

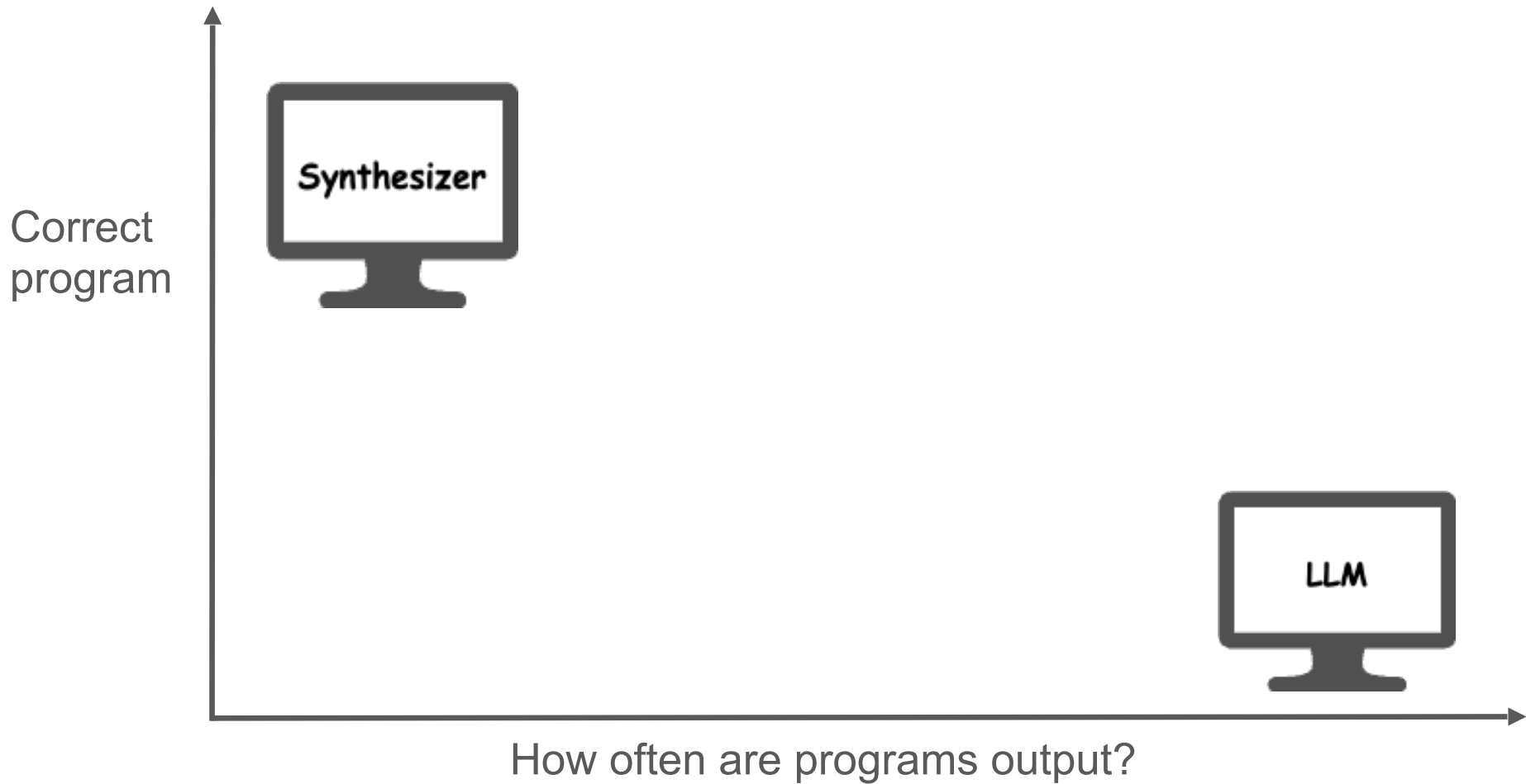
# SYNVER: Towards Automated Verification of LLM-Synthesized C Programs

Prasita Mukherjee and Benjamin Delaware

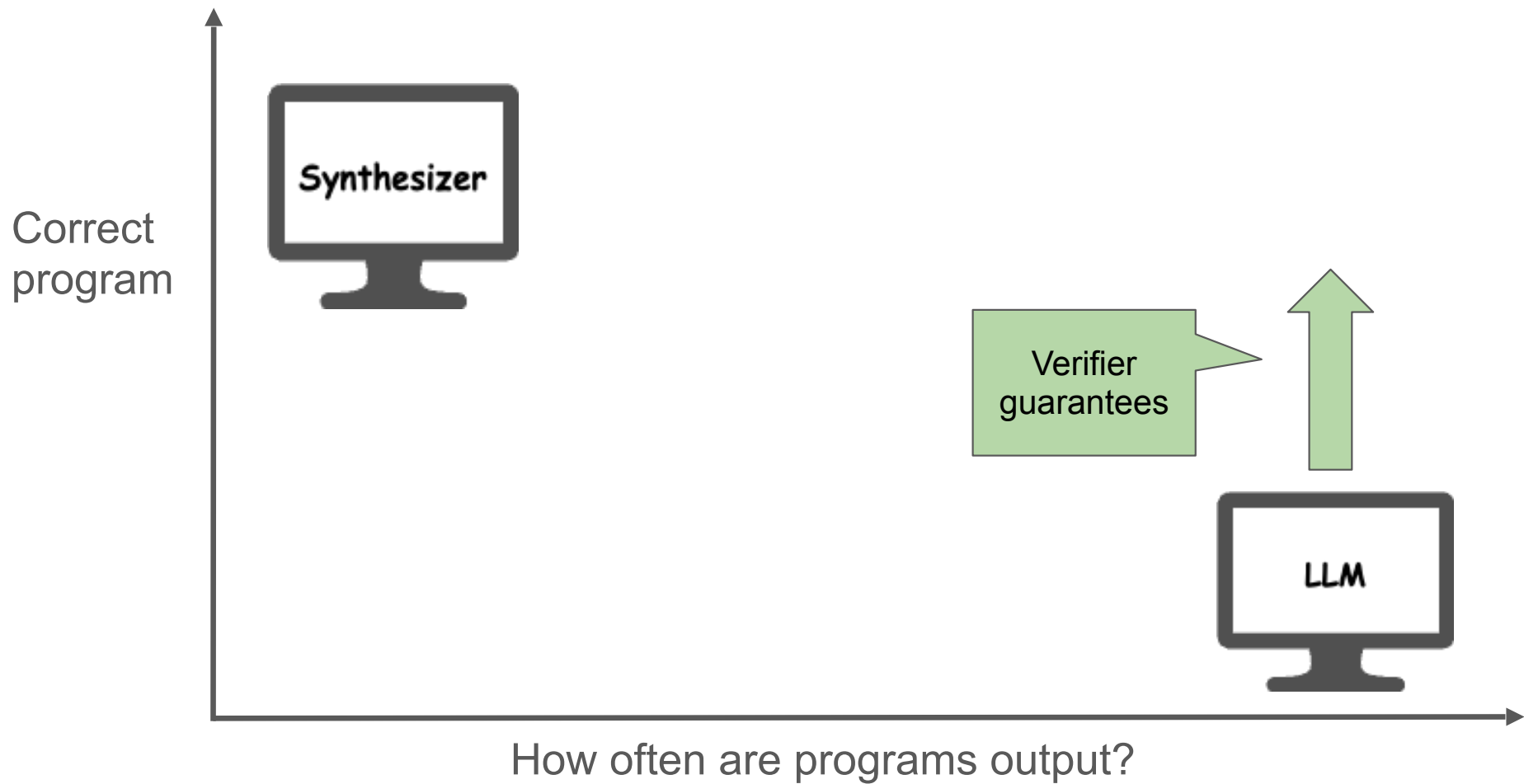
CoqPL 2025



# Program Synthesis: Present Scenario



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# Synthesizing programs: LLMs

$\forall a, b (max(a, b) = a \wedge a \geq b)$ $\vee (max(a, b) = b \wedge b > a)$
<i>Definition</i> $fi2 := fun\ x \Rightarrow x =? 2$ $listrep\ l\ h\ *listrep\ (filter\ fi2\ l)\ a$
.....
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struct sll \*addf(..)  
struct sll \*addb(..)

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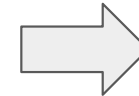
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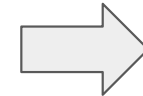
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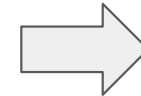
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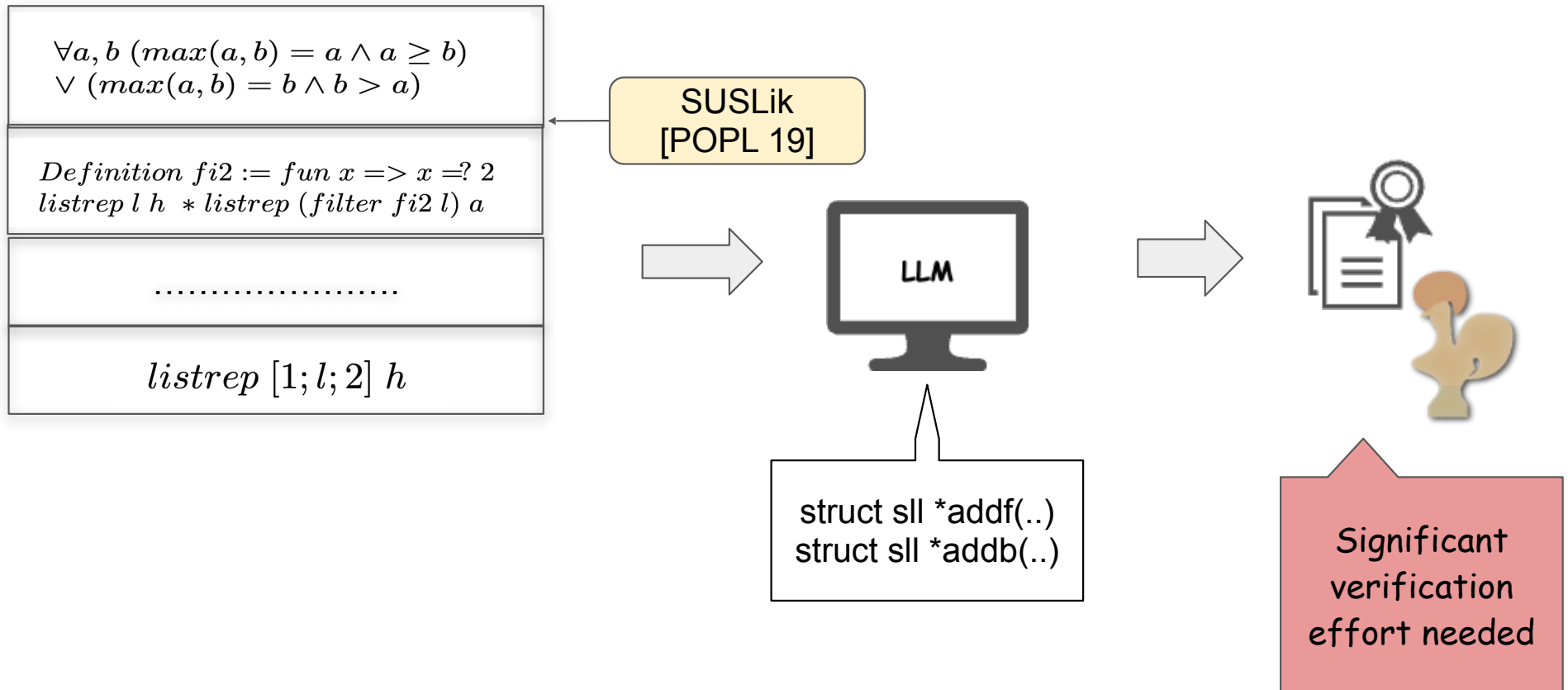
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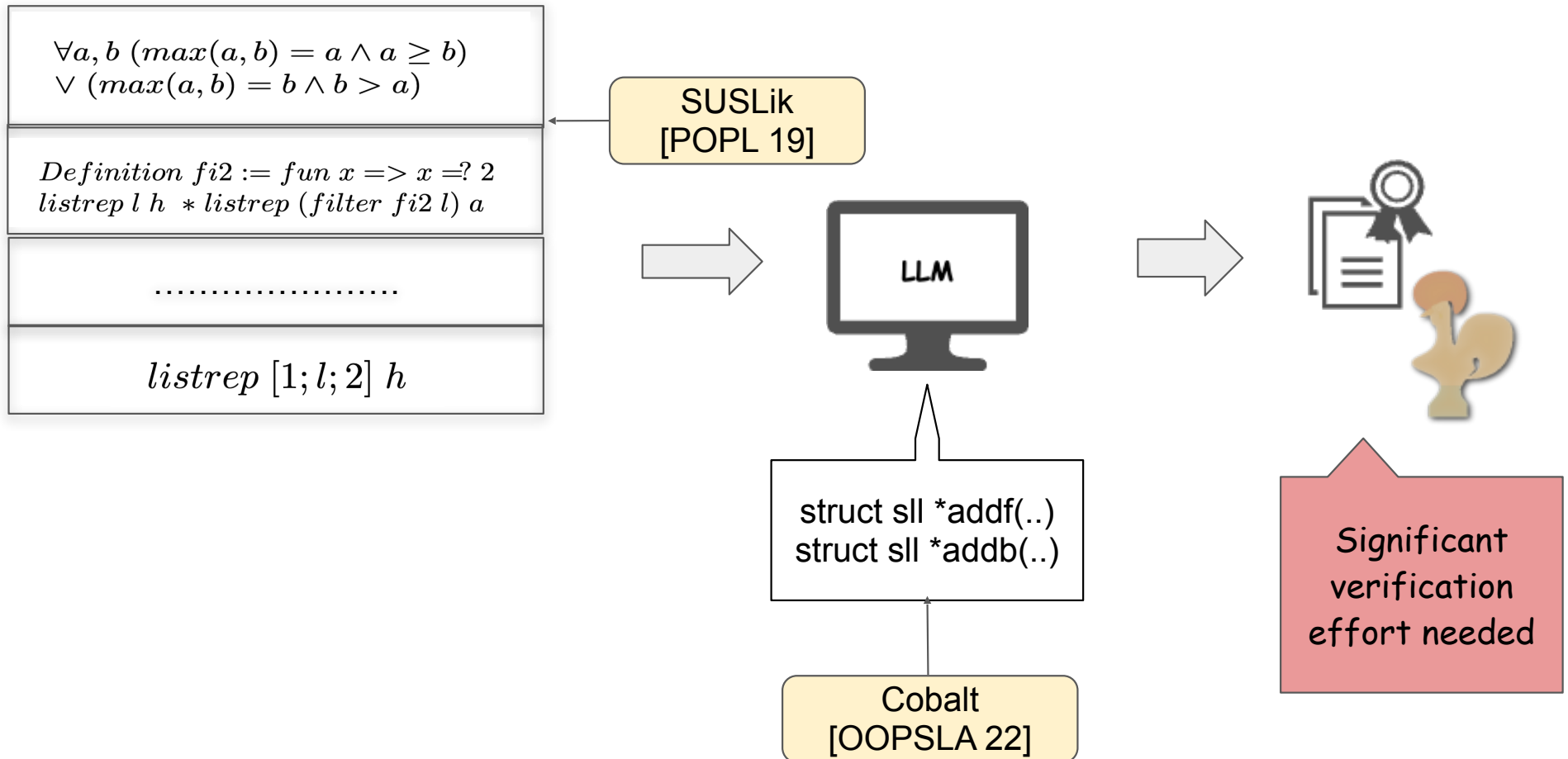
Significant  
verification  
effort needed



