

# VISVESVARAYA TECHNOLOGICAL UNIVERSITY

Jnana Sangama, Belagavi-590018, Karnataka



Project Report  
on

**“A\* Maze Solver”**

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# VISVESVARAYA TECHNOLOGICAL UNIVERSITY

"JnanaSangama", Belagavi-590018, Karnataka

## BANGALORE INSTITUTE OF TECHNOLOGY

Bengaluru-560 004



DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING

### *Certificate*

Certified that the **Mini** Project work entitled "**A\* Maze Solver**" carried out by

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of V semester, Computer Science and Engineering branch as partial fulfillment of the course **Artificial Intelligence (17CS562)** prescribed by **Visvesvaraya Technological University, Belgaum** during the academic year 2019-20. It is certified that all corrections/suggestions indicated for Internal Assessment have been incorporated in the report.

The **Mini** Project report has been approved as it satisfies the academic requirements in respect of project work in Artificial Intelligence.

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# CHAPTER 1

## INTRODUCTION

### 1.1 Introduction

In general, a maze has a complicated path like a labyrinth. To solve a maze is a game or puzzle which connects a path between a start and a goal. There are a lot of studies about shortest path algorithms.

Reaching a destination via a shortest route is a daily activity we all do. A\* is a generic search algorithm used to find solutions for several problems, maze solving basically to be one of them which is the process of plotting an efficiently traversable path between points, called nodes. It is one of the most successful search algorithms. It is an informed search algorithm, as it uses information about path cost and also heuristics to find the solution.

A\* search algorithm outperforms Dijkstra's to solve a maze in terms of search time and is based on the Dijkstra's method. The difference is to use a cost function of searching paths in the A\* search algorithm to solve maze more quickly. This algorithm brings together feature of uniform-cost search and heuristic search. When a location is examined, the A\* algorithm is completed when that location is the goal. In any other case, it helps make note of all that location neighbors for additional exploration. A\* could possibly be the most popular pathfinding algorithm in game AI. The time complexity of this algorithm is depending on heuristic use. A\* achieves *optimality* and *completeness*, two important property of search algorithms

### 1.2 Motivation

In the business world today, not a day passes without advancement in Artificial Intelligence (AI) sector as it has become mainstream. Industry experts are of the opinion that Artificial Intelligence is closely connected to popular culture thus driving the people to have improbable expectations and unrealistic fears about how it will alter the life and workplace in general.

AI is spreading throughout various industries. Some of them include education, healthcare, banking, finance, and many more. Industries that has adopted AI on a bigger level are more from the service industry. Even industries like web design and development being impacted by AI in a huge manner.

## CHAPTER 2

# PROBLEM STATEMENT

### 2.1 Problem Statement

Consider the problem of finding a route across the maze of ROWxCOL grid. The start position is (x,y) and the end position is (p,q). Movement is allowed by one square in any of the four directions. The standard movement cost is 1. There are barriers that occupy certain positions of the grid.

A route with the lowest cost should be found using the A\* search algorithm (there are multiple optimal solutions with the same total cost).

Print the optimal route in text format, as well as the total cost of the route.

Optionally, draw the optimal route and the barrier positions.

### 2.2 Objectives

- The game environment will be handled as a graph. This involves breaking the map into different points or locations, which are usually called nodes. These nodes are used to record the progress of the search.
- In addition to holding the map location, each node has three other attributes that are fitness, goal, and heuristic commonly known as f, g, and h. Different values will be assigned to paths between the nodes.
- Typically, these values would represent the distances between the nodes. The cost between nodes doesn't have to be distance. The cost could be time, if you wanted to find the path that takes the shortest amount of time to traverse. A\* using two lists (open list and closed list). The open list contains all the nodes in the map that have not been totally explored yet. The closed list contains of all the nodes that have been totally explored. In addition to the standard open/closed lists, marker arrays can be used to find out whether a state is in the open or closed list.

