

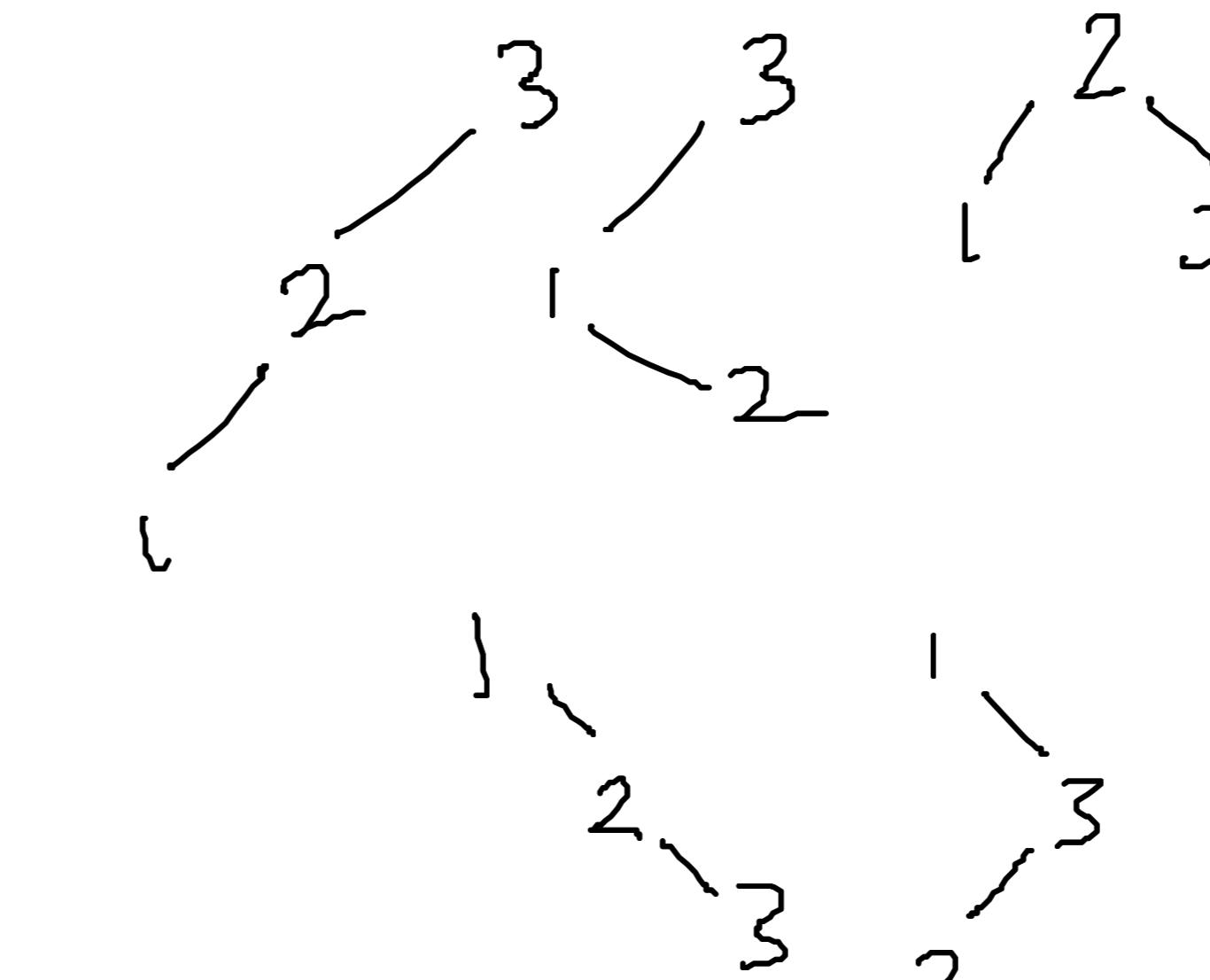
Next Q?