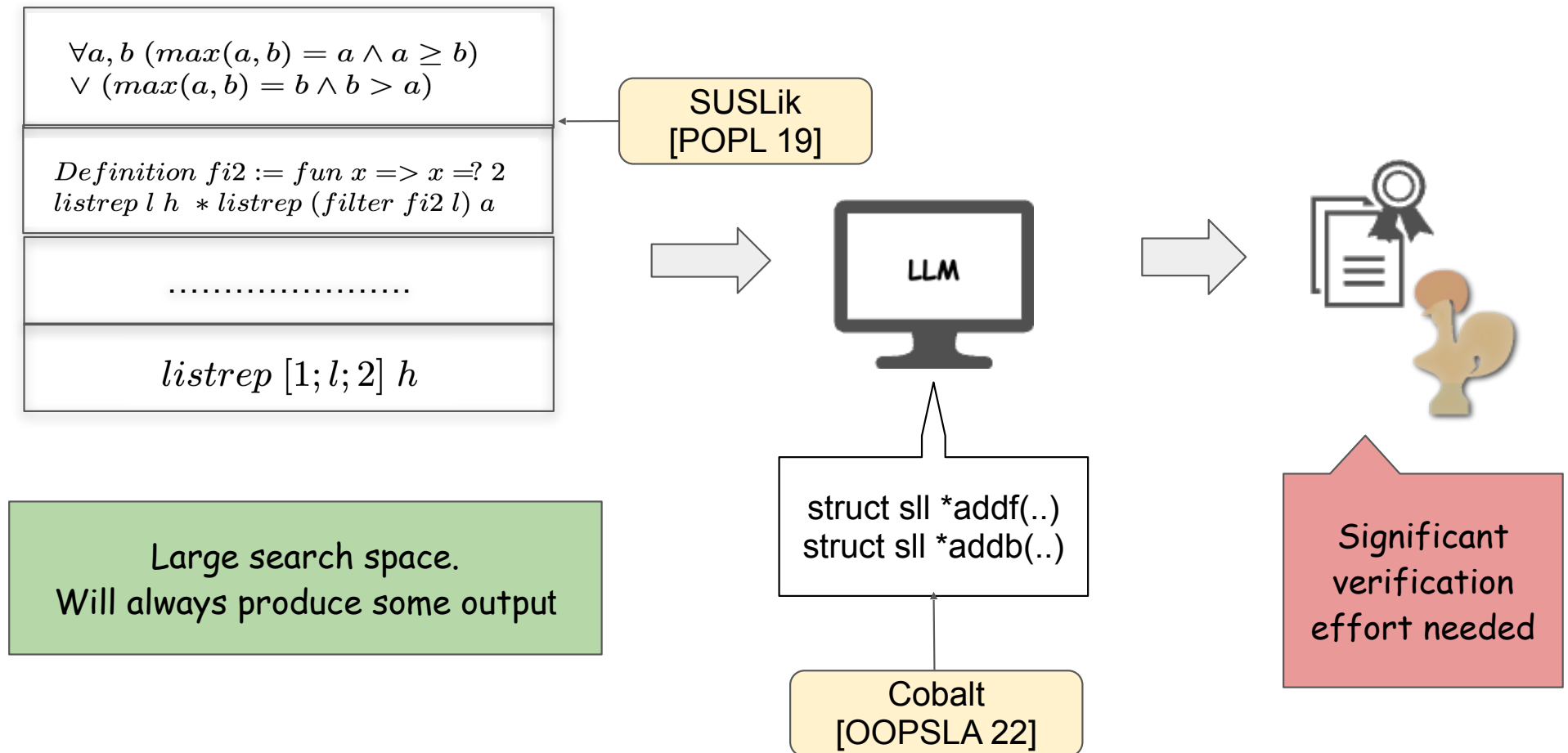
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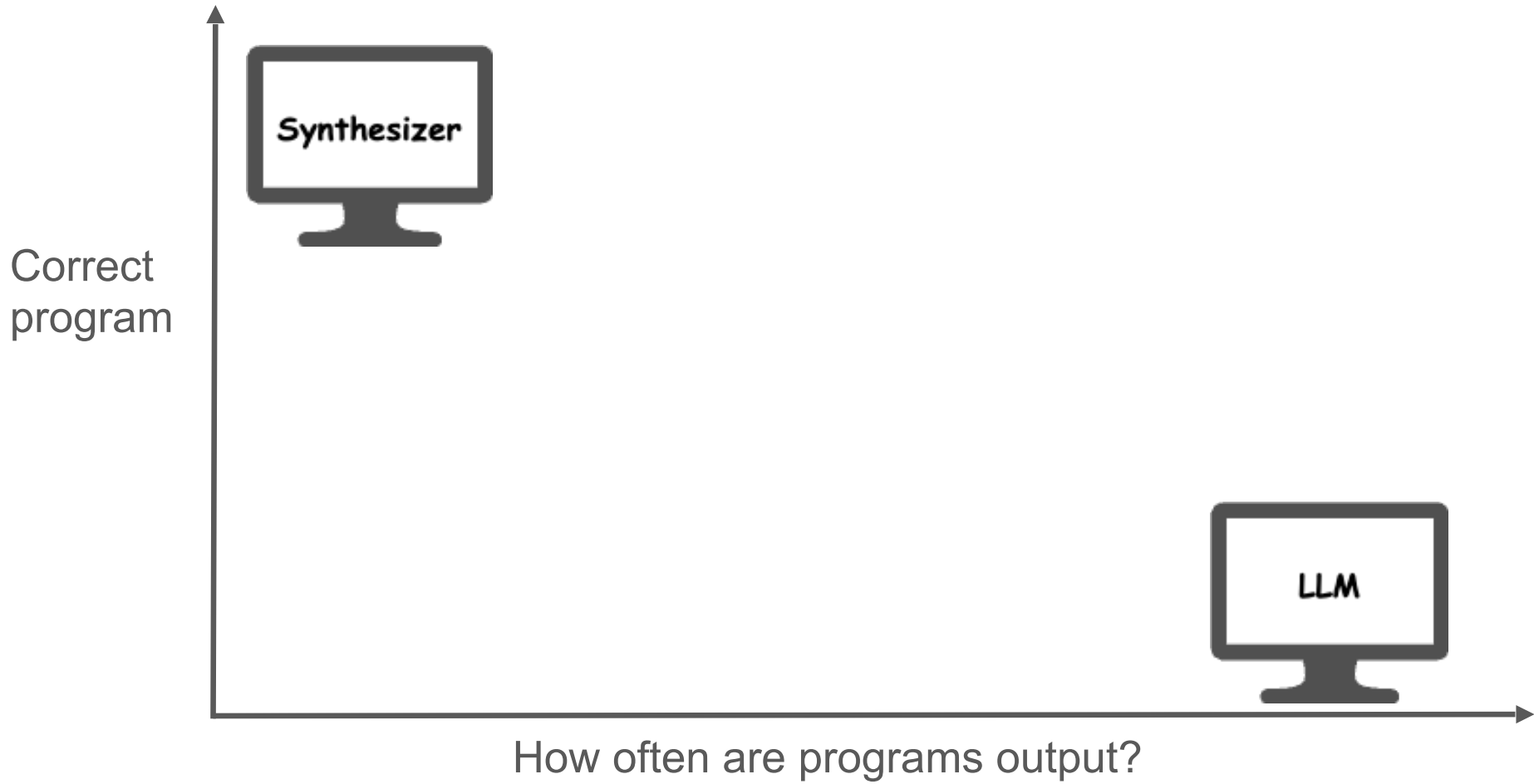
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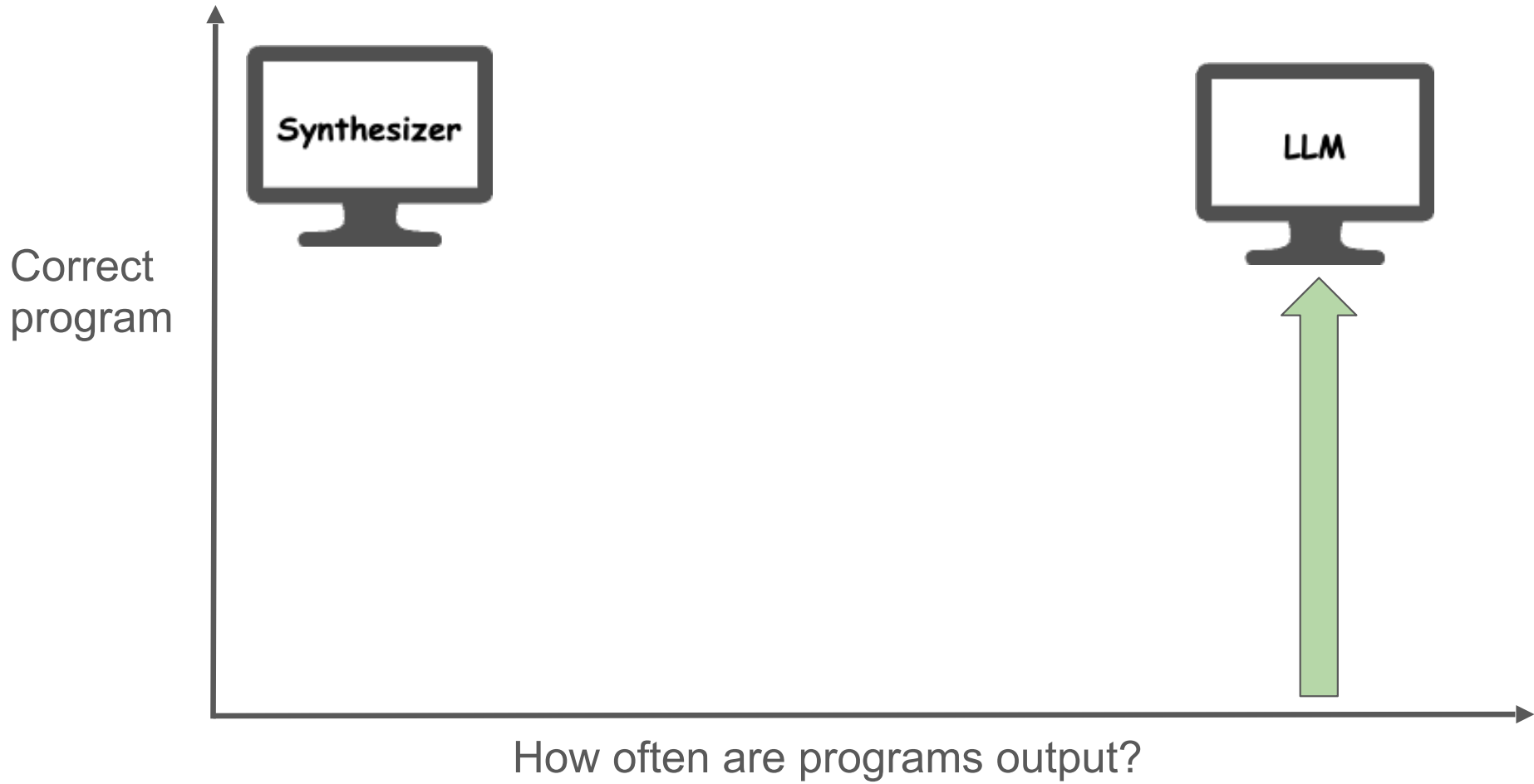
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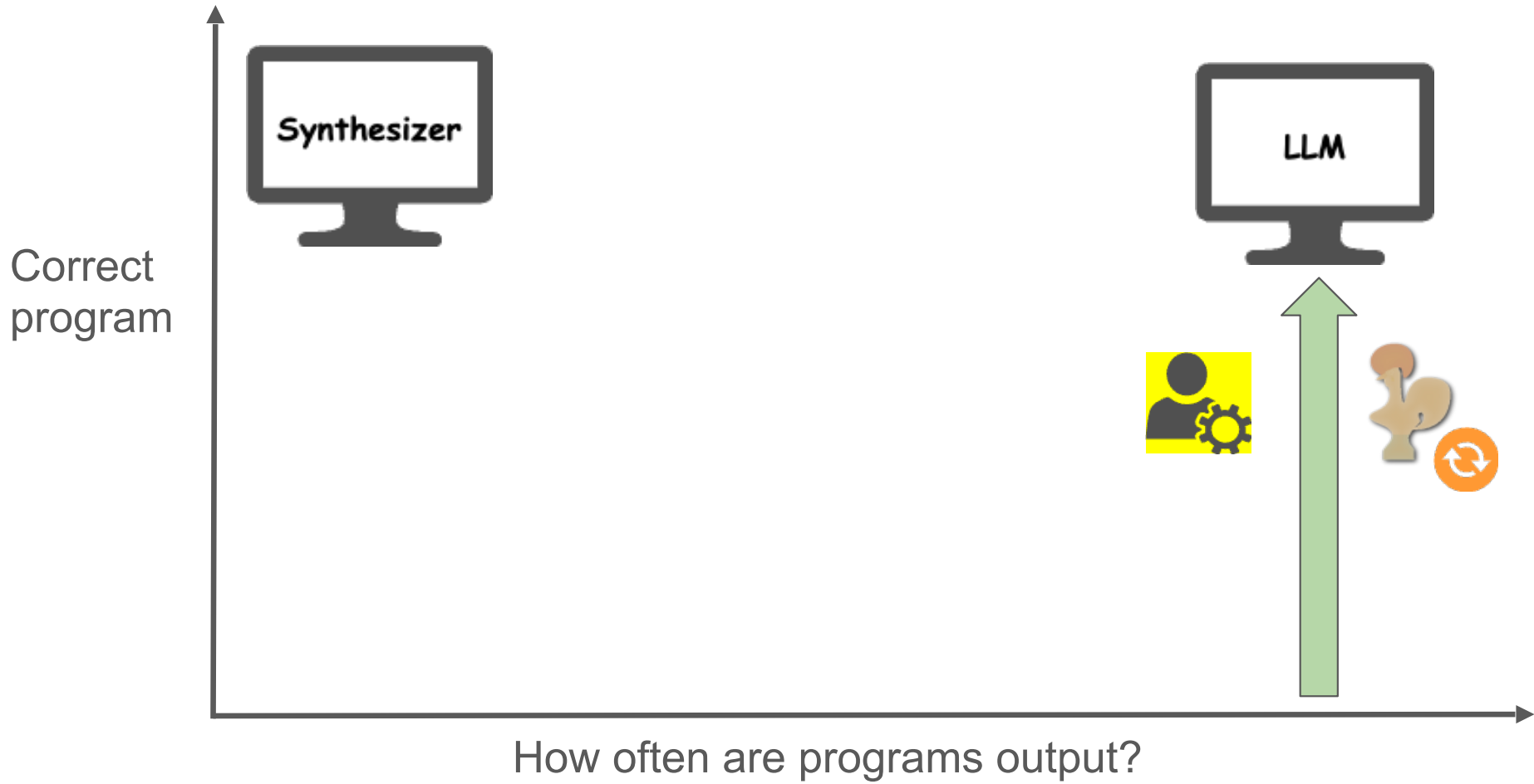
## Goal of SYNVER



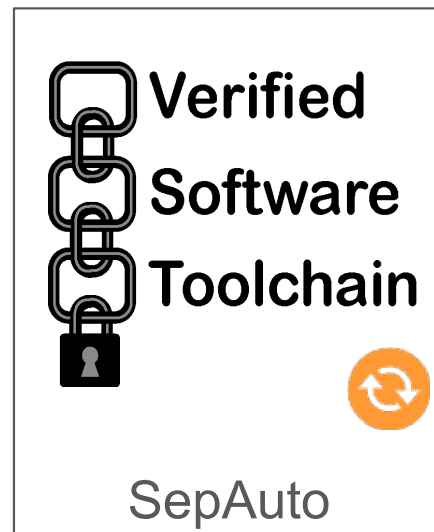
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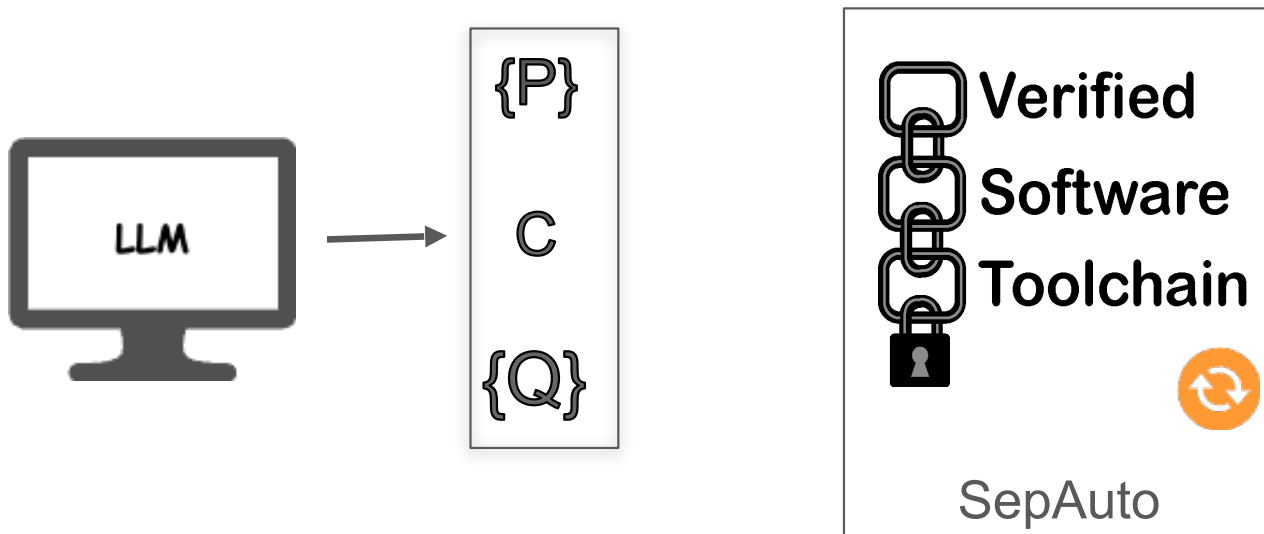
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# Verification Challenges

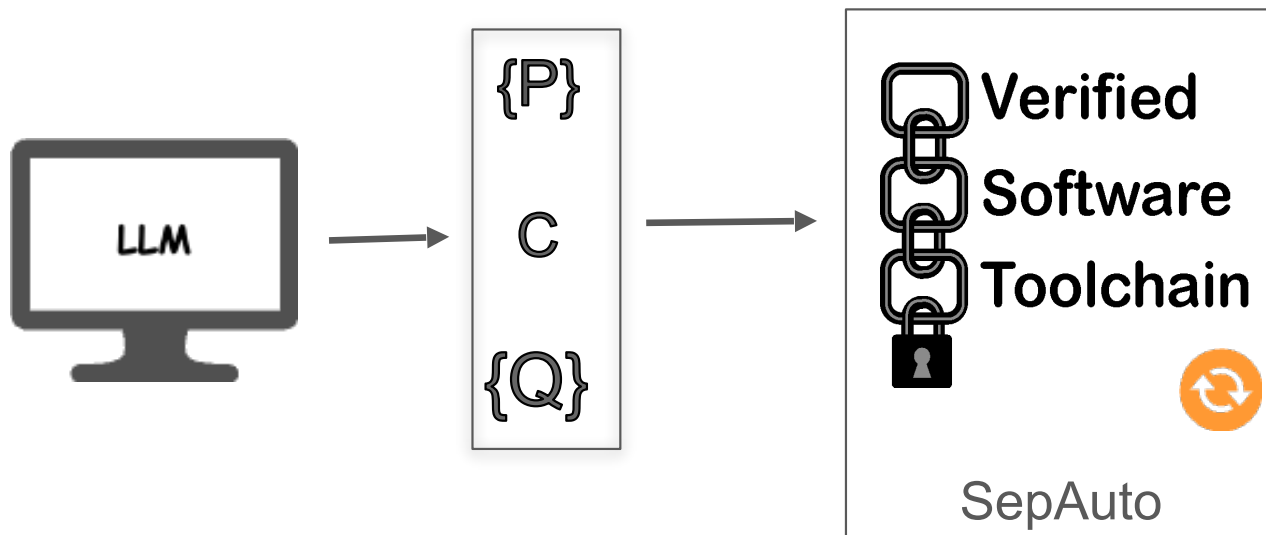


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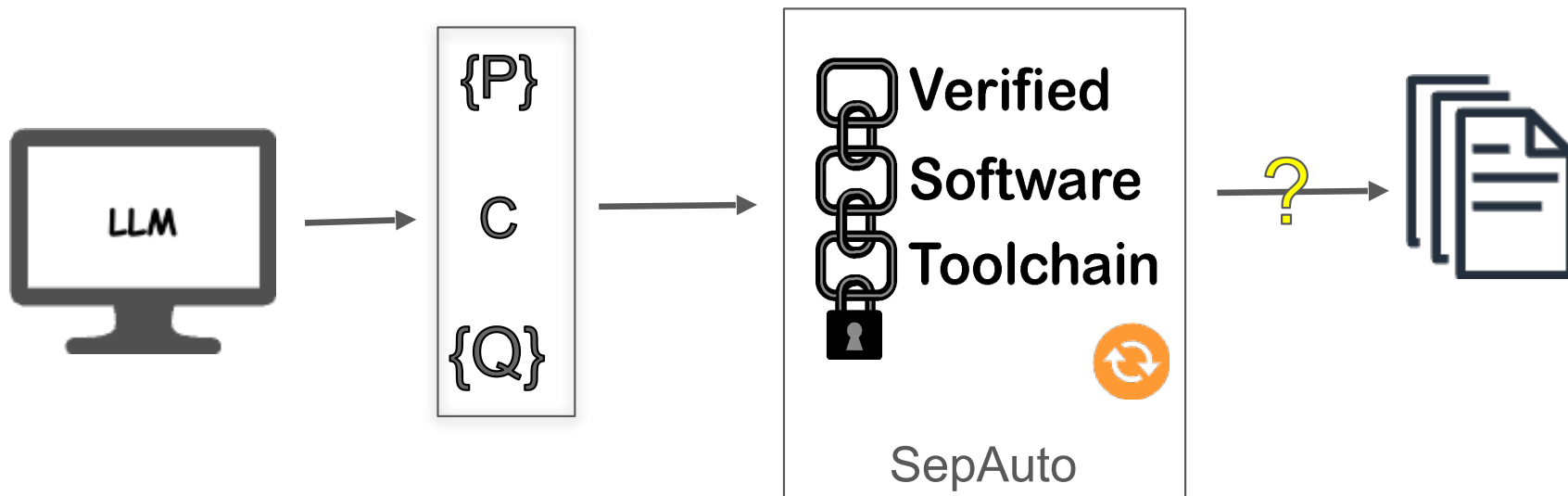