## CHAPTER 3

# SYSTEM REQUIREMENTS

### 3.1 Hardware requirements

PC with Windows/Linux OS/MacOS.

### 3.2 Software Requirements

Windows/Linux OS/MacOS with a C++ IDE.

## CHAPTER 4

# ARCHITECTURE

Now to understand how A\* works, first we need to understand few terminologies:

- **Node** (also called **State**) — All potential position or stops with unique identification.
- **Transition** — The act of moving between states or nodes.
- **Starting Node** — Where to start searching from.
- **Goal Node** — The target to stop searching.
- **Search Space** — A collection of nodes, like all board positions of a board game.
- **Cost** — Numerical value (say distance, time taken or financial expense) for the path from a node to another node.

**$g(n)$**  — this represents the *exact cost* of the path from the **starting node** to any node  **$n$** .

- **$h(n)$**  — this represents the heuristic *estimated cost* from node  **$n$**  to the goal node.
- **$f(n)$**  — lowest cost in the neighboring node  **$n$** .

Each time A\* enters a node, it calculates the cost  $f(n)$  of all of the neighboring nodes, and then enters the node with the lowest value of  $f(n)$ . It is calculated with the formula:  $f(n)=g(n)+h(n)$ . The purpose of  $g(n)$ ,  $h(n)$ , and  $f(n)$  is to quantify how promising a path is.

## 4.1 Algorithm

Initialize open and closed lists

Make the start vertex current

Calculate heuristic distance of start vertex to destination (h)

Calculate f value for start vertex ( $f = g + h$ , where  $g = 0$ )

WHILE current vertex is not the destination

FOR each vertex adjacent to current

IF vertex not in closed list and not in open list THEN

Add vertex to open list

Calculate distance from start (g)

Calculate heuristic distance to destination (h)

Calculate f value ( $f = g + h$ )

IF new f value < existing f value or there is no existing f value THEN

Update f value

Set parent to be the current vertex

END IF

END IF

NEXT adjacent vertex

Add current vertex to closed list

Remove vertex with lowest f value from open list and make it current

END WHILE

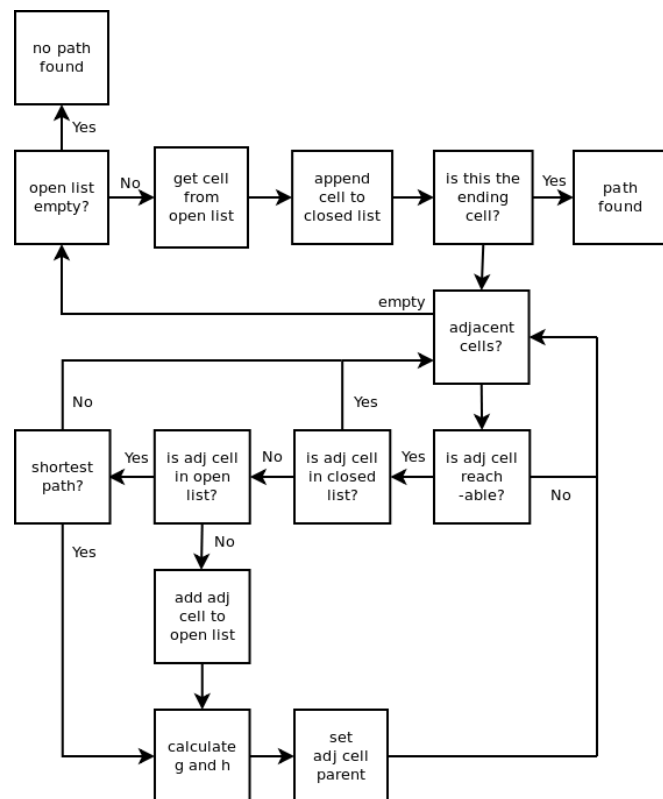


Figure 4.1 – Flowchart of the A\* Algorithm

## CHAPTER 5

### MODULE DESCRIPTIONS

#### 5.1 Module Description

struct cell - A structure is created to hold the necessary parameters  $f(x)$ ,  $g(x)$  and  $h(x)$ .

