



Data Structure & Algorithms

Sunbeam Infotech



Stack and Queue

- Stack & Queue are utility data structures. → temp storage during process.

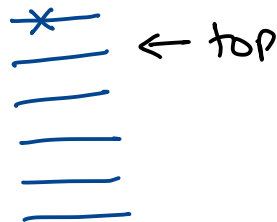
- Can be implemented using array or linked lists.

- Usually time complexity of stack & queue operations is $O(1)$.

- Stack is Last-In-First-Out structure.

- Stack operations → ADT

- ✓ push()
- ✓ pop()
- ✓ peek()
- ✓ isEmpty()
- isFull(*)

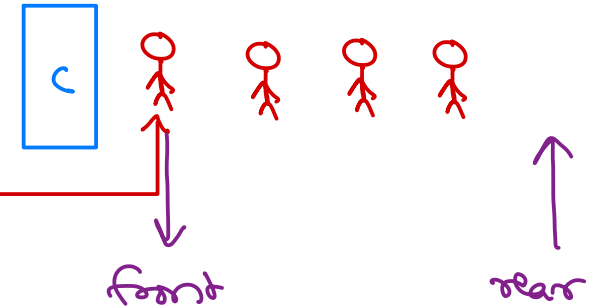


ops
done
from
same
end

- Simple queue is First-In-First-Out structure.

- Queue operations - ADT

- ✓ push() / enqueue()
- ✓ pop() / dequeue()
- ✓ peek()
- ✓ isEmpty()
- isFull(*)



- Queue types

- Linear queue
- Circular queue
- Deque
- Priority queue



Linear Queue

push \rightarrow rear
pop \leftarrow front

init:

$arr = \text{new int}[6];$

$f = -1;$

$r = -1;$

push:

$r++;$

$arr[r] = val;$

pop:

$f++;$

peek:

return $arr[f+1];$

is Full:

$r == \text{MAX} - 1$

is Empty:

$f == r$

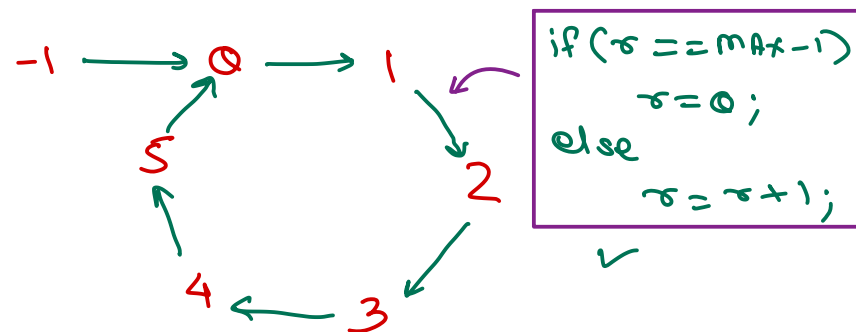
arr
-1

0	1	2	3	4	5
✓	✓	30	40	50	60

f

r

improper memory utilization

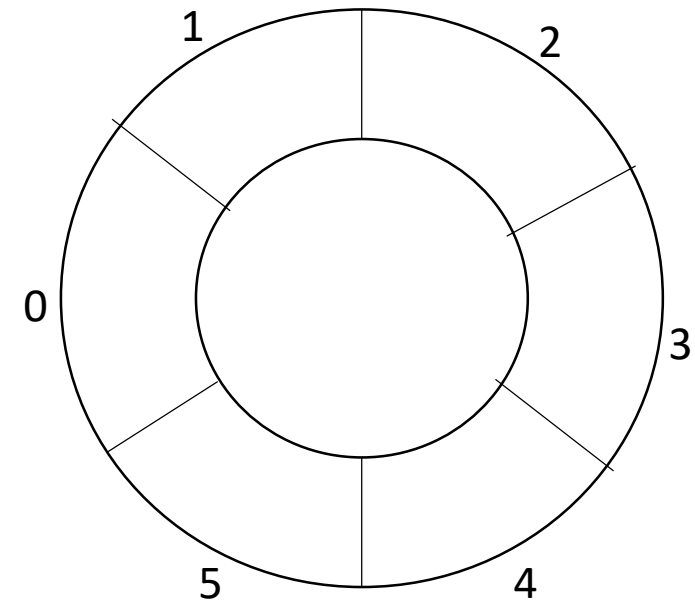


$r = (r + 1) \% \text{MAX};$

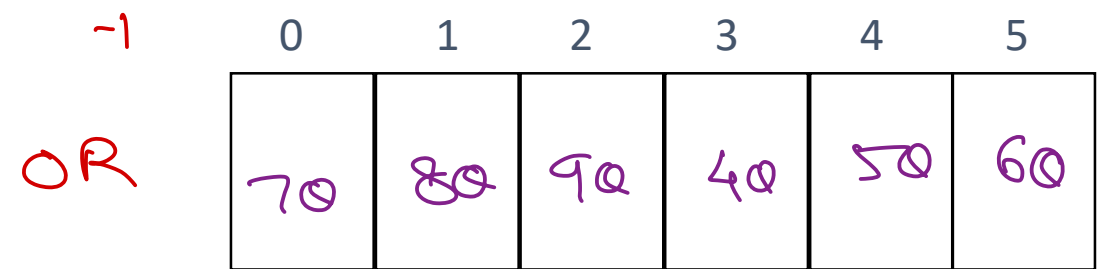
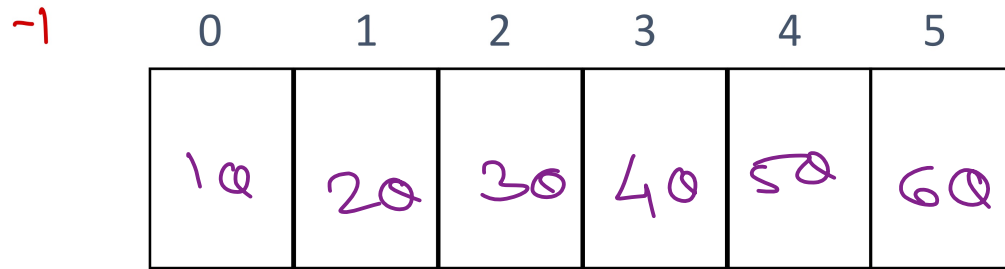


Circular Queue

- In linear queue (using array) when *rear* reaches last index, further elements cannot be added, even if space is available due to deletion of elements from *front*. Thus space utilization is poor.
- Circular queue allows adding elements at the start of array if *rear* reaches last index and space is free at the start of the array.
- Thus *rear* and *front* can be incremented in circular fashion i.e. 0, 1, 2, 3, ..., n-1. So they are said to be circular queue.
- However queue full and empty conditions become tricky.



Circular Queue - full & push



$f \&\& r$

f

$$(f == -1 \&\& r == \text{max} - 1) \quad || \quad (f == r \&\& f \neq -1)$$

push:

$$r = (r + 1) \% \text{MAX};$$

$$\text{arr}[r] = \text{val};$$



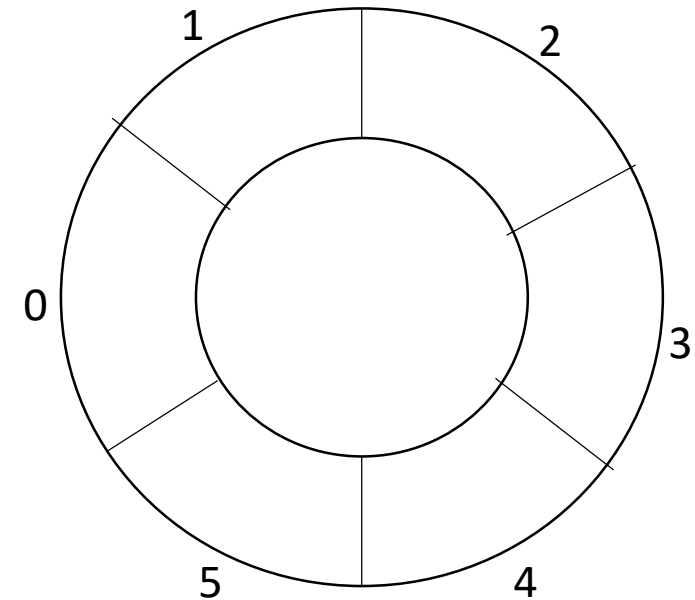
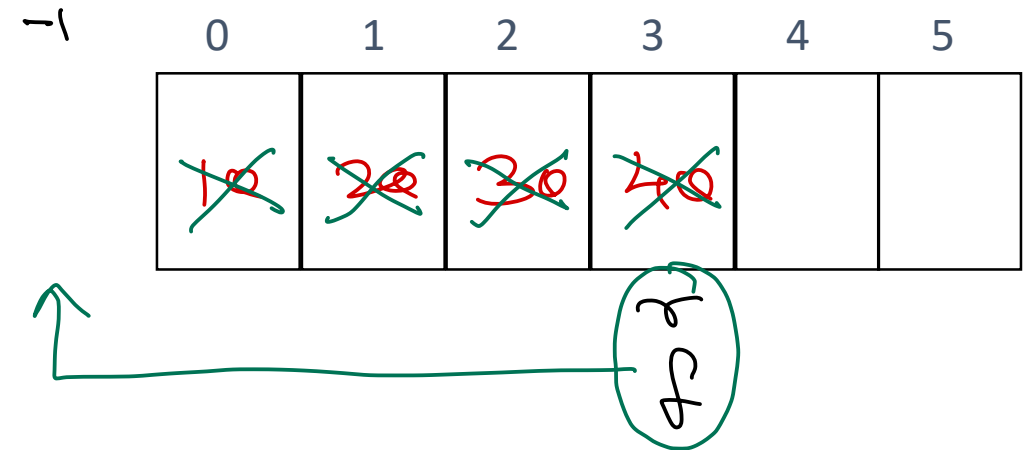
Circular Queue - empty & pop

$(r == f \ \&\& \ f == -1)$

pop:

$f = (f + 1) \% \text{max};$

$\text{if}(f == r)$
 $\{$
 $f = -1$
 $r = -1$
 $\}$

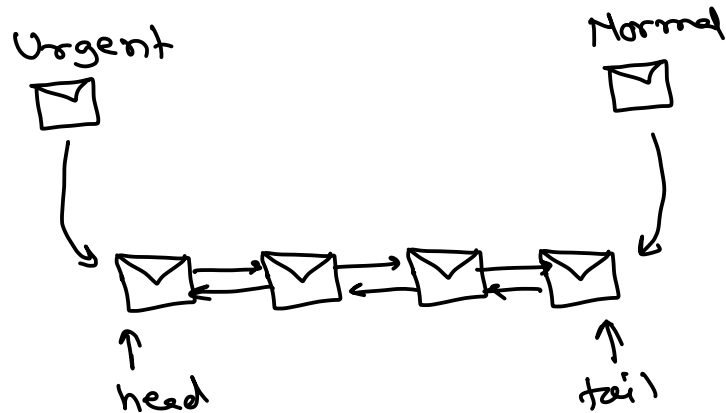


DeQueue

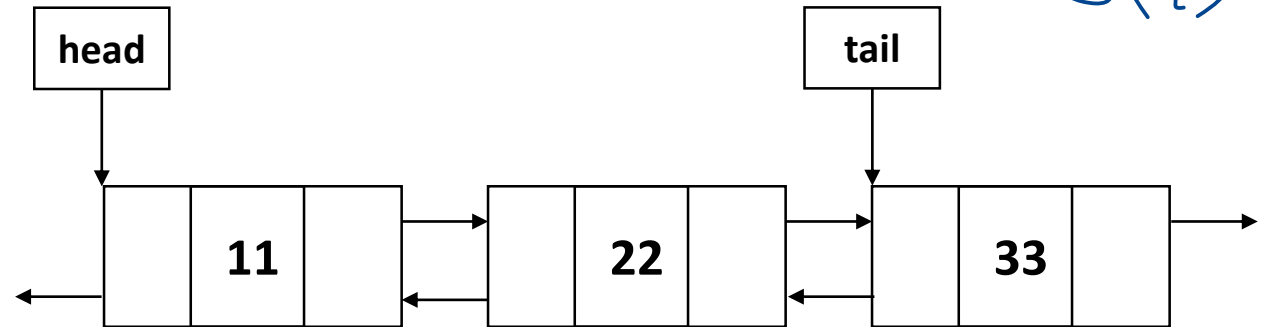
- In double ended queue, values can be added or deleted from front end or rear end.

applications

message queue in RTOS.