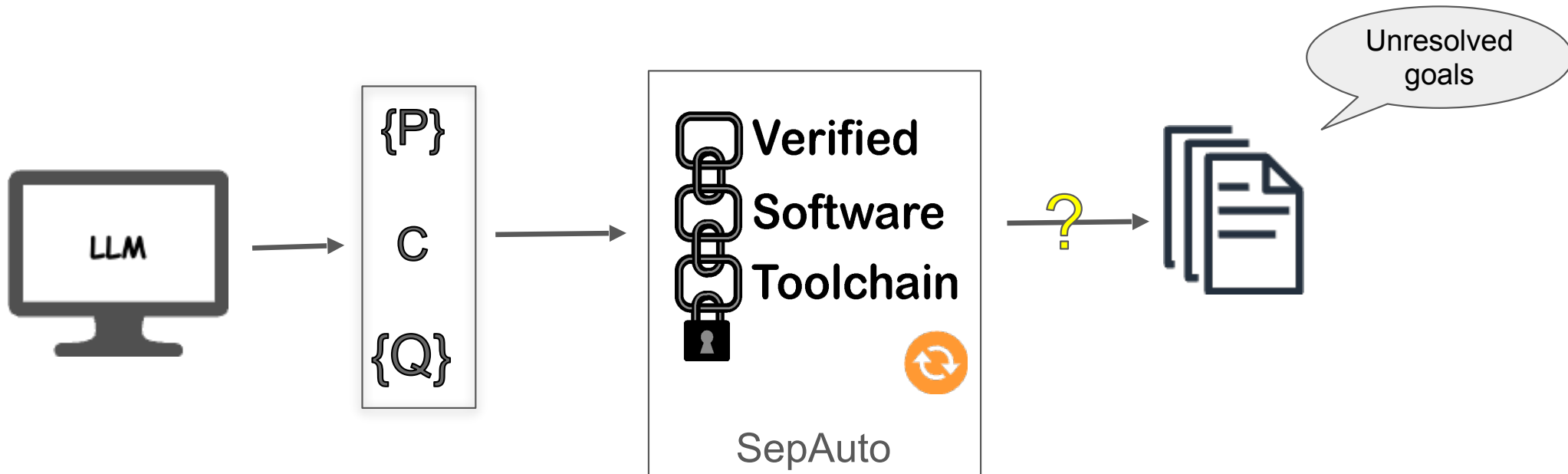
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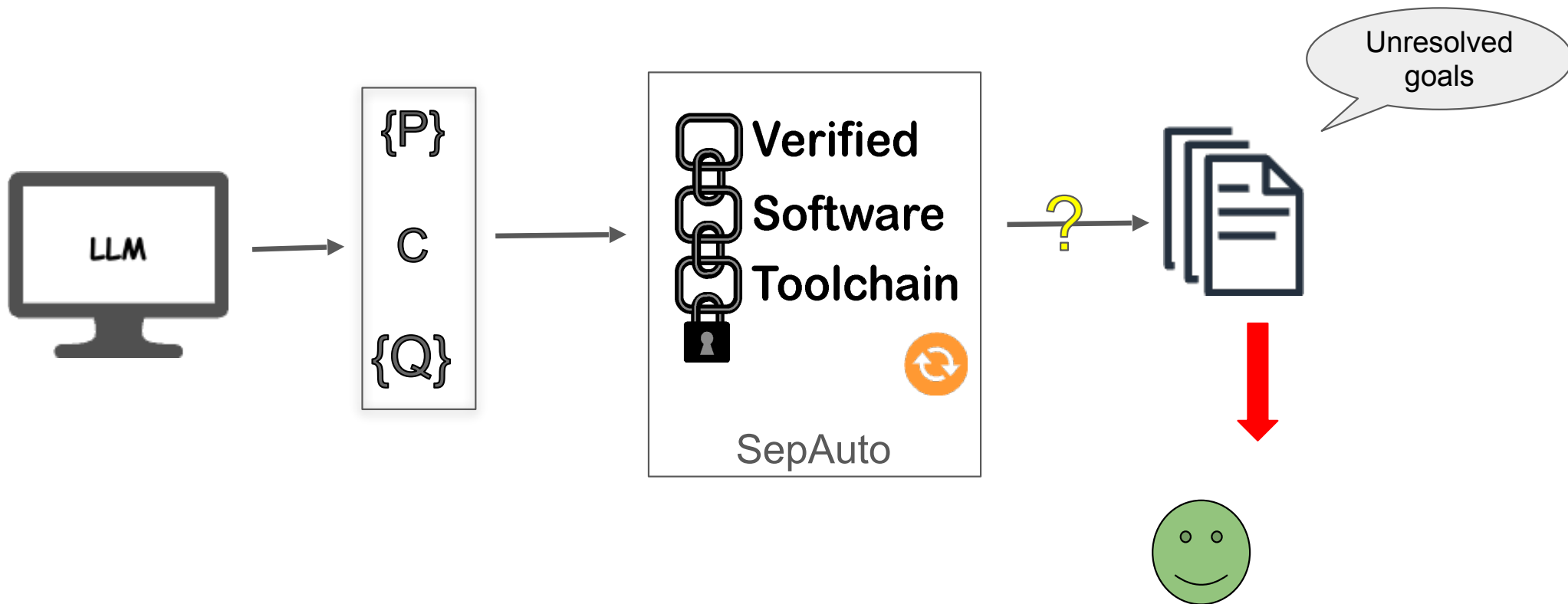
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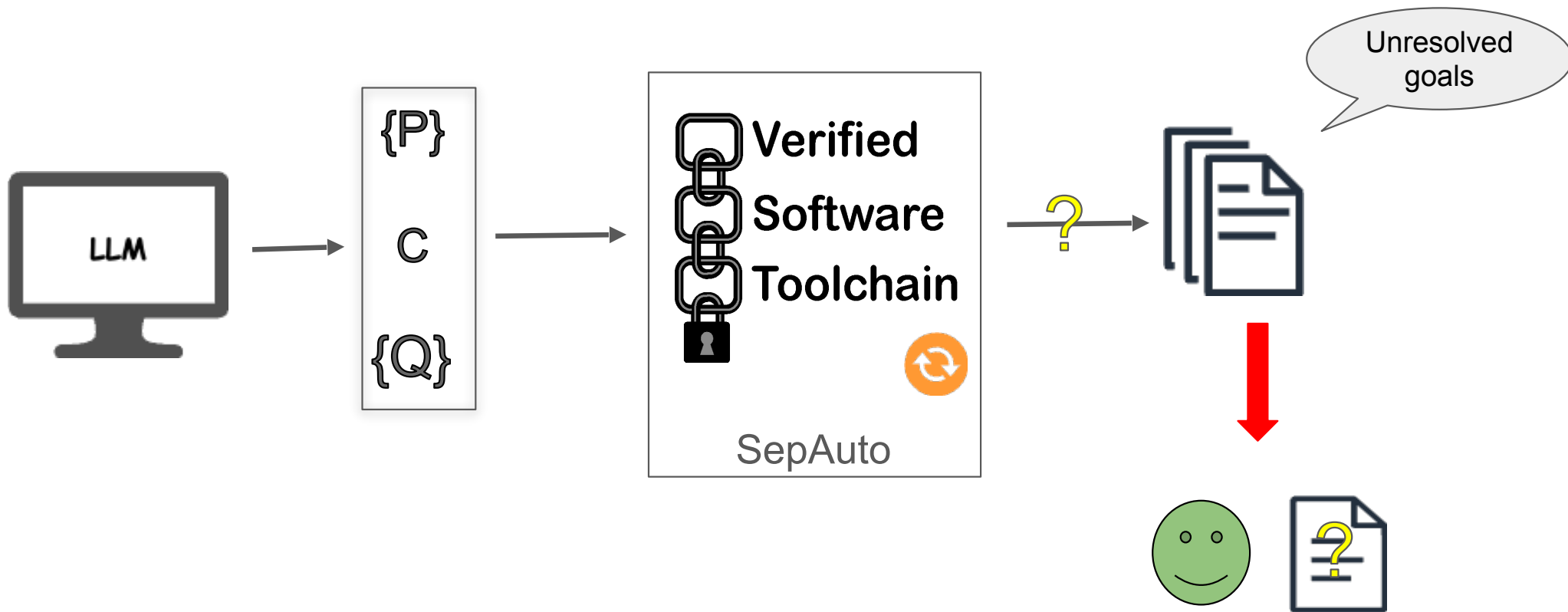
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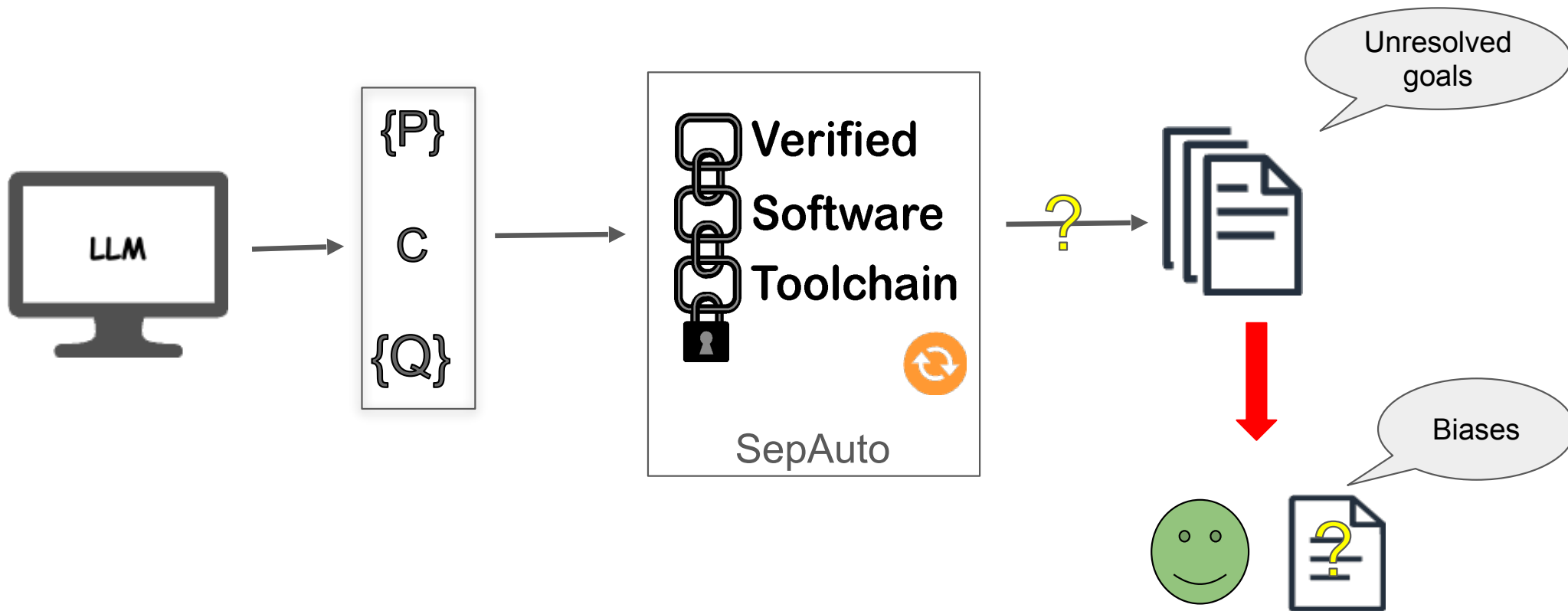
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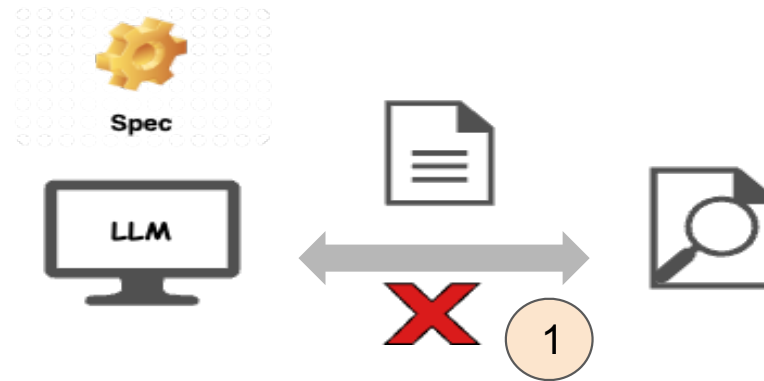
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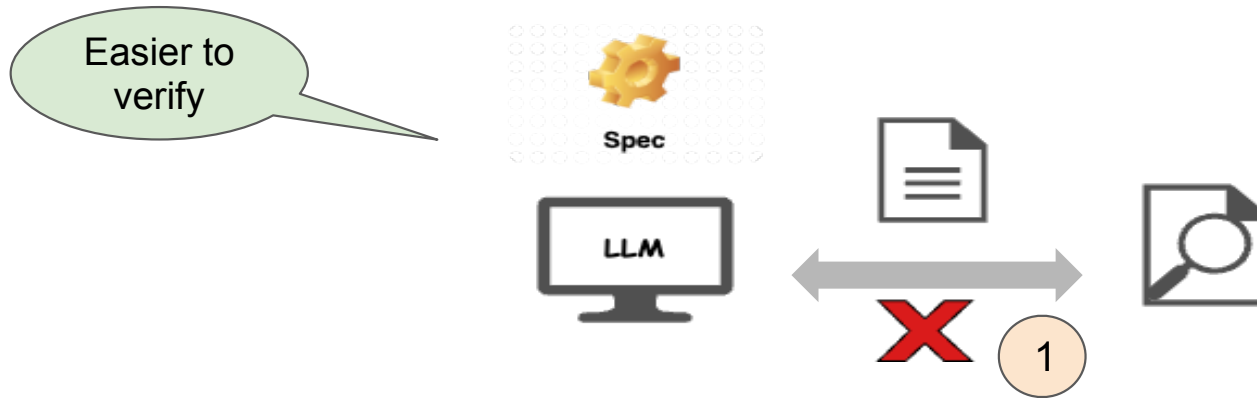
# SYNVER: Workflow



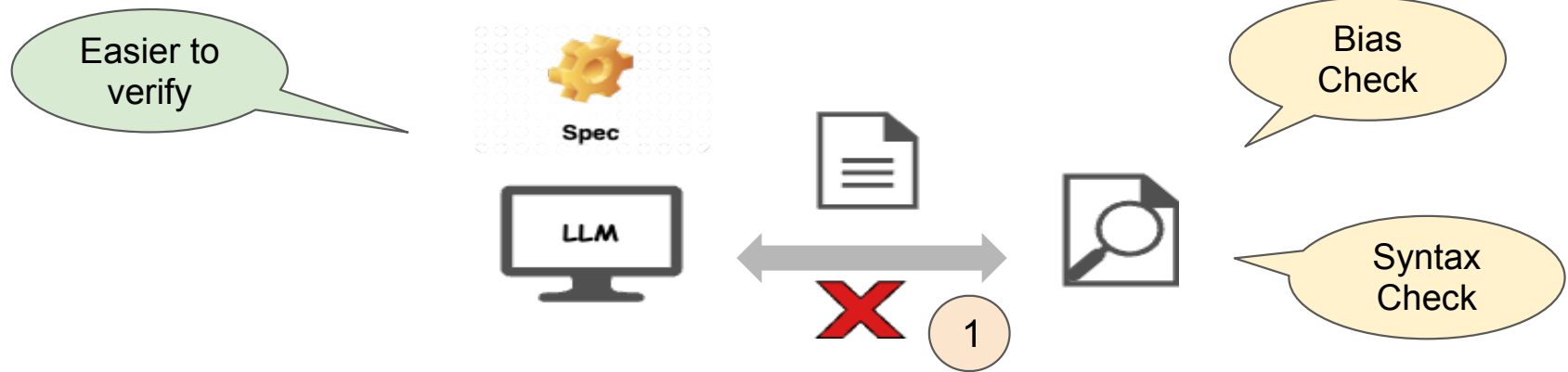
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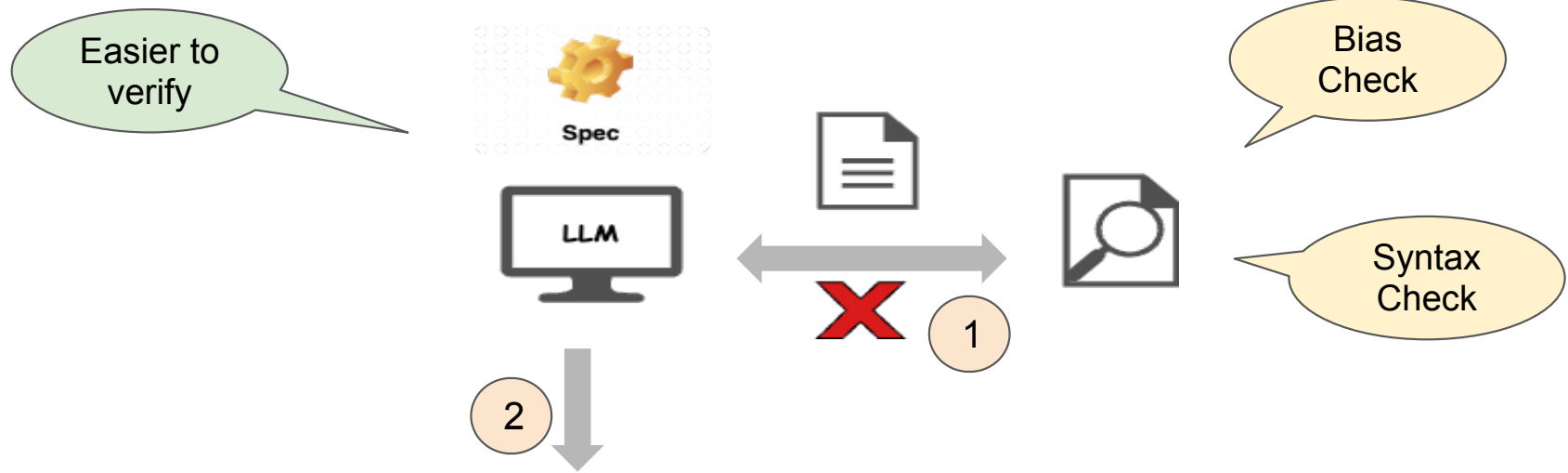
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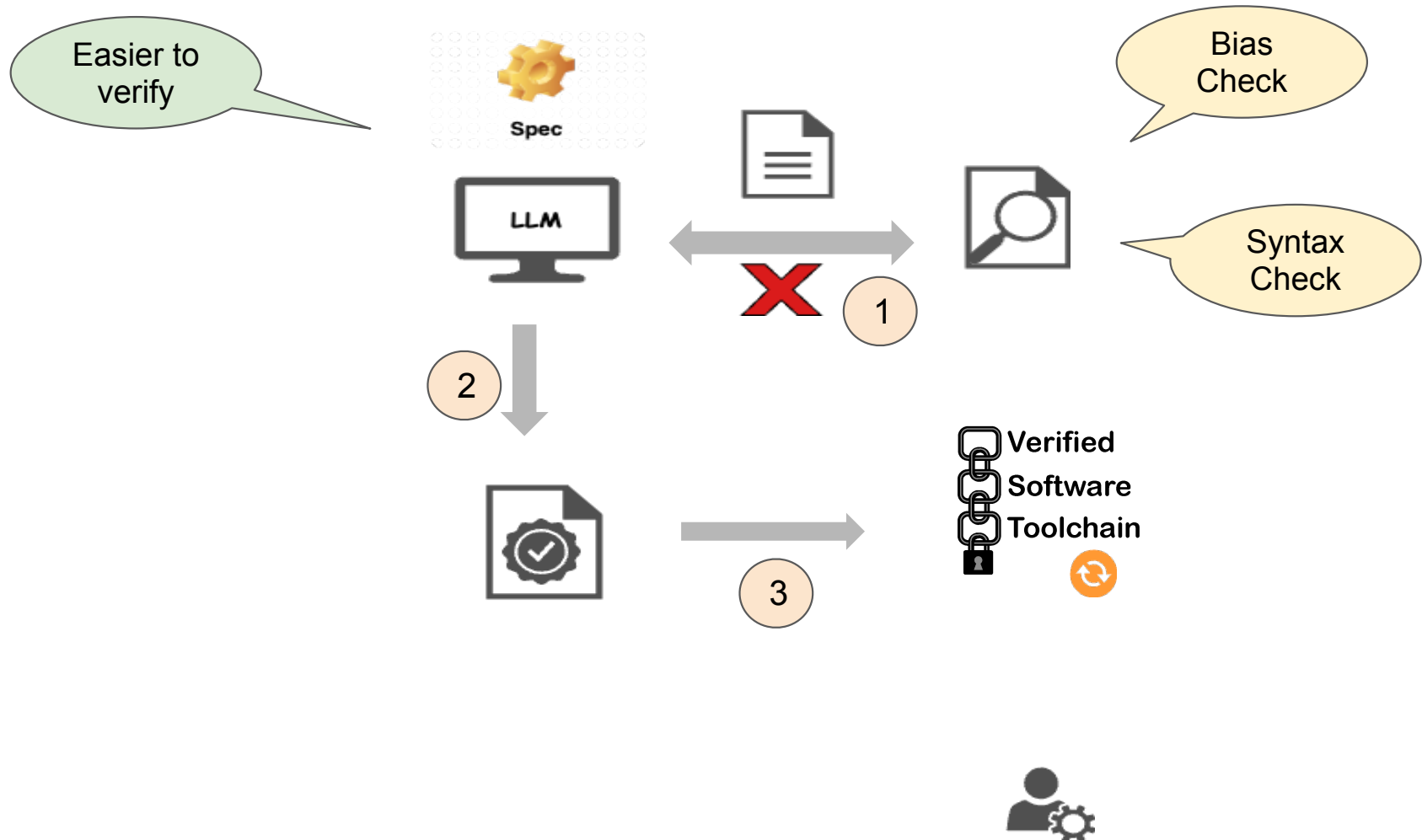
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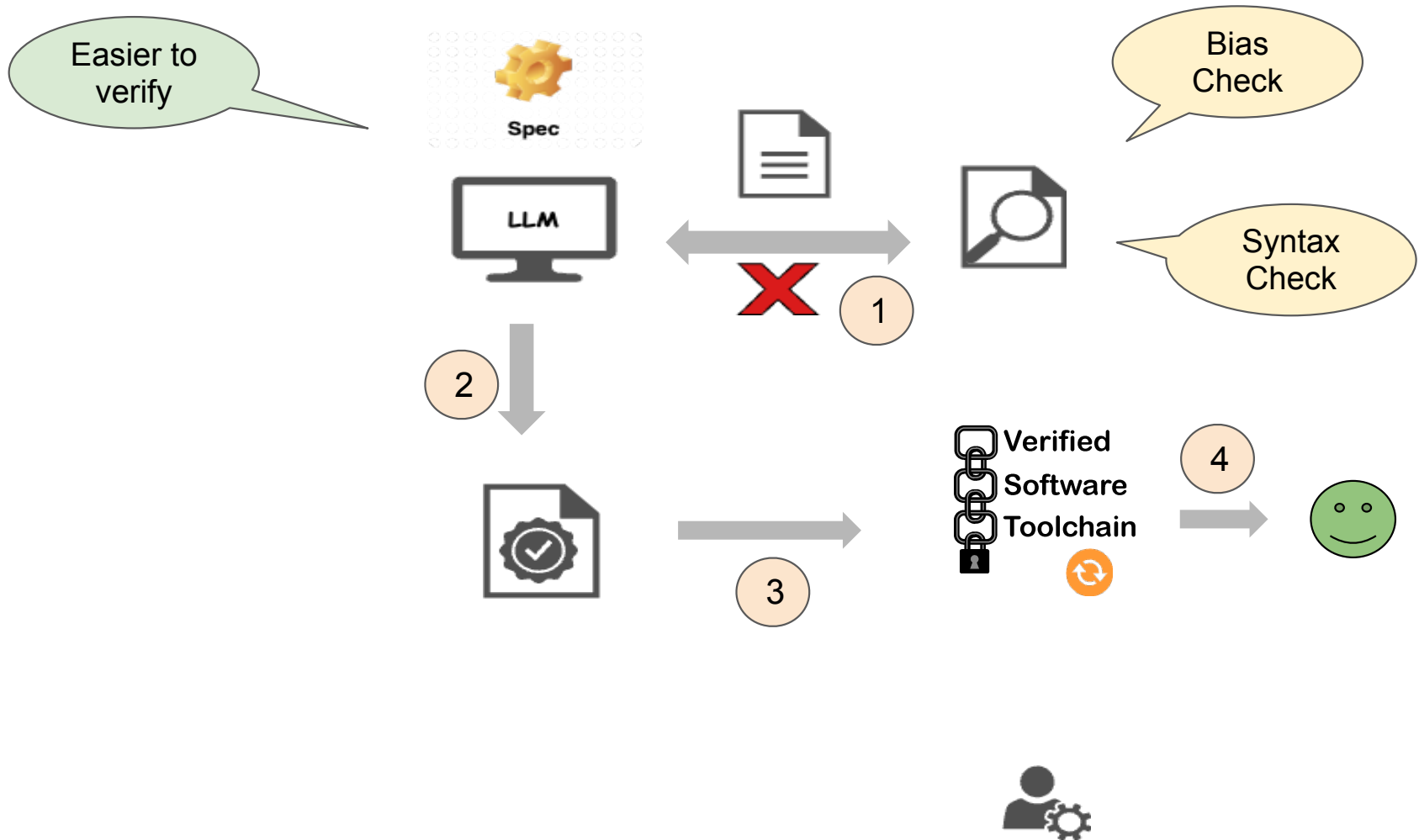
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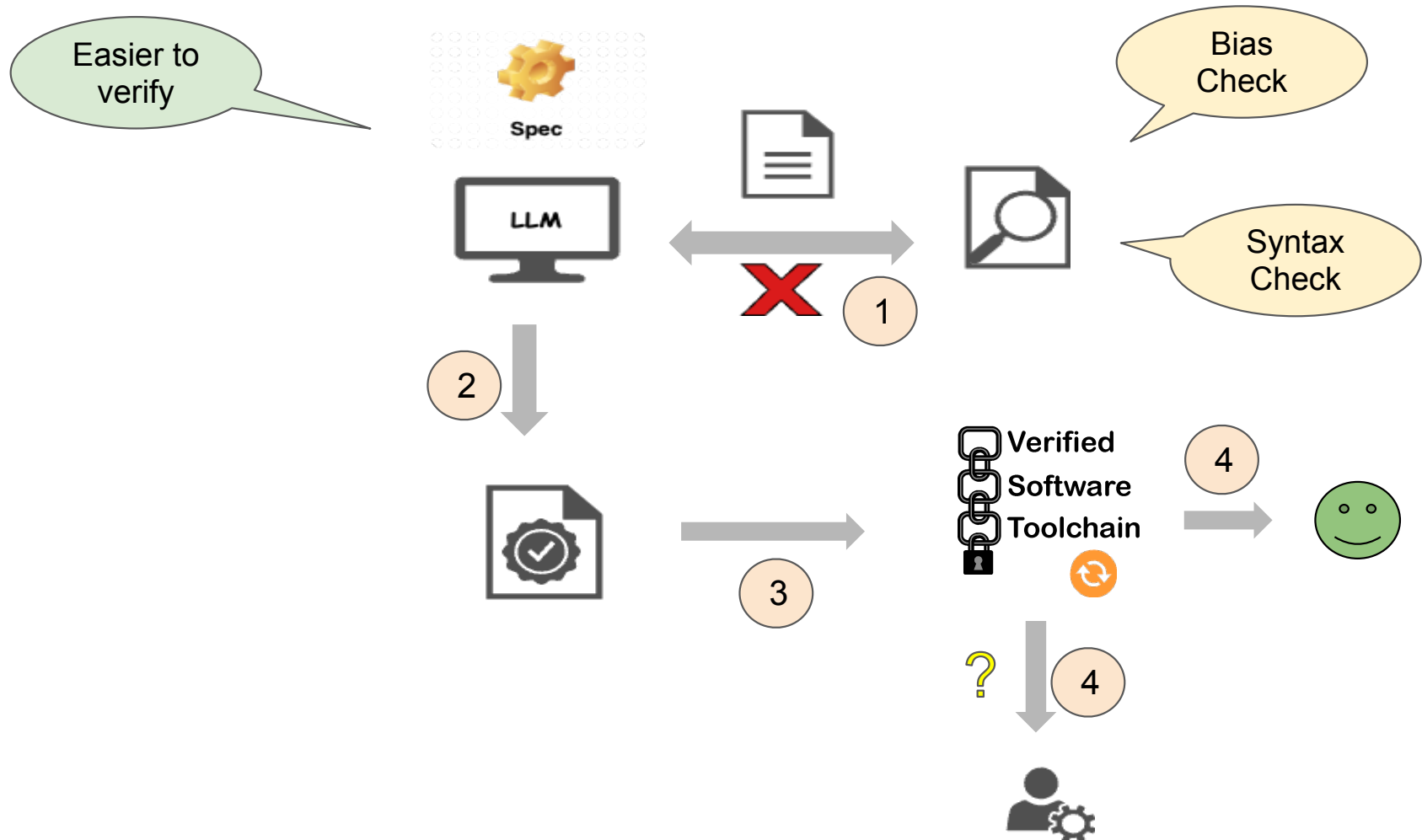
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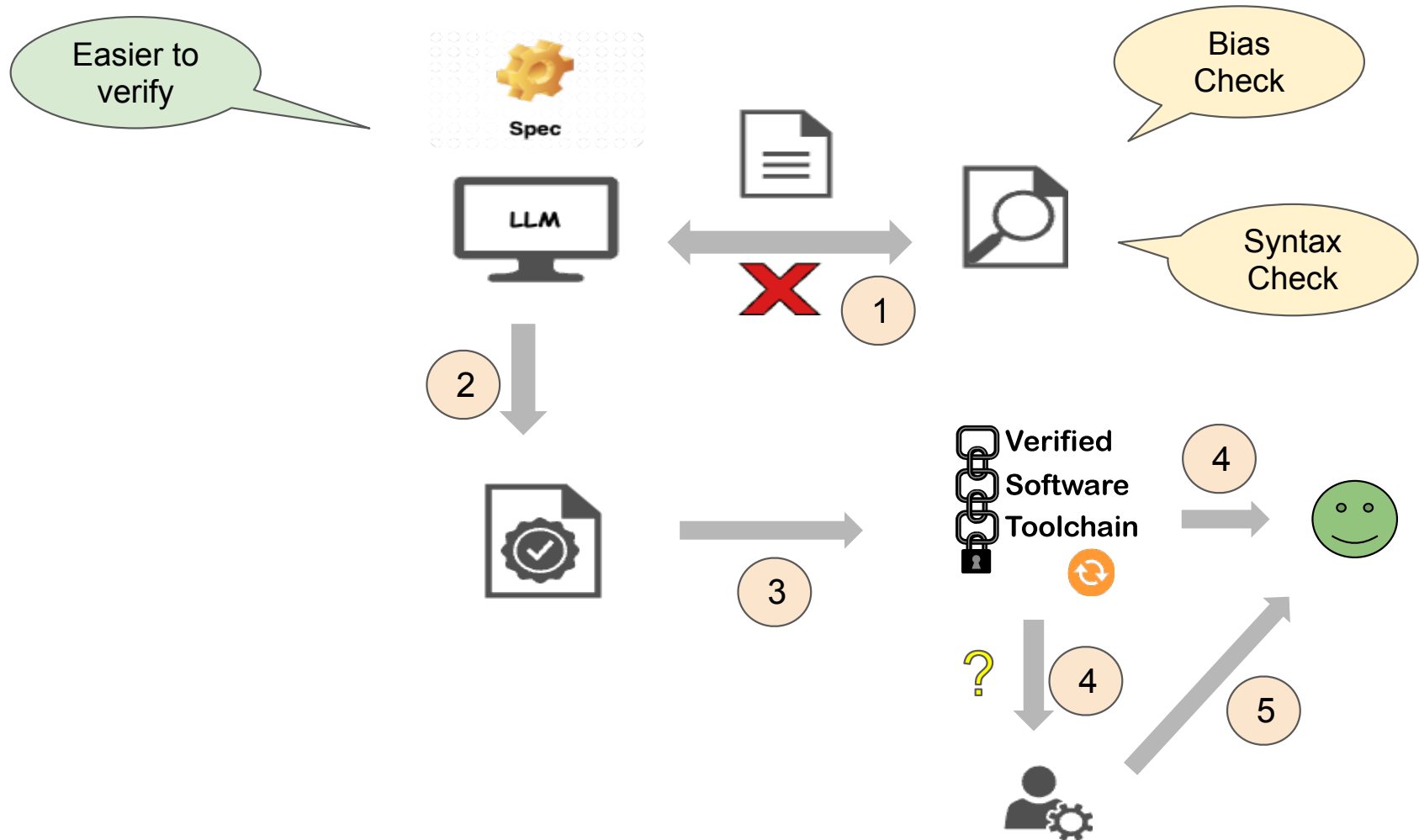
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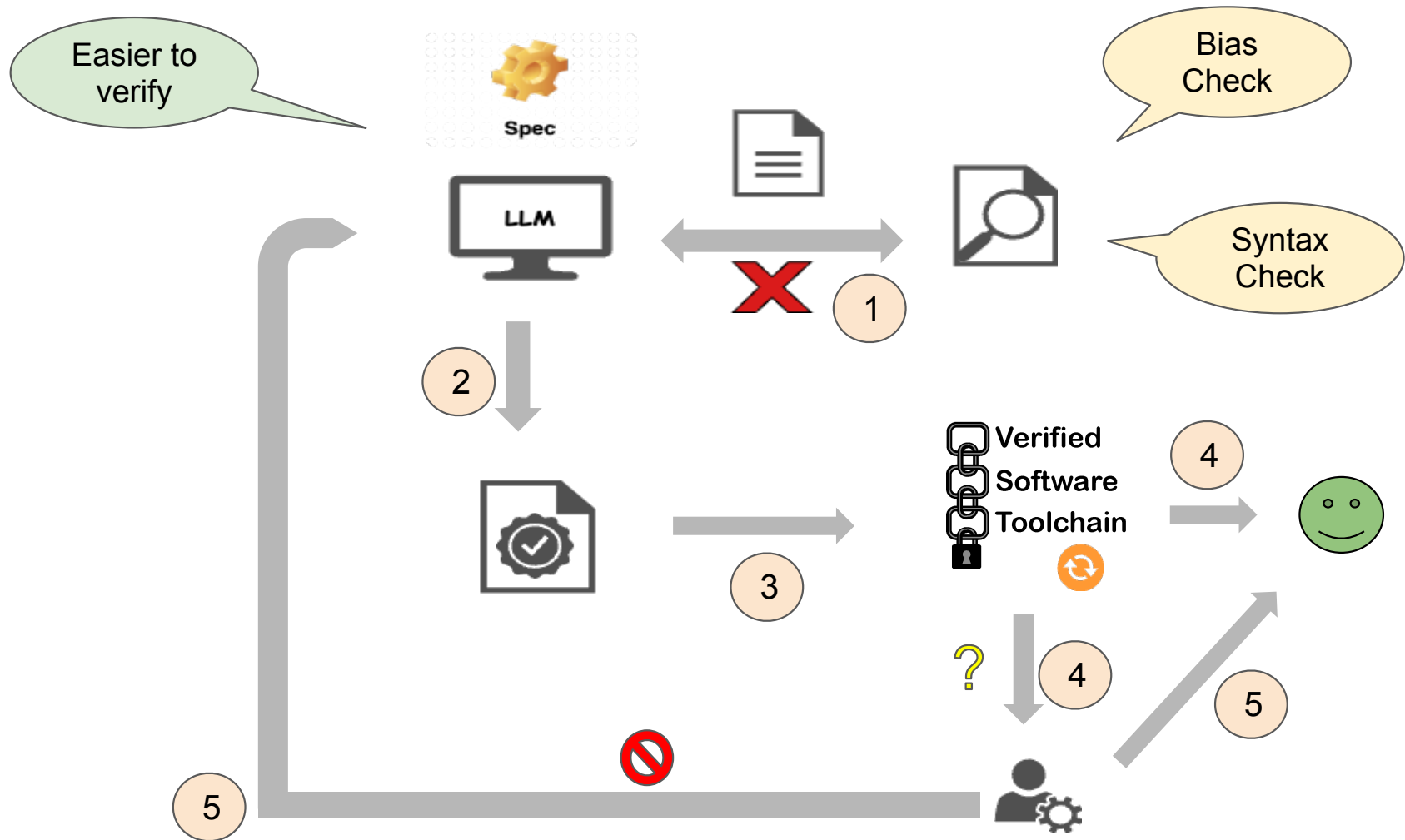