Find No of

$N=3$

- ✓ BST
- BT's
- ✓ Unlabelled BT's

Using $\frac{N \text{ Nodes}}{(1 \dots N)}$



5 BST's

$$BT = \frac{2^n}{(n+1)!} \cancel{\left(\frac{n+1}{2}\right)^n}$$

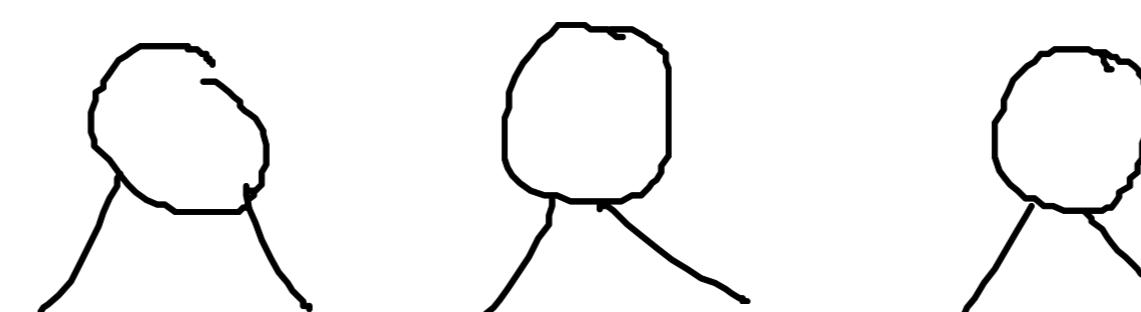
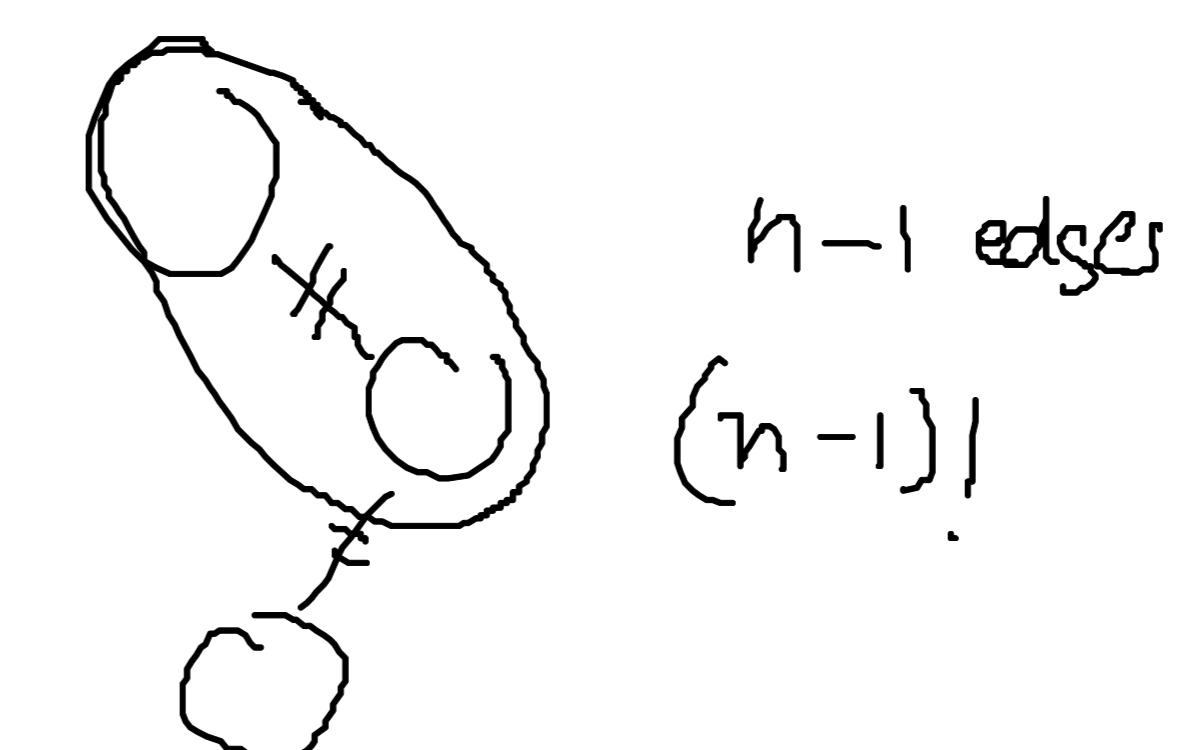
N elements



