### BT5110: Tutorial 3 — Functional Dependencies

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AY25/26 S1



Example Relation, 
$$R = \{A, B, C\}$$
 and  $\Sigma = \{\{A\} \rightarrow \{B\}, \{B\} \rightarrow \{C\}\}$ 

The column in red decides the color of the row: The rows having same values in the red column, have same color.

Α	В	C
a1	b1	c1
a1	b1	c1
a2	b2	c2
a2	b2	c2
а3	b3	c2
a3	b3	c2
a4	b1	c1
a4	b1	c1
a5	b2	c2

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{ <mark>A} −</mark>	→ { <i>A</i> ,	B, C

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a4	b1	c1
a5	b2	c2
{ <b>C</b>	$\rightarrow \{$	<b>C</b> }

### Question 1.a.

Your company, Apasaja Private Limited, is commissioned by *Toko Kopi Luwak* to design the relational schema for coffee beans, drinks, and cafes. We are given

$$R=\{A,B,C,D,E,F,G,H\},$$

$$\Sigma = \{\{A\} \to \{C, E\}, \{A, B\} \to \{D\}, \{F\} \to \{H\}, \{C, E\} \to \{A\}, \{B, C, E\} \to \{D\}, \{A, B, F\} \to \{D, G\}, \{B, C, E, F\} \to \{G\}\}.$$

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### Attribute mapping

Symbol	Meaning
A	name (bean brand)
В	dname (drink name)
C	cultivar
D	price



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### Attribute mapping

Symbol	Meaning
A	name (bean brand)
В	dname (drink name)
C	cultivar
D	price

Symbol	Meaning
E	region
F	bname (branch)

Attribute mapping

G qty (quantity sold)
H address

Notes: A bean is identified by name or by (cultivar, region); a drink name is unique per bean; branches sell drinks, with address and quantity recorded.

### Rules for Functional Dependencies (Armstrong's Axioms)

W, X, Y, Z mentioned here are sets of attributes.

### Sound & complete inference system

- **Reflexivity**: If  $Y \subseteq X$ , then  $X \to Y$ .
- Augmentation: If  $X \to Y$ , then  $XZ \to YZ$  for any Z.
- Transitivity: If  $X \to Y$  and  $Y \to Z$ , then  $X \to Z$ .

### Common derived rules (from the three above)

- Union / Additivity: If  $X \to Y$  and  $X \to Z$ , then  $X \to Y \cup Z$ .
- **Decomposition / Projectivity**: If  $X \to Y \cup Z$ , then  $X \to Y$  and  $X \to Z$ .
- Pseudotransitivity: If  $X \to Y$  and  $W \cup Y \to Z$ , then  $W \cup X \to Z$ .
- Composition: If  $X \to Y$  and  $Z \to W$ , then  $X \cup Z \to Y \cup W$ .
- Self-determination:  $X \to X$ .

### How we use them here

- LHS reduction: test if  $X \setminus \{A\} \to Y$  holds via closure using the axioms.
- Redundancy removal: test if an FD X → Y is implied by the rest (i.e., Y ⊆ X<sup>+</sup> computed from the others).



$$\Sigma = \{ \{A\} \to \{C, E\}, \{A, B\} \to \{D\}, \{F\} \to \{H\}, \{C, E\} \to \{A\}, \{B, C, E\} \to \{D\}, \{A, B, F\} \to \{D, G\}, \{B, C, E, F\} \to \{G\} \}.$$

**Observation:** Since B and F never appear on the RHS of any FD, every key must contain B and F. We enumerate closures starting from supersets of  $\{B, F\}$ .

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```
Size 2 \{B, F\}^+ = \{B, F, H\} (not a key)

Size 3 (supersets of \{B, F\})

\{A, B, F\}^+ = \{A, B, C, D, E, F, G, H\}

\{B, C, F\}^+ = \{B, C, F, H\}

\{B, D, F\}^+ = \{B, D, F, H\}

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Size 4 (not containing \{A, B, F\})
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```

Larger supersets (that are not supersets of keys above) All such closures remain proper subsets of R.

### Conclusion Candidate keys: $\{A, B, F\}$ and $\{B, C, E, F\}.$

Tedious by hand? I made vour life a bit easier: Here

### Question 1.c. Prime attributes

Keys are  $\{A, B, F\}$  and  $\{B, C, E, F\}$ . Prime attributes are those that appear in at least one key.