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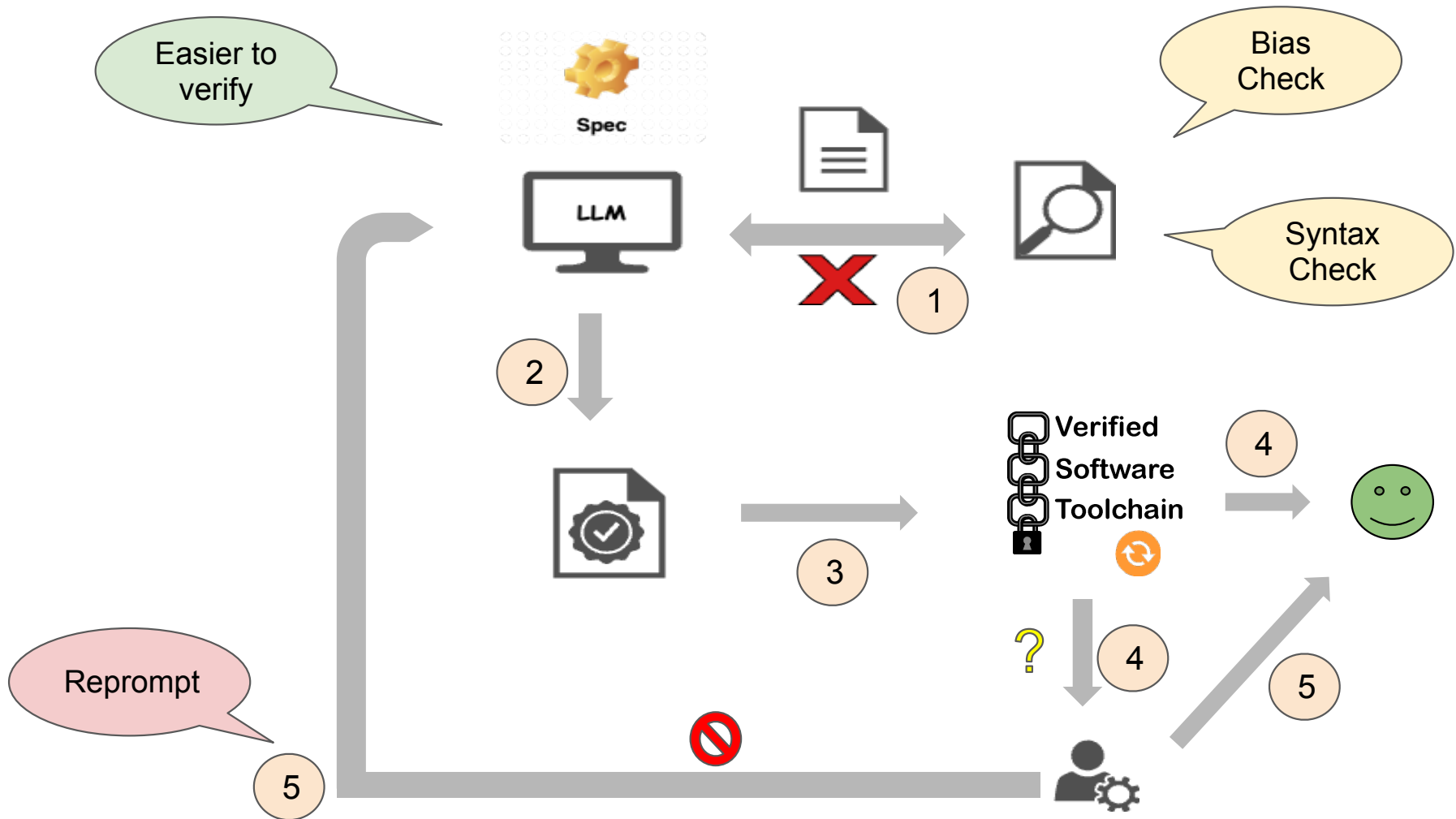




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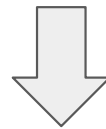
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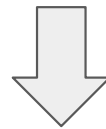
Why incorporate biases?

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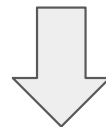


Generate programs that are easier to verify

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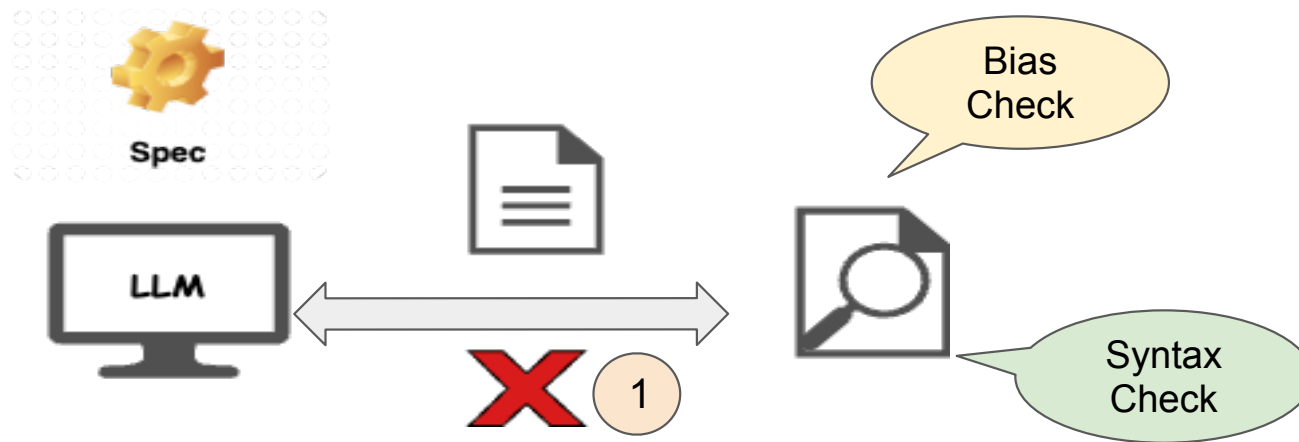
Generate programs that are easier to verify



✗ Infer novel specifications and verify them

Biases: No loops and novel helper functions

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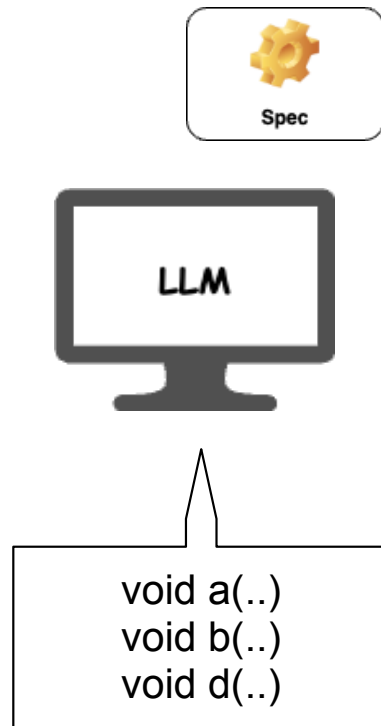


