

## Assignment No. 2.



1. Print elements of an array (forward and backward)

```

int [] arr = new arr[5];
for (int i = 0; i < arr.length; i++) {
    s.o.p(arr[i]);
}
// forward direction

for (int j = arr.length - 1; j >= 0; j--) {
    s.o.p(arr[j]);
}
// backward direction
    
```

e.g. Ex(1) arr = [1, 2, 3, 4] // length = 4

- ①  $i = 0, 0 < 4 \rightarrow arr[0] = 1$
- ②  $i = 1, 1 < 4 \rightarrow arr[1] = 2$
- ③  $i = 2, 2 < 4 \rightarrow arr[2] = 3$
- ④  $i = 3, 3 < 4 \rightarrow arr[3] = 4$
- ⑤  $i = 4, 4 < 4 \rightarrow break;$

- ①  $j = 4 - 1 = 3, 3 >= 0 \rightarrow arr[3] = 4$
- ②  $j = 4 - 2 = 2, 2 >= 0 \rightarrow arr[2] = 3$
- ③  $j = 4 - 3 = 1, 1 >= 0 \rightarrow arr[1] = 2$
- ④  $j = 4 - 4 = 0, 0 >= 0 \rightarrow arr[0] = 1$
- ⑤  $j = 4 - 5 = -1, -1 >= 0 \rightarrow break;$

\* // Using recursion

```

main () {
    fun(arr, 0);
}

void fun(int[] arr, int index) {
    if (index < arr.length) return;
    s.o.p(arr[index]);
    fun(arr, index + 1);
}
    
```

e.g. arr = [1, 2, 3, 4]

fun(arr, 0) → arr[0] = 1  
↳ fun(arr, 1) → arr[1] = 2  
↳ fun(arr, 2) → arr[2] = 3  
↳ fun(arr, 3) → arr[3] = 4  
↳ fun(arr, 4) → break;

\* // element in reverse

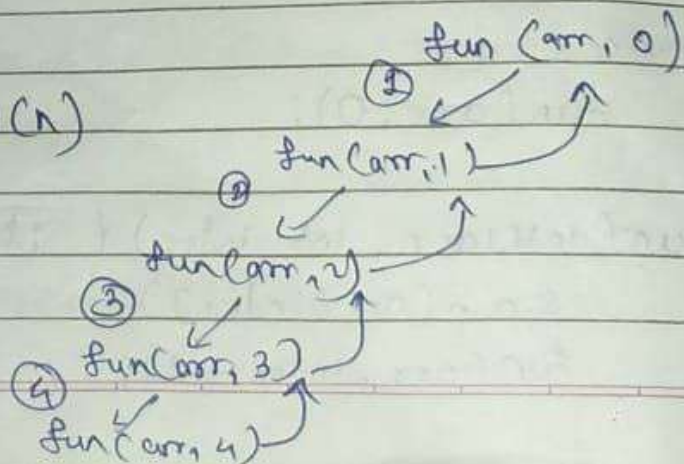
```
main() {  
    fun(arr, arr.length-1);  
}
```

```
void fun(arr, index) {  
    if (arr.length - index <= 0) return;  
    swap(arr[index]);  
    fun(arr, index-1);  
}
```

e.g. arr = [1, 2, 3, 4]

fun(arr, 3) → arr[3] = 4  
↳ fun(arr, 2) → arr[2] = 3  
↳ fun(arr, 1) → arr[1] = 2  
↳ fun(arr, 0) → arr[0] = 1  
↳ fun(arr, -1) → arr[-1] → break;

Time Complexity =  $O(n)$





2. Find the sum of digits of a number.

```
main() {  
    s.o.p(sum(num));  
}  
int sum(num) {  
    if (num <= 0) return 0;  
    num = num / 10; rem = num % 10;  
    return rem + sum(num / 10);  
}
```

\* eg.  $N = 1234$

①  $\text{sum}(1234) \rightarrow \text{num} \neq 0 \rightarrow 1234 \% 10 = 4 + \text{sum}(123)$

②  $\text{sum}(123) \rightarrow \text{num} \neq 0 \rightarrow 123 \% 10 + \text{sum}(12) = 3 + \text{sum}(12)$

