IT5008: Tutorial 7 — Functional Dependencies

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AY25/26 S1

Identifying Functional Dependencies

Example Relation, $R = \{A, B, C\}$ and $\Sigma = \{\{A\} \rightarrow \{B\}, \{B\} \rightarrow \{C\}\}$

The column in red decides the color of the row: The rows having same values in the red column, have same color.

Task: Check which other columns are following "same value if same color" for each deciding (red) column. The columns for which this holds, are implied by the red column.

A	B	C
a1	b1	c1
a1	b1	c1
a2	b2	c2
a2	b2	c2
a3	b3	c2
a3	b3	c2
a4	b1	c1
a4	b1	c1
a5	b2	c2

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$\{A\} \rightarrow \{A, B, C\}$

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$\{A\} \rightarrow \{A, B, C\}$

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a2	b2	c2
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a3	b3	c2
a4	b1	c1
a4	b1	c1
a5	b2	c2

$\{B\} \rightarrow \{B, C\}$

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$\{A\} \rightarrow \{A, B, C\}$

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a2	b2	c2
a3	b3	c2
a3	b3	c2
a4	b1	c1
a4	b1	c1
a5	b2	c2

$\{B\} \rightarrow \{B, C\}$

A	B	C
a1	b1	c1
a1	b1	c1
a2	b2	c2
a2	b2	c2
a3	b3	c2
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$\{A\} \rightarrow \{A, B, C\}$

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a1	b1	c1
a2	b2	c2
a2	b2	c2
a3	b3	c2
a3	b3	c2
a4	b1	c1
a4	b1	c1
a5	b2	c2

$\{B\} \rightarrow \{B, C\}$

A	B	C
a1	b1	c1
a1	b1	c1
a2	b2	c2
a2	b2	c2
a3	b3	c2
a3	b3	c2
a4	b1	c1
a4	b1	c1
a5	b2	c2

$\{C\} \rightarrow \{C\}$

Rules for Functional Dependencies (Armstrong's Axioms)

W, X, Y, Z mentioned here are sets of attributes.

Sound & complete inference system

- **Reflexivity**: If $Y \subseteq X$, then $X \to Y$.
- **Augmentation**: If $X \to Y$, then $X \cup Z \to Y \cup Z$ for any Z.
- **Transitivity**: If $X \to Y$ and $Y \to Z$, then $X \to Z$.

Common derived rules (from the three above)

- **Union / Additivity**: If $X \to Y$ and $X \to Z$, then $X \to Y \cup Z$.
- **Decomposition / Projectivity**: If $X \to Y \cup Z$, then $X \to Y$ and $X \to Z$.
- **Pseudotransitivity**: If $X \to Y$ and $W \cup Y \to Z$, then $W \cup X \to Z$.
- **Composition**: If $X \to Y$ and $Z \to W$, then $X \cup Z \to Y \cup W$.
- **Self-determination**: $X \to X$.

How we use them here

- *LHS reduction*: test if $X \setminus \{A\} \to Y$ holds via closure using the axioms.
- *Redundancy removal*: test if an FD $X \to Y$ is implied by the rest (i.e., $Y \subseteq X^+$ computed from the others).

Question 1.a.

Your company, Apasaja Private Limited, is commissioned by *Toko Kopi Luwak* to design the relational schema for coffee beans, drinks, and cafes. We are given

$$R = \{A, B, C, D, E, F, G, H\},$$

$\Sigma = \{\{A\} \rightarrow \{C, E\}, \{A, B\} \rightarrow \{D\}, \{F\} \rightarrow \{H\}, \{C, E\} \rightarrow \{A\}, \{B, C, E\} \rightarrow \{D\}, \{A, B, F\} \rightarrow \{D, G\}, \{B, C, E, F\} \rightarrow \{G\}\}$.

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Attribute mapping

Symbol	Meaning
A	`name` (bean brand)
B	`dname` (drink name)
C	`cultivar`
D	`price`

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Attribute mapping

Symbol	Meaning
A	name (bean brand)
B	dname (drink name)
C	cultivar
D	price

Attribute mapping

Symbol	Meaning
E	region
F	bname (branch)
G	qty (quantity sold)
H	address

Notes: A bean is identified by name or by (cultivar, region); a drink name is unique *per bean*; branches sell drinks, with address and quantity recorded.