$2n$ edges

$n-1$ edges choose

$$2^n C_{n-1} \times (n-1)!$$



$$2^N (2N-1)(2N-2)\dots(2N-N+1)$$

$$\approx 2^N \underline{1 \dots (N+1)^N}$$

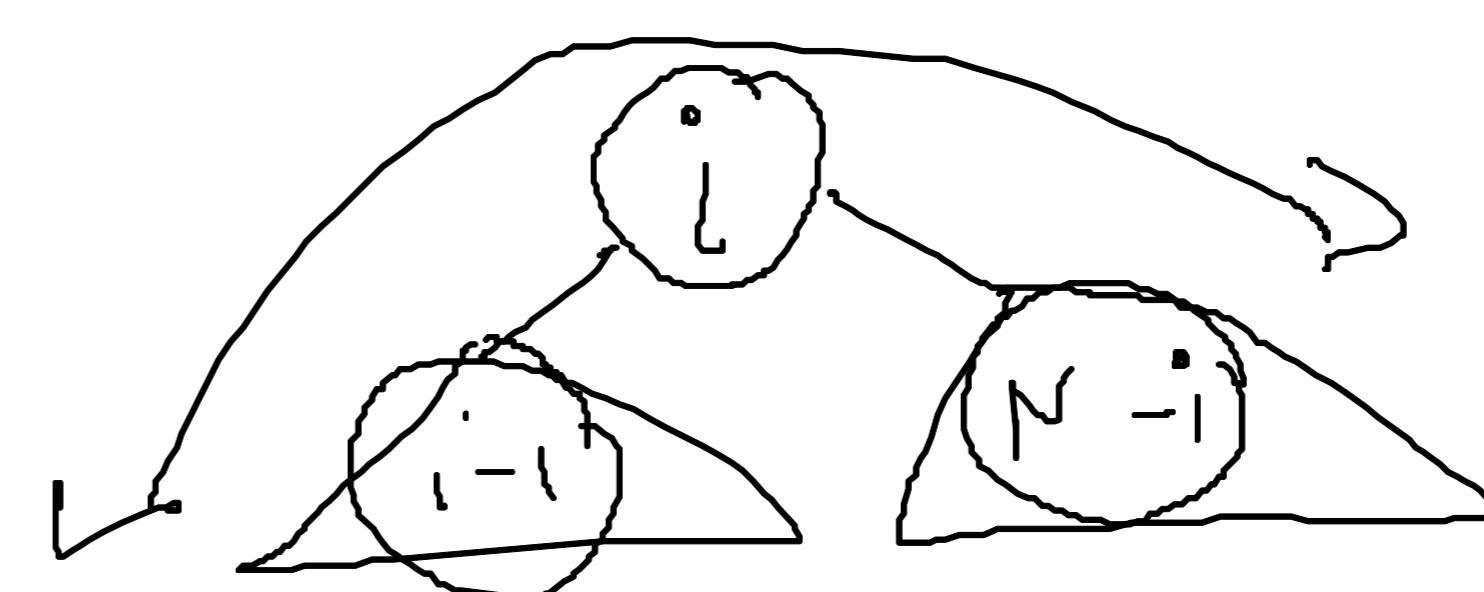
$$\frac{2^N}{N!}$$

$$\frac{2^N}{N!}$$

N elements

\textcircled{N} choices

1, 2, 3, 4, 5, ..., \textcircled{i} , ..., N



$$f(2) = \sum_{i=1}^2 f(i) * f(N-i)$$

$$= 1 * 1 + 1 * 1 = \boxed{2}$$

x ways and y ways

$$x_c_1 \cdot y_c_1$$

$$= xy$$

N=3

$$i=1 \quad f(0) * f(2) \\ 1 * 2 = \begin{array}{c} 1 \\ | \\ 2 \end{array}$$

$$i=2 \quad f(1) * f(1) = \begin{array}{c} 1 \\ | \\ 1 \end{array}$$

$$i=3 \quad f(2) * f(0) = \begin{array}{c} 1 \\ | \\ 2 \end{array}$$

$$\begin{array}{c} 1 \\ | \\ 2 \\ | \\ 3 \end{array}$$

$$\text{Catalan No} = \frac{2^N C_N}{(N+1)} = \frac{6 C_3}{4} = \frac{6 \cdot 5 \cdot 4}{1 \cdot 2 \cdot 3 \cdot 4} = \boxed{5}$$

$$\frac{2^N C_N}{N+1} * N!$$

Unlabelled Tree

$f(N) = \text{No of BST's using } N \text{ nodes}$

$$x = f(i) = \checkmark \text{ } i \text{ nodes}$$

$$y = f(N-i) = \checkmark \text{ } N-i \text{ nodes}$$

$$f(N) = \sum_{i=1}^N f(i-1) * f(N-i)$$

Recursion

[optimize: DP]