③  $\text{sum}(12) \rightarrow \text{num} \neq 0 \rightarrow 12 \% 10 + \text{sum}(1) = 2 + \text{sum}(1)$

④  $\text{sum}(1) \rightarrow \text{num} \neq 0 \rightarrow 1 \% 10 + \text{sum}(0) = 1 + \text{sum}(0)$

⑤  $\text{sum}(0) \rightarrow \text{num} \leq 0 \rightarrow \text{return};$

$\text{s.o.p}(\text{sum}) = \underline{\underline{10}}$

Time complexity =  $O(\log n)$

3. Find the product of digits of a number.

```
main() {  
    s.o.p(product(num));  
}  
product(int num) {  
    if (product <= 1) return 1;  
    int rem = num % 10;  
    return (rem * product(num / 10));  
}
```

eg.  $N = 231$

① product(231)

↳  $231 > 1$

↳  $rem = 231 \% 10 = 1 \rightarrow rem * product(23)$

② product(23)

↳  $23 > 1$

↳  $rem = 23 \% 10 = 3 \rightarrow rem * product(2)$

③ product(2)

↳  $2 > 1$

↳  $rem = 2 \% 10 = 2 \rightarrow rem * product(0)$

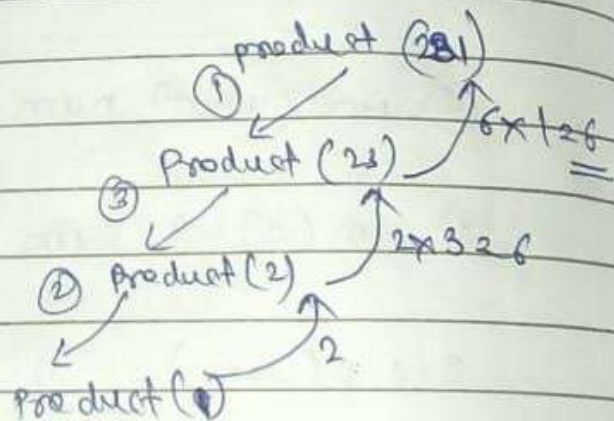
④ product(0)

↳  $0 \leq 1$

↳ return 1;

$$\therefore \text{product} = 1 * 3 * 2 = \underline{6}$$

$\therefore \text{Time Complexity} = O(\log n)$





4. Count number of digits in a number

```
main() {
```

```
    s.o.p (count(num));
```

```
}
```

```
count(int num) {
```

```
    if (num <= 0) return 0;
```

```
    return 1 + count(num/10);
```

```
}
```

eg. 98765

① count(98765)  $\rightarrow$  num  $> 0 \rightarrow 1 + \text{count}(9876)$

② count(9876)  $\rightarrow$  num  $> 0 \rightarrow 1 + 1 + \text{count}(987)$

③ count(987)  $\rightarrow$  num  $> 0 \rightarrow 2 + 1 + \text{count}(98)$

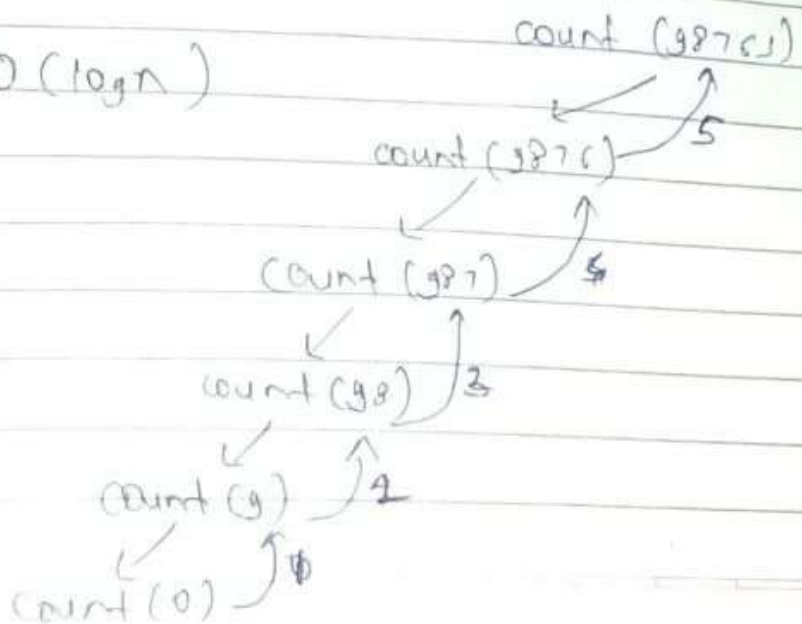
④ count(98)  $\rightarrow$  num  $> 0 \rightarrow 3 + 1 + \text{count}(9)$