bool isValid(int row, int col) - A Utility Function to check whether given cell (row, col) is a valid cell or not.

bool isDestination(int row, int col, Pair dest) -  
A Utility Function to check whether destination cell has been reached or not.

double calculateHValue(int row, int col, Pair dest) - A Utility Function to calculate the 'h' heuristics when prompted.

void tracePath(cell cellDetails[][COL], Pair dest) - A Utility Function to trace the path from the source to destination after successfully tracing the shortest path.

void aStarSearch(int grid[][COL], Pair src, Pair dest) - A Function to find the shortest path between a given source cell to a destination cell according to A\* Search Algorithm. Initializes the parameters of the starting node and creates an open list and closed list. Then Searches through them while removing and adding from these lists as required.

int main () - Driver program to test above function. Description of the Grid is given by 1 for the cells that are not blocked and 0 for the cells that are blocked.

In the main driver program the Source and Destination are declared.



## 5.2 Example

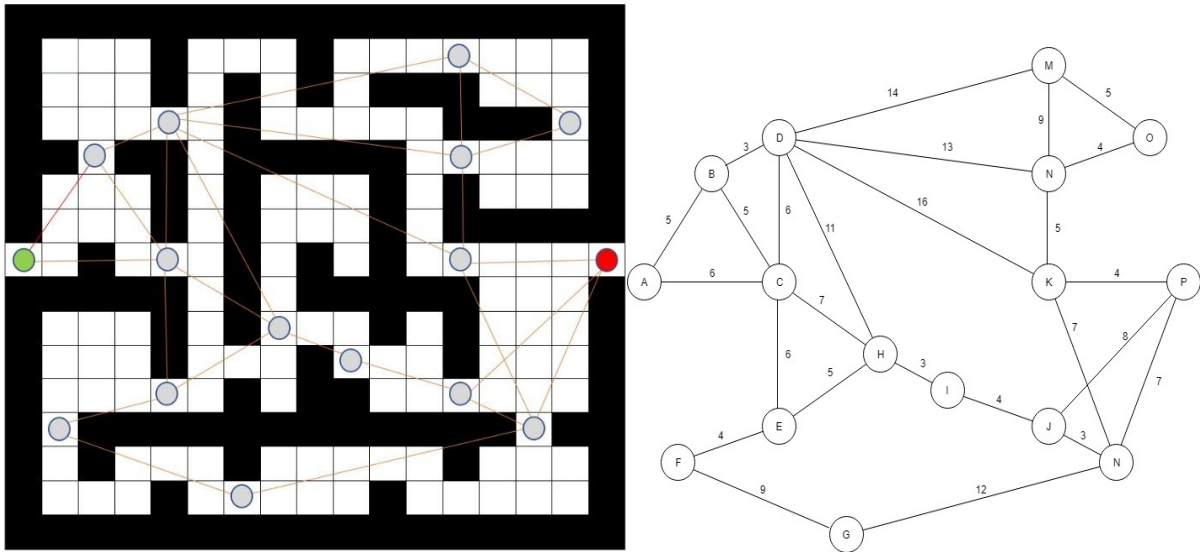


Figure 5.1 – Maze taken as an example and Maze represented as a Graph

Table 5.1 – Values of  $g(x)$ ,  $h(x)$  and  $f(x)$  with their parent nodes

Node	Distance from A (g)	Heuristic Distance (h)	$f = g + h$	Previous Node
A	0	16	16	
B	5	17	22	A
C	5	13	18	A
D	8	16	24	B
E	12	16	28	C
F		20		
G		17		
H	13	11	24	C
I	16	10	26	H
J	20	8	28	I
K	24	4	28	D
L	21	7	28	D
M	22	10	32	D
N	23	7	30	J
O		5		
P	28	0	28	J

