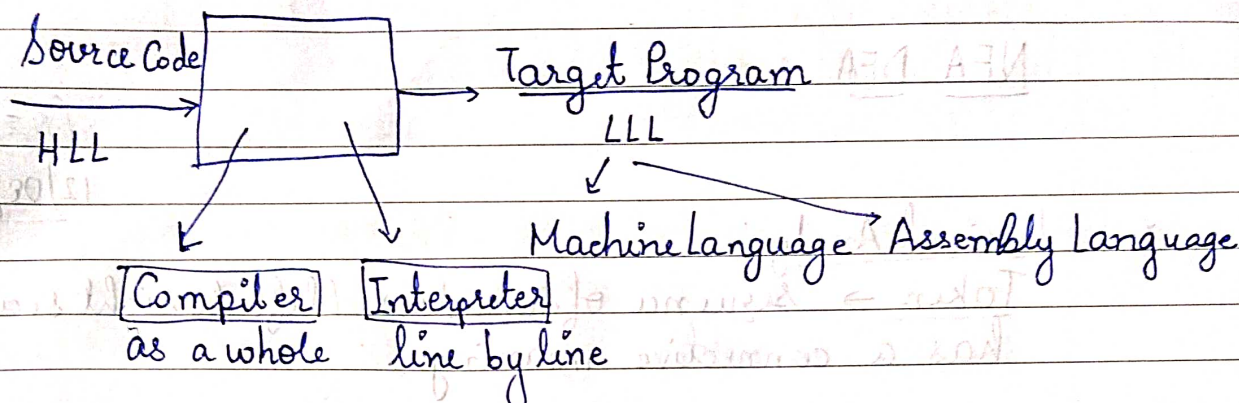


Compiler Design

Interpreter



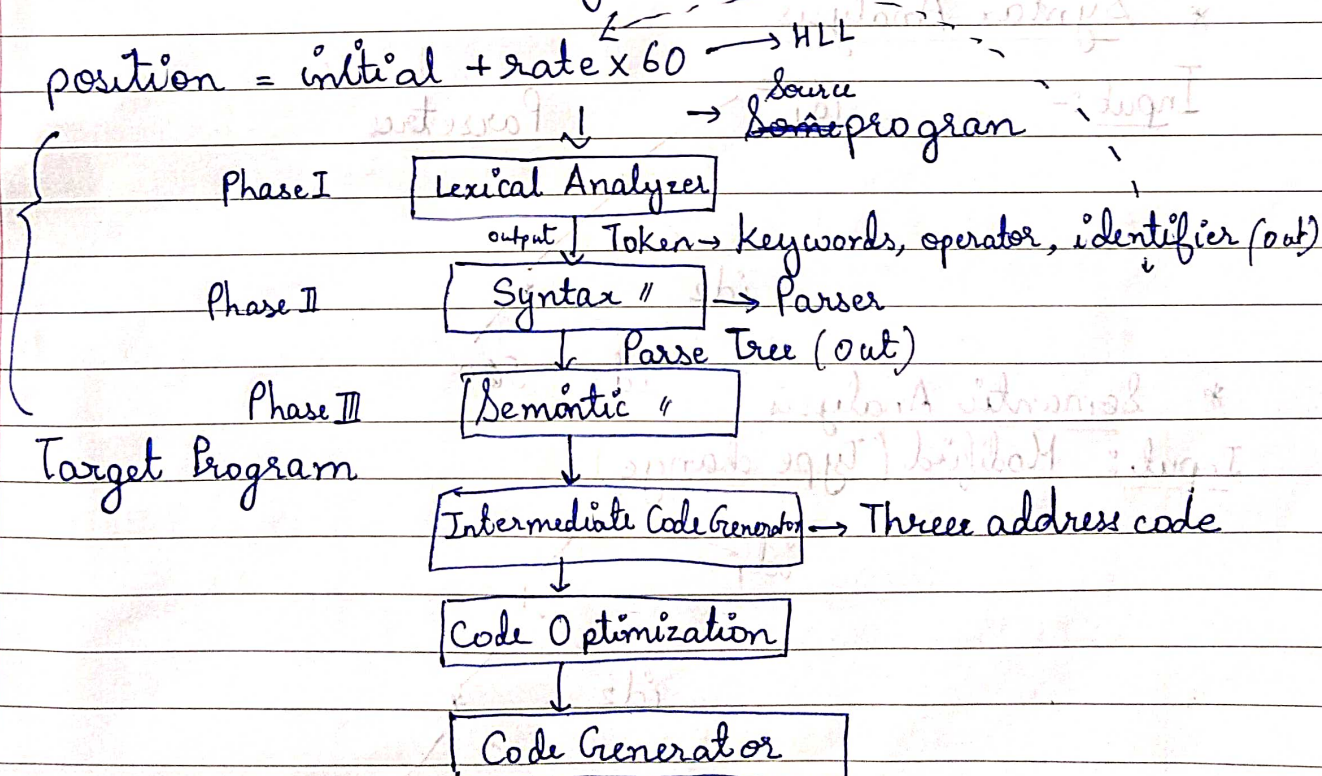
Advantage: Debugging is easier in Interpreter.

Disadvantage: Time consuming → Interpreter.