$O(1)$
push-front
pop-front
 $O(1)$



$O(1)$
push-back
pop-back
 $O(1)$



Priority queue → Not a FIFO queue

- In priority queue, element with highest priority is removed first.

↳ internally elements are stored in sorted manner.

Can be implemented

① using array → insertion logic.

② using linked list → insertion logic.

$$\text{push} = O(n)$$

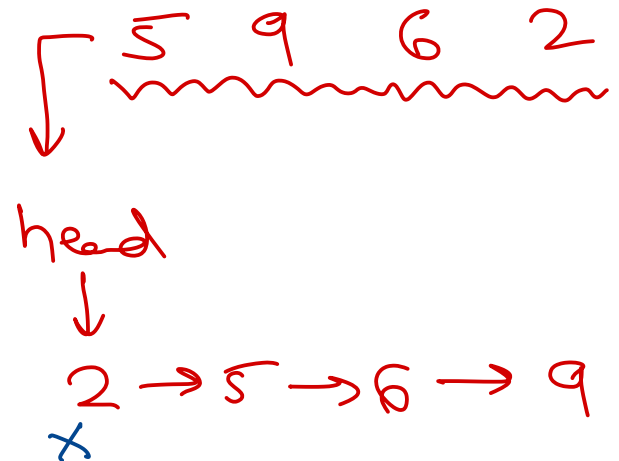
$$\text{pop} = O(1)$$

③ using heap (array representation of bin tree).

efficient

$$\text{push} = O(\log n)$$

$$\text{pop} = O(\log n)$$



Stack / Queue in Java collections

- class java.util.Stack<E>

- ✓ E push(E);
- ✓ E pop();
- ✓ E peek();
- ✓ boolean isEmpty();

- interface java.util.Queue<E>

- ✓ boolean offer(E e); - push
- ✓ E poll(); - POP
- ✓ E peek();
- ✓ boolean isEmpty();

ArrayDeque<> ✓

LinkedList<> ✓

PriorityQueue<>



Infix to Postfix

• $5 + 9 - 4 * (8 - 6 / 2) + 1 \$ (7 - 3)$

- ① traverse infix from left to right.
- ② if operand found, added to post fix.
- ③ if operator found, push to stack ④.
 - ④ if priority of topmost operator on stack \geq priority of current operator, pop it from stack & append to post fix.
- ④ if opening (found, push it on stack.
- ⑤ if closing) found, pop operators from stack & append to post fix until opening is found. also pop & discard opening (from stack.
- ⑥ if all syms from infix are completed, pop one by one & append to post fix.



Infix to Prefix

Stack

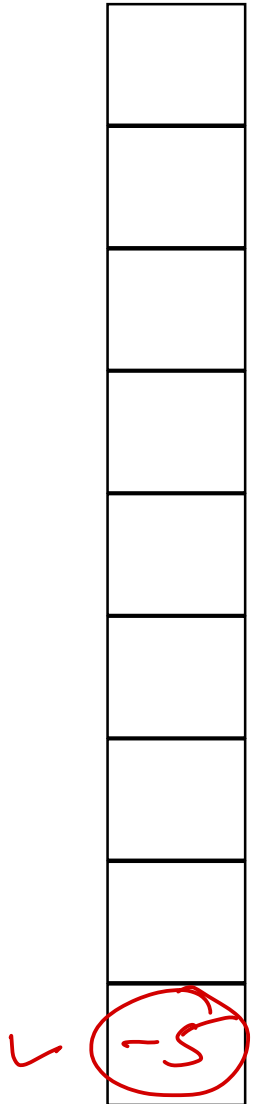
- $5 + 9 - 4 * (8 - 6 / 2) + 1 * (7 - 3)$