## CHAPTER 6

### IMPLEMENTATION DETAILS

#### 6.1 Source Code

```
// A C++ Program to implement A* Search Algorithm

#include<bits/stdc++.h>

using namespace std;

#define ROW 16

#define COL 17

// Creating a shortcut for int, int pair type
typedef pair<int, int> Pair;

// Creating a shortcut for pair<int, pair<int, int>> type
typedef pair<double, pair<int, int> > pPair;

// A structure to hold the necessary parameters
struct cell{

    // Row and Column index of its parent

    // Note that 0 <= i <= ROW-1 & 0 <= j <= COL-1

    int parent_i, parent_j;

    // f = g + h

    double f, g, h;

};

// A Utility Function to check whether given cell (row, col) is a valid cell or not.
bool isValid(int row, int col){

    // Returns true if row number and column number is in range

    return (row >= 0) && (row < ROW) &&

        (col >= 0) && (col < COL);

}
```

```

// A Utility Function to check whether the given cell is blocked or not
bool isUnBlocked(int grid[][COL], int row, int col){

    // Returns true if the cell is not blocked else false

    if (grid[row][col] == 1)

        return (true);

    else

        return (false);

}

// A Utility Function to check whether destination cell has been reached or not
bool isDestination(int row, int col, Pair dest){

    if (row == dest.first && col == dest.second)

        return (true);

    else

        return (false);

}

// A Utility Function to calculate the 'h' heuristics.
double calculateHValue(int row, int col, Pair dest){

    // Return using the distance formula

    return(((double)sqrt((row-dest.first)*(row-dest.first)+(col-dest.second)*(col-dest.second))));

}

// A Utility Function to trace the path from the source to destination
void tracePath(cell cellDetails[][COL], Pair dest){

    printf ("\nThe Path is ");

    int row = dest.first;

    int col = dest.second;

    stack<Pair> Path;

```

```

while (!(cellDetails[row][col].parent_i == row
        && cellDetails[row][col].parent_j == col )){

    Path.push (make_pair (row, col));

    int temp_row = cellDetails[row][col].parent_i;
    int temp_col = cellDetails[row][col].parent_j;
    row = temp_row;
    col = temp_col;

}

Path.push (make_pair (row, col));

while (!Path.empty()){

    pair<int,int> p = Path.top();

    Path.pop();

    printf("-> (%d,%d) ",p.first,p.second);

}

return;

}

// A Function to find the shortest path between a given source cell to a destination cell
//according to A* Search Algorithm

void aStarSearch(int grid[][COL], Pair src, Pair dest){

    // If the source is out of range

    if (isValid (src.first, src.second) == false){

        printf ("Source is invalid\n");

        return;

    }

    // If the destination is out of range

    if (isValid (dest.first, dest.second) == false){

        printf ("Destination is invalid\n");

```

```

        return;
    }

    // Either the source or the destination is blocked
    if (isUnBlocked(grid, src.first, src.second) == false ||
        isUnBlocked(grid, dest.first, dest.second) == false){
        printf ("Source or the destination is blocked\n");
        return;
    }

    // If the destination cell is the same as source cell
    if (isDestination(src.first, src.second, dest) == true){
        printf ("We are already at the destination\n");
        return;
    }

    // Create a closed list and initialise it to false which means that no cell has been included
    //yet. This closed list is implemented as a boolean 2D array
    bool closedList[ROW][COL];

    memset(closedList, false, sizeof (closedList));

    // Declare a 2D array of structure to hold the details of that cell
    cell cellDetails[ROW][COL];

    int i, j;
    for (i=0; i<ROW; i++){
        for (j=0; j<COL; j++){
            cellDetails[i][j].f = FLT_MAX;
            cellDetails[i][j].g = FLT_MAX;
            cellDetails[i][j].h = FLT_MAX;
            cellDetails[i][j].parent_i = -1;
            cellDetails[i][j].parent_j = -1;
        }
    }

