

Timestamp Methods for Concurrency Control

- Timestamp is a unique identifier created by the DBMS to identify the relative starting time of a transaction.
- Typically, timestamp values are assigned in the order in which the transactions are submitted to the system. So, a timestamp can be thought of as the transaction start time. Therefore, time stamping is a method of concurrency control in which each transaction is assigned a transaction timestamp. A transaction timestamp is a monotonically increasing number, which is often based on the system clock. The transactions are managed so that they appear to run in a timestamp order. Timestamps can also be generated by incrementing a logical counter every time a new transaction starts.
- The schedule is equivalent to the particular Serial Order corresponding to the order of the Transaction timestamps. An algorithm must ensure that, for each item accessed by Conflicting Operations in the schedule, the order in which the item is accessed does not violate the ordering.
- To ensure this, use two Timestamp Values relating to each database item X.
 - $W_TS(X)$ is the largest timestamp of any transaction that executed $write(X)$ successfully.
 - $R_TS(X)$ is the largest timestamp of any transaction that executed $read(X)$ successfully.

Timestamp Ordering Protocol

- The Timestamp Ordering Protocol is used to order the transactions based on their Timestamps. The order of transaction is nothing but the ascending order of the transaction creation.
- The priority of the older transaction is higher that's why it executes first. To determine the timestamp of the transaction, this protocol uses system time or logical counter.
- The lock-based protocol is used to manage the order between conflicting pairs among transactions at the execution time. But Timestamp based protocols start working as soon as a transaction is created.
- Let's assume there are two transactions T1 and T2. Suppose the transaction T1 has entered the system at 007 times and transaction T2 has entered the system at 009 times. T1 has the higher priority, so it executes first as it is entered the system first.
- The timestamp ordering protocol also maintains the timestamp of last 'read' and 'write' operation on a data.

Basic Timestamp ordering protocol works as follows:

1. Check the following condition whenever a transaction T_i issues a **Read (X)** operation:
 - If $W_TS(X) > TS(T_i)$ then the operation is rejected.
 - If $W_TS(X) \leq TS(T_i)$ then the operation is executed.
 - Timestamps of all the data items are updated.

2. Check the following condition whenever a transaction T_i issues a **Write(X)** operation:

- If $TS(T_i) < R_TS(X)$ then the operation is rejected.
- If $TS(T_i) < W_TS(X)$ then the operation is rejected and T_i is rolled back otherwise the operation is executed.

Where,

TS(Ti) denotes the timestamp of the transaction T_i .

R_TS(X) denotes the Read time-stamp of data-item X .

W_TS(X) denotes the Write time-stamp of data-item X .

Advantages and Disadvantages of TO protocol:

- TO protocol ensures serializability since the precedence graph is as follows:



Image: Precedence Graph for TS ordering

- TS protocol ensures freedom from deadlock that means no transaction ever waits.
- But the schedule may not be recoverable and may not even be cascade- free.