⑤ count(9)  $\rightarrow$  num  $> 0 \rightarrow 4 + 1 + \text{count}(0)$

⑥ count(0)  $\rightarrow$  return 0;

$\therefore$  count = 5

Time complexity =  $O(\log n)$



5. Find the maximum element in an array.

```

for main() {
    solve(max(num, 0index)); // ([2, 5, 9, 1, 6], 0)
    // int max = 0;
    int max(int num, int index) {
        if (index == num.length) return 0;
        return Math.max(num[index], num[index+1, num]);
    }
}

```

eg. [2, 5, 9, 1, 6]

①  $\text{max}(\text{num}, 0) \rightarrow 0 \neq 5 \rightarrow \text{Math.max}(2, \text{max}(\text{num}, 1))$  <sup>⑨</sup>

②  $\text{max}(\text{num}, 1) \rightarrow 1 \neq 5 \rightarrow \text{Math.max}(5, \text{max}(\text{num}, 2))$  <sup>⑧</sup>

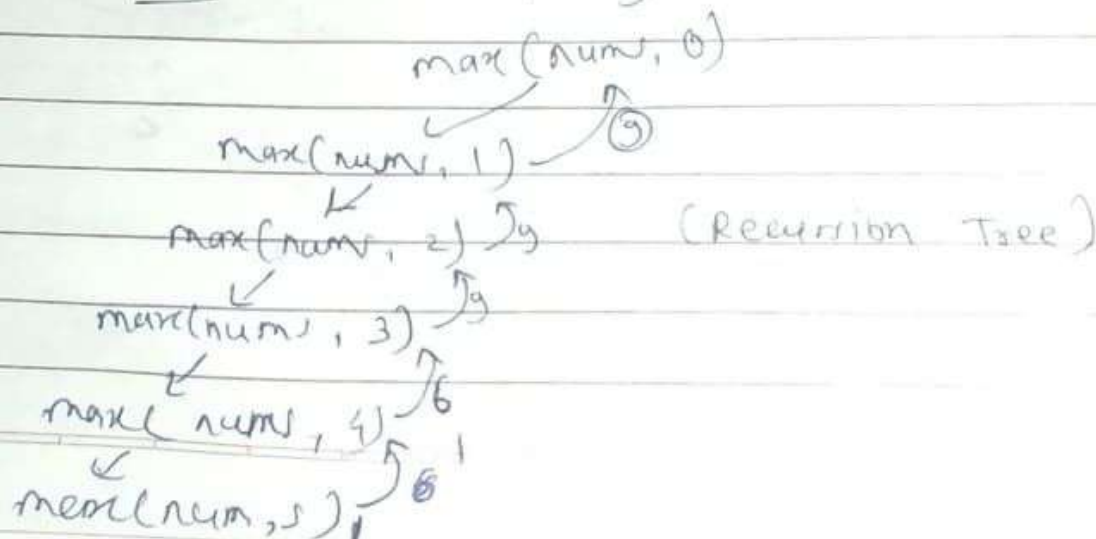
③  $\text{max}(\text{num}, 2) \rightarrow 2 \neq 5 \rightarrow \text{Math.max}(9, \text{max}(\text{num}, 3))$  <sup>⑦</sup>

④  $\text{max}(\text{num}, 3) \rightarrow 3 \neq 5 \rightarrow \text{Math.max}(1, \text{max}(\text{num}, 4))$  <sup>⑥</sup>

⑤  $\text{max}(\text{num}, 4) \rightarrow 4 \neq 5 \rightarrow \text{Math.max}(6, \text{max}(\text{num}, 5))$  <sup>⑤</sup>