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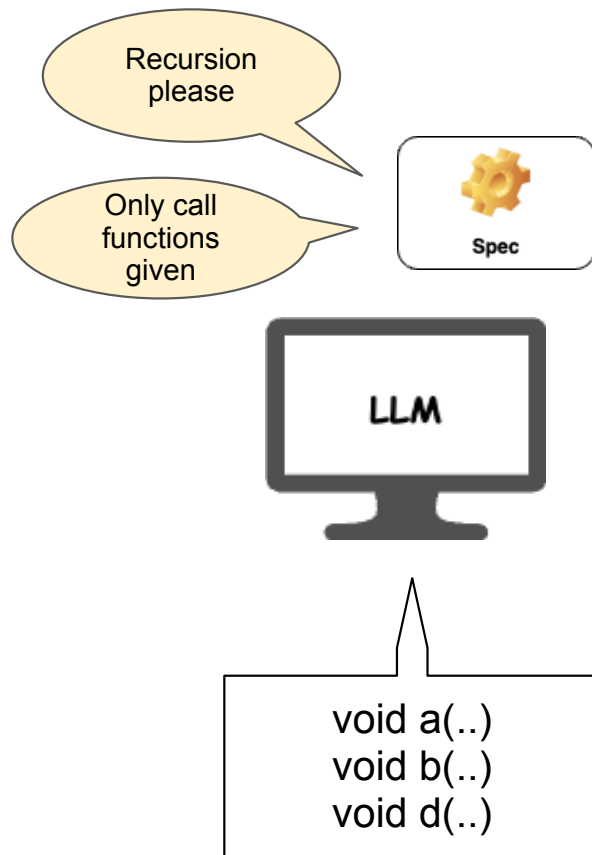
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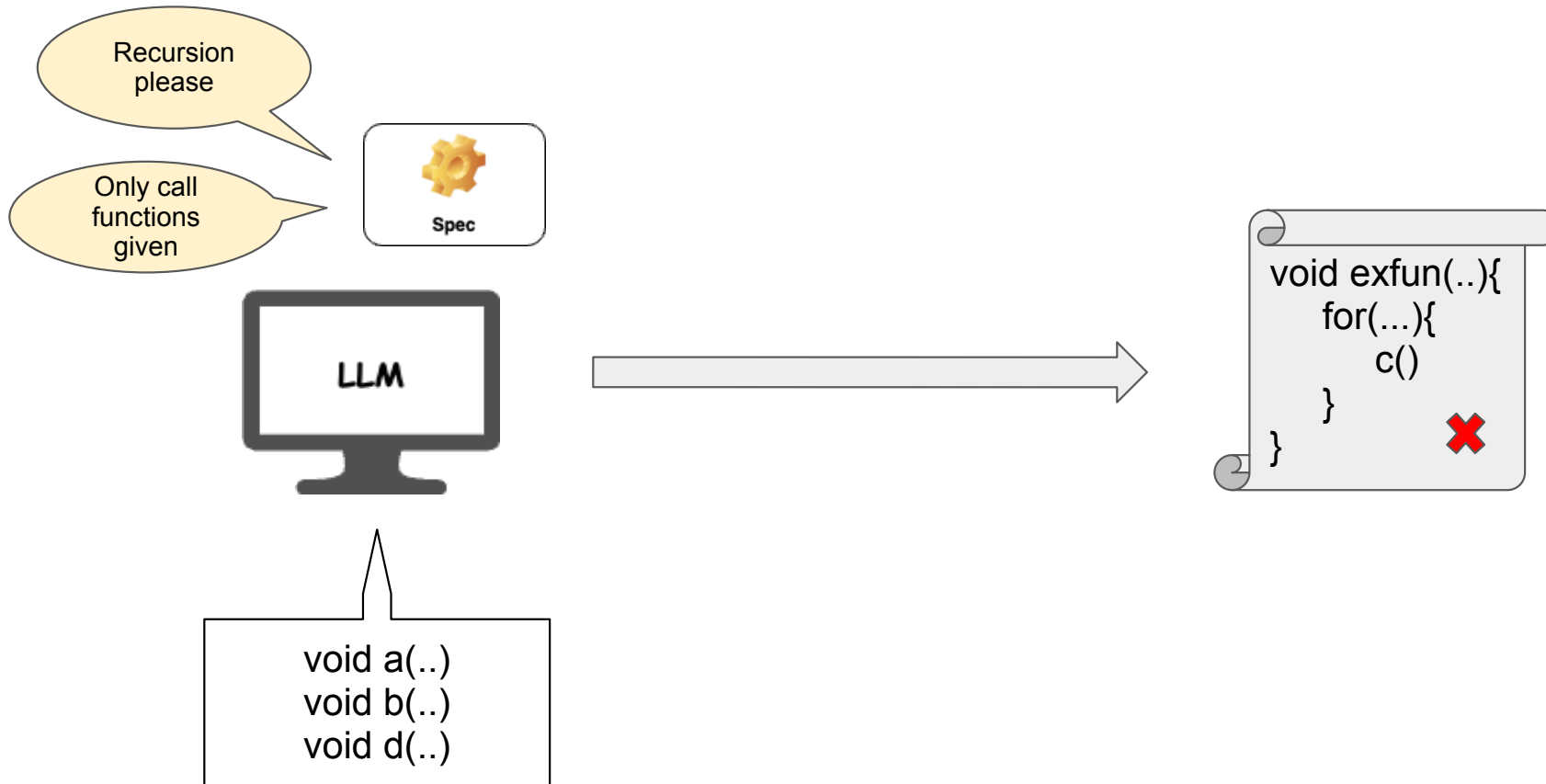
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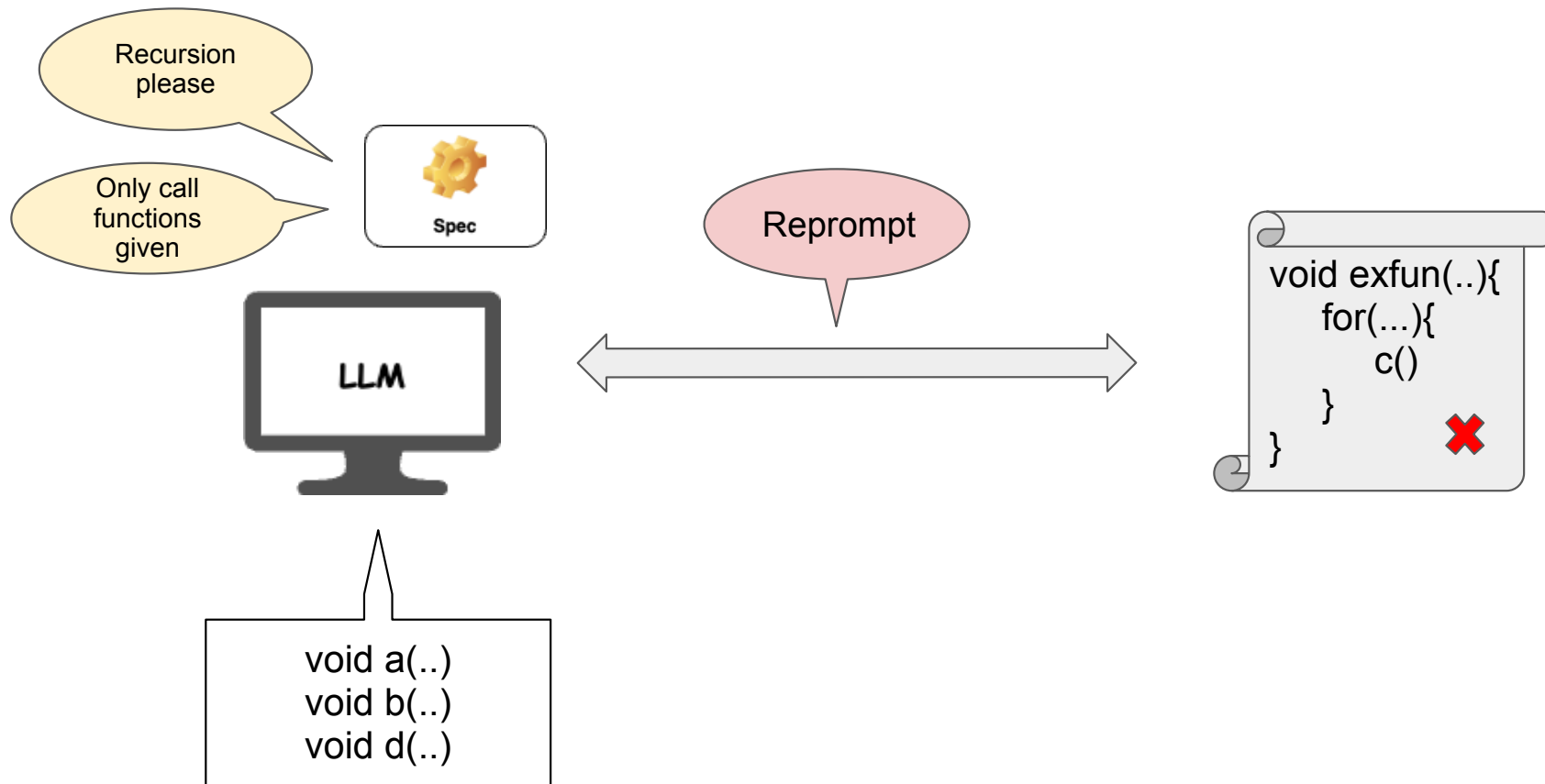
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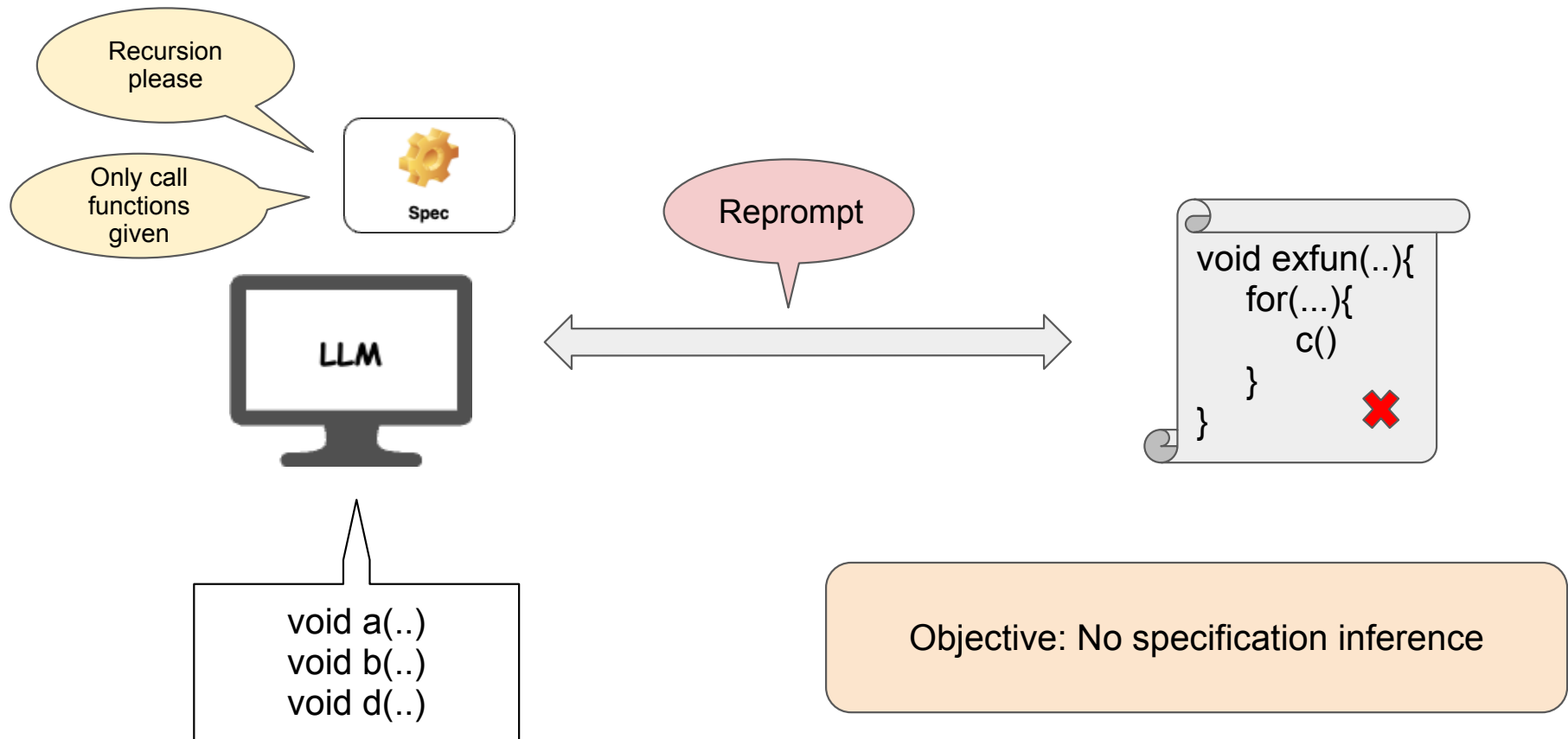
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# Writing automation friendly specifications

- Equivalent specifications amenable to automated reasoning
- Express with predicates like `Forall`, `Forall2`
- Check if an array is sorted in ascending order

$$\forall i, j \ 0 \leq i < len \rightarrow i < j < len \rightarrow a[i] \leq a[j]$$



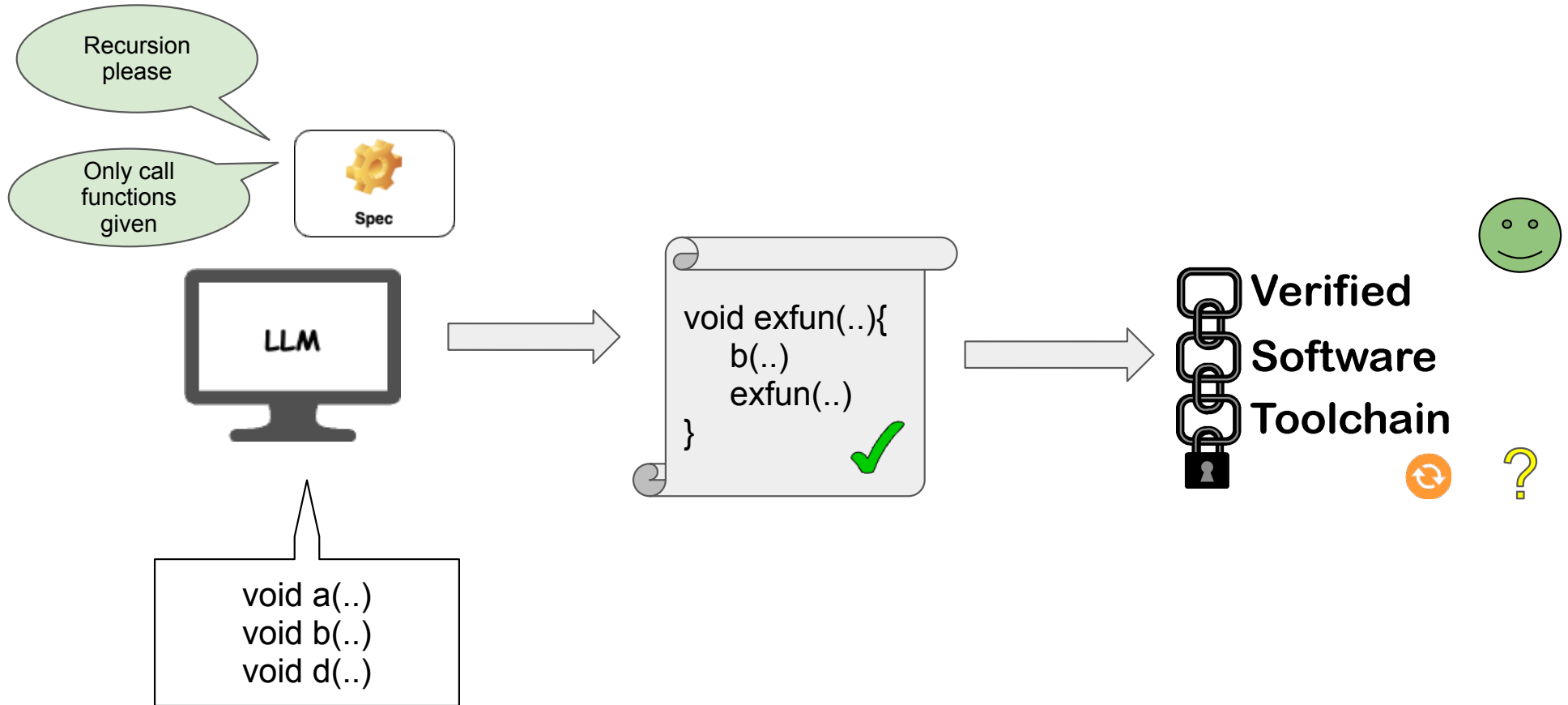
$$\text{Forall2} \leq (\text{sublist } 0 \ (len - 1) \ a) \ (\text{sublist } 1 \ len \ a)$$



list\_solve

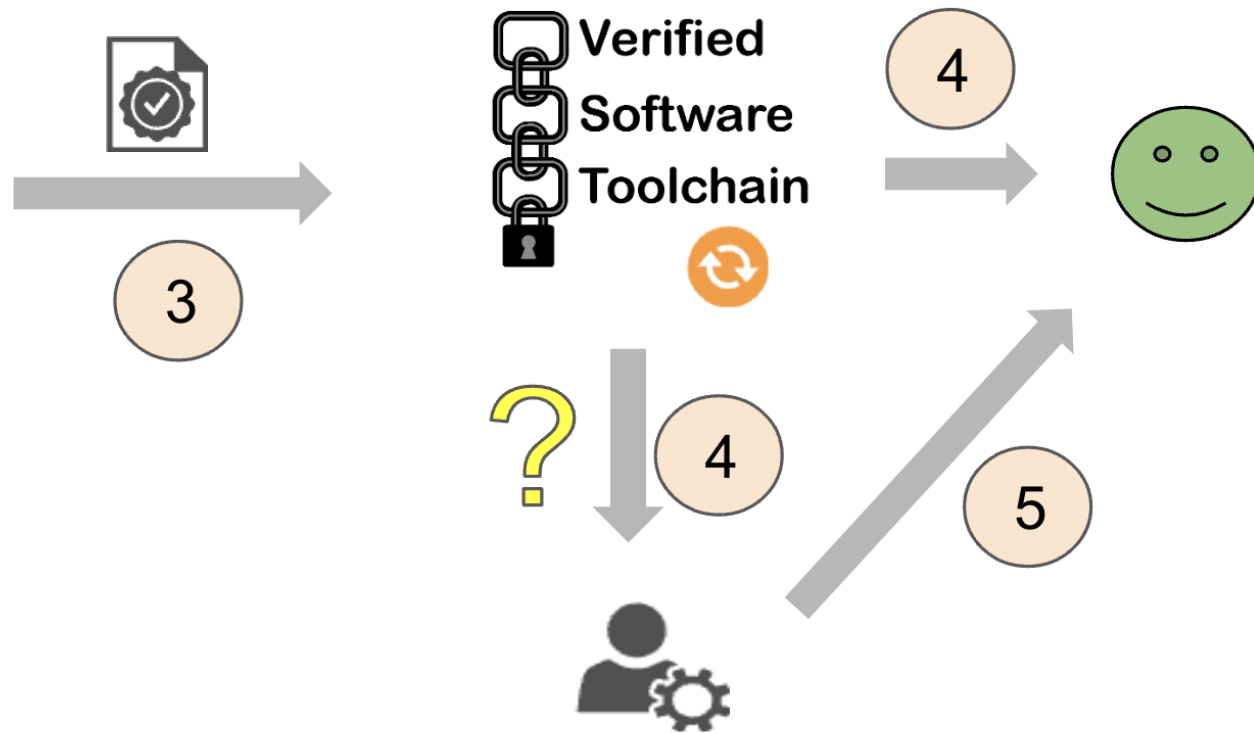


# Biases checked



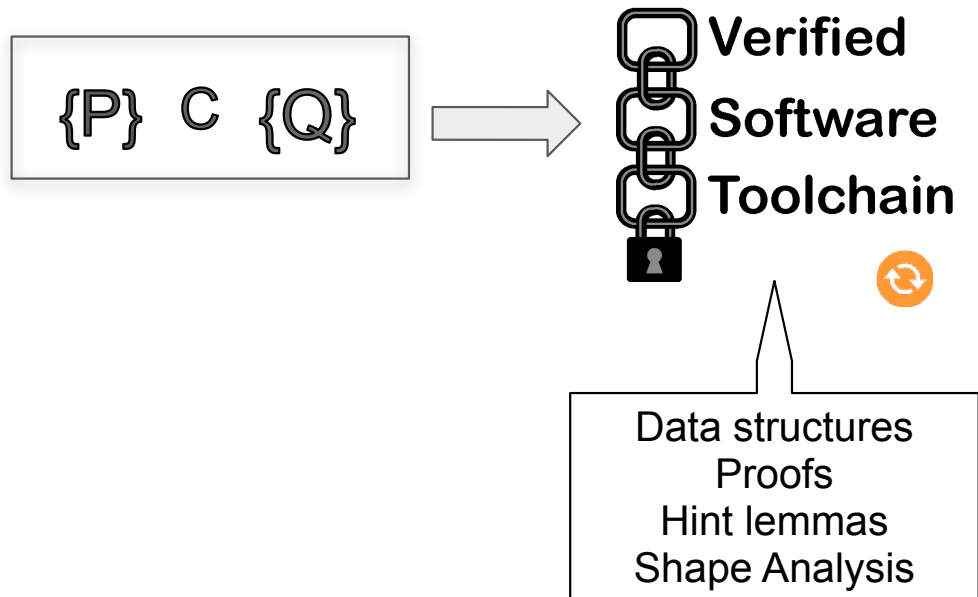
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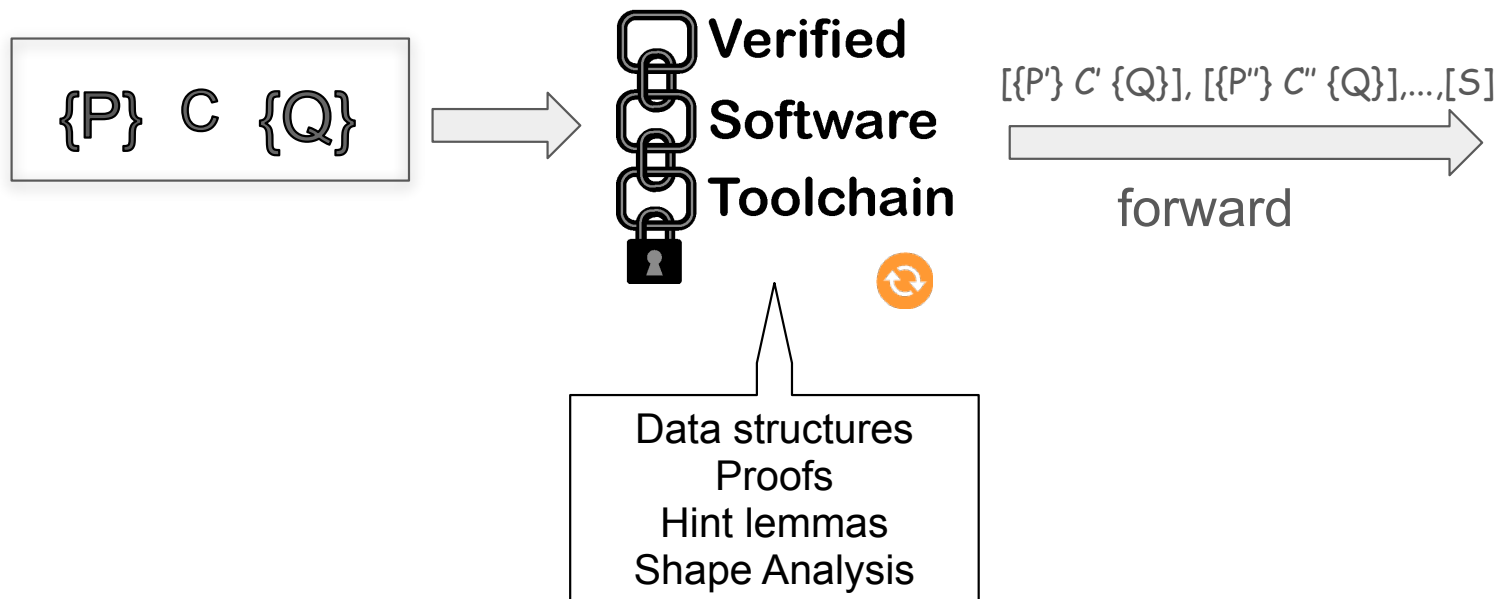


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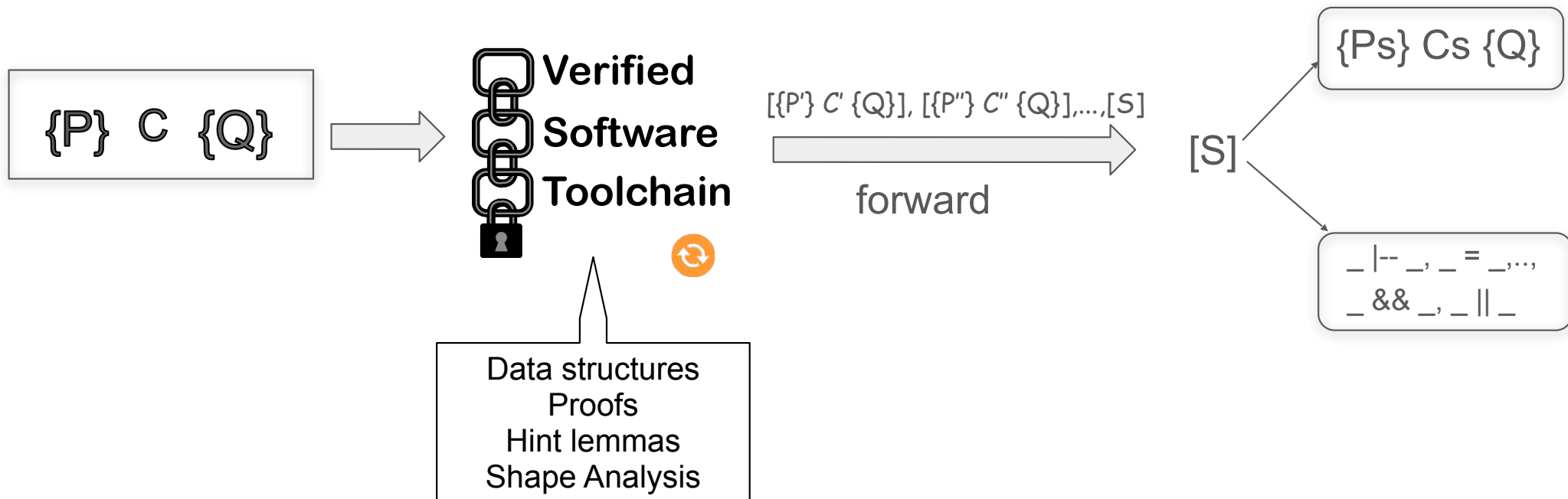
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SepAuto: if statements



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$\{Ps\}$  if  $b$  then  $c$  else  $d$ ;  $e$   $\{Q\}$

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forward\_if ✖

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Requires a joint post-condition

