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CS 402 - ADVANCED STUDIES - II
HOMEWORK 2

1.5

[4] < S 1.6 > Consider 3 different processors P₁, P₂ & P₃ executing the same instruction set. P₁ has a 3 GHz clock rate & a CPI of 1.5. P₂ has a 2.5 GHz clock rate & a CPI of 1.0. P₃ has a 4.0 GHz clock rate & has a CPI of 2.2

- (a) Which processor has the highest performance expressed in instructions per second?
- (b) If the processors each execute a program in 10 sec, find the no. of cycles & the no. of instructions.
- (c) We are trying to reduce the execution time by 30% but this leads to an increase of 20% in the CPI. What clock rate should we have to get this time reduction?

⇒

(a)

$$\text{CPU time}(t) = (\text{Instruction count}) * \frac{(\text{clock cycle})}{\text{time}} * (\text{CPI})$$
$$= \frac{\text{Instruction count}}{\text{clock rate}} * \text{CPI}$$

$$t_1 = \frac{1 * 1.5}{3 * 10^9} = 5 * 10^{-10}$$

$$t_2 = \frac{1 * 1}{2.5 * 10^9} = 4 * 10^{-10}$$

$$t_3 = \frac{1 * 2.2}{4.0 * 10^9} = 5.5 * 10^{-10}$$

For performance - take inverse of the CPU time.

$$P_1 = \frac{1}{t_1} = \frac{1}{5 \times 10^{-10}} = 2 \times 10^9 \text{ instructions/sec.}$$

$$P_2 = \frac{1}{t_2} = \frac{1}{4 \times 10^{-10}} = 2.5 \times 10^9 \text{ instructions/sec}$$

$$P_3 = \frac{1}{t_3} = \frac{1}{5.5 \times 10^{-10}} \approx 1.82 \times 10^9 \text{ instructions/sec}$$

$$P_2 > P_1 > P_3$$

∴ P_2 has the highest performance.

(b) $P_{cycles} = \text{Clock Rate} \times t$

$$P_{ins} = \frac{P_{cycles}}{\text{CPI}}$$

$$P_1 \text{ cycles} = 3 \times 10^9 \times 10 = 3 \times 10^{10}$$

$$P_2 \text{ cycles} = 2.5 \times 10^9 \times 10 = 2.5 \times 10^{10}$$

$$P_3 \text{ cycles} = 4.0 \times 10^9 \times 10 = 4.0 \times 10^{10}$$

$$P_1 \text{ ins} = \frac{3 \times 10^{10}}{1.5} = 2 \times 10^{10}$$

$$P_2 \text{ ins} = \frac{2.5 \times 10^{10}}{1} = 2.5 \times 10^{10}$$

$$P_3 \text{ ins} = \frac{4.0 \times 10^{10}}{2.2} \approx 1.82 \times 10^{10}$$

(c) $P_{CR} = \frac{P_{ins} \times 1.2 \times \text{CPI}}{0.7 \times t}$

$$P_{CR} = \frac{2 \times 10^{10} \times 1.2 \times 1.5}{0.7 \times 10} \approx 5.14 \times 10^9$$

$$P_{2CR} = \frac{2.5 \times 10^{10} \times 1.2 \times 1}{0.7 \times 10}$$

$$\approx 4.29 \times 10^9$$

$$P_{3CR} = \frac{1.82 \times 10^{10} \times 1.2 \times 2.2}{0.7 \times 10}$$

$$\approx 6.86 \times 10^9$$

3.6 Consider two different implementations of the same instruction set architecture. The instructions can be divided into four classes according to their CPI (class A, B, C & D).

P₁ with a clock rate of 2.5 GHz & CPIs of 1, 2, 3 & 3 & P₂ with a clock rate of 3 GHz & CPIs of 2, 2, 2 & 2.

Given a program with a dynamic instruction count of 1.0E6 instructions divided into classes as follows : 10% class A, 20% class B, 50% class C & 20% class D, which implementation is faster?

- (a) What is the global CPI for each implementation?
- (b) Find the clock cycles required in both cases.