⑥  $\text{max}(\text{num}, 5) \rightarrow 5 = 5 \rightarrow \text{return } 1$

$\therefore \underline{\underline{\text{max} = 9}}$ , Time complexity =  $O(n)$ ,





6. check if an array is sorted (strictly increasing)  
→ main() {  
    isSorted(nums, 1);  
}

boolean isSorted (int[] nums, index) {  
    if (nums[index-1] >= nums[index]) return false;  
    else if (index == nums.length) return true;  
    return isSorted(nums, index+1);  
}

① isSorted(nums = [1, 2, 3, 4], 1)  
    ↳ (nums[0] < nums[1])  
        ↳ isSorted(nums, 2)

② isSorted(nums, 2)  
    ↳ (nums[1] < nums[2])  
        ↳ isSorted(nums, 3)

③ isSorted(nums, 3)  
    ↳ (nums[2] < nums[3])  
        ↳ isSorted(nums, 4)

④ isSorted(nums, 4)  
    ↳ ~~nums~~ (index == 4 == nums.length);  
        ↳ true;

∴ true, array is sorted.

Time complexity =  $O(n)$

7. check if a number is prime

eg. 13  $\rightarrow$   $13 \times 1 = 13$   
 $1 \times 13 = 13$

```
int counter = 0;  
for (int i = 1; i <= sqrt(n); i++) {
```

```
    if (n % i == 0 && n % (n/i) != 0) {
```

```
        counter++;
```

```
    }
```

```
    if (counter > 1) {
```

```
        break;
```

```
    }
```

```
}
```

```
if (counter == 1) {
```

```
    s.o.p("Prime");
```

```
}
```

time complexity =  $O(\sqrt{n})$

eg. 36

36  $\rightarrow$

1  $\times$  36

2  $\times$  18

3  $\times$  12

4  $\times$  9

6  $\times$  6

```
for (i = 1; i <= 6; i++) {
```

①  $i = 1$

$\rightarrow 36 \% 1 \neq 0$  &  $36 \% 36 = 0$

$\rightarrow$  counter = 1;

②  $i = 2$

$\rightarrow 36 \% 2 = 0$  &  $36 \% 18 = 0$

$\rightarrow$  counter = 2;

break;

counter = 2, ~~is prime~~ Not Prime

eg. 29

$\rightarrow$

1  $\times$  29

```
for (i = 1; i <= 5; i++) {
```

①  $i = 1$

$\rightarrow 29 \% 1 \neq 0$  &  $29 \% 29 = 0$  ✓

$\rightarrow$  counter = 1;

$\rightarrow$  for,  $i = 2, 3, 4, 5$

$\rightarrow 29 \% i \neq 0$

② .

counter = 1, so, it is Prime



8. Find the first index of an element in an array.

```

main() {
    firstOcc(arr, 0);
    if (index == arr.length) return -1;
    key = 7;
    int firstOcc(arr, index) {
        if (index == arr.length) return -1;
        if (arr[index] == key)
            return index;
        return firstOcc(arr, index + 1);
    }
}

```

① firstOcc([4, 2, 7, 2], 0) {  
 ↳ arr[index] ≠ key  
 firstOcc(arr, 1) ← index

② firstOcc(arr, 1)  
 ↳ arr[1] ≠ key;  
 firstOcc(arr, 2) ← index

③ firstOcc(arr, 2)  
 ↳ arr[2] = key;  
 return 2; // index

∴ Time complexity =  $O(n)$ ;

```

graph TD
    A[firstOcc(arr, 0)] -- 2 --> B[firstOcc(arr, 1)]
    B -- 2 --> C[firstOcc(arr, 2)]

```

9. Find the last index of an element in an array.

```

main() {
    lastOcc (arr, arr.length-1);
}

int lastOcc (arr, index) {
    if (index <= -1) return -1;
    else if (arr[index] == key)
        return index;

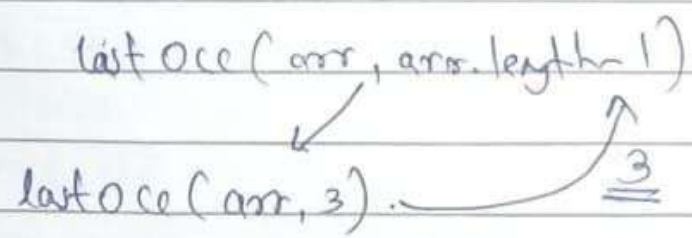
    return lastOcc (arr, index-1);
}
    
```

e.g. [4, 2, 7, 7, 9]