Assembly Language: converts assembly language to machine level language. (Low-level → Low level)

We will learn compiling here (not execution).

$$\text{position} = \text{initial} + \text{rate} \times 60$$



Compiler Design
Aho, Ullman, Sethi

Compiler
Sridha

NFA DFA

(Q, Σ , δ , F)
12/08/2022

* Lexical Analysis

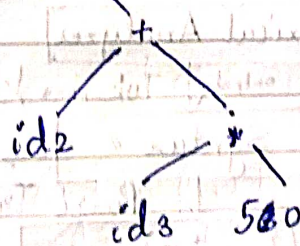
Token \rightarrow Sequence of characters (left to right scan) that has a connective meaning.

Inputs: Total = num + num1 * 50
 $\downarrow \quad \downarrow \quad \downarrow \quad \downarrow \quad \downarrow \quad \downarrow$
 id1 id2 op1 id3 op2 50

Outputs: id1 = id2 + id3 * 50 (presented in tokens)
 (stored in symbol table)

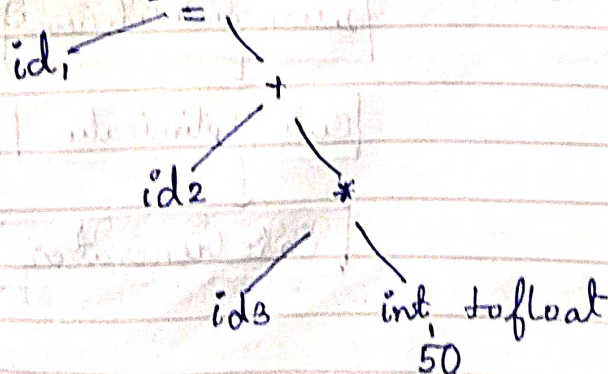
* Syntax Analysis

Input:- id1 = Parse tree



* Semantic Analysis

Input: Modified (Type change)



Date ____ / ____ / ____

* Intermediate Code Generator

Using temporary variable store $t_1 = (\text{int to float}) 50$.

Another temporary variable store, $t_2 = id_3 * 50 t_1$

$t_3 = id_2 + t_2$

$id_1 = t_3$

~~MUL id3, t1, t2~~
~~ADD id2, t2~~

MUL t2, id3, t1
STORE t2
ADD t3, id2, t2
STORE t3
MOV id1, t3

* Code optimization

$t_3 = id_3 * 50.0$

$id_1 = id_2 + t_3$

Low LL \rightarrow Binary, Assembly

MOV F R2, #50

MUL F R2, id3

MOV F R1, R2

ADD F R1, id2

~~MOV~~

Q) Position = Initial + rate * 60

~~1.~~ 1. Lexical Analysis

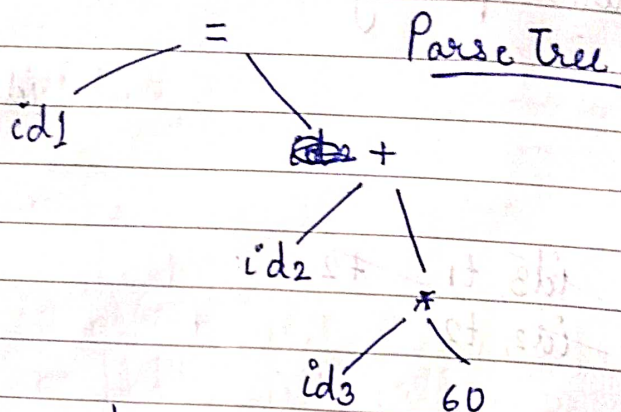
Input: Position = Initial + rate * 60

↓ ↓ ↓ ↓ ↓ ↓
id1 op1 id2 op2 id3 60

Output: $id1 = id2 + id3 * 60$ (sent to Syntax Analyzer)

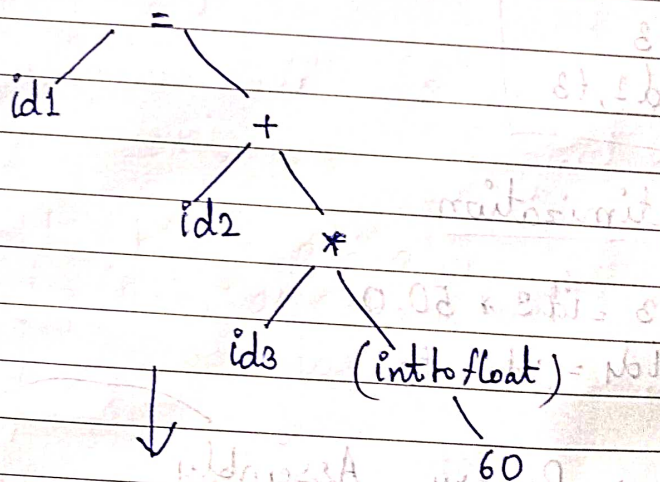
* Syntax Analyzer

Input:



* Semantic Analyzer

Input:



* Intermediate Code Generator

$t1 = (\text{int to float}) 60$

$t2 = id3 * t1$

$t3 = id2 + t2$

$id1 = id2$

* Code Optimization

$t3 = id3 * 50.0$

$id1 = id2 + t3$

Date ____ / ____ / ____

Low LL

MOVF R1, #60

MULF R0, id3

~~MOVF R1, R2~~

ADD F R1, id2