Step 3: Soot each sublists using any sooting algorithm

step 4: Recursively find P as median of median of each sublists

Steps: Pastition S into Sq & Sz such that Si contains all elements less then or equal to p & Si contains all elements greater than P

Step 6: Now there can be 3 cases:

(a) If $|S_1| = |K|$ then $S_1(K)$ will be 1cth smaller

(b) if $|S_1| > |K|$ then |K| smallest must be

present in $|S_1|$ (c) otherwise |K| smallest must be in $|S_2|$

· Proof of the complexity

Step 2 takes o(n/5) time, step 3 takes
o(n)

Now we claim that 3/10 out of theorem

Orc Ep

Step 4 takes 4n/10 time

'-T(n) \equiv cn + T(n/5) + T(7n/10)

now we know that

T(n)= Q(n)
T(n)= Q(n)