Prime attributes = 
$$\{A, B, F\} \cup \{B, C, E, F\} = \{A, B, C, E, F\}$$
.

### Start from $\Sigma$ : $\{A\} \rightarrow \{C, E\},$ $\{A,B\} \rightarrow \{D\},$ $\{F\} \rightarrow \{H\},$ $\{C,E\} \rightarrow \{A\},$ $\{B,C,E\} \rightarrow \{D\},$ $\{A,B,F\} \rightarrow \{D,G\},$ $\{B,C,E,F\} \rightarrow \{G\}.$

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## $\begin{array}{c} \text{Step 1: Split RHS} \\ \{A\} \to \{C\}, \\ \{A\} \to \{E\}, \\ \{A,B\} \to \{D\}, \\ \{F\} \to \{H\}, \\ \{C,E\} \to \{A\}, \\ \{B,C,E\} \to \{D\}, \\ \\ \{A,B,F\} \to \{D\}, \\ \\ \{A,B,F\} \to \{G\}, \\ \{B,C,E,F\} \to \{G\}. \end{array}$

### Start from Σ: $\{A\} \rightarrow \{C, E\}, \\ \{A,B\} \rightarrow \{D\}, \\ \{F\} \rightarrow \{H\}, \\ \{C,E\} \rightarrow \{A\}, \\ \{B,C,E\} \rightarrow \{D\}, \\ \{A,B,F\} \rightarrow \{D,G\}, \\ \{B,C,E,F\} \rightarrow \{G\}.$

### Step 1: Split RHS $\{A\} \rightarrow \{C\},$ $\{A\} \rightarrow \{E\},$ $\{A,B\} \rightarrow \{D\},$ $\{F\} \rightarrow \{H\},$ $\{C,E\} \rightarrow \{A\},$ $\{B,C,E\} \rightarrow \{D\},$ $\{A,B,F\} \rightarrow \{D\},$ $\{A,B,F\} \rightarrow \{G\},$ $\{B,C,E,F\} \rightarrow \{G\}.$

# Step 2: Reduce LHS where possible $\{A\} \rightarrow \{C\},$ $\{A\} \rightarrow \{E\},$ $\{A,B\} \rightarrow \{D\},$ $\{F\} \rightarrow \{H\},$ $\{C,E\} \rightarrow \{A\},$ $\{B,C,E\} \rightarrow \{D\},$ $\{A,B,F\} \rightarrow \{G\},$ $\{B,C,E,F\} \rightarrow \{G\},$ $\{B,C,E,F\} \rightarrow \{G\},$

### Step 3: Remove redundancies

### Minimal Cover

$$\{A\} \rightarrow \{C\}, 
 \{A\} \rightarrow \{E\}, 
 \{F\} \rightarrow \{H\}, 
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### Step 3: Remove redundancies

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### Minimal Cover

$${A} \rightarrow {C},$$
  
 ${A} \rightarrow {E},$   
 ${F} \rightarrow {H},$   
 ${C,E} \rightarrow {A},$   
 ${B,C,E} \rightarrow {D},$   
 ${B,C,E,F} \rightarrow {G}.$ 

(Alternative minimal covers exist, e.g., replace 
$$\{B,C,E\} \to \{D\}$$
 by  $\{A,B\} \to \{D\}$  and  $\{B,C,E,F\} \to \{G\}$  by  $\{A,B,F\} \to \{G\}$ .)

### Question 2.b. Canonical cover (merged by LHS)

One canonical cover (aka Compact Minimal Cover) (grouping by identical LHS) is:  $\{A\} \rightarrow \{C, E\},\$   $\{F\} \rightarrow \{H\},\$   $\{C,E\} \rightarrow \{A\},\$   $\{B,C,E\} \rightarrow \{D\},\$ 

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 ${F} \rightarrow {H},$   
 ${C,E} \rightarrow {A},$   
 ${B,C,E} \rightarrow {D},$   
 ${B,C,E,F} \rightarrow {G}$ 

Other canonical covers are possible depending on which equivalent minimal cover you merge.

Questions?

Drop a mail at: pratik.karmakar@u.nus.edu