Catalan Series

$$T_n = \frac{2^n C_n}{n+1}$$

BST's

Applications
of
Catalan number

```

int f(n) {
    n == 0
    return 1;
}

```

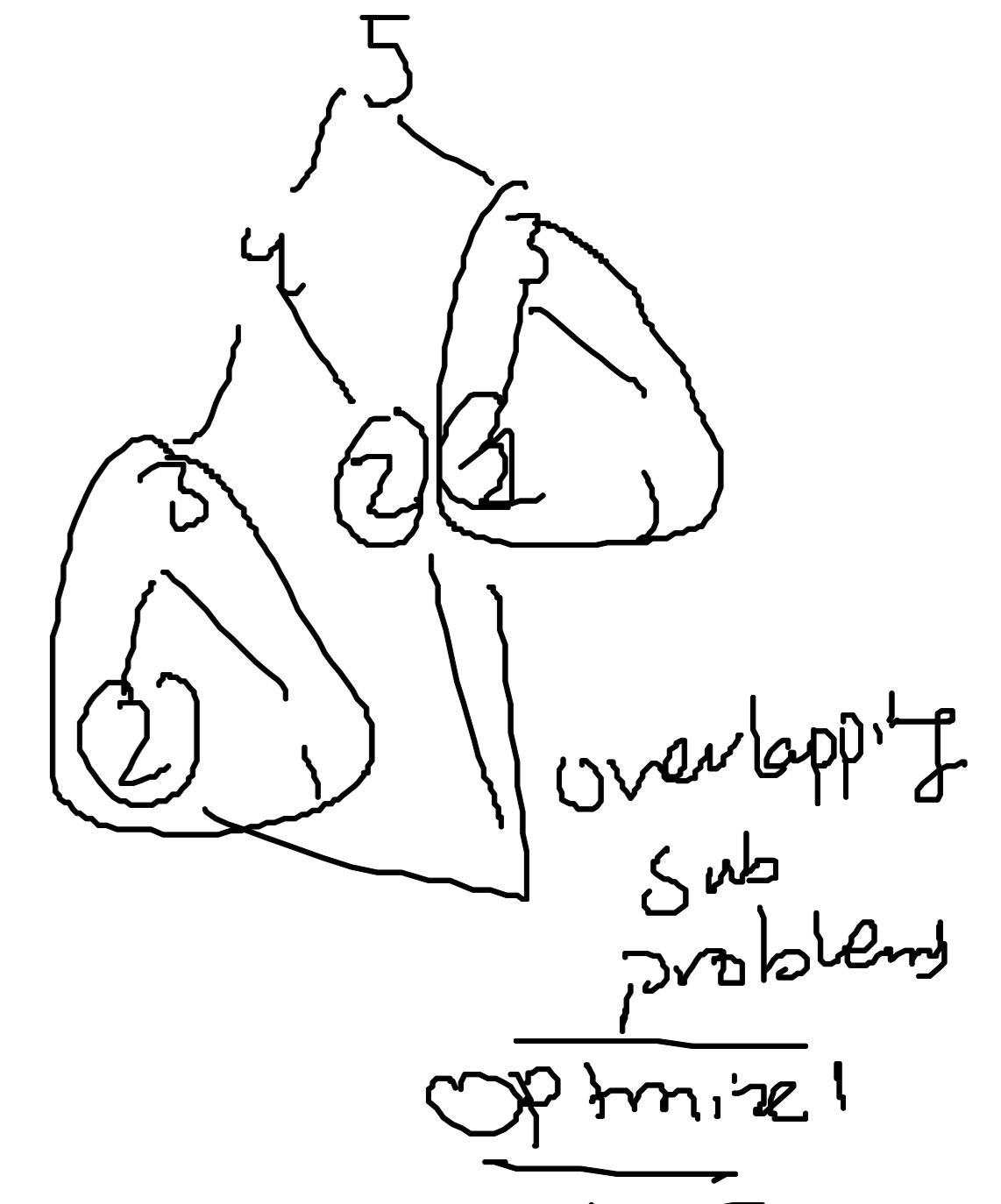
```

int ans = 0;
for(i=1; i <= n; i++) {
    ans += f(n-i) * f(i-1);
}
return ans;
}

