

# **CSE 5331- DBMS Models and Implementation**

## **Project 2: Implementation of a Transaction Manager**

**Team no.:** 30

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## 1. Overall status:

Following is the status on implementation tasks assigned in this project:

1. **zgt\_tx.C** - In this class, we have completed the implementation of the four functions (excluding begintx function) assigned.
  - a. readtx
  - b. writetx
  - c. aborttx
  - d. committx
2. **zgt\_tm.C** - In this class, we have completed the implementation of the following functions.
  - a. TxRead
  - b. TxWrite
  - c. CommitTx
  - d. AbortTx
3. We have successfully combined and ran using the test data files given.

Below is more information on how we have implemented the above mentioned classes and functions.

**1.1. Transaction manager class (zgt\_tm.C):** In this class, we first have set the team\_no to our team no. i.e. 30. This class basically creates threads for each operation of a transaction as per mentioned in the input text file. There are 6 operations that are basically performed.

**1. openlog(string lfile):** This function creates a log file and writes the initial headers in the log file. In all the other functions, we have implemented a similar procedure to be followed.

**2. BeginTx(long tid, int thrNum, char type):** This function basically creates a thread corresponding to the begin transaction operation mentioned in the input text file. For doing so it creates a node with the required information such as tid (transaction id), Txtype (transaction type), object number (obno = -1) and count (= 0). This node is then passed as arguments to the begintx function in zgt\_tx.C (discussed later).

**3. TxRead(long tid, long obno, int thrNum):** This function basically performs the same thing as BeginTx but here it sets the object number to the required obno and count is the sequence number of the tid decremented by 1. And then this node is then passed as arguments to the readtx function in zgt\_tx.C (discussed later).

**4. TxWrite(long tid, long obno, int thrNum):** Similar to the TxRead function. The difference is just that the node is passed as arguments to the writetx function in zgt\_tx.C (discussed later).

**5. CommitTx(long tid, int thrNum):** The arguments passed are tid and count only to the committx function in zgt\_tx.C (discussed later).

**6. AbortTx(long tid, int thrNum):** Same as in the CommitTx function.

**7. endTx(int thrNum):** This function is called when end all is encountered in the input text file. This function was already implemented in the given code files.

**1.2. Transactions class (zgt\_tx.C):** This class basically performs the transaction operations by setting the required locks for different threads.

**1. begintx(void \*arg):** This function is called from the BeginTx function in the transaction manager class. It starts the operation by calling start\_operation function. The start\_operation function makes all the threads belonging to the same transaction mutually exclusive from each other depending upon the condition set. Then adds the transaction node to the transaction manager data structure at LASTR (while modifying the transaction manager values, we had to lock the transaction manager using semaphore 0). Once this is done, finish the operation by unlocking the threads.

**2. readtx(void \*arg):** This function calls the start\_operation() function and then gets the transaction from the transaction manager using the get\_tx() function. Then we send the transaction values i.e, tid, obno and count in a shared lockmode('S') to the set\_lock() function (discussed below in detail). Finally, it calls the finish\_operation() function and exits the assigned thread using pthread\_exit(NULL).

**3. writetx(void \*arg):** This function performs the same operations as in the readtx function above. The only difference is that it sets an exclusive lock.

**4. aborttx(void \*arg):** This function first starts the operation using the start\_operation function and then calls the do\_commit\_abort() function to perform the abort and commit operations. While calling the function do\_commit\_abort() function, the operation (abort or commit) is indicated by the status argument. If the status argument is passed as 'A' then the abort operation is performed and if it is 'C' then a commit operation is performed. After that, the operation is completed by calling finish\_operation function.

**5. committx(void \*arg):** This function behaves the same way as the aborttx() function but sends a 'C' character to the do\_commit\_abort() function as the status to indicate a commit of the transaction.

**6. do\_commit\_abort(long t, char status):** This function performs the actual commit and abort of the transactions. It first gets the transaction from the transaction manager using the get\_tx() function and then checks if the status is 'C' or 'A'. If the transaction status is 'A', it writes the abort of the transaction to the logfile. It frees all the locks that the transaction is holding using the free\_locks() function and then removes the transaction from the Transaction manager using the remove\_tx() function so that none of the operations done by the transaction are made permanent. If the transaction status is 'C', it does the same thing that it does for 'A' with one key difference: it calls the end\_tx() function instead of the remove\_tx() function so that the operations done by the transaction are not removed.