① lastOcc (arr, 4)  
 ↳ arr[4] ≠ 7  
 ↳ index = 4-1 = 3 // lastOcc (arr, 3)

② lastOcc (arr, 3)  
 ↳ arr[3] = 7  
 ↳ return 3;

∴ Time Complexity =  $O(n)$





10. Reverse a number using Recursion.

```
main() {  
    s.o.p (reverseNum(num));  
}  
int reverseNum (num) {  
    if (num < 20) return 0;  
    return ((num % 10) * 10 + reverseNum (num / 10));  
}
```

e.g.  $N = 1234$

① reverseNum(1234) {  
 ↳ num > 0  
 ↳  $4 \times 10 + \text{reverseNum}(123)$

↓  
② reverseNum(123)  
 ↳ num > 0  
 ↳  $3 \times 10 + \text{reverseNum}(12)$

↓  
③ reverseNum(12)  
 ↳ num > 0  
 ↳  $2 \times 10 + \text{reverseNum}(1)$

↓  
④ reverseNum(1)  
 ↳ num > 0  
 ↳  $1 \times 10 + \text{reverseNum}(0)$

↓  
⑤ return 0; // reverseNum = 4321

∴ Time Complexity =  $O(\log n)$

11. Count how many times a digit appears in a number.

```
→ main () {
    int count = 0;
    s.o.p (countNum (717237));
}
```

```
int countNum (int num) {
    if (num <= 0) return 0; rem = num % 10;
    num = num / 10;
```

```
    if (rem == digit) {
        count++;
    }
```

```
    countNum (num);
    return count;
}
```

① countNum (717237)  
↳ rem = 7, count = 1, num = 71723

② countNum (71723)  
↳ rem = 3, num = 7172

③ countNum (7172)  
↳ rem = 2, num = 717

④ countNum (717)  
↳ rem = 7, count = 2, num = 71

⑤ countNum (71)  
↳ rem = 1, num = 7

⑥ count (7)  
↳ rem = 7, count = 3

Time complexity =  $O(\log n)$



return (num % 10);

12. check if a number is palindromic.

→ main() {

int reverse = reverseNum(num);

s.o.p (isPalindrom(num, reverseNum));

}

int reverseNum(num) {

if (num <= 0) return 0;

return (num % 10 \* 10 + reverseNum(num / 10));

}

boolean isPalindrom(num, reverseNum) {

if (num == reverseNum) {

return true;

} else {

return false;

}

}

- eg. N = 121

① reverse(121)

↳ 1 \* 10 + reverse(12)

↓

② reverse(12)

↳ 12 \* 10 + reverse(1)

↓

③ reverse(1)

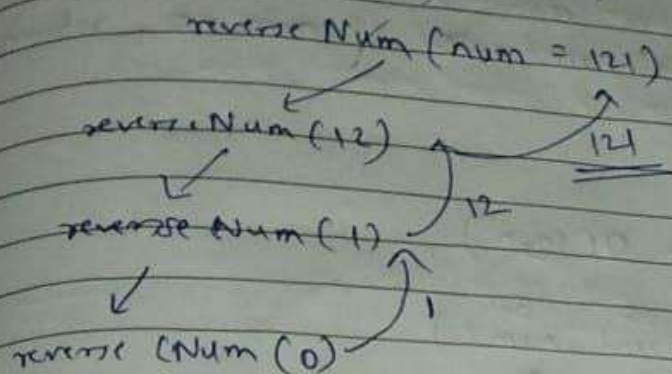
↳ 120 + 1; reverse = 121

1. Palindrome (121, 121)

↳  $121 = 121$

↳ Therefore, 121 is palindrome.