```

3

$$a[3] = -1$$



Heaps (Priority Queue)

Insert $O(\log N)$

Get Min/Max $O(1)$

Remove Min/Max $O(\log N)$

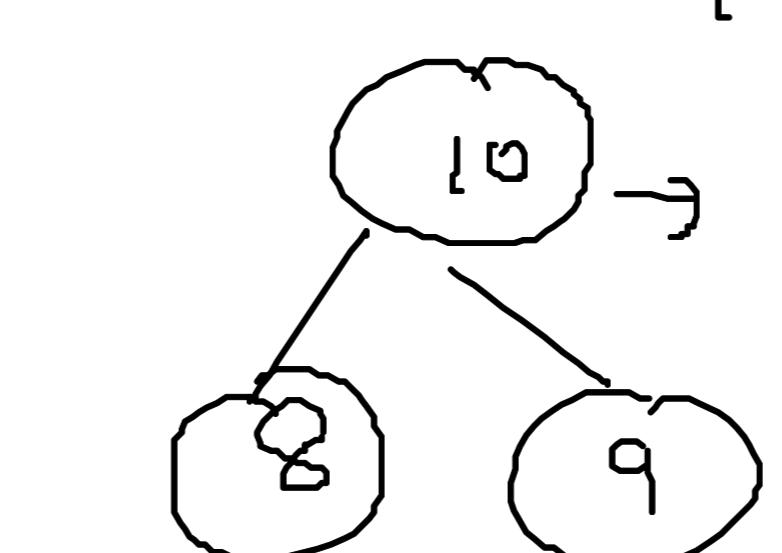
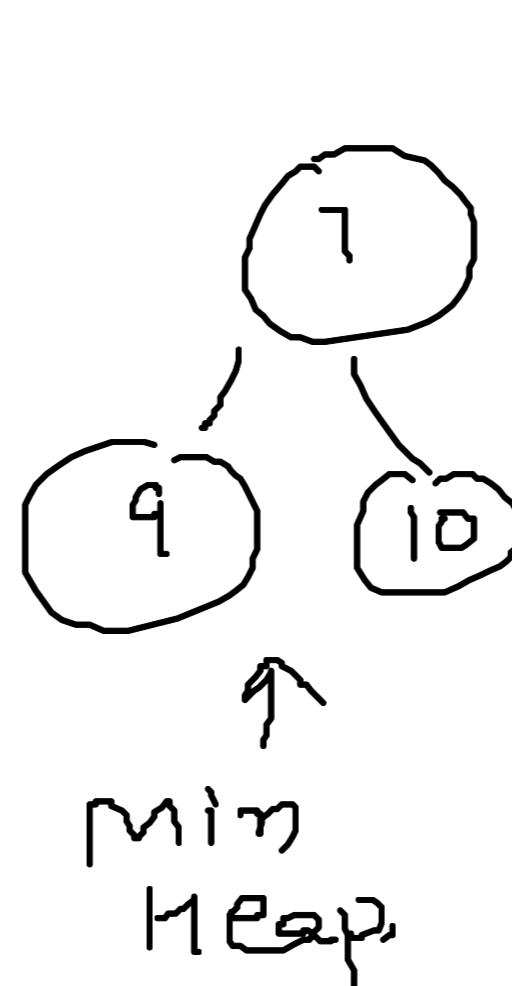
10 sort