**7. remove\_tx () and end\_tx():** These two functions are used by do\_commit\_abort() functions. The remove\_tx() function is used to remove the transaction from the transaction list in the transaction manager without committing the changes to the disk. Whereas the end\_tx() function is used to remove transactions as well as committing the changes to the disk. Thus the remove\_tx() function is used to perform the abort operation and end\_tx() function is used to perform the commit operation.

**8. free\_locks():** This function simply removes all the locks that a transaction holds. It does this by using the head pointer of the transaction to locate all the objects locked by the transaction i.e, all the objects that the transaction put into the hash table. It removes all these objects from the hash table, thus unlocking them.

**9. set\_lock(long tid1, long sgno1, long obno1, int count, char lockmode1):** As mentioned in the above points, the readTx and writeTx functions use the set\_lock function. The set\_lock function basically handles the locks (shared and exclusive) for all objects. Following are the steps that we have followed for locking a particular object:

- a. Check the hash table (lock table) in the transaction manager data structure to find if there is any existing lock on the object with obno1. Note: Before reading or making any changes in the lock table, we need to lock the semaphore 0 (semaphore for locking transaction manager).
- b. If there is no lock (shared or exclusive) on the required object then grant the access. On granting the access, perform read/ write by calling the perform\_read\_write() function (discussed below).
- c. If there is a lock on the required object then check if the same current transaction is having the access then perform read/ write by calling the perform\_read\_write() function as there is no need to grant access.
- d. If there is a lock on the required object but the transaction is different then find the transaction that is holding the current lock. And then,
  - i. If the current transaction requires shared lock and the transaction holding the current lock is a read-only transaction then grant the shared lock on the required object and perform read write.
  - ii. If the transaction currently holding the object is a read/ write transaction or if current transaction requires an exclusive lock or if current transaction requires shared lock and transaction holding the lock is read-only transaction but there are other transactions that are waiting on the required object then do the following: Set a semaphore for the current transaction and make it wait till the current transaction performs its operations. This is done using the setTx\_semno() function and then zgt\_p() and zgt\_v() functions. Setting the value of zgt\_p(x) allows the transaction x to complete its critical section and it is unset by using zgt\_v(x). While the transaction having the lock on the required object executes its critical section, the current transaction is made to wait by setting its values to default values as mentioned in the project description file (i.e. obno = -1, lockmode = ' ' and status = TR\_ACTIVE).

**10. perform\_read\_write(long tid, long obno, char lockmode):** This function performs the actual read and write operations on the objects in a simulated manner. First, it checks the lockmode of the object. If it is a shared lock denoting a read operation, it accesses the object in the object array using the obno and decrements its value by 1 and writes this to the log file. After this, to simulate the amount of time required for an actual disk read operation, the control sleeps for the given amount of time using the sleep() function. The time for which the operation is made to sleep depends on the optime array which is generated randomly using a large prime number as seed.

If the lockmode is exclusive denoting a write operation, the function does the same thing as a read operation but instead of decrementing the value, it increments the value of the object in the object array by 1 and then sleeps for the required amount of time.

## 2. Input/output, log files and overall status of outputs:

### 2.1. Input and log files:

#### 1. no\_conflicts\_2Txs.txt

```
// serializable history
// 2 transactions (no conflicts)
// same object accessed
// multiple times
Log no_conflicts_2Txs.log
BeginTx 1 R
Read  1 1
Read  1 2
BeginTx 2 W
Read  2 8
Read  2 7
Write 2 6
Write 2 5
Commit 2
read  1 3
read  1 4
Commit 1
end all
```

```
[axp3623@omega src]$ cat no_conflicts_2Txs.log
```

