

$$Q:- T(n) = T\left(\frac{3n}{8}\right) + T\left(\frac{5n}{8}\right) + n^2 + C$$

$$\rightarrow T(n) = 3T\left(\frac{n}{8}\right) + 5T\left(\frac{n}{8}\right) + n^2 + C$$

$$= 8T\left(\frac{n}{8}\right) + n^2 + C$$

$$= 8\left(8T\left(\frac{n}{8^2}\right) + n^2\right) + C + n^2$$

$$= 8^2 T\left(\frac{n}{8^2}\right) + 8n^2 + C + n^2$$

after h iterations

$$= 8^h T\left(\frac{n}{8^h}\right) + 8h n^2 + C + n^2$$

$$\text{Let } \frac{n}{8^h} = 1$$

$$n = 8^h$$

$$h = \log_8 n$$

$$T(n) = 8^{\log_8 n} T(1) + 8 \log_8 n n^2 + C + n^2$$

$$O(n^2 \log_8 n)$$

$$Q:- T(n) = T\left(\frac{n}{9}\right) + T\left(\frac{8n}{9}\right) + \log_2 n + C$$

→ For simplicity we can assume

$$T\left(\frac{n}{9}\right) = T\left(\frac{8n}{9}\right) = T\left(\frac{n}{2}\right)$$

as both equal to  $n$  i.e.

$$\frac{n}{9} + \frac{8n}{9} = n$$

$$\text{and } \frac{n}{2} + \frac{n}{2} = n$$

and in complexity we need only have an approximate

$$T(n) = 2^h T\left(\frac{n}{2}\right) + \log_2 n + C$$

$$= 2^h \left( 2^{h-1} T\left(\frac{n}{2^2}\right) + \log_2 \frac{n}{2} \right) + \log_2 n + C$$

$$= 2^2 T\left(\frac{n}{2^2}\right) + 2 \log_2 \frac{n}{2} + \log_2 n + C$$

After  $h$  iteration

$$= 2^h T\left(\frac{n}{2^h}\right) + 2^{h-1} \log_2 \frac{n}{2^{h-1}} + 2^{h-2} \log_2 \frac{n}{2^{h-2}}$$

$$+ \dots + 2 \log_2 \frac{n}{2} + \log_2 n + C$$

$$= 2^h T\left(\frac{n}{2^h}\right) + 2^{h-1} \log_2 n + 2^{h-2} \log_2 n + \dots + 2 \log_2 n$$

$$= 2^{h-1} \log_2 2^{h-1} + 2^{h-2} \log_2 2^{h-2} + \dots + 2 \log_2 2 + \log_2 n + C$$

$$\text{Put } \frac{n}{2^h} = 1$$

$$\log_2 n = h$$

$$\begin{aligned} T(n) &= 2^h T\left(\frac{n}{2^h}\right) + (2^{h-1} + 2^{h-2} + 2^{h-3} + \dots + 2) \log_2 n \\ &\neq (2^{h-1}(h-1) + 2^{h-2}(h-2) + \dots + 2 \log 2) \\ &\quad + \log_2 n + C \end{aligned}$$

$$= 2^{\log_2 n} (1) + \frac{2^{h-1} - 2}{2 - 1} \log_2 n \neq (2^{h-1}(h-1)$$

$$+ (2^{h-2}(h-2) - \dots + 2 \log 2) + \log_2 n + C$$

$$= n + n \log n + \log n + C$$

$$\neq (2^{h-2}(h-2) + \dots + 2 \log 2)$$

we assume

$$2^{h-2}(h-2) + 2^{h-1}(h-1) - \dots - 2(\log 2)$$

is approximately

$$n (\log n)^2$$

log sum is approximately log n & we know the multiplication is  $\log n$  & log n (power)

So;

$$n (\log n)^2$$