```

```

    }
}

// Initialising the parameters of the starting node

i = src.first, j = src.second;

cellDetails[i][j].f = 0.0;
cellDetails[i][j].g = 0.0;
cellDetails[i][j].h = 0.0;
cellDetails[i][j].parent_i = i;
cellDetails[i][j].parent_j = j;

/* Create an open list having information as- <f, <i, j>>
where f = g + h, and i, j are the row and column index of that cell
Note that 0 <= i <= ROW-1 & 0 <= j <= COL-1
This open list is implemented as a set of pair of pair.*/
set<pPair> openList;

// Put the starting cell on the open list and set its 'f' as 0
openList.insert(make_pair (0.0, make_pair (i, j)));

// We set this boolean value as false as initially the destination is not reached.
bool foundDest = false;

while (!openList.empty()){
    pPair p = *openList.begin();

    // Remove this vertex from the open list
    openList.erase(openList.begin());

    // Add this vertex to the closed list

    i = p.second.first;
    j = p.second.second;

```

```

        closedList[i][j] = true;

/*Generating all the 4 successor of this cell

        N

        W--Cell--E

        S

Cell-->Popped Cell (i, j)

N --> North    (i-1, j)

S --> South    (i+1, j)

E --> East     (i, j+1)

W --> West     (i, j-1) */

// To store the 'g', 'h' and 'f' of the 4 successors

double gNew, hNew, fNew;

//----- 1st Successor (North) -----

// Only process this cell if this is a valid one
if (isValid(i-1, j) == true) {

    // If the destination cell is the same as the current successor

    if (isDestination(i-1, j, dest) == true) {

        // Set the Parent of the destination cell

        cellDetails[i-1][j].parent_i = i;

        cellDetails[i-1][j].parent_j = j;

        printf ("The destination cell is found\n");

        tracePath (cellDetails, dest);

        foundDest = true;

        return;

    }

    // If the successor is already on the closed list or if it is blocked, then
    //ignore it. Else do the following

```

```

else if (closedList[i-1][j] == false &&isUnBlocked(grid, i-1, j) == true){

    gNew = cellDetails[i][j].g + 1.0;

    hNew = calculateHValue (i-1, j, dest);

    fNew = gNew + hNew;

    //If it isn't on the open list, add it to the open list. Make the
    //current square the parent of this square. Record f, g, and h costs
    //of the square cell OR If it is on the open list already, check to
    //see if this path to that square is better, using 'f' cost as measure.

    if(cellDetails[i-1][j].f==FLT_MAX||cellDetails[i-1][j].f>fNew){

        openList.insert( make_pair(fNew,make_pair(i-1, j)));

        // Update the details of this cell

        cellDetails[i-1][j].f = fNew;

        cellDetails[i-1][j].g = gNew;

        cellDetails[i-1][j].h = hNew;

        cellDetails[i-1][j].parent_i = i;

        cellDetails[i-1][j].parent_j = j;

    }

}

}

//----- 2nd Successor (South) -----

// Only process this cell if this is a valid one

if (isValid(i+1, j) == true) {

    // If the destination cell is the same as the current successor

    if (isDestination(i+1, j, dest) == true) {

        // Set the Parent of the destination cell

        cellDetails[i+1][j].parent_i = i;

        cellDetails[i+1][j].parent_j = j;

```



```

        printf("The destination cell is found\n");

        tracePath(cellDetails, dest);

        foundDest = true;

        return;
    }

    // If the successor is already on the closed list or if it is blocked, then
    //ignore it. Else do the following
    else if (closedList[i+1][j] == false && isUnBlocked(grid, i+1, j) == true){

        gNew = cellDetails[i][j].g + 1.0;

        hNew = calculateHValue(i+1, j, dest);

        fNew = gNew + hNew;

        //If it isn't on the open list, add it to the open list. Make the
        //current square the parent of this square. Record f, g, and h costs
        //of the square cell OR If it is on the open list already, check to
        //see if this path to that square is better, using 'f' cost as measure

        if(cellDetails[i+1][j].f == FLT_MAX || cellDetails[i+1][j].f > fNew) {

            openList.insert( make_pair (fNew, make_pair (i+1, j)));

            // Update the details of this cell

            cellDetails[i+1][j].f = fNew;

            cellDetails[i+1][j].g = gNew;

            cellDetails[i+1][j].h = hNew;

            cellDetails[i+1][j].parent_i = i;

            cellDetails[i+1][j].parent_j = j;

        }

    }

}

//----- 3rd Successor (East) -----

// Only process this cell if this is a valid one