TxId	Txtype	Operation	Obld:Obvalue:optime	LockType	Status	TxStatus
T1	R	BeginTx				
T1		readTx	1:-1:29614	ReadLock	Granted	P
T2	W	BeginTx				
T2		readTx	8:-1:3675	ReadLock	Granted	P
T2		readTx	7:-1:3675	ReadLock	Granted	P
T2		writeTx	6:1:3675	WriteLock	Granted	P
T2		writeTx	5:1:3675	WriteLock	Granted	P
T2		CommitTx				
T1		readTx	2:-1:29614	ReadLock	Granted	P
T1		readTx	3:-1:29614	ReadLock	Granted	P
T1		readTx	4:-1:29614	ReadLock	Granted	P
T1		CommitTx				

## 2. interleaved\_RW.txt

log interleaved\_RW.log

BeginTx 1 W

Read 1 1

Write 1 2

Read 1 3

Write 1 4

BeginTx 2 W

Write 2 8

Write 2 2

Read 2 1

BeginTx 3 W

Write 3 1

Write 3 2

Write 3 9

Read 3 8

Commit 1

commit 2

Commit 3

end all

[axp3623@omega src]\$ cat interleaved\_RW.log

TxId	Txtype	Operation	ObId:Obvalue:optime	LockType	Status	TxStatus
T1	W	BeginTx				
T1		readTx 1:-1:29614		ReadLock	Granted	P
T2	W	BeginTx				
T3	W	BeginTx				
T2		writeTx8:1:3675		WriteLock	Granted	P
T2		writeTx2:1:3675		WriteLock	Granted	P



### 3. Multi\_ROTxs.txt

```
// Multiple RO Txs test case
// read only transactions
log Multi_ROTxs.log
// op   Tx#   type
// op   Tx#   Obj
BeginTx 1 R
Read 1 1
Read 1 2
Read 1 3
Read 1 8
BeginTx 2 R
Read 2 1
Read 2 8
Read 2 5
BeginTx 3 R
Read 3 1
Read 3 5
Read 3 3
read 3 7
Commit 2
commit 3
Commit 1
end all
```

```
[axp3623@omega src]$ cat Multi_ROTxs.log
```

TxId	Txtype	Operation	Obld:Obvalue:optime	LockType	Status	TxStatus
T1	R	BeginTx				
T1		readTx	1:-1:29614	ReadLock	Granted	P
T2	R	BeginTx				
T2		readTx	1:-2:3675	ReadLock	Granted	P
T2		readTx	8:-1:3675	ReadLock	Granted	P
T2		readTx	5:-1:3675	ReadLock	Granted	P
T3	R	BeginTx				
T3		readTx	1:-3:19271	ReadLock	Granted	P
T2		CommitTx				
T3		readTx	5:-2:19271	ReadLock	Granted	P
T1		readTx	2:-1:29614	ReadLock	Granted	P
T3		readTx	3:-1:19271	ReadLock	Granted	P
T3		readTx	7:-1:19271	ReadLock	Granted	P
T1		readTx	3:-2:29614	ReadLock	Granted	P
T3		CommitTx				

T1	readTx	8:-2:29614	ReadLock	Granted	P
T1	CommitTx				

#### 4. disj\_multi\_accesses.txt

```
// serial history
// 2 transactions
// same disjoint objects accessed
// multiple times
Log disj_multi_accesses.log
BeginTx 1 W
Read  1 1
Read  1 2
Write 1 3
Write 1 4
read  1 1
write 1 2
write 1 4
write 1 4
commit 1
begintx 2 W
read  2 5
write 2 5
write 2 6
read  2 6
commit 2
end all
```

[axp3623@omega src]\$ cat disj\_multi\_accesses.log

TxId	Txtype	Operation	ObId:Obvalue:optime	LockType	Status	TxStatus
T1	W	BeginTx				
T1		readTx	1:-1:29614	ReadLock	Granted	P
T2	W	BeginTx				
T2		readTx	5:-1:3675	ReadLock	Granted	P
T2		writeTx	5:0:3675	WriteLock	Granted	P
T2		writeTx	6:1:3675	WriteLock	Granted	P
T2		readTx	6:0:3675	ReadLock	Granted	P
T2		CommitTx				
T1		readTx	2:-1:29614	ReadLock	Granted	P
T1		writeTx	3:1:29614	WriteLock	Granted	P
T1		writeTx	4:1:29614	WriteLock	Granted	P
T1		readTx	1:-2:29614	ReadLock	Granted	P
T1		writeTx	2:0:29614	WriteLock	Granted	P