N

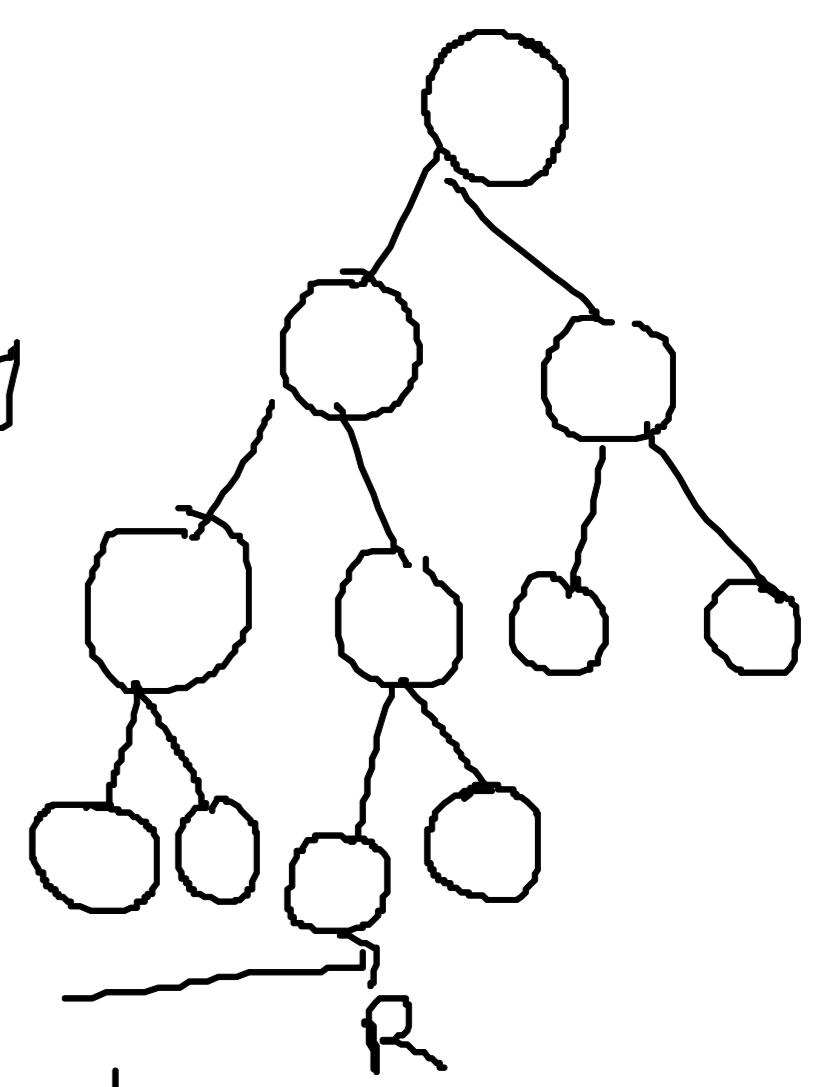
$N \log N$

$K \log N$