Prefix Evaluation

Stack

• + - + 5 9 * 4 - 8 / 6 2 \$ 1 - 7 3
↑

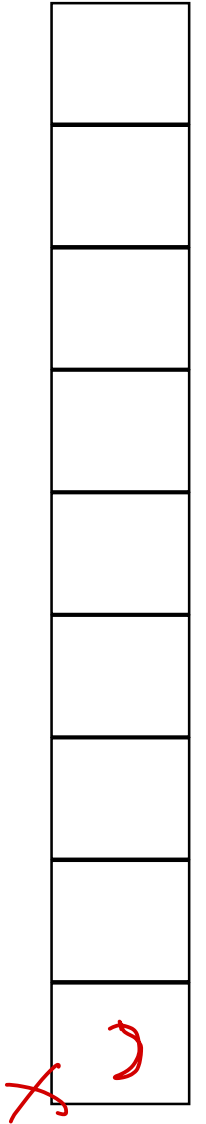


Parenthesis Balancing

• $5 + ([9 - 4] * (8 - \{6 / 2\}))$ ~~X~~
↑

	0	1	2
open	([{
close)]	}

0	0
()



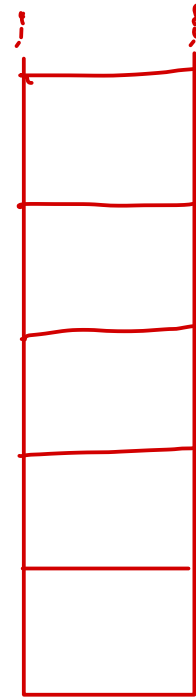
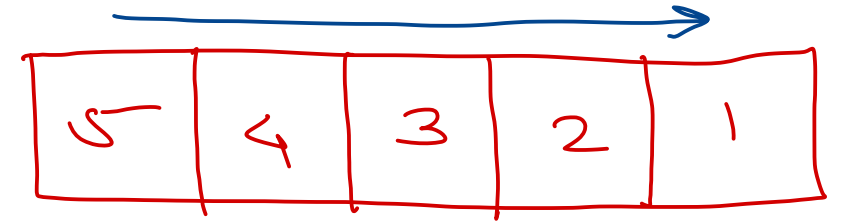
Stack / Queue – Competitive Programming

- Reverse array, string or linked list.

$\text{for}(i=0; i < n; i++)$
 $\quad s.\text{push}(a[i]);$ $\left| O(n) \right.$

$\text{for}(i=0; i < n; i++)$
 $\quad a[i] = s.\text{pop}();$ $\left| O(n) \right.$

$$O(\underline{2n}) = \underline{O(n)}$$



Stack / Queue – Competitive Programming

- Create two stacks in single array in efficient manner.

init 1:

$top1 = -1$

is Empty 1:

$top1 == -1$

push 1():

$top1++$;

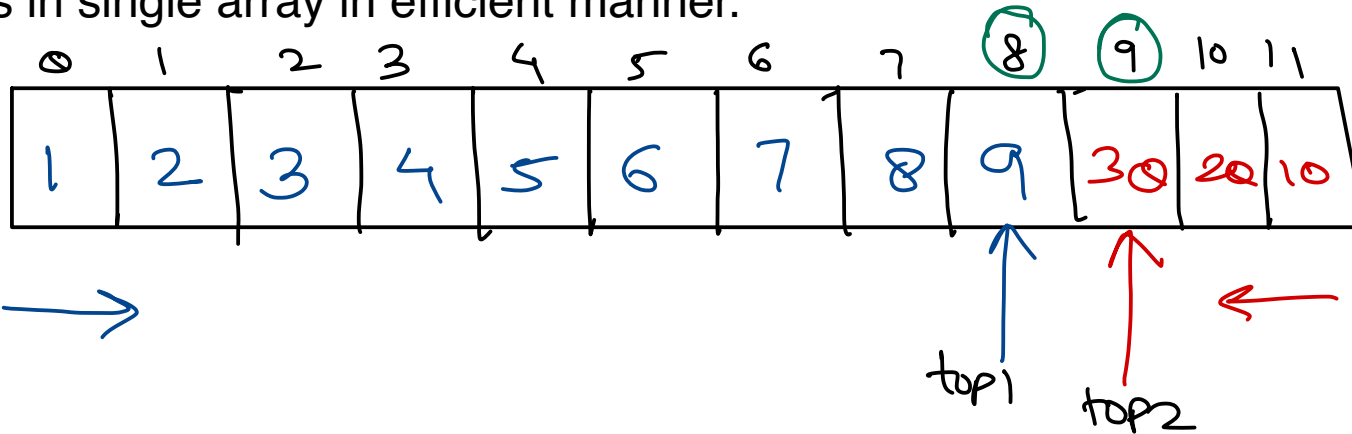
$arr[top1] = val$;