⇒ Proportioning the instructions with their respective percentages

$$I_A = (0.1) \times (1 \times 10^6) = 1 \times 10^5$$

$$I_B = (0.2) \times (1 \times 10^6) = 2 \times 10^5$$

$$I_C = (0.5) \times (1 \times 10^6) = 5 \times 10^5$$

$$I_D = (0.2) \times (1 \times 10^6) = 2 \times 10^5$$

(a) To find global CPI for each processor, first we will have to calculate its time

$$t_N = \frac{I_N \times CPI_N}{CR_n}$$

$N \rightarrow$ various class

$n \rightarrow$ processor number

for 1st processor:

$$t_A = \frac{1 \times 10^5 \times 1}{2.5 \times 10^9} = 4 \times 10^{-5} \text{ sec}$$

$$t_B = \frac{2.5 \times 10^5 \times 2}{2.5 \times 10^9} = 1.6 \times 10^{-4} \text{ sec}$$

$$t_C = \frac{5 \times 10^5 \times 3}{2.5 \times 10^9} = 6 \times 10^{-4} \text{ sec}$$

$$t_D = \frac{2 \times 10^5 \times 3}{2.5 \times 10^9} = 2.4 \times 10^{-4} \text{ sec}$$

$$CPI_{\text{global}} = \frac{CR_n \cdot \sum_{i=0}^N t_i}{I_{\text{total}}}$$

$$CPI_{\text{global}} = \frac{(0.4 + 1.6 + 6 + 2.4) \times 10^{-4} \times 2.5 \times 10^9}{1 \times 10^6}$$

$$CPI_{\text{global}} = 2.6$$

for 2nd processor:

$$t_A = \frac{1 \times 10^5 \times 2}{3 \times 10^9} = \frac{2}{3} \times 10^{-4} \text{ sec}$$

$$t_B = \frac{2 \times 10^5 \times 2}{3 \times 10^9} = \frac{4}{3} \times 10^{-4} \text{ sec}$$

$$t_c = \frac{5 \times 10^5 \times 2}{3 \times 10^9} = \frac{10}{3} \times 10^{-4} \text{ sec}$$

$$t_D = \frac{2 \times 10^5 \times 2}{3 \times 10^9} = \frac{4}{3} \times 10^{-4} \text{ sec}$$

$$CPI_{2 \text{ global}} = \frac{(2/3 + 4/3 + 10/3 + 4/3) \times 10^{-4} \times 3 \times 10^9}{1 \times 10^6}$$

$$CPI_{2 \text{ global}} = 2$$

$$(b) P_{\text{clock cycles}} = \sum_{i=0}^N I_{N,i} \cdot CPI_{N,i}$$

$$P_1 \text{ clock cycles} = (1 \times 10^5)(1) + (2 \times 10^5)(2) + (5 \times 10^5)(3) \\ + (2 \times 10^5)(3)$$

$$P_1 \text{ clock cycles} = 2.6 \times 10^6$$

$$P_2 \text{ clock cycles} = (1 \times 10^5)(2) + (2 \times 10^5)(2) + (5 \times 10^5)(2) \\ + (2 \times 10^5)(2)$$

$$P_2 \text{ clock cycles} = 2 \times 10^6$$

1.7 Compilers can have a profound impact on the performance of an application. Assume that for a program, compiler A results in a dynamic instruction count of 1.0×10^9 & has an execution time of 1.1 sec, while compiler B results in a dynamic instruction count of 1.2×10^9 & an execution time of 1.5 sec.

(a) Find the avg. CPI for each program given that the processor has a clock cycle time of 1 ns .

- (b) Assume the compiled programs run on two different processors. If the execution time on the 2 processors are the same, how much faster is the clock of the processor running compiler A's code versus the clock of the processor running compiler B's code?
- (c) A new compiler is developed that uses only 6.0E8 instructions & has an avg. CPI of 1.1. What is the speed up of using this new compiler versus using compiler A or B on the original processor?

⇒ (a)

$$CPI = t_{exec} \cdot$$

clock cycles × Instruction count

$$CPI_A = 1.1 \times \frac{1}{10^{-9} \times 10^9} = 1.1$$

$$CPI_B = 1.5 \times \frac{1}{1.2 \times 10^{-9} \times 10^9} = 1.25$$

(b) To compare two compilers

$$\frac{1.2 \times 10^9 \times 1.25}{10^9 \times 1.1} \approx 1.37$$

Compiler B is approximately 1.37 times faster than compiler A.

e (c) To compare with a new compiler

$$\frac{T_A}{T_{\text{new}}} = \frac{1.1 \times 1 \times 10^9}{1.1 \times 6 \times 10^8} \approx 1.67$$

$$\frac{T_B}{T_{\text{new}}} = \frac{1.025 \times 1.2 \times 10^9}{1.1 \times 6 \times 10^8} \approx 2.27$$

The speed up with regards to compiler A is approximately 1.67 & for compiler B is approximately 2.27.