```

```

if (isValid (i, j+1) == true) {

    // If the destination cell is the same as the current successor

    if (isDestination(i, j+1, dest) == true) {

        // Set the Parent of the destination cell

        cellDetails[i][j+1].parent_i = i;

        cellDetails[i][j+1].parent_j = j;

        printf("The destination cell is found\n");

        tracePath(cellDetails, dest);

        foundDest = true;

        return;

    }

    // If the successor is already on the closed list or if it is blocked, then
    //ignore it. Else do the following

    else if (closedList[i][j+1]==false&&isUnBlocked(grid, i, j+1) == true){

        gNew = cellDetails[i][j].g + 1.0;

        hNew = calculateHValue (i, j+1, dest);

        fNew = gNew + hNew;

        //If it isn't on the open list, add it to the open list. Make the
        //current square the parent of this square. Record f, g, and h costs
        //of the square cell OR If it is on the open list already, check to
        //see if this path to that square is better, using 'f' cost as measure.

        if(cellDetails[i][j+1].f==FLT_MAX||cellDetails[i][j+1].f>fNew){

            openList.insert( make_pair(fNew,make_pair(i,j+1)));

            // Update the details of this cell

            cellDetails[i][j+1].f = fNew;

            cellDetails[i][j+1].g = gNew;

            cellDetails[i][j+1].h = hNew;

            cellDetails[i][j+1].parent_i = i;

```

```

        cellDetails[i][j+1].parent_j = j;
    }
}

//----- 4th Successor (West) -----
// Only process this cell if this is a valid one
if (isValid(i, j-1) == true){
    // If the destination cell is the same as the current successor
    if (isDestination(i, j-1, dest) == true){
        // Set the Parent of the destination cell
        cellDetails[i][j-1].parent_i = i;
        cellDetails[i][j-1].parent_j = j;
        printf("The destination cell is found\n");
        tracePath(cellDetails, dest);
        foundDest = true;
        return;
    }
    // If the successor is already on the closed list or if it is blocked, then
    //ignore it. Else do the following
    else if (closedList[i][j-1] == false && isUnBlocked(grid, i, j-1) == true)
    {
        gNew = cellDetails[i][j].g + 1.0;
        hNew = calculateHValue(i, j-1, dest);
        fNew = gNew + hNew;

        //If it isn't on the open list, add it to open list. Make the current
        //square the parent of this square. Record the f, g, and h costs of
        //the square cell OR If it is on the open list already, check to see
        //if this path to that square is better, using 'f' cost as the measure.
    }
}

```

```

        if(cellDetails[i][j-1].f==FLT_MAX||cellDetails[i][j-1].f>fNew){
            openList.insert( make_pair (fNew,make_pair(i,j-1)));
            // Update the details of this cell
            cellDetails[i][j-1].f = fNew;
            cellDetails[i][j-1].g = gNew;
            cellDetails[i][j-1].h = hNew;
            cellDetails[i][j-1].parent_i = i;
            cellDetails[i][j-1].parent_j = j;
        }
    }
}

// When destination cell is not found and the open list is empty, we conclude that we
//failed to reach the destination cell. This may happen when there is no way to reach
//destination cell (due to blockages)

if (foundDest == false)

    printf("Failed to find the Destination Cell\n");

return;
}

int main(){

    // 1--> The cell is not blocked

    // 0--> The cell is blocked

    int grid[ROW][COL] ={

        {0,0,0,0,0,0,0,0,0,0,0,0,0,0,0},

        {0,1,1,1,0,1,1,1,0,1,1,1,1,1,0},

        {0,1,1,1,0,1,0,1,0,1,0,0,0,1,1,0},

        {0,1,1,1,1,1,0,1,1,1,1,0,0,0,1,0},
    }
}