pop 1():

$top1--$;

is Full 1:

$top1 + 1 == top2$



✓ $top1 + 1 == top2$

✓ $top1 == top2 - 1$

✓ $top2 - top1 == 1$

init 2:

$top2 = arr.length$

is Empty 2:

$top2 == arr.length$;

push 2():

$top2--$;

$arr[top2] = val$;

pop 2():

$top2++$

is Full 2:

$top1 + 1 == top2$



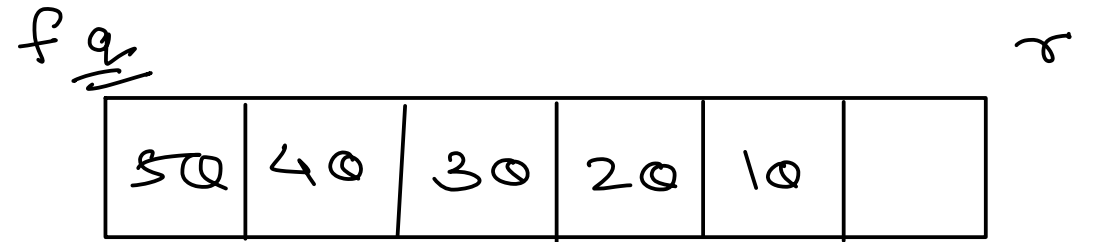
Stack / Queue – Competitive Programming

- Create stack using queue.

↓
LIFO

↓
FIFO

sq →



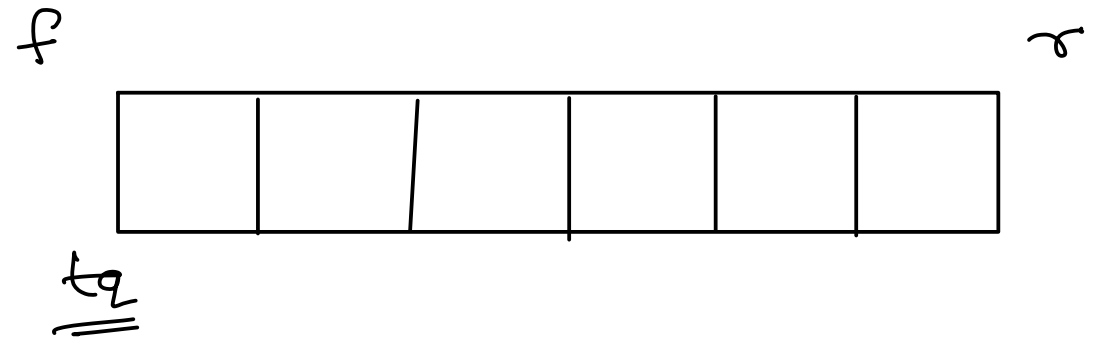
✓ push → $O(n)$

(tq)

```
while (!q.isEmpty()) {  
    tq.push(q.pop());  
}
```

```
q.push(val);
```

```
while (!tq.isEmpty()) {  
    q.push(tq.pop());  
}
```



pop → $O(1)$
q.pop()



Stack / Queue – Competitive Programming

- Create queue using stack.





Thank you!

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