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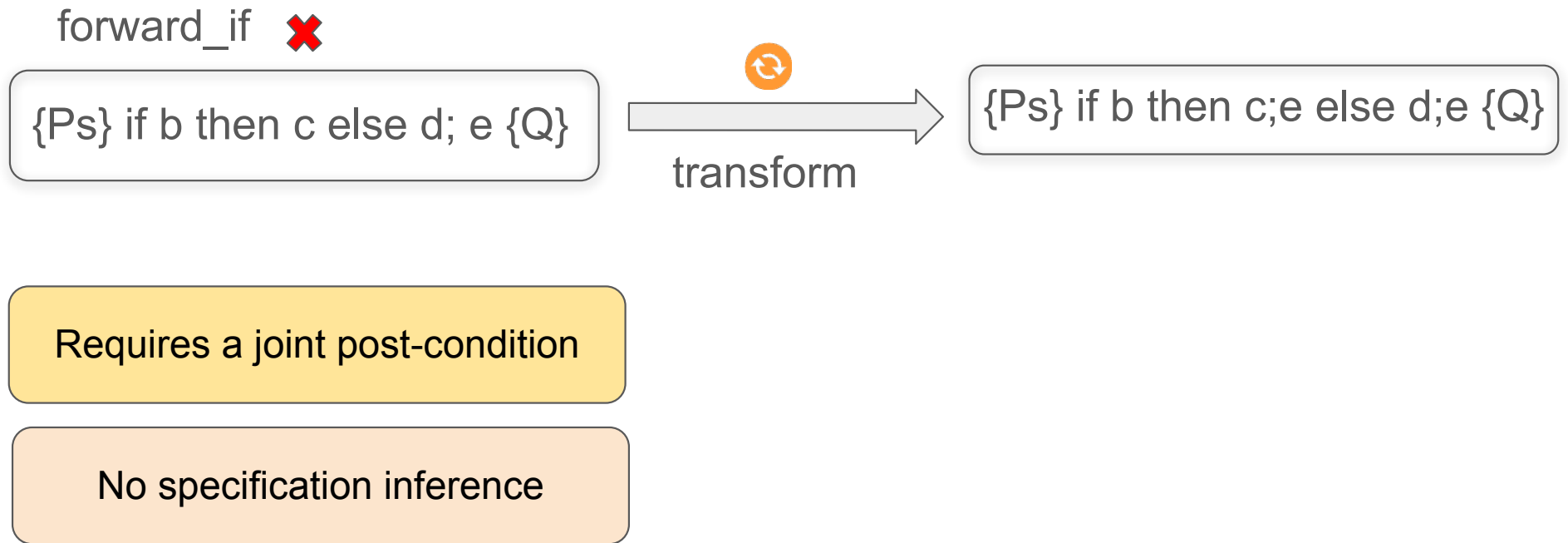
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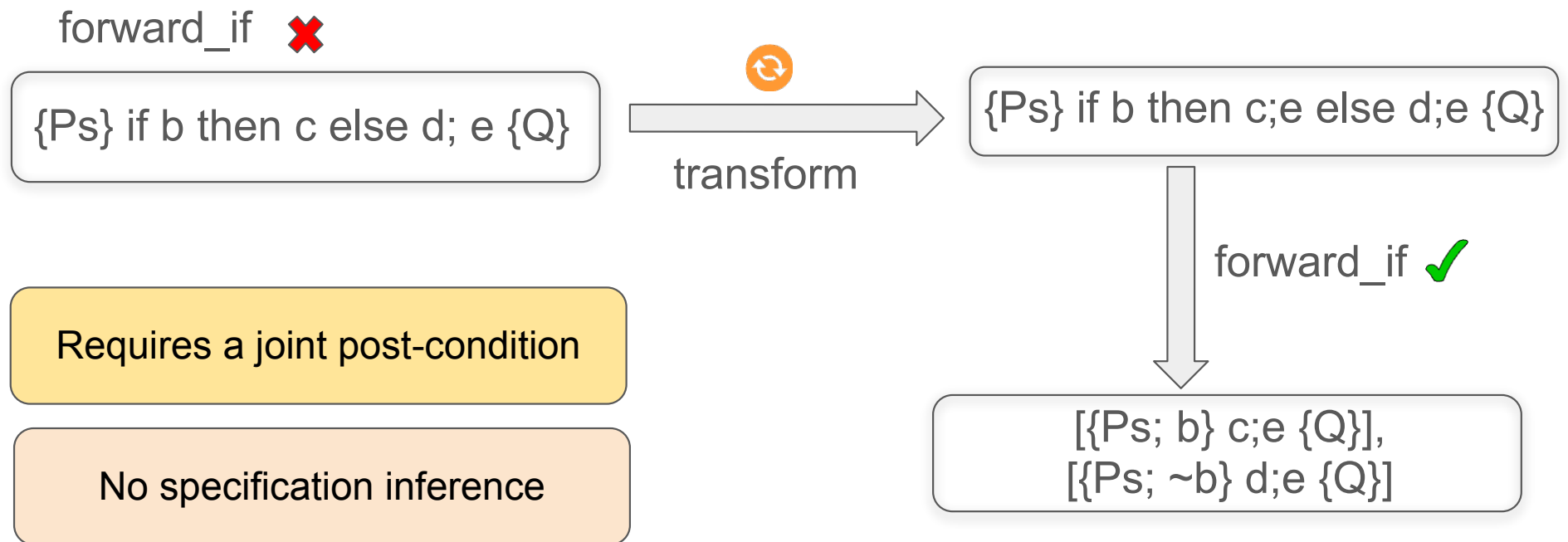
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No specification inference

## SepAuto: if statements



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# SepAuto: Function calls

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$\{Ps\} f(\dots); e \{Q\}$



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forward\_call(....)

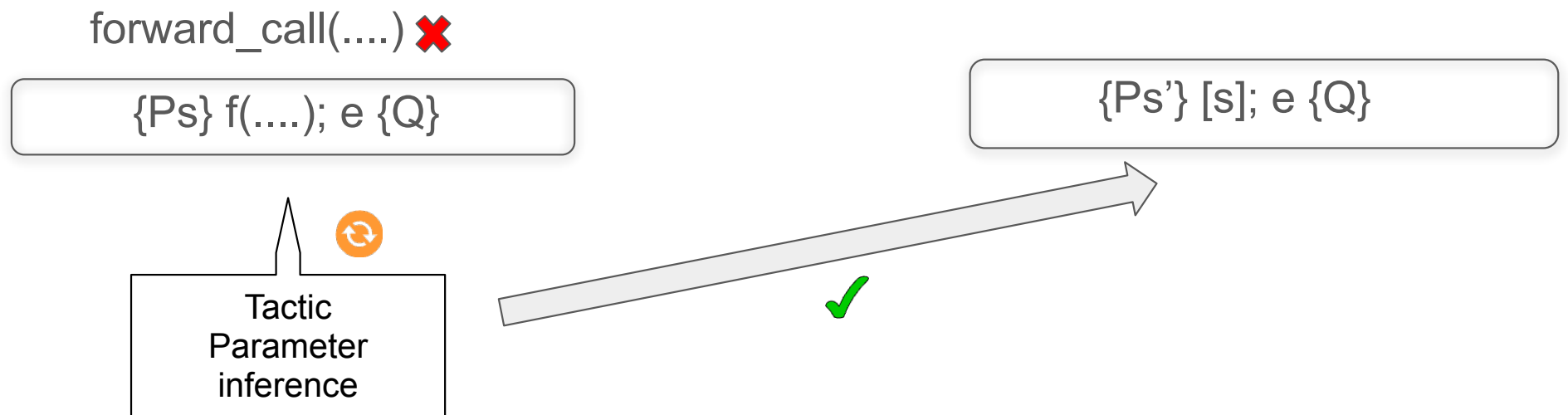
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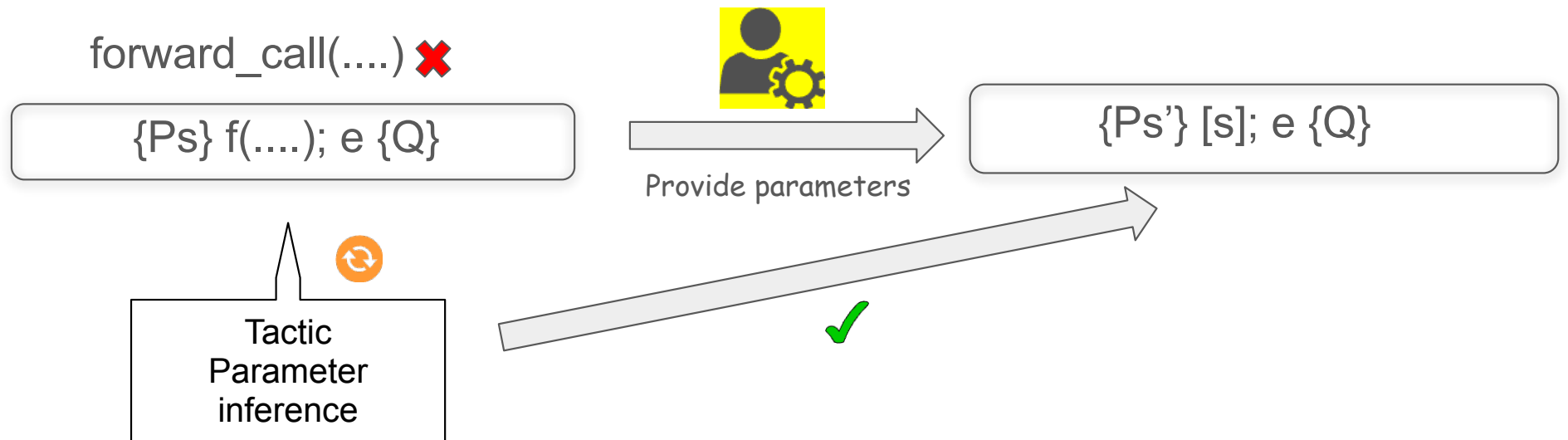
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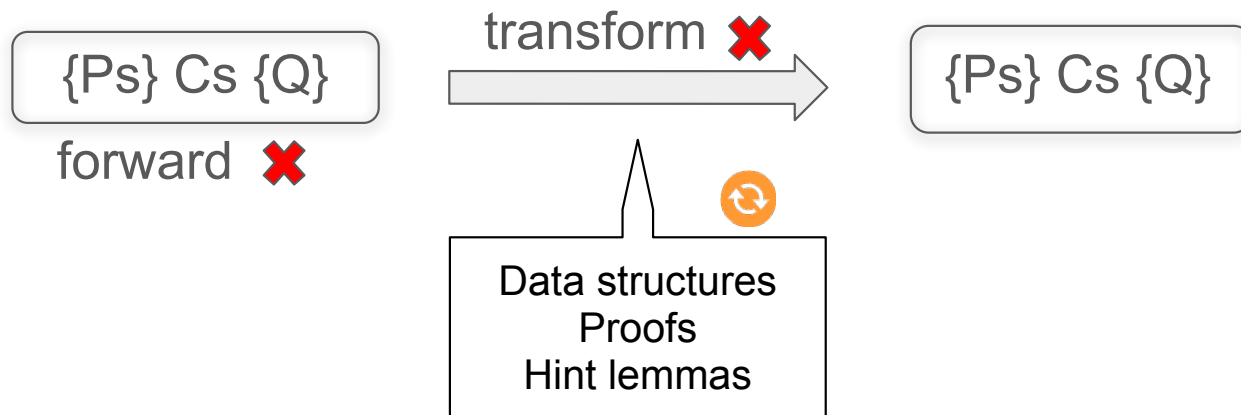
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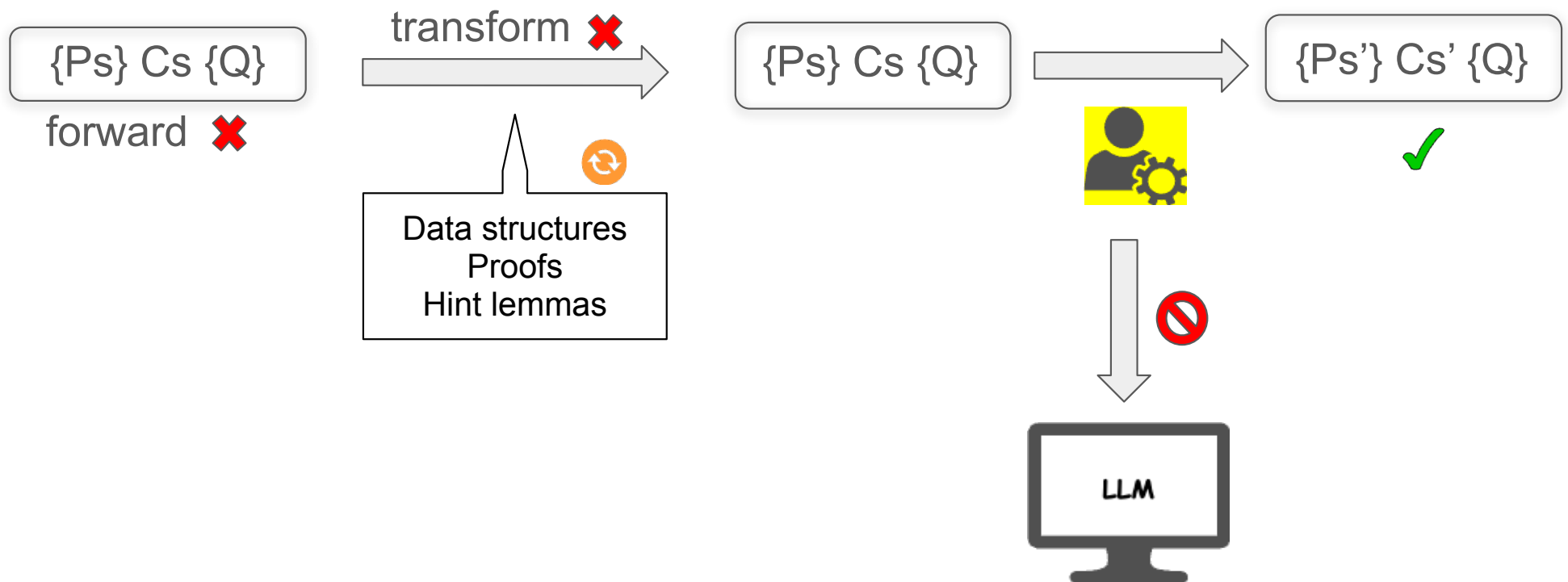
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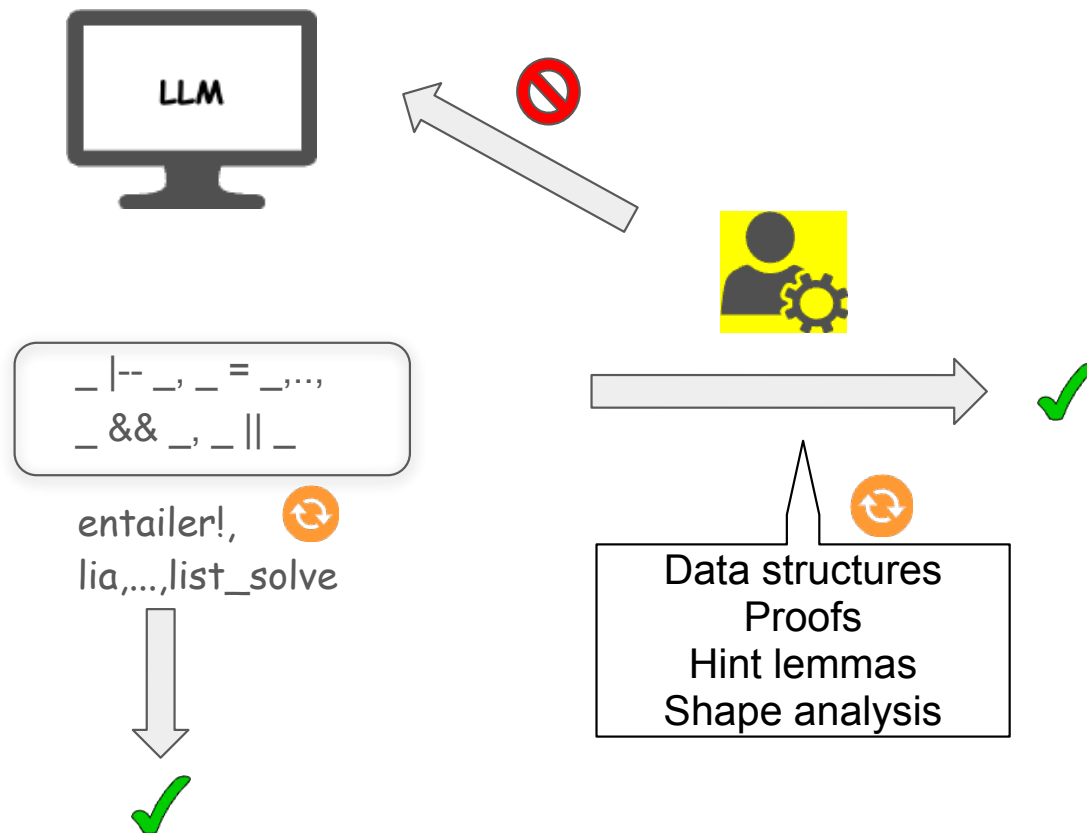
## SepAuto: When forward fails



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# SepAuto: Non-Hoare Triples



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# Experimental Setup

- Benchmark of 50 programs across basic [1], SL [2] and API [3]
- Multiple types supported for programs
- GPT-4o as the LLM
- Prompted with five components

[1] Md Rakib Hossain Misu, Cristina V. Lopes, Iris Ma, and James Noble. FSE (2024), 812–835. <https://doi.org/10.1145/3643763>

[2] Nadia Polikarpova and Ilya Sergey. 2019. Structuring the synthesis of heap-manipulating programs. POPL (2019), 72:1–72:30. <https://doi.org/10.1145/3290385>

[3] Cormen, T.H., Leiserson, C.E., Rivest, R.L., Stein, C.: Introduction to Algorithms, Third Edition. The MIT Press, 3rd edn. (2009)

# Results

- GPT-4o produced correct code in the first try for 49/50 programs
- Incorrectly generated program(s) required detailed information about the algorithm

Benchmark	Bias	Types	Recursion	SC	Attempts
arrayLookUpTable	✓	{Z, S}	✗	0	1
completegraphAdjMatrix	✓	{Z, S}	✗	0	1
deleteBST	✓	{Z}	✗	0	1
generateSkewedBST	✓	{Z}	✗	2	3
addLastVoid	✓	{Z, S}	✗	0	1
stackAPILinkedList	✓	{Z, S}	✗	1	1

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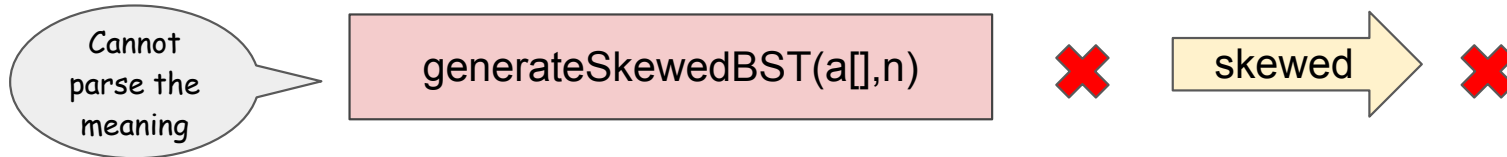


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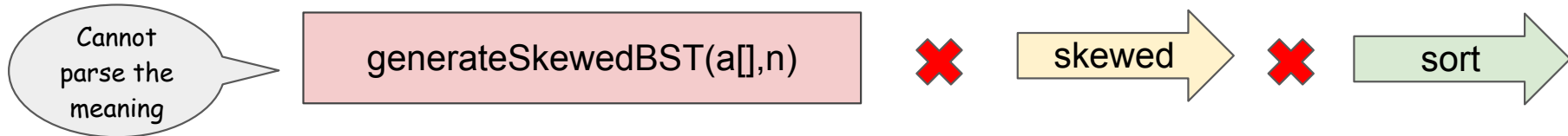


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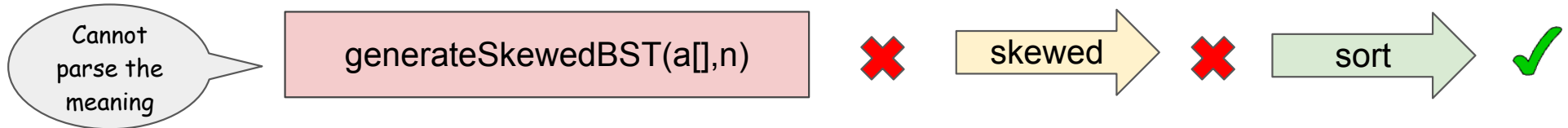
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# Results

- Verifies mostly automatically
- Side conditions primarily generated for complex specifications

SL						Basic					
listDelEnd	✓	{Z, S}	✓	1	1	checkZ	✓	{Z, S}	✓	1	1
listLookup	✓	{Z, S}	✓	0	1	consecNumbers	✓	{Z, S}	✓	0	1
listCopy	✓	{Z, S}	✓	0	1	nIsGreater	✓	{Z, S}	✓	0	1
listInsertEnd	✓	{Z, S}	✓	0	1	firstodd	✓	{Z}	✓	3	1
listFilter	✓	{Z, S}	✓	0	1	arrayMemberValue	✓	{Z, S}	✓	0	1
listAndArraySame	✓	{Z, S}	✓	0	1	isOddAtIndexOdd	✓	{Z}	✓	3	1
listAdd1	✓	{Z}	✓	0	1	lastelempos	✓	{Z, S}	✓	2	1
bstInit	✓	{Z}	✗	0	1	arrayModifyWithOneValue	✓	{Z, R, S}	✓	0	1
bstFree	✓	{Z}	✓	0	1	compareTwoArrays	✓	{Z, S}	✓	0	1
bstLookup	✓	{Z}	✓	0	1	addArrayElementsBy1	✓	{Z}	✓	0	1
bstInsert	✓	{Z}	✓	0	1						
bstMinNode	✓	{Z}	✓	2	1						

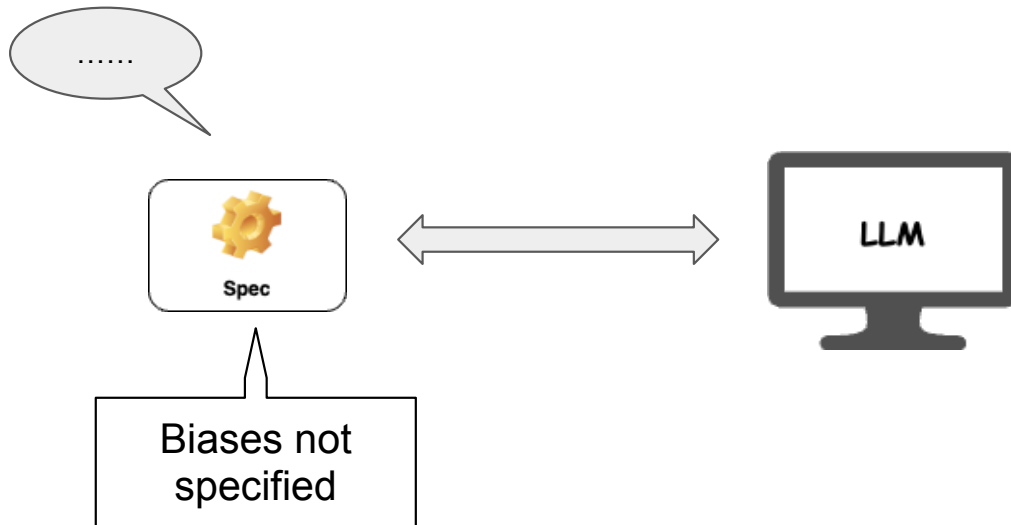
RQ1: Are the biases necessary?

