

# **N-Queens Problem Report**

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**N-Queens Problem Solution using Backtracking**

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**Course:** [Introduction To AI]

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## 1. Introduction:

The N-Queens problem is a well-known combinatorial problem where the goal is to place N queens on an  $N \times N$  chessboard such that no two queens threaten each other. This means no two queens can be in the same row, column, or diagonal. The problem has both theoretical and practical applications in constraint satisfaction problems, artificial intelligence, and optimization.

The most straightforward approach to solve the N-Queens problem is through backtracking, a depth-first search algorithm that tries to place queens one by one on the board while checking for conflicts. Once all queens are placed successfully, a solution is found.

This report presents a solution to the N-Queens problem using the backtracking method and discusses the code, the results, and the significance of the solution.

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## 2. Problem Statement:

The N-Queens problem asks us to place N queens on an  $N \times N$  chessboard such that no two queens can attack each other. A queen can attack another queen if they are on the same row, column, or diagonal. We are required to find all possible arrangements of queens on the board that satisfy these conditions.

The objective of the report is to:

1. Develop a solution to the N-Queens problem using backtracking.
  2. Implement the solution in Python.
  3. Analyze the output generated from the solution.
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## 3. Methodology:

To solve the N-Queens problem, we use the **backtracking algorithm**, which is a common technique for solving constraint satisfaction problems.

### Backtracking Approach:

- **Step 1:** Place a queen in the first row, and check all the columns for possible positions.
- **Step 2:** For each valid position in the row, place the queen and move to the next row.
- **Step 3:** If a queen cannot be placed in any column of the current row (because of conflicts), backtrack by removing the queen from the previous row and try the next valid position.
- **Step 4:** Continue this process until all queens are placed on the board without conflicts.

The algorithm will explore all possible configurations of queens on the chessboard but prune those that violate the constraints.

### Time Complexity:

The time complexity of this approach is exponential, as we explore every possible configuration. However, the backtracking helps reduce the number of invalid configurations, improving efficiency compared to brute-force methods.

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## 4. Implementation (Code):

```
python
Copy
def is_safe(board, row, col, N):
    # Check for queen in the same column
    for i in range(row):
        if board[i] == col or abs(board[i] - col) ==
abs(i - row):
            return False
    return True

def solve_n_queens_util(board, row, N, solutions):
    if row == N:
        # A solution is found, add the board
configuration
        solutions.append(board.copy())
        return

    for col in range(N):
```

```

        if is_safe(board, row, col, N):
            board[row] = col
            solve_n_queens_util(board, row + 1, N,
solutions)
            board[row] = -1 # Backtrack

def solve_n_queens(N):
    board = [-1] * N # Initialize the board
    solutions = []
    solve_n_queens_util(board, 0, N, solutions)
    return solutions

def print_solution(solutions, N):
    for sol in solutions:
        for row in range(N):
            line = ['Q' if col == sol[row] else '.' for
col in range(N)]
            print(" ".join(line))
        print("\n")

# Driver code to test the solution
N = int(input("Enter Queens")) # For example, 8-Queens
problem
solutions = solve_n_queens(N)
print(f"Total solutions for {N}-Queens:
{len(solutions)}\n")
print_solution(solutions, N)

```

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## 5. Results and Output:

For an 8-Queens problem, the backtracking algorithm finds 92 valid solutions. Here's a sample output for the 8-Queens problem:

Total solutions for 8-Queens: 92

```

Q . . . . . . .
. . . Q . . . .
. . . . . Q . .
. . Q . . . . .
. . . . Q . . .

```

.	Q	.	.	.	.	.	.
.	.	.	.	.	.	Q	.
.	.	.	.	.	Q	.	.
Q	.	.	.	.	.	.	.
.	.	.	.	.	.	Q	.
.	.	Q	.	.	.	.	.
.	.	.	.	.	.	.	Q
.	.	.	.	Q	.	.	.
.	Q	.	.	.	.	.	.
.	.	.	.	.	Q	.	.
.	.	.	Q	.	.	.	.

Each solution represents a configuration of queens on the chessboard, where Q indicates a queen and . indicates an empty square.

The total number of solutions for 8 queens on an 8×8 board is 92, and the code will display all 92 valid configurations.

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## 6. Conclusion:

The backtracking algorithm successfully solves the N-Queens problem by exploring all possible configurations and pruning invalid solutions as soon as conflicts are detected. For  $N = 8$ , the algorithm finds 92 valid configurations, demonstrating the efficiency of backtracking in solving constraint satisfaction problems.

Further improvements or alternatives, such as using optimization techniques like simulated annealing or genetic algorithms, could provide more efficient solutions for larger values of  $N$ . However, the backtracking approach remains a robust and widely used method for solving the N-Queens problem.

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