Time Complexity =  $O(\log n)$



13. Find GCD of two numbers using recursion.

→ main() {

    s.o.f (gcd(num1, num2));

}

int gcd(num1, num2) {

    if (num1 == 0 || num2 == 0) return Math.max(num1, num2);

    max = Math.max(num1, num2);

    min = Math.min(num1, num2);

    return gcd(max % min, min);

}

e.g.  $A = 24, B = 36$

① gcd(24, 36)

↳  $\max = 36, \min = 24$

↳  $\text{gcd}(36 \% 24, 24) = \text{gcd}(12, 24)$



②  $\text{gcd}(12, 24)$

↳  $\text{max} = 24, \text{min} = 12$

↳  $\text{gcd}(24 \% 12, 12) = \text{gcd}(0, 12)$

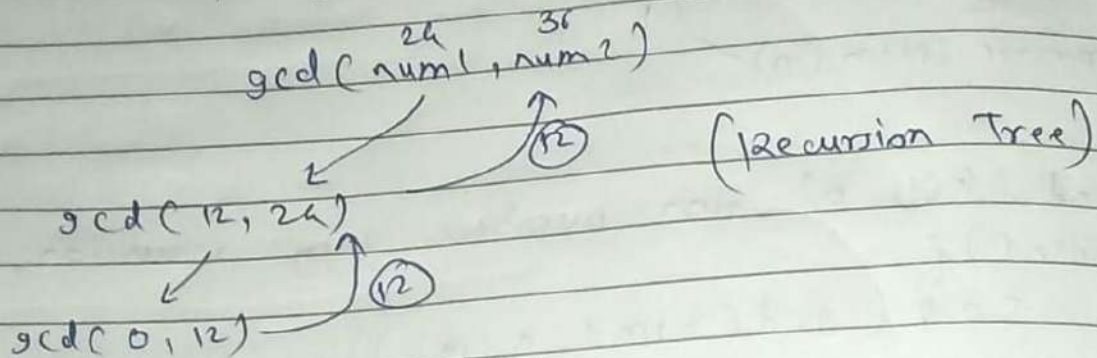
③  $\text{gcd}(0, 12)$

↳  $\text{max} = 12, \text{min} = 0$

↳  $\text{num1} = 24 \parallel \text{num} = 0, \text{return } \text{max}$

∴  $\text{gcd} = \underline{\underline{12}}$

1. Time complexity =  $O(\log n)$



14 Print all numbers from 1 to N divisible by 3.

main () {

divisible By Three (num, 1);

}

void divisible By Three (num, index) {

if (~~num~~ \* index > num)

return;

s.o.p (3 \* index);

divisible By Three (num, index + 1);

}

① divisible By Three (10, 1)

↳  $3 \times 1 < \text{num}$

↳ s.o.p (3); divisible By Three (10, 2)

② divisible By Three (10, 2)

↳  $3 \times 2 < \text{num}$

↳ s.o.p (6); divisible By Three (10, 3)

③ divisible By Three (10, 3)

↳  $3 \times 3 < \text{num}$

↳ s.o.p (9); divisible By Three (10, 4)

④ divisible By Three (10, 4)

↳  $3 \times 4 > \text{num} \rightarrow \text{return};$

∴ Time Complexity =  $O(\frac{n}{3})$

divisible By Three (10, 1)

③ ✓

divisible By Three (10, 2)

⑥ ✓

divisible By Three (10, 3)

⑨ ✓

divisible By Three (10, 4)



15. Find power of a number using recursion.

→

```
main() {
```

```
    s.o.p(pow(int A, int B));
```

```
}
```

```
int pow(num1, num2) {
```

```
    if (num2 <= 0) return 1;
```

```
    return num1 * pow(num1, num2-1);
```

```
}
```

eg.  $\text{pow}(2, 4)$

↳  $4 > 0$

↳ return  $2 * \text{pow}(2, 3)$

$8 * 2 = 16$

$\text{pow}(2, 3)$

↳  $3 > 0$

↳ return  $2 * \text{pow}(2, 2)$

$4 * 2 = 8$

$\text{pow}(2, 2)$

↳  $2 > 0$

↳ return  $2 * \text{pow}(2, 1)$

$2 * 2 = 4$

$\text{pow}(2, 1)$

↳  $1 > 0$

↳ return  $2 * \text{pow}(2, 0)$

$2 * 1 = 2$

return 1;

Power = 16, Time Complexity =  $O(\text{num2})$ ;