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Without additional biases, what programs do LLMs tend to generate?

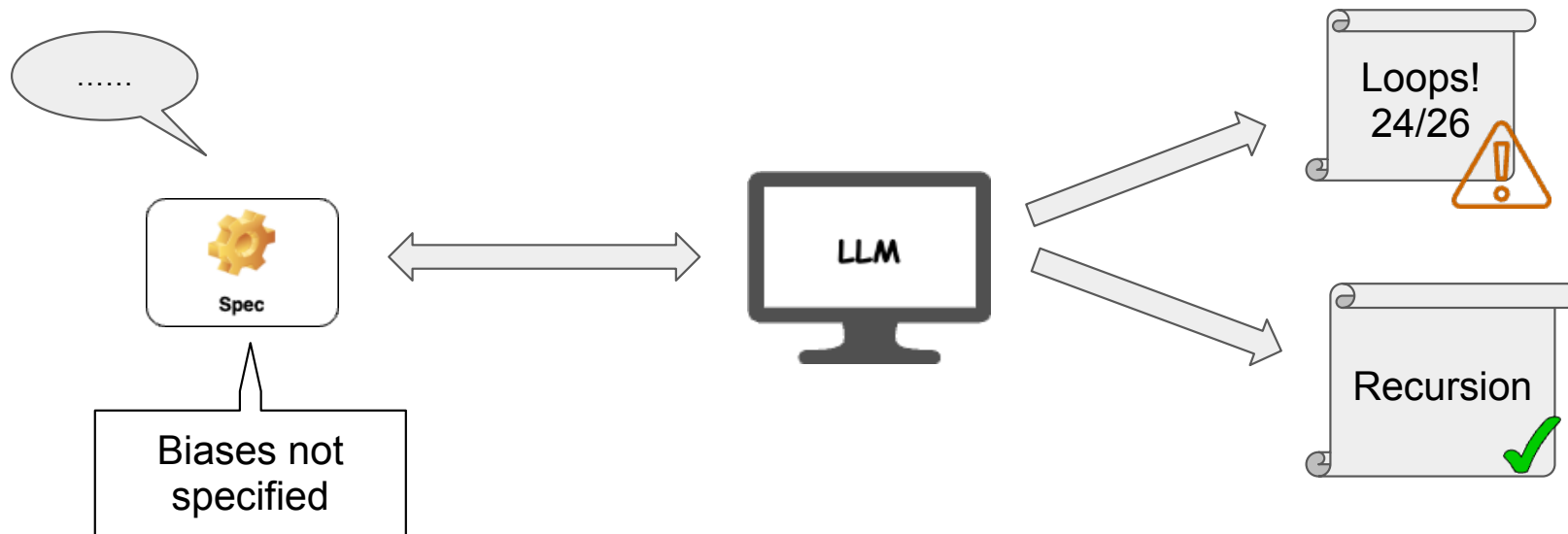
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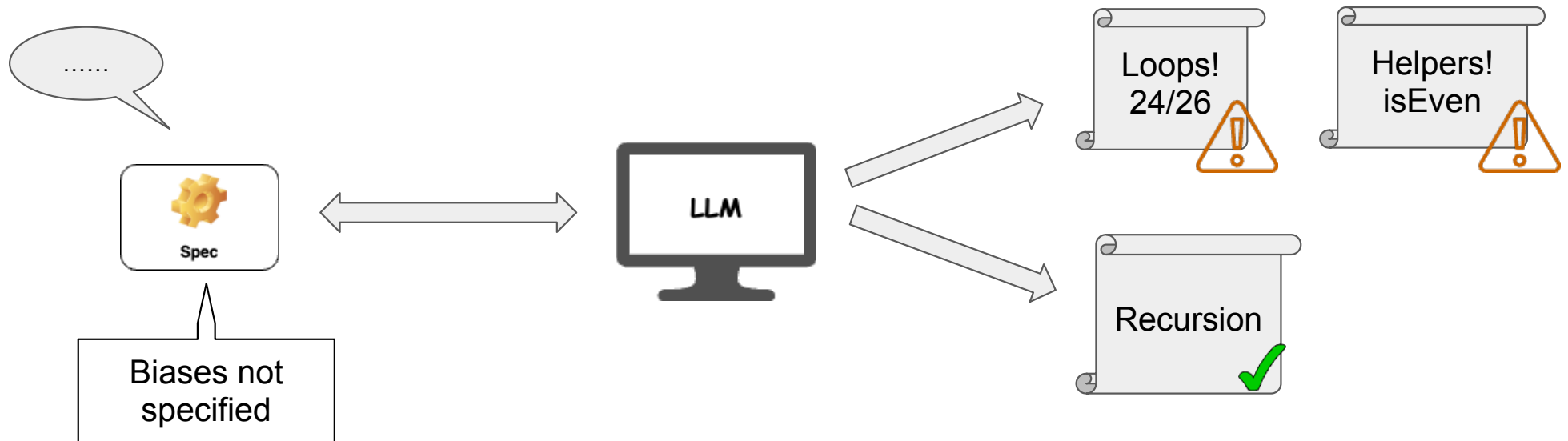
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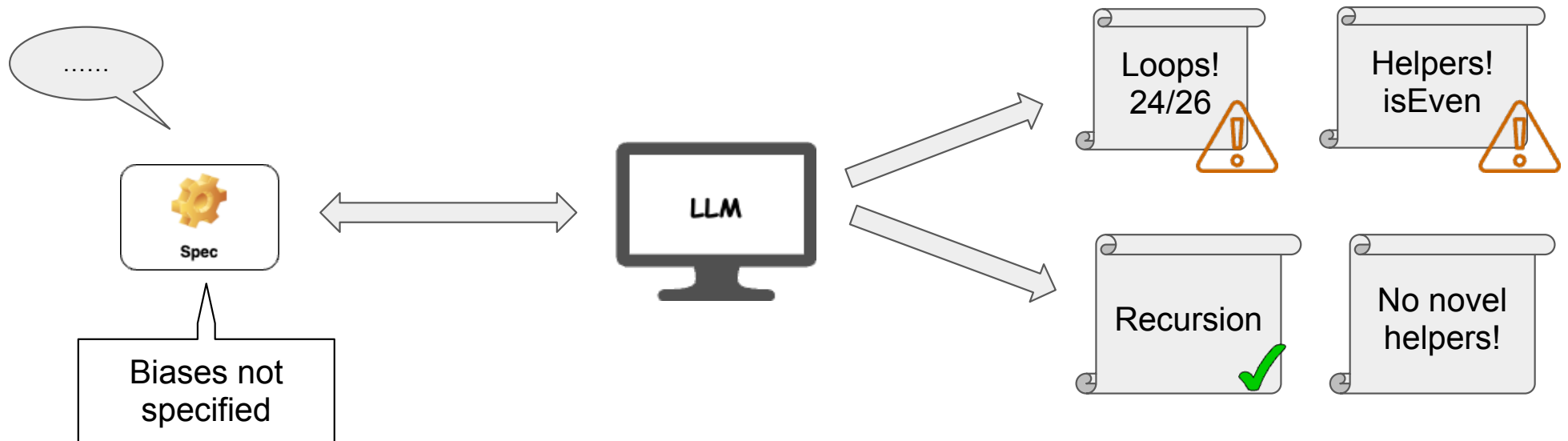
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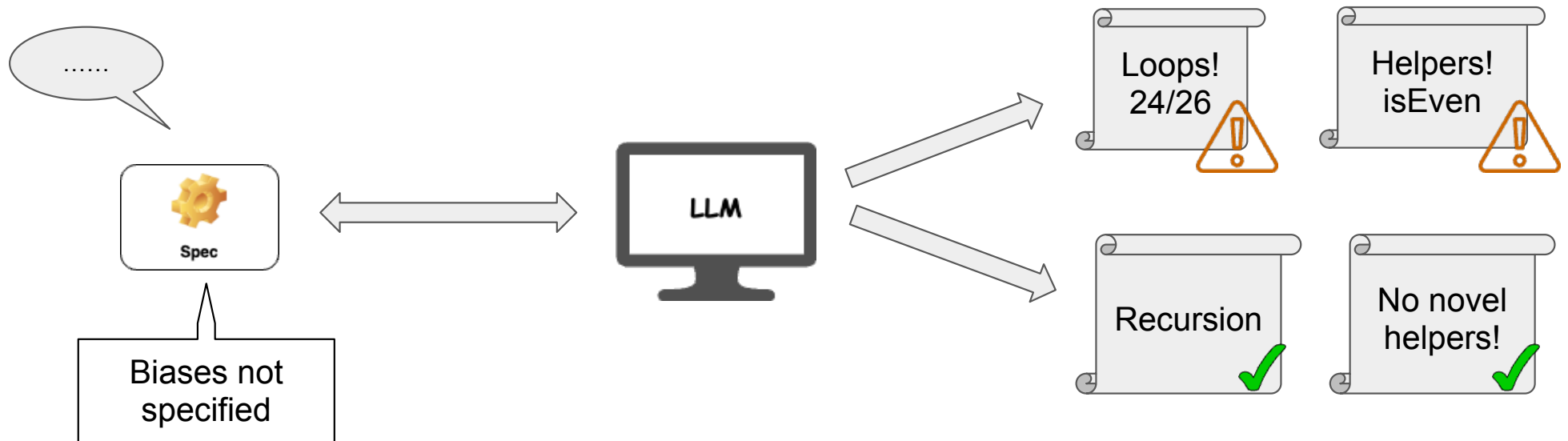
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## RQ2: Where does SepAuto Fail?

Unable to infer relation between two variables

```
H8 : ind >= 0
H6 : ind < Zlength contents
H9 : Znth ind contents <> ele
Hform : Znth vret contents = ele /\
        Forall (fun x : Z => x <> ele)
              (sublist (vret + 1) (ind - 1 + 1) contents)
```

---

(1/1)

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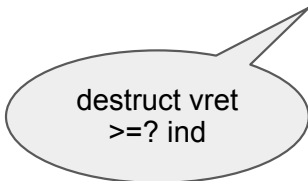
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>=? ind

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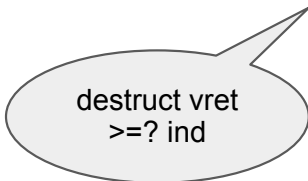
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destruct vret  
>=? ind

sublist\_split

## RQ2: Where does SepAuto Fail?

Unable to pose and prove non-trivial premises

```
H : p <> nullval
```

```
H4 : y = nullval <--> hs = []
```

---

```
(1/1)
```

```
listrep hs y
```

```
|-- listrep (sublist 1 (Zlength (h0 :: hs)) (h0 :: hs)) y
```



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---

(1/1)

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## RQ2: Where does SepAuto Fail?

Automation is unsupported for the wand operator

```
H : p <> nullval
H6 : field_compatible t_struct_tree [] p
H7 : value_fits t_struct_tree
      (Vint (Int.repr k),
       (Vint (Int.repr v), (nullval, pb)))
PNpb : is_pointer_or_null pb
```

---

```
(1/1)
data_at Tsh t_struct_tree
  (Vint (Int.repr k),
   (Vint (Int.repr v), (nullval, pb))) p
* tree_rep b pb
|-- EX a0 : val,
   !! (a0 <> nullval /\ p = a0) &&
   (tree_rep (inorderSuccessor (T E k v b)) a0
    * (tree_rep (inorderSuccessor (T E k v b)) a0 -* 
      tree_rep (T E k v b) p))
```











RQ2.5: SepAuto versus other verifiers?

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Where do other verifiers fail in comparison to SepAuto?











## RQ2.5: SepAuto versus other verifiers?

Where do other verifiers fail in comparison to SepAuto?

Frama C		VeriFast		RefinedC	
Basic		Expressive		Expressive	
Dynamic allocation		Annotation inference		SL implication	
Separation Logic		Precise messaging		Precise messaging	
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```

void addHelper(struct node* n, int x)
  //@ requires n!= 0 &*lseg(n, 0, ?vs);
  //@ ensures lseg(n, 0, append(vs, cons(x, nil)));
{
  //@ open lseg(n, 0, vs);
  if(n->next == 0) {
    // @ open lseg(0, 0, _);
    struct node* nn = create_node(0, x);
    n->next = nn;
  } else {
    addHelper(n->next, x);
  }
}

```

## RQ2.5: SepAuto versus other verifiers?











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<pre>void addHelper(struct node* n, int x)   //@ requires n!= 0 &amp;*&amp; lseg(n, 0, ?vs);   //@ ensures lseg(n, 0, append(vs, cons(x, nil))); {   //@ open lseg(n, 0, vs);   if(n-&gt;next == 0) {     // @ open lseg(0, 0, _);     struct node* nn = create_node(0, x);     n-&gt;next = nn;   } else {     addHelper(n-&gt;next, x);   } }</pre>				Hint Lemmas	

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}

```

Postcondition





## RQ2.5: SepAuto versus other verifiers?

```
[[rc::parameters("l : {list Z}", "p : loc", "n : Z" )]]
[[rc::args("p @ &own<l @ list_t>", "n @ int<i32>")]]
[[rc::exists("b : bool")]]
[[rc::returns("b @ builtin_boolean")]]
[[rc::ensures("own p : l @ list_t", "{b ↔ n ∈ l}")]]
//[[rc::tactics("all: try set_solver.")]]
bool member_rec (list_t *p, int k) {
    if (*p == NULL) {
        return false;
    }
    int head = (*p)->val;
    if (head == k) {
        return true;
    }
    return member_rec(&(*p)->next, k);
}
```

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```



```
HCASE : CASE_DISTINCTION_INFO
        (optional == ... : n :: x1 ≠ [])
        [loc_38; loc_38]
HCASE0 : CASE_DISTINCTION_INFO
        (if false, false) [loc_38]
H1 : n ≤ max_int i32
H0, H2 : min_int i32 ≤ n
H3 : n ≤ max_int i32
HCASE1 : CASE_DISTINCTION_INFO
        (if bool_decide (n = n), true)
        [loc_21]
```

---

```
(1/3)
True ↔ n ∈ n :: x1
```

---

```
(2/3)
x' ↔ n ∈ x0 :: x1
```

---

```
(3/3)
False ↔ n ∈ []
```

## RQ2.5: SepAuto versus other verifiers?

```
[[rc::parameters("l : {list Z}", "p : loc", "n : Z" )]]
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    return true;
  }
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}
```



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        [loc_21]
```

---

(1/3)  
True ↔ n ∈ n :: x1

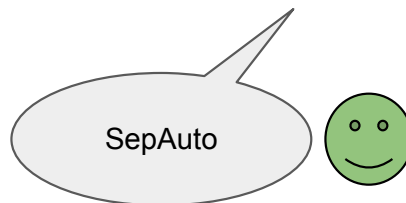
---

(2/3)  
x' ↔ n ∈ x0 :: x1

---

(3/3)  
False ↔ n ∈ []

---



RQ3: Do the components of the prompt matter?

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English Description



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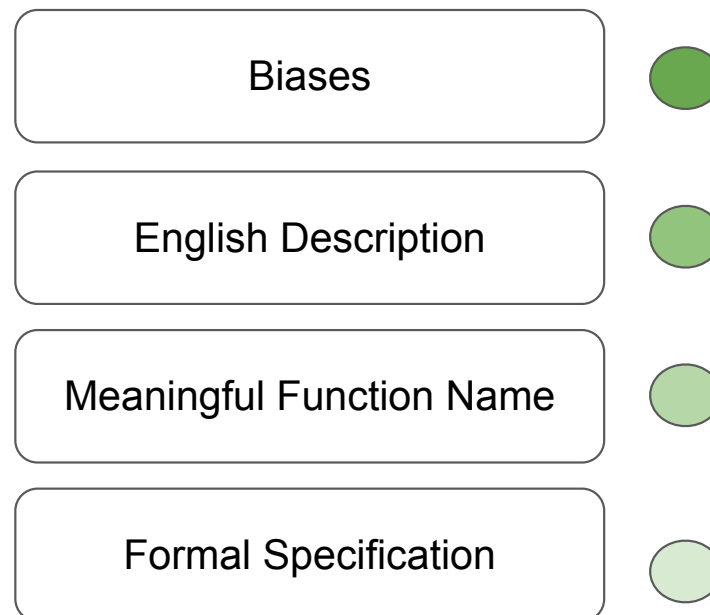


Meaningful Function Name








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Biases	
English Description	
Meaningful Function Name	
Formal Specification	
I/O Examples	

## RQ3: Prompt Components

Biases



English



Name



Formal Spec.



I/O Examples



*Definition*  $fi2 := fun\ x \Rightarrow x =? 2$   
 $listrep\ l\ h\ * listrep\ (filter\ fi2\ l)\ a$

## RQ3: Prompt Components

Biases



English



Name



Formal Spec.



I/O Examples



*Definition*  $fi2 := fun\ x \Rightarrow x ==? 2$   
 $listrep\ l\ h\ * listrep\ (filter\ fi2\ l)\ a$

```
#include "treelistdef.c"
struct sll* abc(struct sll* h) {
    if (h == NULL) {
        return NULL;
    }
    if (h->key == 2) {
        h->next = abc(h->next);
        return h;
    }
    else {
        return abc(h->next);
    }
}
```

## RQ3: Prompt Components

Biases



English



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Formal Spec.



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```

$listrep\ (filter\ fi2\ l)\ a$

## RQ3: Prompt Components

Biases



English



Name



Formal Spec.



I/O Examples



Basic: 17/19

SL: 12/23

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```

$listrep\ (filter\ fi2\ l)\ a$



## RQ3: Prompt Components

Biases 

English 

Name 

Formal Spec. 

I/O Examples 

*Definition*  $fi2 := fun\ x \Rightarrow x \leq 2$   
 $listrep\ l\ h\ * listrep\ (filter\ fi2\ l)\ a$

# RQ3: Prompt Components

Biases



English



Name



Formal Spec.



I/O Examples



