

Assignment 7

Title: Symbol Table Hashing

Problem Statement:

It is generated by compiler. From this perspective a symbol table is a set of name-attribute pairs. Perform the following operations on symbol table:

- 1) Determine if particular name is in symbol table.
- 2) Retrieve the attributes of that name.
- 3) Modify the attributes of that name.
- 4) Insert a new name and its attributes.
- 5) Delete a name & its attributes.

Objective:

To understand the concept of chaining used in Hash tables.

Outcome:

It will be able to implement symbol table structure using hashing & chaining with & without replacement.

Requirements:

- 1) 64 bit machine, operating system.
- 2) Editor, Compiler (g++)
- 3) CPU, RAM

Theory:

In linear probing technique if a large number of keys are being mapped to the same location then searching time complexity no longer remains $O(1)$ but it can take $O(n)$ time for search. Chaining technique tries to reduce time. In chaining all keys which get the same location on applying hash function are stored in chain.

• Pseudo Code:

• add (symbol a)
 if (replac)

 if $\text{hash}(s[b].k) \neq b$ and $s[b]$ non Empty
 swap($s[b]$, a)

 else if $s[b]$ non empty

 for $c = b$ until $c = -1$

$c = s[c].\text{chain}$

l :- First empty loc after b

 if (no loc.)

 error

 else

$s[l] := a$

$s[c].\text{chain} := l$

else from $l := b$

while ($s[l]$ not empty $\{ \text{hash } s[x].k \neq b \} \{$

if (all loc. non empty)

print 'error' & return

$\}$

if $s[l] = \text{empty}$

$s[l] = a$ ret

if $c := l$ until $s[c].\text{chain} = -1$

$c := s[c].\text{chain}$

$m = \text{next empty after } c$

if (not found)

print 'error', return

else

$s[m] := a$

$s[c].\text{chain} := m$

• retrieve (symbol a)

$b := \text{find}(a.k)$

if ($b = -1$)

print "not found"

else

Read attr. of $s[b]$

print attr.

return.

- update (symbol a)

b := find (a.key)

if b = -1

print "Not found" & return

else

read attr of s[b] from user

return.

- delete (symbol a)

b := find (a.key)

if (b = -1)

print (not found) ret

else

l = loc. s.t (s[l].chain = b)

if (loc = found)

s[l].chain = s[b].chain

make s[b] = empty

return.

* Problems. faced

1. Probing without replacement
2. Collision resolution.

- Test cases:

	Input	Operation	Output	Result
1)	Abc, Ade, Bcd, Bdc, Abcd, Def (without Replacement)	find (Abc)	loc: 0	Success
		find (Bcd)	loc: 2	
		find (Def)	loc: 5	
		find (Abcd)	loc: 4	
		delete (Ade)	deleted	

- Conclusion:

Symbol table is implemented using hashing using chaining with & without replacement.