Question 1.b. Candidate keys via closures

$\Sigma = \{\{A\} \rightarrow \{C, E\}, \{A, B\} \rightarrow \{D\}, \{F\} \rightarrow \{H\}, \{C, E\} \rightarrow \{A\}, \{B, C, E\} \rightarrow \{D\}, \{A, B, F\} \rightarrow \{D, G\}, \{B, C, E, F\} \rightarrow \{G\}\}$.

Observation: Since B and F never appear on the RHS of any FD, every key must contain B and F. We enumerate closures starting from supersets of $\{B, F\}$.

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Size 2

$\{B, F\}^+ = \{B, F, H\}$ (not a key)

Size 3 (supersets of $\{B, F\}$)

$\{A, B, F\}^+ =$
$\{A, B, C, D, E, F, G, H\}$
$\{B, C, F\}^+ = \{B, C, F, H\}$
$\{B, D, F\}^+ = \{B, D, F, H\}$
$\{B, E, F\}^+ = \{B, E, F, H\}$
$\{B, F, G\}^+ = \{B, F, G, H\}$
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Question 1.b. Candidate keys via closures

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$\{A, B, F\}^+ = \{A, B, C, D, E, F, G, H\}$

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$\{B, D, F\}^+ = \{B, D, F, H\}$

$\{B, E, F\}^+ = \{B, E, F, H\}$

$\{B, F, G\}^+ = \{B, F, G, H\}$

$\{B, F, H\}^+ = \{B, F, H\}$

Size 4 (not containing $\{A, B, F\}$)

$\{B, C, D, F\}^+ = \{B, C, D, F, H\}$

$\{B, C, E, F\}^+ = \{A, B, C, D, E, F, G, H\}$

$\{B, C, F, G\}^+ = \{B, C, F, G, H\}$

$\{B, C, F, H\}^+ = \{B, C, F, H\}$

$\{B, D, E, F\}^+ = \{B, D, E, F, H\}$

$\{B, D, F, G\}^+ = \{B, D, F, G, H\}$

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Question 1.b. Candidate keys via closures

$\Sigma = \{\{A\} \to \{C, E\}, \{A, B\} \to \{D\}, \{F\} \to \{H\}, \{C, E\} \to \{A\}, \{B, C, E\} \to \{D\}, \{A, B, F\} \to \{D, G\}, \{B, C, E, F\} \to \{G\}\}$.

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$\{A, B, F\}^+ = \{A, B, C, D, E, F, G, H\}$
$\{B, C, F\}^+ = \{B, C, F, H\}$
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Larger supersets (that are not supersets of keys above)
All such closures remain proper subsets of R.

Conclusion

Candidate keys: $\{A, B, F\}$ and $\{B, C, E, F\}$.

Tedious by hand?
I made your life a bit easier: Here

Question 1.c. Prime attributes

Keys are $\{A, B, F\}$ and $\{B, C, E, F\}$. Prime attributes are those that appear in at least one key.

$$\text{Prime attributes} = \{A, B, F\} \cup \{B, C, E, F\} = \{A, B, C, E, F\}.$$

Question 2.a. Minimal cover of R with Σ

Start from Σ:
$\{A\} \rightarrow \{C, E\}$,
{A,B} $\rightarrow \{D\}$,
{F} $\rightarrow \{H\}$,
{C,E} $\rightarrow \{A\}$,
{B,C,E} $\rightarrow \{D\}$,
{A,B,F} $\rightarrow \{D, G\}$,
{B,C,E,F} $\rightarrow \{G\}$.

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{B,C,E,F} $\rightarrow \{G\}$.

Step 1: Split RHS
$\{A\} \rightarrow \{C\}$,
{A} $\rightarrow \{E\}$,
{A,B} $\rightarrow \{D\}$,
{F} $\rightarrow \{H\}$,
{C,E} $\rightarrow \{A\}$,
{B,C,E} $\rightarrow \{D\}$,
{A,B,F} $\rightarrow \{D\}$,
{A,B,F} $\rightarrow \{G\}$,
{B,C,E,F} $\rightarrow \{G\}$.