```
#include "treelistdef.c"
struct sll* copyfilteredListWhereElementis2AndReturnItsHead(struct sll *h)
{
    if (h == NULL) {
        return NULL;
    }
    if (h->key == 2)
    {
        struct sll* new_node = (struct sll*)malloc(sizeof(struct sll));
        new_node->key = 2;
        new_node->next = copyfilteredListWhereElementis2AndReturnItsHead(h->next);
        return new_node;
    }
    else
    {
        return copyfilteredListWhereElementis2AndReturnItsHead(h->next);
    }
}
```

*Definition*  $fi2 := \text{fun } x \Rightarrow x ==? 2$   
 $\text{listrep } l \ h \ * \ \text{listrep } (\text{filter } fi2 \ l) \ a$

# RQ3: Prompt Components

Biases



English



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Formal Spec.



I/O Examples




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        return new_node;
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    {
        return copyfilteredListWhereElementis2AndReturnItsHead(h->next);
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```

*Definition*  $fi2 := fun\ x \Rightarrow x ==? 2$   
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# RQ3: Prompt Components

Biases 

English 

Name 

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I/O Examples 

Basic: 17/19

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```

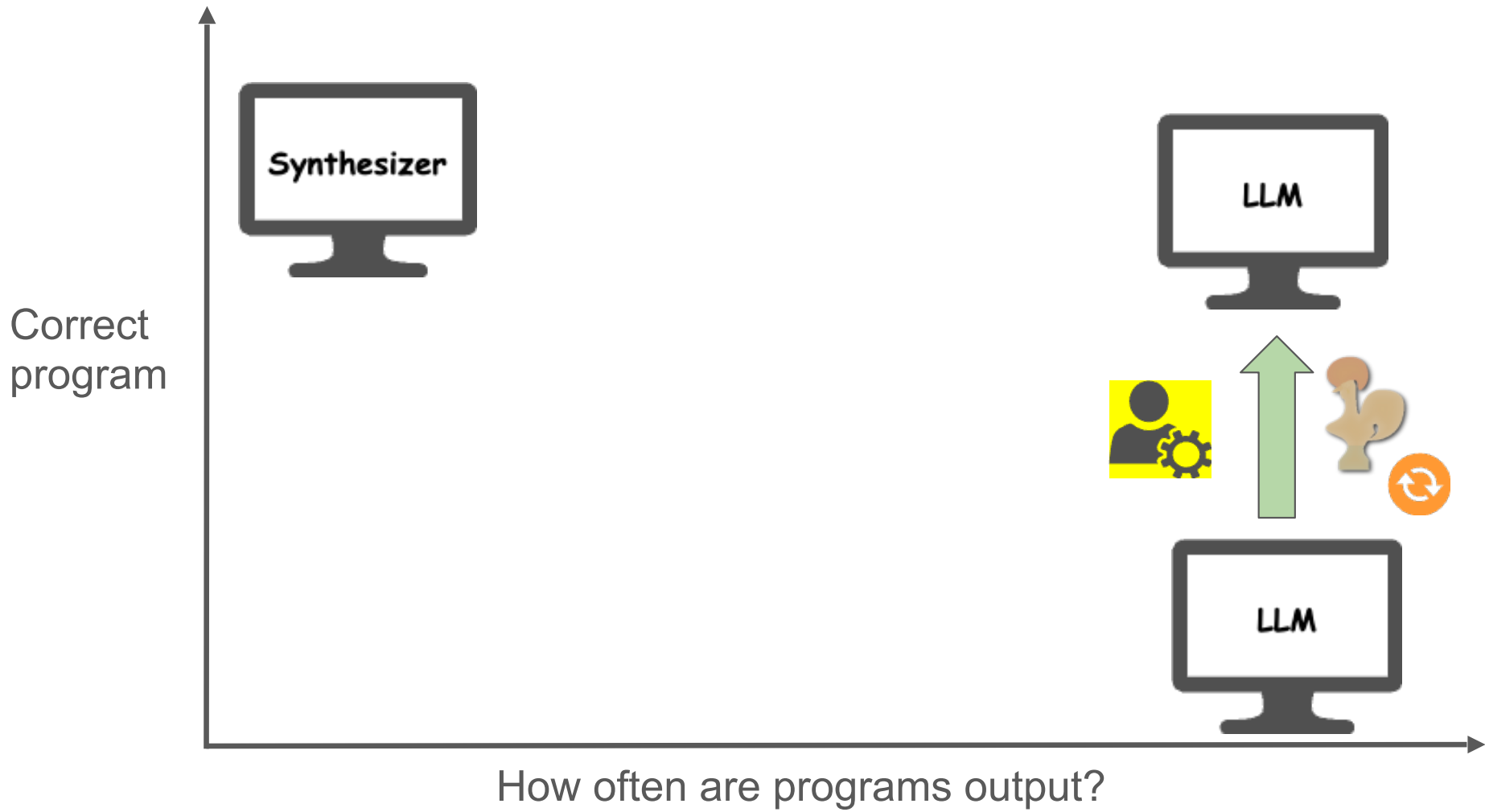
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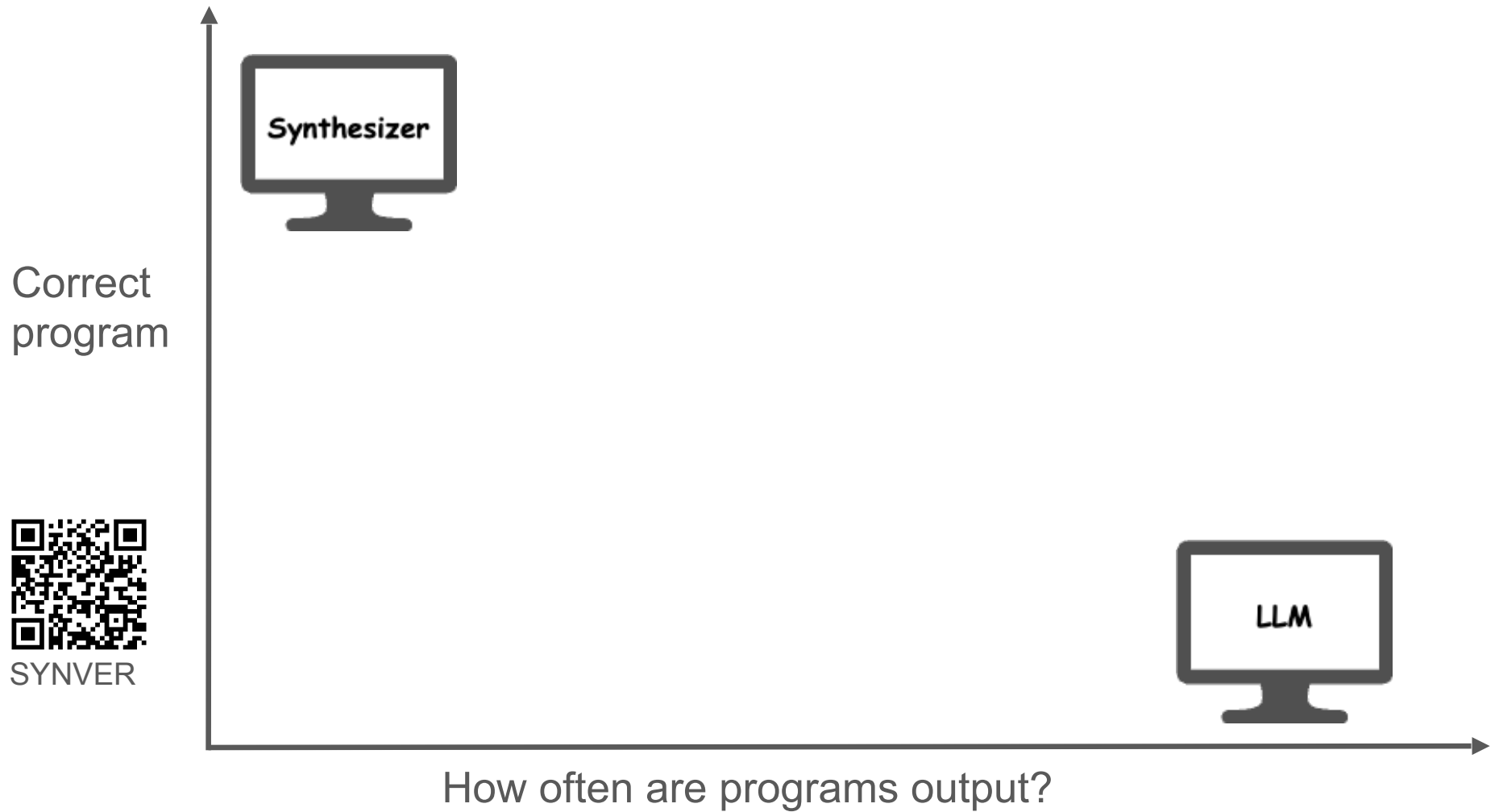
# Conclusions

- LLMs with foundational verifiers as general purpose synthesizers
- Large search space well suited for a popular, mainstream language like C
- LLMs respond positively to biases and natural language
- Automation is well suited to the biases
- SepAuto is not complete
- Data leakage may have required less re prompting

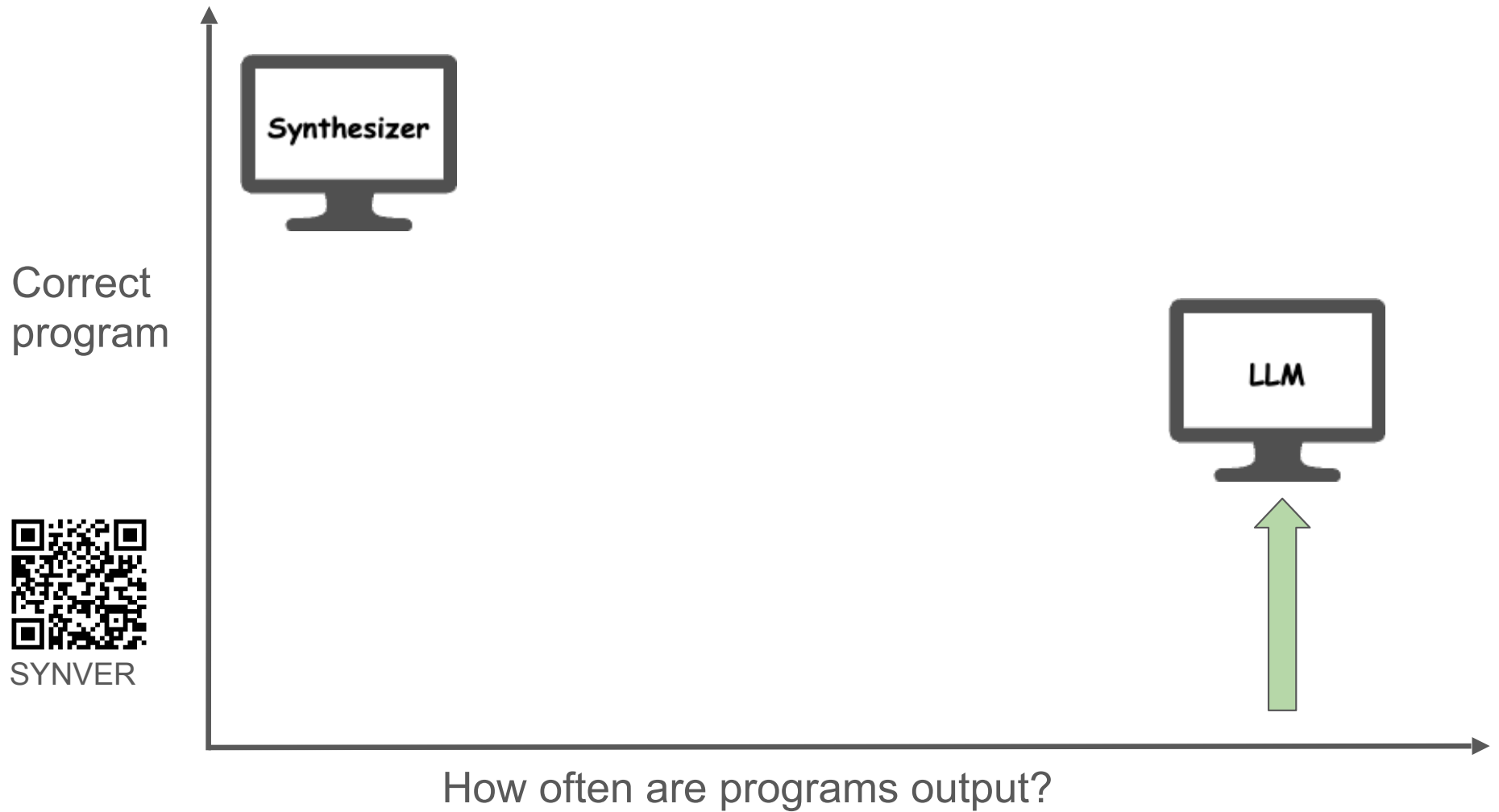
## Current State



# What's next?

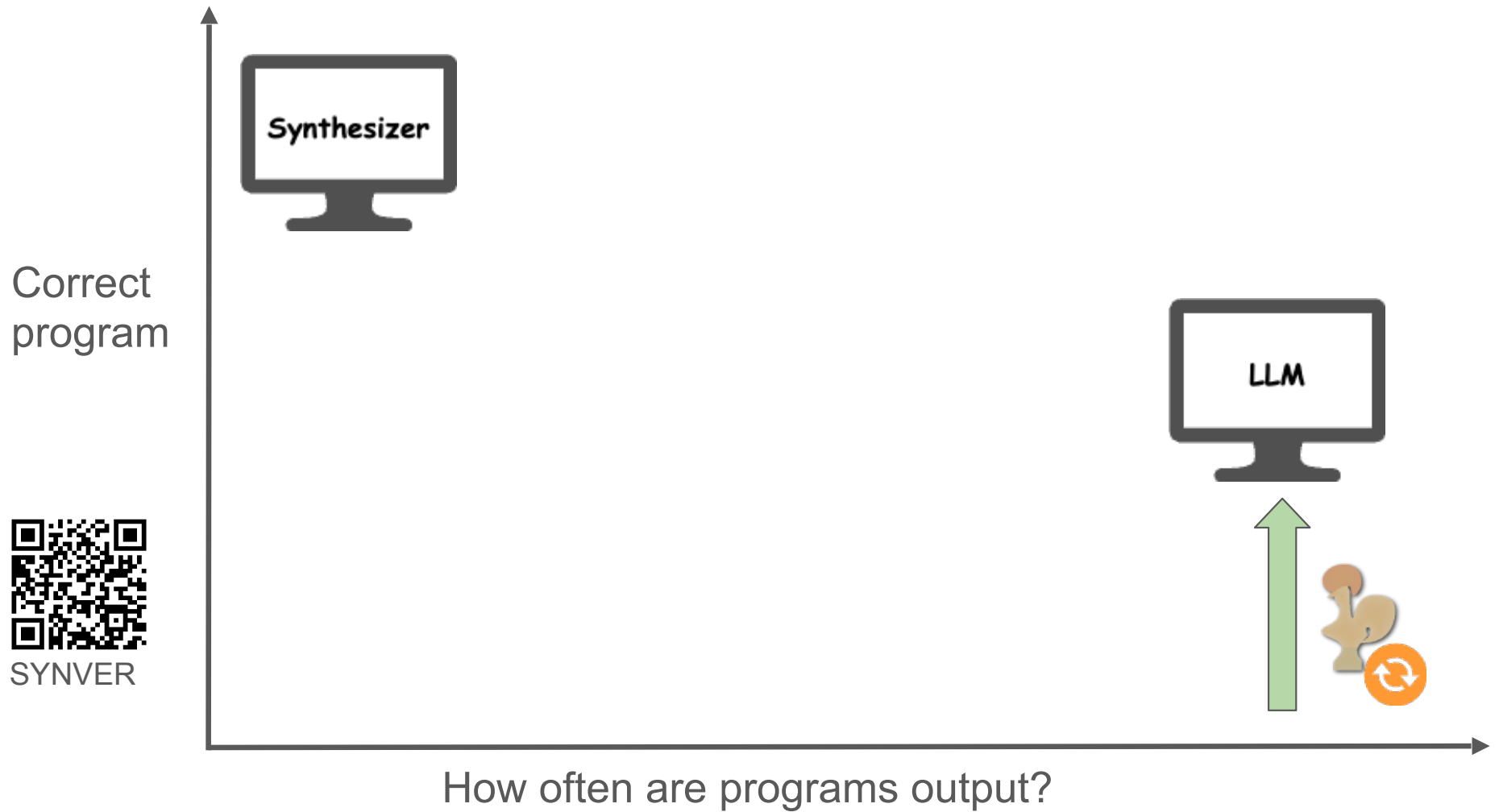


# What's next?

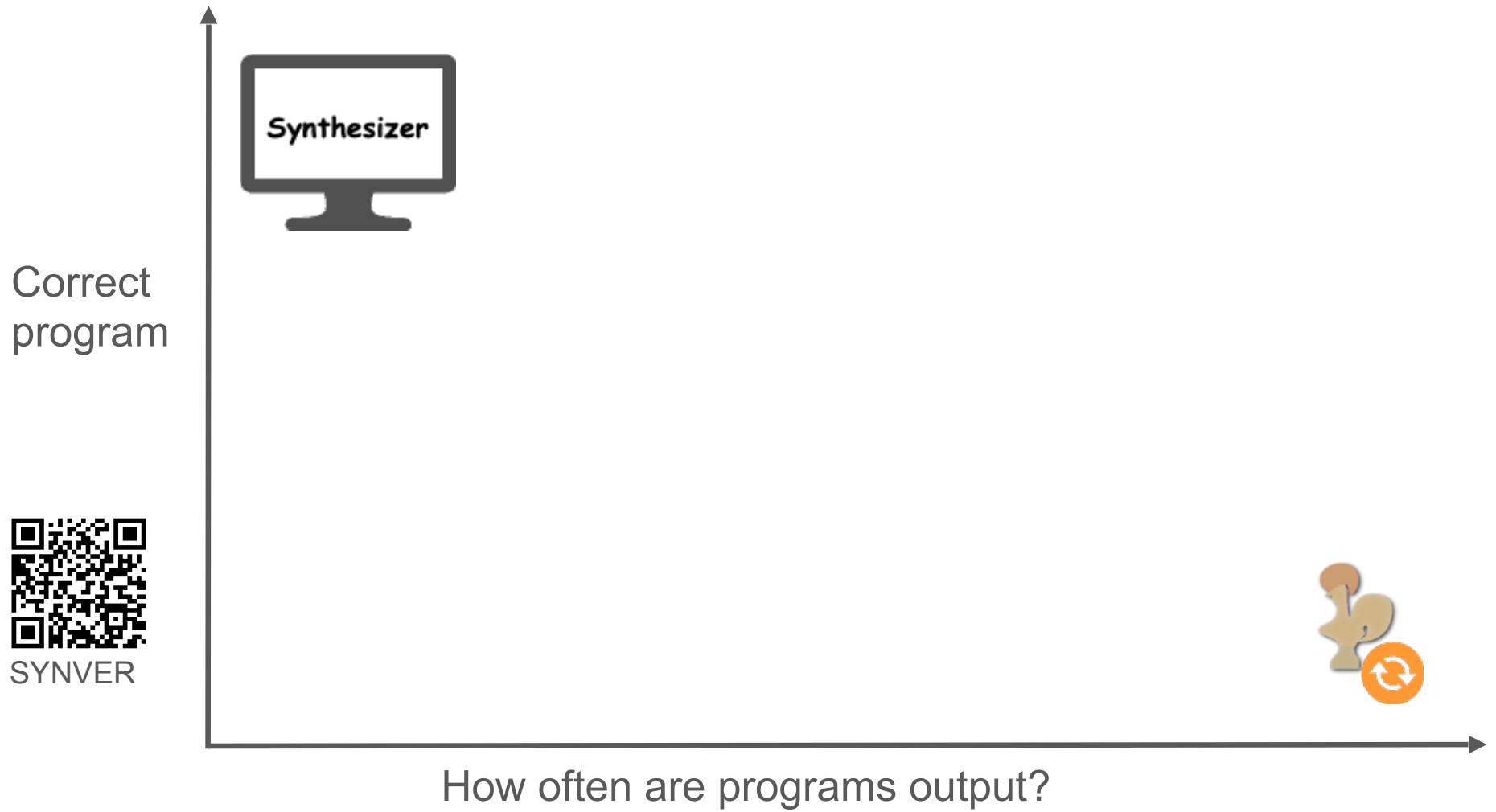




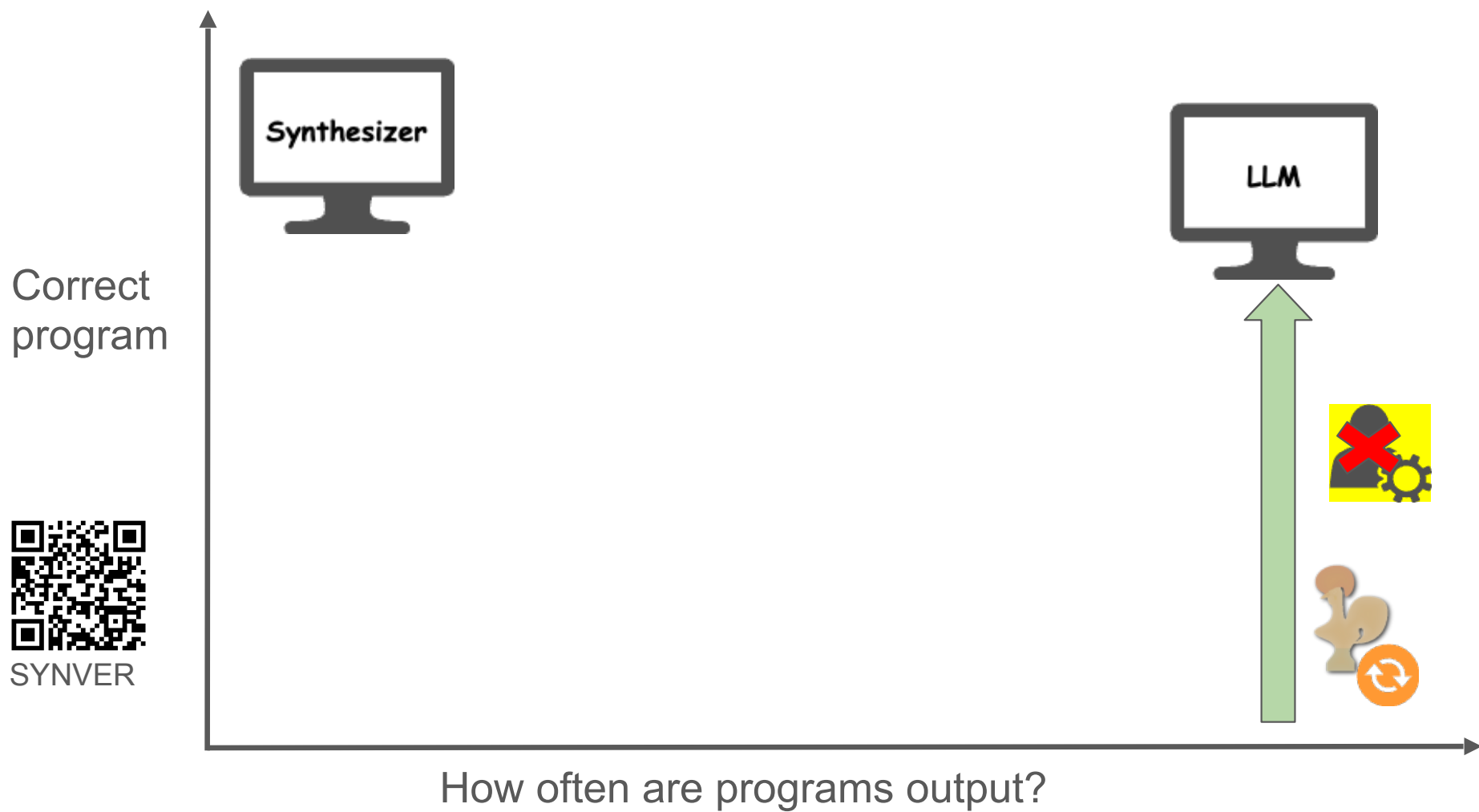
# What's next?



# What's next?



# What's next?



# What's next?