T1	writeTx	4:2:29614	WriteLock	Granted	P
T1	writeTx	4:3:29614	WriteLock	Granted	P
T1	CommitTx				

## 5. ddlk\_3Tx.txt

// possible deadlock test case

// Two write transactions

log ddlk\_3Tx.log

// op Tx# type

BeginTx 1 W

// op Tx# Obj

Read 1 1

Write 1 2

Read 1 6

BeginTx 2 W

Read 2 2

Write 2 1

Read 2 7

commit 2

Commit 1

beginTx 3 R

read 3 2

write 3 1

read 3 2

end all

[axp3623@omega src]\$ cat ddlk\_3Tx.log

TxId	Txtype	Operation	ObId:Obvalue:optime	LockType	Status	TxStatus
T2	W	BeginTx				
T1	W	BeginTx				
T1		readTx	1:-1:29614	ReadLock	Granted	P
T2		readTx	2:-1:3675	ReadLock	Granted	P
T3	R	BeginTx				
T3		readTx	2:-2:19271	ReadLock	Granted	P

## 6. ddlk\_2Txs.txt

```
// 2 transactions
// classic deadlock
// will hang w/o deadlock resolution
Log ddlk_2Tx.log
BeginTx 1 W
BeginTx 2 W
Read  1 1
Read  2 2
Write 1 2
Write 2 1
Commit 1
commit 2
end all
```

```
[axp3623@omega src]$ cat ddlk_2Tx.log
```

TxId	Txtype	Operation	Obld:Obvalue:optime	LockType	Status	TxStatus
T1	W	BeginTx				
T2	W	BeginTx				
T1		readTx	1:-1:29614	ReadLock	Granted	P
T2		readTx	2:-1:3675	ReadLock	Granted	P

## 7. test\_abort.txt

```
// simple deadlock test case
// Two write transactions
log test_abort.log
// op   Tx#   type
// op   Tx#   Obj
BeginTx 1 W
read 1 6
write 1 7
write 1 7
read 1 6
beginTx 2 W
read 2 8
write 2 7
abort 2
beginTx 3 R
read 3 4
write 3 5
read 3 9
commit 3
commit 1
end all
```

[axp3623@omega src]\$ cat test\_abort.log

TxId	Txtype	Operation	Obld:Obvalue:optime	LockType	Status	TxStatus
T1	W	BeginTx				
T1		readTx	6:-1:29614	ReadLock	Granted	P
T2	W	BeginTx				
T2		readTx	8:-1:3675	ReadLock	Granted	P
T2		writeTx	7:1:3675	WriteLock	Granted	P
T2		AbortTx				
T3	R	BeginTx				
T3		readTx	4:-1:19271	ReadLock	Granted	P
T3		writeTx	5:1:19271	WriteLock	Granted	P
T3		readTx	9:-1:19271	ReadLock	Granted	P
T1		writeTx	7:2:29614	WriteLock	Granted	P
T3		CommitTx				
T1		writeTx	7:3:29614	WriteLock	Granted	P
T1		readTx	6:-2:29614	ReadLock	Granted	P
T1		CommitTx				

## 8. RW\_disjoint.txt

```
// Multiple RW Txs test case with no deadlock
log RW_disjoint.log
// op   Tx#   type
// op   Tx#   Obj
BeginTx 1 W
Read 1 1
Write 1 2
Read 1 3
BeginTx 2 W
Write 2 4
Write 2 5
BeginTx 3 W
Write 3 6
Write 3 7
read 3 8
Commit 3
commit 2
Commit 1
begintx 5 R
read 5 9
read 5 10
read 5 11
read 5 12
read 5 13
read 5 1
commit 5
end all
```

```
[xap3623@omega src]$ cat RW_disjoint.log
```

TxId	Txtype	Operation	ObId:Obvalue:optime	LockType	Status	TxStatus
T1	W	BeginTx				
T1		readTx	1:-1:29614	ReadLock	Granted	P
T2	W	BeginTx				
T2		writeTx	4:1:3675	WriteLock	Granted	P
T2		writeTx	5:1:3675	WriteLock	Granted	P
T3	W	BeginTx				
T3		writeTx	6:1:19271	WriteLock	Granted	P
T2		CommitTx				
T5	R	BeginTx				
T5		readTx	9:-1:16371	ReadLock	Granted	P
T3		writeTx	7:1:19271	WriteLock	Granted	P