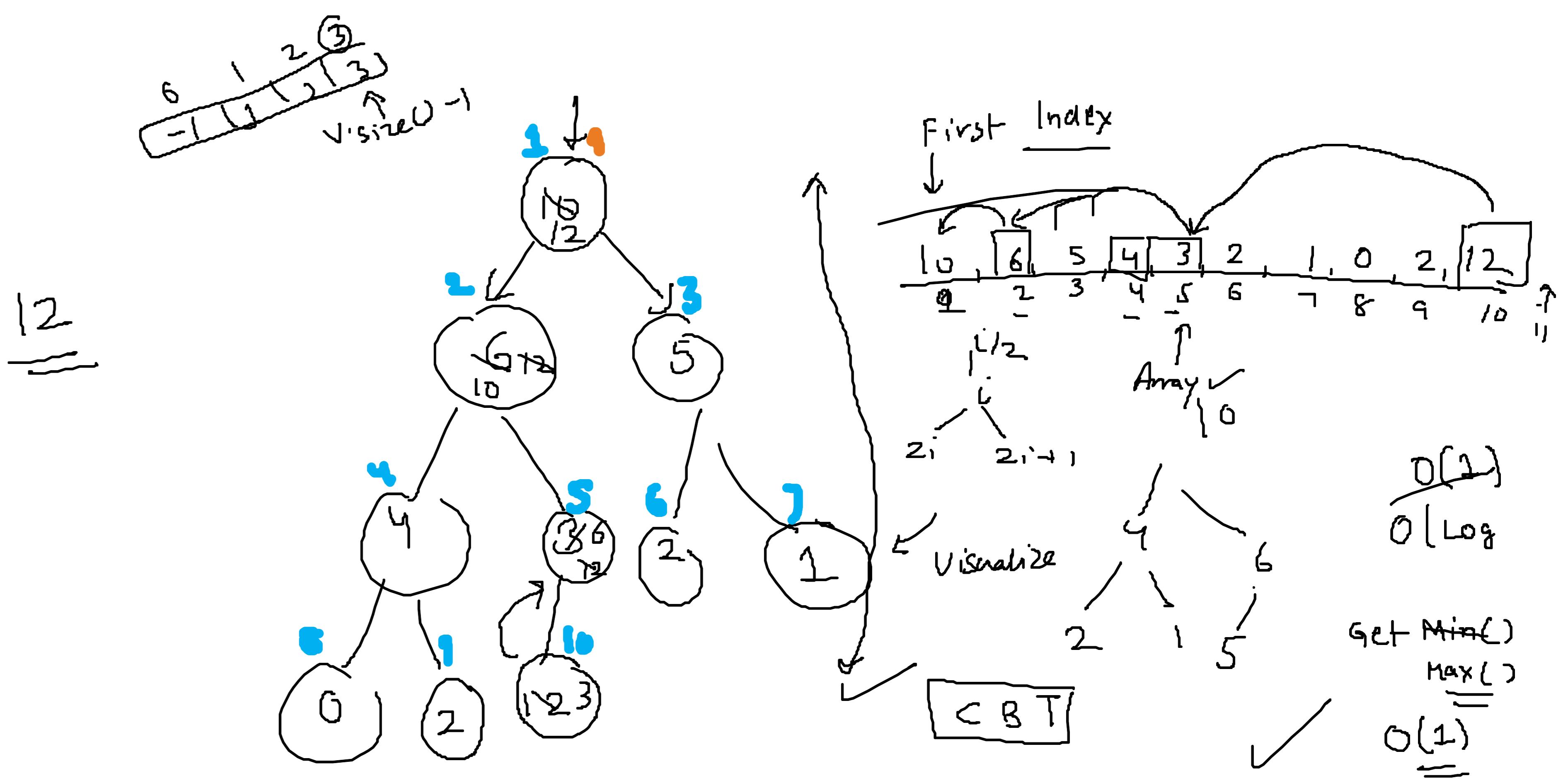
\rightarrow Insert
 \rightarrow Delete
 \rightarrow Search