1.9

1.9.1

$$t_N = I_N \cdot CPI_N$$

CR_n

$$t_1 = \frac{\left(\frac{2.56 \times 10^9}{0.7} \times 1 \right) + \left(\frac{1.28 \times 10^9}{0.7} \times 12 \right) + \left(2.56 \times 10^8 \times 5 \right)}{2 \times 10^9}$$

$$t_1 = 13.44 \text{ sec}$$

$$t_2 = \frac{\left(\frac{2.56 \times 10^9}{0.7 \times 2} \times 1 \right) + \left(\frac{1.28 \times 10^9}{0.7 \times 2} \times 12 \right) + \left(2.56 \times 10^8 \times 5 \right)}{2 \times 10^9}$$

$$t_2 = 7.04 \text{ sec}$$

$$t_4 = \frac{\left(\frac{2.56 \times 10^9}{0.7 \times 4} \times 1 \right) + \left(\frac{1.28 \times 10^9}{0.7 \times 4} \times 12 \right) + \left(2.56 \times 10^8 \times 5 \right)}{2 \times 10^9}$$

$$t_4 = 3.84 \text{ sec}$$

$$t_8 = \frac{\left(\frac{2.56 \times 10^9}{0.4 \times 8} \times 1 \right) + \left(\frac{1.28 \times 10^9}{0.4 \times 8} \times 12 \right) + \left(2.56 \times 10^8 \times 5 \right)}{2 \times 10^9}$$

$$t_8 = 2.24 \text{ sec}$$

$$S_1 = 1$$

$$S_2 = 1.91$$

$$S_4 = 3.5$$

$$S_8 = 6$$

1.9.2 Doubling the CPI for the arithmetic value,

$$t_1 = 15.27 \text{ sec}$$

$$t_2 = 7.95 \text{ sec}$$

$$t_4 = 4.30 \text{ sec}$$

$$t_8 = 2.47 \text{ sec}$$

1.9.3 It should be reduced to 3.

1.11

$$1.11.1 \text{ Clock Rate} = \frac{1}{\text{Cycle Time}} = \frac{1}{0.333 \times 10^{-9}} \\ \approx \frac{1}{3 \times 10^9}$$

$$\text{CPI (bit 2)} = 3 \times 10^9 \times \frac{750}{2389 \times 10^9} \\ \approx 0.94$$

\therefore Required CPI is approx. 0.94

$$1.11.2 \text{ SPECratio} = \frac{t_r}{J_N} = \frac{9650}{750} \approx 12.87$$

\therefore SPECratio is approx. 12.87

1.11.3 We know that,

$$\begin{aligned} \text{CPU time} &= (\underbrace{\text{Instruction count}}_{\text{Instruction}}) \times (\underbrace{\text{CPI}}_{\text{time}}) \times (\underbrace{\text{Clock cycle time}}_{\text{Clock Rate}}) \\ &= \frac{\text{Instruction count} \times \text{CPI}}{\text{Clock Rate}} \end{aligned}$$

If we assume CPI & clock rate to be constant, there should be a 10% increase in CPU time.

$$\begin{aligned} 1.11.4 \frac{\text{CPU time (after)}}{\text{CPU time (before)}} &= \left(\frac{1.1 \times \text{Instruction count} \times 1.05 \times \text{CPI}}{\text{Clock rate}} \right) \\ &\quad \times \left(\frac{\text{Clock Rate}}{\text{Instruction count} \times \text{CPI}} \right) \\ &= 1.1 \times 1.05 \\ &= 1.155 \end{aligned}$$

∴ The CPU time increases by 15.5%

$$\begin{aligned} 1.11.5 \frac{\text{SPEC ratio (after)}}{\text{SPEC ratio (before)}} &= \left(\frac{t_r}{\text{CPU time (after)}} \right) \times \left(\frac{\text{CPU time (before)}}{t_r} \right) \\ &= \frac{\text{CPU time (after)}}{\text{CPU time (before)}} \\ &= \frac{1}{1.155} \\ &\approx 0.87 \end{aligned}$$

∴ The SPEC ratio decreased by 13%

1.11.6

$$CPI = \frac{CPU\ time \times Clock\ rate}{Instruction\ count}$$

$$= \frac{700 \times 4 \times 10^9}{0.85 \times 2389 \times 10^9} \approx 1.38$$

∴ CPI is approx. 1.38

1.11.7

$$\frac{4}{3} \approx 1.33$$

$$\frac{1.38}{0.94} \approx 1.47$$

These 2 values are different due to the fact that although the number of instructions were reduced by 15%, the CPU time was reduced by a lower percentage.

