?id=5HYTP1dIAp4C&printsec=frontcover&hl=de&source=gbv\\_ViewAPI](https://books.google.com/books?id=5HYTP1dIAp4C&printsec=frontcover&hl=de&source=gbv_ViewAPI)
- Prautzsch, Boehm & Paluszny: *Bézier and B-Spline techniques*, [https://books.google.com/books?id=xP7A8F6NZGQC&printsec=frontcover&hl=de&source=gbv\\_ViewAPI](https://books.google.com/books?id=xP7A8F6NZGQC&printsec=frontcover&hl=de&source=gbv_ViewAPI)
- Foley, van Dam, Feiner, Hughes: *Computer Graphics: Principles and Practice*, p. 478–491 [https://books.google.com/books?id=-4ngT05gmAQC&printsec=frontcover&hl=de&source=gbv\\_ViewAPI](https://books.google.com/books?id=-4ngT05gmAQC&printsec=frontcover&hl=de&source=gbv_ViewAPI)
- de Boor, B-form basics, SIAM, 1987, <ftp://ftp.cs.wisc.edu/Approx/BBform.pdf>
- Farin *Triangular Bernstein-Bézier patches*, J. CAGD 3 (1986), <https://www.ljll.math.upmc.fr/frey/ftp/M5105/Farin%20G.,%20Triangular%20Bernstein%20Bezier%20patches.pdf>