```

```

        {0,0,1,0,0,1,0,0,0,0,0,1,1,1,1,0},
        {0,1,1,1,0,1,0,1,1,1,0,1,0,1,1,0},
        {0,1,1,1,0,1,0,1,1,1,0,1,0,0,0,0},
        {1,1,0,1,1,1,0,1,0,1,0,1,1,1,1,1},
        {0,0,0,0,0,1,0,1,0,0,0,0,0,1,1,0},
        {0,1,1,1,0,1,0,1,1,1,0,1,0,1,1,0},
        {0,1,1,1,0,1,1,1,0,1,1,1,0,1,1,0},
        {0,1,1,1,1,1,0,1,0,0,1,1,1,1,1,0},
        {0,1,0,0,0,0,0,0,0,0,0,0,0,0,1,0},
        {0,1,0,1,1,1,0,1,1,1,1,1,0,1,1,0},
        {0,1,1,1,0,1,1,1,1,1,0,1,1,1,1,0},
        {0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0}
    };

    Pair src = make_pair(7, 0);
    Pair dest = make_pair(7, 16);
    aStarSearch(grid, src, dest);
    return(0);
}

```

## CHAPTER 7

### RESULTS

#### 7.1. Output

A custom maze is created and solved with the help of A\* algorithm and is implemented with the help of a C++ code.

The destination cell is found

The Path is -> (7,0) -> (7,1) -> (6,1) -> (6,2) -> (6,3) -> (7,3) -> (7,4) -> (7,5) -> (8,5) -> (9,5) -> (10,5) -> (10,6) -> (10,7) -> (9,7) -> (9,8) -> (9,9) -> (10,9) -> (10,10) -> (10,11) -> (11,11) -> (11,12) -> (11,13) -> (10,13) -> (9,13) -> (9,14) -> (8,14) -> (8,15) -> (7,15) -> (7,16)

## CHAPTER 8

### APPLICATIONS

- Games - Games with bots use a modified version of the A\* algorithm to find their path to the player/object.
- Town Planning - Used to optimize the placement of essential utilities from a home unit.
- Space Exploration – Used in rovers to calculate the most efficient path around a crater/rock.
- flight path optimization – Optimum use of airspace and fuel by a whole fleet is maintained by planning the most effective paths for flights to take.
- Rescue Robots - Robots are widely used in many areas, where Human life cannot be compromised. In Rescue operations, the robots need to search for the trapped humans and find the efficient path between the robot and the trapped person. This path is planned using a Hybrid A\*
- NLP - Parsing using stochastic grammars.

## CHAPTER 9

# CONCLUSION AND FUTURE ENHANCEMENT

### 9.1 Conclusion

A-star (A\*) is a mighty algorithm in Artificial Intelligence with a wide range of usage. However, it is only as good as its heuristic function (which can be highly variable considering the nature of a problem). A\* is the most popular choice for pathfinding because it's reasonably flexible.

### 9.2 Future Enhancement

A\* has a low efficiency in the path selection process of very complex networks, if the search space is large and the entity is irregular, this algorithm is often unfavorable to search the optimal path, expanding a lot of redundant nodes and consuming much time. The presence of local minima deviates original A\* from the best path.

A\* algorithm considers only the local costs rather than the global costs and results in some substandard results, we may say that it has erroneously moved in the wrong direction. Further examination of nodes, existing in the linked list, by generating their children or the subsequent nodes could provide greater information regarding the cost and may lead to the best strategy for query optimization.

A closed list need not be maintained and that a simple onList flag can be set for each node, thereby removing the need to search a closed list repeatedly.

The A-star algorithm and its variations like Iterative Deepening A-star (IDA-star), Jump Point Search (JPS), etc. are the most popular algorithm in the field of Path-finding in Maze and are enhanced versions of A\* that are optimized for their usage.

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