1.11.8

$$\frac{700}{750} \approx 0.933$$

750

CPU time reduced by 6.4%

1.11.9

$$Instruction\ count = \frac{CPU\ time \times Clock\ Rate}{CPI}$$

$$= \frac{960 \times 0.9 \times 4 \times 10^9}{1.61}$$

$$\approx 2.147 \times 10^{12}$$

∴ Instruction count is approx. 2.147×10^{12}

1.11.10 $\text{Clock Rate}_{(\text{new})} = \frac{\text{Instruction count} \times \text{CPI}}{\text{CPU time}}$

Rearranging,

$$\text{Clock Rate}_{(\text{new})} = \frac{\text{Clock Rate}_{(\text{old})}}{0.9}$$
$$\approx 3.33 \times 10^9$$

\therefore New Clock Rate is 3.33 GHz

1.11.11 $\text{Clock Rate}_{(\text{new})} = \frac{0.85}{0.80} \times \text{Clock Rate}_{(\text{old})}$

$$\approx 3.19 \times 10^9$$

\therefore New Clock Rate is approx. 3.19 GHz

~~1.11.12~~

1.12

1.12.1 $\text{Clock Rate} = \frac{\text{Instruction count} \times \text{CPI}}{\text{CPU time}}$

$$t_{P1} = \frac{5 \times 10^9 \times 0.9}{4 \times 10^9} = 1.125 \text{ sec}$$

$$t_{P2} = \frac{1 \times 10^9 \times 0.45}{3 \times 10^9} = 0.25 \text{ sec}$$

Clock rate for processor 1 is larger than processor 2 but processor 2 has better performance than processor 1.

\therefore Statement isn't true

$$1.12.2 \quad t_{P_1} = t_{P_2}$$

$$t_{P_1} = \frac{\text{No. of Instructions}}{\text{Clock rate}} \times CPI$$

$$t_{P_1} = 2.53 \times 3$$

$$t_{P_1} = \frac{1 \times 10^9 \times 0.9}{4 \times 10^9} = 0.225$$

Number of instructions executed in this time frame from processor 2.

$$N_2 = \frac{0.225 \times 3 \times 10^9}{0.75} = 9 \times 10^8$$

$$\therefore N_2 = 9 \times 10^8$$

1.12.3

$$\text{MIPS} = \frac{\text{Clock rate}}{\text{CPI}} \times 10^{-6}$$

$$\text{MIPS}_1 = \frac{4 \times 10^9 \times 10^{-6}}{0.9} \approx 4.44 \times 10^3$$

$$\text{MIPS}_2 = \frac{3 \times 10^9 \times 10^{-6}}{0.75} = 4 \times 10^3$$

\therefore The MIPS value is greater in processor 1, yet we still have the performance to be better for processor 2.

1.12.4

$$\text{MFLOPS} = \frac{\text{No. of FP operations}}{\text{texec} \times 1 \times 10^6}$$

$$\text{MFLOPS}_1 = \frac{0.4 \times 5 \times 10^9}{1.125 \times 10^6} \\ \approx 1.78 \times 10^3$$

$$\text{MFLOPS}_2 = \frac{0.4 \times 1 \times 10^9}{0.25 \times 10^6} \\ = 1.6 \times 10^3$$

∴ MFLOPS₁ is greater than MFLOPS₂, yet we still have the performance to be better for processor 2.

1.13

1.13.1

Reducing FP operations by 20%, we get

$$t_{FP} = 70 \times 0.8 = 56 \text{ sec}$$

$$t_{\text{total}} = 56 + 85 + 40 + 55 = 236 \text{ sec}$$

∴ We get a 5.6% reduction from reducing the FP operations by 20%

1.13.2

Taking total time & reducing it by 20%

$$t_{\text{new}} = 236 \times 0.8 \\ = 200 \text{ sec}$$

$$t_{\text{new}} - (t_{\text{FP}} + t_{\text{L/S}} + t_{\text{branch}}) = 200 - 165 \\ = 35 \text{ sec}$$

\therefore We get a 58.8% reduction for the INT operations.