T5	readTx	10:-1:16371	ReadLock	Granted	P
T1	writeTx	2:1:29614	WriteLock	Granted	P
T3	readTx	8:-1:19271	ReadLock	Granted	P
T5	readTx	11:-1:16371	ReadLock	Granted	P
T3	CommitTx				
T5	readTx	12:-1:16371	ReadLock	Granted	P
T1	readTx	3:-1:29614	ReadLock	Granted	P
T5	readTx	13:-1:16371	ReadLock	Granted	P
T5	readTx	1:-2:16371	ReadLock	Granted	P
T1	CommitTx				
T5	CommitTx				

## 9. RW\_pot\_ddlk.txt

// Multiple RW Txs test case with no deadlock

log RW\_pot\_ddlk.log

// op Tx# type

// op Tx# Obj

BeginTx 1 W

Read 1 1

Write 1 2

Read 1 3

Write 1 8

BeginTx 2 W

Write 2 4

Write 2 5

BeginTx 3 W

Write 3 6

Write 3 7

Read 3 9

Commit 3

commit 2

Commit 1

begintx 5 R

read 5 1

read 5 2

read 5 3

read 5 8

read 5 6

read 5 7

commit 5

end all

[axp3623@omega src]\$ cat RW\_pot\_ddlk.log

TxId	Txtype	Operation	Obld:Obvalue:optime	LockType	Status	TxStatus
T1	W	BeginTx				
T1		readTx	1:-1:29614	ReadLock	Granted	P
T2	W	BeginTx				
T2		writeTx	4:1:3675	WriteLock	Granted	P
T2		writeTx	5:1:3675	WriteLock	Granted	P
T3	W	BeginTx				
T3		writeTx	6:1:19271	WriteLock	Granted	P
T2		CommitTx				
T5	R	BeginTx				
T5		readTx	1:-2:16371	ReadLock	Granted	P
T3		writeTx	7:1:19271	WriteLock	Granted	P
T5		readTx	2:-1:16371	ReadLock	Granted	P
T3		readTx	9:-1:19271	ReadLock	Granted	P
T5		readTx	3:-1:16371	ReadLock	Granted	P
T3		CommitTx				
T5		readTx	8:-1:16371	ReadLock	Granted	P
T5		readTx	6:0:16371	ReadLock	Granted	P
T5		readTx	7:0:16371	ReadLock	Granted	P
T5		CommitTx				
T1		writeTx	2:0:29614	WriteLock	Granted	P
T1		readTx	3:-2:29614	ReadLock	Granted	P
T1		writeTx	8:0:29614	WriteLock	Granted	P
T1		CommitTx				



## 10. unlikely\_ddlk.txt

```
// serializable history
// ddlk unlikely
log unlikely_ddlk.log
BeginTx 1 W
Read  1 3
Read  1 2
BeginTx 2 W
Read  2 1
Write 2 3
Write 1 3
Write 1 2
Write 2 2
Commit 1
commit 2
end all
```

```
[axp3623@omega src]$ cat unlikely_ddlk.log
```

TxId	Txtype	Operation	Obld:Obvalue:optime	LockType	Status	TxStatus
T1	W	BeginTx				
T1		readTx	3:-1:29614	ReadLock	Granted	P
T2	W	BeginTx				
T2		readTx	1:-1:3675	ReadLock	Granted	P
T1		readTx	2:-1:29614	ReadLock	Granted	P
T1		writeTx	3:0:29614	WriteLock	Granted	P
T1		writeTx	2:0:29614	WriteLock	Granted	P
T1		CommitTx				
T2		writeTx	3:1:3675	WriteLock	Granted	P
T2		writeTx	2:1:3675	WriteLock	Granted	P
T2		CommitTx				