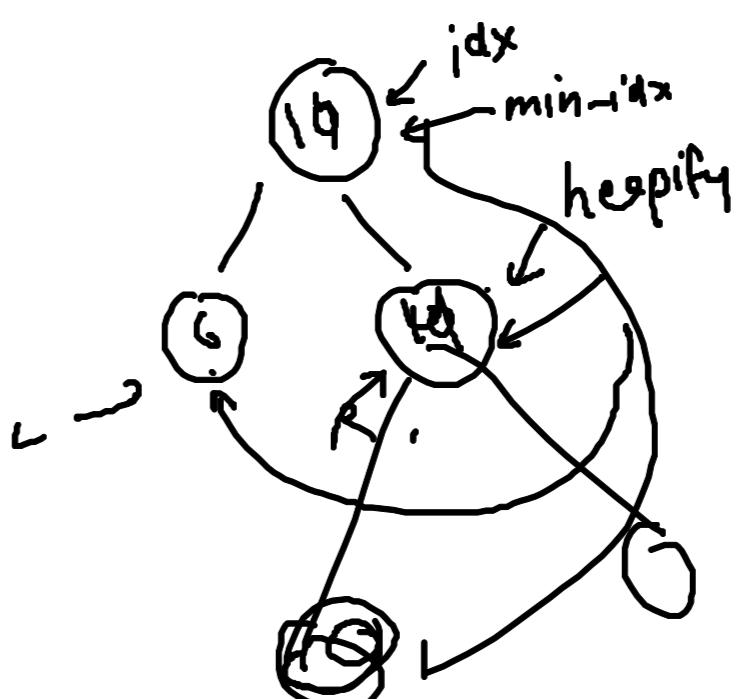
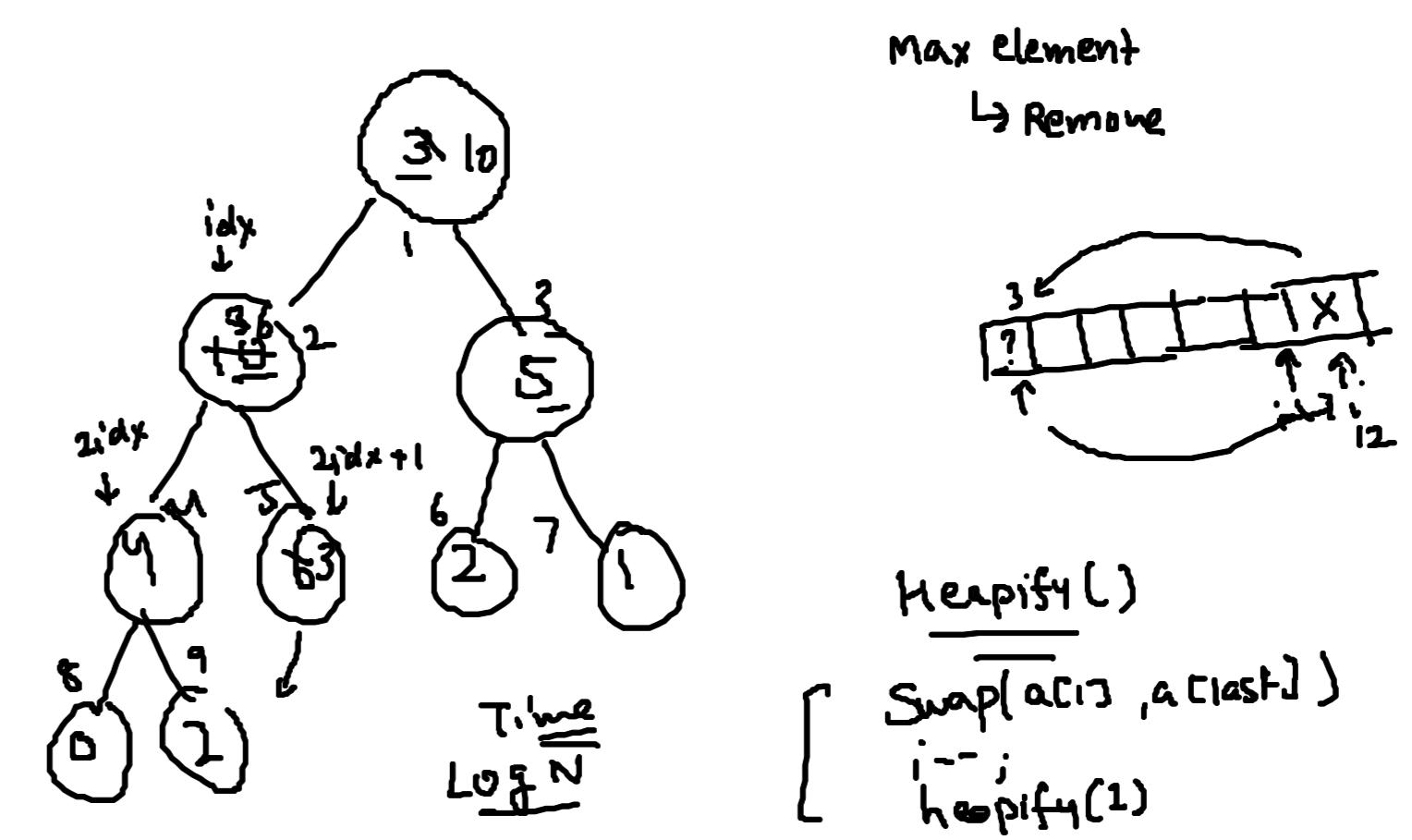
Heap ?
Parent \rightarrow Data Structure
Children \rightarrow Complete Binary Tree
for all nodes
All levels are completely filled except last level
which may be partially filled
but must be by Left to Right order.



\hookrightarrow Heaps Order Property



remove Min() / Max()



$\rightarrow \frac{N \log N}{\text{with}} + k \log N$

ARRAY \rightarrow head (Build) = $\frac{\sigma(N)}{2} + k \log N$

How to Build a heap from Array in $O(n)$ time.