Question 2.a. Minimal cover of R with Σ

Start from Σ:
{A} → {C, E},
{A,B} → {D},
{F} → {H},
{C,E} → {A},
{B,C,E} → {D},
{A,B,F} → {D, G},
{B,C,E,F} → {G}.

Step 1: Split RHS
{A} → {C},
{A} → {E},
{A,B} → {D},
{F} → {H},
{C,E} → {A},
{B,C,E} → {D},
{A,B,F} → {D},
{A,B,F} → {G},
{B,C,E,F} → {G}.

Step 2: Reduce LHS where possible
{A} → {C},
{A} → {E},
{A,B} → {D},
{F} → {H},
{C,E} → {A},
{B,C,E} → {D},
{A,B,$\cancel{F}$} → {D} because {A, B} → {D}
{A,B,F} → {G},
{B,C,E,F} → {G}.

Question 2.a. Minimal cover of R with Σ

Step 3: Remove redundancies

$\{A\} \rightarrow \{C\}$,
$\{\text{A}\} \rightarrow \{E\}$,
$\cancel{\{\text{A,B}\} \rightarrow \{D\}}$
Already exists (repetition),
$\{\text{F}\} \rightarrow \{H\}$,
$\{\text{C,E}\} \rightarrow \{A\}$,
$\{\text{B,C,E}\} \rightarrow \{D\}$,
$\cancel{\{\text{A,B}\} \rightarrow \{D\}}$
because $\{\text{A}\} \rightarrow \{C, E\}$ and $\{B, C, E\} \rightarrow \{D\}$,
$\cancel{\{\text{A,B,F}\} \rightarrow \{G\}}$
because $\{\text{A}\} \rightarrow \{C, E\}$ and $\{B, C, E, F\} \rightarrow \{G\}$,
$\{\text{B,C,E,F}\} \rightarrow \{G\}$.

Canonical Cover

$\{A\} \rightarrow \{C\}$,
$\{\text{A}\} \rightarrow \{E\}$,
$\{\text{F}\} \rightarrow \{H\}$,
$\{\text{C,E}\} \rightarrow \{A\}$,
$\{\text{B,C,E}\} \rightarrow \{D\}$,
$\{\text{B,C,E,F}\} \rightarrow \{G\}$.

Question 2.a. Minimal cover of R with Σ

Step 3: Remove redundancies

$\{A\} \rightarrow \{C\}$,

$\{A\} \rightarrow \{E\}$,

$\cancel{\{A,B\} \rightarrow \{D\}}$

Already exists (repetition),

$\{F\} \rightarrow \{H\}$,

$\{C,E\} \rightarrow \{A\}$,

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because $\{A\} \rightarrow \{C, E\}$ and $\{B, C, E\} \rightarrow \{D\}$,

$\cancel{\{A,B,F\} \rightarrow \{G\}}$

because $\{A\} \rightarrow \{C, E\}$ and $\{B, C, E, F\} \rightarrow \{G\}$,

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Canonical Cover

$\{A\} \rightarrow \{C\}$,

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$\{F\} \rightarrow \{H\}$,

$\{C,E\} \rightarrow \{A\}$,

$\{B,C,E\} \rightarrow \{D\}$,

$\{B,C,E,F\} \rightarrow \{G\}$.

(Alternative canonical covers exist, e.g., replace $\{B, C, E\} \rightarrow \{D\}$ by $\{A, B\} \rightarrow \{D\}$ and $\{B, C, E, F\} \rightarrow \{G\}$ by $\{A, B, F\} \rightarrow \{G\}$.)

Question 2.b. Compact Minimal Cover (merged by LHS)

One Compact Minimal Cover (grouping by identical LHS) is:

$\{A\} \rightarrow \{C, E\}$,
{F} $\rightarrow \{H\}$,
{C,E} $\rightarrow \{A\}$,
{B,C,E} $\rightarrow \{D\}$,
{B,C,E,F} $\rightarrow \{G\}$

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Other minimal covers are possible depending on which equivalent canonical cover you merge.

Questions?

Drop a mail at: pratik.karmakar@u.nus.edu