1.13.3

$$t_{\text{new}} = 200 \text{ sec}$$

$$t_{\text{branch}} = 0$$

$$t_{\text{FP}} + t_{\text{L/S}} + t_{\text{INT}} = 210 \text{ sec} > 200 \text{ sec}$$

Since this is greater than the threshold this is not possible.

1.14

$$\begin{aligned} 1.14.1 \text{ Clock cycles} &= (\text{CPI}_{\text{FP}} \times N_{\text{FP}}) + (\text{CPI}_{\text{INT}} \times N_{\text{INT}}) \\ &\quad + (\text{CPI}_{\text{L/S}} \times N_{\text{L/S}}) + (\text{CPI}_{\text{branch}} \times N_{\text{branch}}) \end{aligned}$$

$$t_{\text{cpu}} = \frac{512 \times 10^6}{2 \times 10^9} = 0.256 \text{ sec}$$

Do we want to run the program twice as fast with only changing the CPI of the FP instructions.

$$\text{CPI}_{\text{new FP}} = \frac{256 - 462}{50} = -4.12$$

\therefore Since the number would be negative, this is impossible.

$$1.14.2 \quad CPI_{new LS} = \frac{256 - 198}{80} \\ = 0.725$$

To run the program twice as fast, we would need L/S CPI to be 0.725

$$1.14.3 \quad CPI_{IFP} = CPI_{INT} = 0.6 \times 1 = 0.6$$

$$CPI_{LS} = 0.7 \times 4 = 2.8$$

$$CPI_{branch} = 0.7 \times 2 = 1.4$$

$$t_{cpu} = 0.171 \text{ sec}$$

\therefore we get 1.5 times speed up from the CPI improvements

1.15

Data should be below, where primes are the ideal number

$$\begin{aligned} t_1 &= 100 \text{ sec} \\ t_2 &= 54 \text{ sec} \\ t_4 &= 29 \text{ sec} \\ t_8 &= 16.5 \text{ sec} \\ t_{16} &= 10.25 \text{ sec} \\ t_{32} &= 7.125 \text{ sec} \\ t_{64} &= 5.5625 \text{ sec} \\ t_{128} &= 4.78125 \text{ sec} \end{aligned}$$

$$\begin{aligned} t'_1 &= 100 \text{ sec} \\ t'_2 &= 50 \text{ sec} \\ t'_4 &= 25 \text{ sec} \\ t'_8 &= 12.5 \text{ sec} \\ t'_{16} &= 6.25 \text{ sec} \\ t'_{32} &= 3.125 \text{ sec} \\ t'_{64} &= 1.5625 \text{ sec} \\ t'_{128} &= 0.78125 \text{ sec} \end{aligned}$$

Speed ups for these values would be

$$S_1 = \frac{100}{100} = 1$$

$$S'_1 = \frac{100}{100} = 1$$

$$S_2 = \frac{100}{54} = 1.85$$

$$S'_2 = \frac{100}{SD} = \underline{0.93} 2$$

$$S_4 = \frac{100}{29} = 3.44$$

$$S'_4 = \frac{100}{25} = 4$$

$$S_8 = \frac{100}{16.5} = 6.06$$

$$S'_8 = \frac{100}{12.5} = 8$$

$$S_{16} = \frac{100}{10.25} = 9.76$$

$$S'_{16} = \frac{100}{6.25} = 16$$

$$S_{32} = \frac{100}{7.125} = 14.03$$

$$S'_{32} = \frac{100}{3.125} = 31.007 \approx 32$$

$$S_{64} = \frac{100}{5.5625} = 17.98$$

$$S'_{64} = \frac{100}{1.5625} = 64$$

$$S_{128} = \frac{100}{4.78125} = 20.91$$

$$S'_{128} = \frac{100}{0.78125} = 128$$

∴ Required Ratio would be

$$\frac{S_1}{S'_1} = 1$$

$$\frac{S_{16}}{S'_{16}} = 0.61$$

$$\frac{S_2}{S'_2} = 0.93$$

$$\frac{S_{32}}{S'_{32}} = 0.44$$

$$\frac{S_4}{S'_4} = 0.86$$

$$\frac{S_{64}}{S'_{64}} = 0.28$$

$$\frac{S_8}{S'_8} = 0.75$$

$$\frac{S_{128}}{S'_{128}} = 0.16$$