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## HandsOn Assignment

$F_1, \dots, F_k$  set of subjects having final exam.

$S_1, \dots, S_n$  set of students giving them

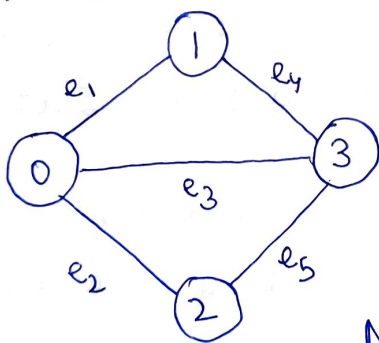
$h$  is the number of slots they have to be scheduled

Constraint: Student cannot give two exams in single slot

A proposed solution can be to take  $F_1, \dots, F_k$  as nodes of graph and the edges connecting them be  $S_1, \dots, S_n$ . So the problem reduces to find the chromatic number of the graph which would be equal to  $h$  slots together with the constraint mentioned above.

Random Graph of  $n$  nodes can be generated

Take  $n=4$



The edges  $e_1, e_2, e_3, e_4, e_5$  represent the students

Now each vertex can be represented by  $h$  colors

So Vertex Constraints: let's take  $V_i$  it would be

Assign color:  $x_1 \vee x_2 \vee \dots \vee x_h$

Exactly one color:  $(\neg x_1 \vee \neg x_2) \wedge (\neg x_1 \vee \neg x_3) \dots \wedge (\neg x_k \vee \neg x_h)$

Edge Constraints Take  $V_i$

So the assigning & formulation is on next page.

Color 1 (Slot 1) :  $(\neg x_1 \vee \neg x_{h+1})$

Color 2 (Slot 2) :  $(\neg x_2 \vee \neg x_{h+2})$

Color h (Slot h) :  $(\neg x_h \vee \neg x_{2h})$

Similarly for all other edges  $e_2, e_3, \dots, e_g$  constraints can be calculated

With all the constraints (Vertex & Edge) of Graph given to SAT Solver will give the optimal no of slots (colors) to be used for scheduling the final exam subject allocation to respective Students.

Note: The graph has no adjacent vertex (final exam subjects) with same color as the student (represented by graph edge) cannot give two exams (nodes) in a single slot (color).

So for the SAT solver, the constraints can be ordered as Vertex and Edge constraints in .cnf file accordingly specifying the number of clauses ( $x$ ; variables) and constraints in the header of it.

The not in cnf file will be treated as - so our constraints can be updated with the respective clauses accordingly.

The following problem is an allocation problem which can be reduced to a NP complete problem which can be passed to the SAT Solver with appropriate clauses & constraints which gives the solution for the same.