

## Lecture ÷ Contest Discussion I

### Agenda

- Longest strictly increasing subarray
- Negative numbers in a range
- Time complexity problem.

Q1

## Negative numbers in a range

Given an  $arr[n]$  and  $Q$  queries. for every query return the negative numbers in given range.

Examples:  $arr[] = [-2^0, 1^1, -8^2, 3^3, 9^4, -6^5]$

| $l$ | $r$ | count |
|-----|-----|-------|
| 0   | 3   | 2     |
| 3   | 5   | 1     |
| 1   | 3   | 1     |
| 2   | 4   | 1     |

Brute force:

for every query —  $O(Q)$

Iterate the array for given range

and count negative no. —  $O(n)$

TC:  $O(Q * n)$

constraint:

$$1 \leq n \leq 10^5$$

$$1 \leq Q \leq 10^5$$

$$Q * n = 10^{10} \quad [O(Q * n) \text{ won't work}]$$

### Optimised code

$$arr[] = \begin{bmatrix} -2 & 1 & -8 & 3 & 9 & -6 \end{bmatrix}$$

$$pf[] = [1 \quad 1 \quad 2 \quad 2 \quad 2 \quad 3]$$

$$pf[0] = 1 \quad \text{if } arr[0] < 0 \rightarrow 1$$

$$\text{else} \rightarrow 0.$$

$$\begin{aligned} \text{pf}[1] &= \text{count of negative no } [0, 1] \\ &= \text{pf}[0] + \text{if } \text{arr}[1] < 0 \rightarrow 1 \\ &\quad \text{else} \rightarrow 0 \end{aligned}$$

$$\text{count of negative no}[l:r] = \text{pf}[r] - \text{pf}[l-1]$$

Edge case  $\div$   $l = 0$

└ pf[r].

Qn longest strictly increasing subarray.

1.  $arr[] = [6, 9, 9, 8]$

$arr[] = [8, 4, 5, 7, 6, 11, 15, 19, 1, 0, 5]$

Brute force:

Go to all subarray —  $O(n^2)$

check if it strictly increasing or not —  $O(n)$ .

Yes — update your ans

No — ignore.

constraint:  $1 \leq n \leq 10^5$

$O(n^3) \approx 10^{15} \rightarrow$  not work

$O(n^2) \approx 10^{10} \rightarrow$  not work

$arr[] = [8, 4, 5, 7, 6, 11, 15, 19, 1, 0, 5]$

$i = 0, j = 1$

$arr[1] < arr[0]$

$i = 1$

$j = 2$

$arr[j] > arr[j-1] \rightarrow$  update the ans.  $j++$

$j = 3$

$arr[j] > arr[j-1] \rightarrow$  update your ans.  $j++$

$j = 4$

$arr[j] > arr[j-1] \rightarrow$  No

arr[] = [ 0 1 2 3 4 5 6 7 8 9 10 ]  
 8 4 5 7 6 11 15 19 1 0 5  
 ↑ i ↑ j

i = 4

|                                 |                       |
|---------------------------------|-----------------------|
| j = 5<br>arr[j] > arr[j-1]      | j++ → j - i + 1 = 2 ✓ |
| j = 6<br>arr[6] > arr[5]        | j++ → j - i + 1 = 3 ✓ |
| j = 7<br>arr[7] > arr[6]        | j++ → j - i + 1 = 4 ✓ |
| <u>j = 8</u><br>arr[8] < arr[7] |                       |

i = 4 and j = 8 [ ans = j - i ]

i = 8  
 j = 9  
 ...

```

while (j < n) {
    if (arr[j] > arr[j-1]) {
        j++;
    } else {
        break;
    }
}
  
```

Q. Time complexity

```

→ for (i = 1; i <= n; i++) {
    for (j = 1; j <= n; j++) {
        print("Priyanka");
        → break;
    }
    if (i * i > n) {
        break;
    }
}

```

8 \* 8 = 64 > 50  
break;

| n = 50 |          |     |
|--------|----------|-----|
| i      | j [1, n] | itr |
| 1      | 1        | 1   |
| 2      | 1        | 1   |
| 3      | 1        | 1   |
| 4      | 1        | 1   |
| ⋮      |          |     |
| 7      | 1        | 1   |
| 8      | 1        | 1   |

n = 50

$\sqrt{n} = 7$

$$O(\sqrt{n} + 1) \approx O(\sqrt{n})$$

## Doubts

```
for(i=0; i<n; i++) {
```

$$j = l + 1.$$

while ( $j < n$ ) {

$$j++;$$
$$i = j - 1$$

$i = j$