## 11. multiple\_aborts.txt

```
// Multiple RW Txs test case with no deadlock
// 23 operations
log multiple_aborts.log
// op   Tx#   type
// op   Tx#   Obj
BeginTx 1 W
Read 1 1
Write 1 2
Read 1 3
Write 1 8
BeginTx 2 W
Write 2 4
Write 2 5
BeginTx 3 W
Write 3 6
Write 3 7
Read 3 9
abort 3
commit 2
abort 1
beginTx 5 R
read 5 1
read 5 2
read 5 3
read 5 8
read 5 6
read 5 7
abort 5
end all
```

```
[axp3623@omega src]$ cat multiple_aborts.log
```

TxId	Txtype	Operation	Obld:Obvalue:optime	LockType	Status	TxStatus
T1	W	BeginTx				
T1		readTx	1:-1:29614	ReadLock	Granted	P
T2	W	BeginTx				
T2		writeTx	4:1:3675	WriteLock	Granted	P
T2		writeTx	5:1:3675	WriteLock	Granted	P
T3	W	BeginTx				
T3		writeTx	6:1:19271	WriteLock	Granted	P
T2		CommitTx				
T5	R	BeginTx				

T5	readTx	1:-2:16371	ReadLock	Granted	P
T5	readTx	2:-1:16371	ReadLock	Granted	P
T3	writeTx	7:1:19271	WriteLock	Granted	P
T3	readTx	9:-1:19271	ReadLock	Granted	P
T5	readTx	3:-1:16371	ReadLock	Granted	P
T3	AbortTx				
T5	readTx	8:-1:16371	ReadLock	Granted	P
T5	readTx	6:0:16371	ReadLock	Granted	P
T5	readTx	7:0:16371	ReadLock	Granted	P
T5	AbortTx				
T1	writeTx	2:0:29614	WriteLock	Granted	P
T1	readTx	3:-2:29614	ReadLock	Granted	P
T1	writeTx	8:0:29614	WriteLock	Granted	P
T1	AbortTx				

## 2.2. Overall status of outputs:

	test file	zgt_test	tmtest 100	tmtest 1000	max tmtest value < 50000 for which it works correctly
1	no_conflict_2Txs.txt	ok	works	works	50000
2	interleaved_RW.txt	hangs (as it should be)	hangs	hangs	hangs
3	Multi_ROTxs.txt	ok	works	works	50000
4	disj_multi_accesses.txt	ok	works	works	50000
5	ddlk_3Txs.txt	hangs	hangs	hangs	hangs
6	ddlk_2Txs.txt	hangs	hangs	hangs	hangs
7	test_abort.txt	aborts T2 correctly	works	works	50000
8	RW_disjoint.txt	ok	works	works	50000
9	RW_pot_ddlk.txt	ok	works	works	50000
10	unlikely_ddlk.txt	ok	works	works	50000
11	multiple_aborts.txt	ok	works	works	50000

### 3. Difficulties and logical errors:

#### 3.1. Where we encountered difficulty:

The difficulties we faced were mostly centered around the `set_lock()` function. It took a lot of time for us to correctly implement all the possible cases. We faced some problems with making the transaction wait using a semaphore when another transaction had the lock.

We also encountered some difficulties with the `perform_read_write()` function because we were not sure how the `sleep()` function corresponded to the execution. When the deadlock test case hung, we were unsure what was causing it: a deadlock or some error in the `sleep()` function.

**3.2. File descriptions:** There are no new files, data structures or additional test cases that we have added in the submission zip file.

#### 3.3 Division of Labor:

Group member name	Functions implemented	Hours spent on project on weekdays	Hours spent on project on weekends
Pratik Antoni Patekar (1001937948)	Read and abort related functions in both the classes	Utmost 3 hours	4 to 5 hours
Anurag Reddy Pingili(1001863623)	Write and Commit functions in both the classes.	Utmost 3 hours	4 to 5 hours

#### 3.4. Logical Errors:

- Some test cases were leading to a segmentation fault. We weren't sure what was causing this but it went away after we implemented our `set_lock()` function correctly.
- Race condition: Our test cases were hanging because we didn't check the number of threads waiting for a lock and granted the lock to a transaction immediately. This went away after we checked the `wait_tx_no > 0`.
- Locking and unlocking TM: We were not doing this before so we were not sure why our test cases kept hanging but we realized our error once we properly understood the `beginTx` function.