The Wavelet Trie: Maintaining an Indexed Sequence of Strings in Compressed Space CSI 5335 Paper presentation

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Use Case

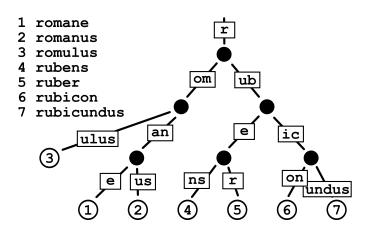
- A lot of things are string sequences.
- Column databases store and index string sequences.
- Great example: access logs

```
188.26.52.117 - - [24/Apr/2013:03:35:48 - 0500] \ "GET / img/welcome/corner.png \\ 188.26.52.117 - - [24/Apr/2013:03:35:49 - 0500] \ "GET / img/welcome/arrowDown.gif \\ 188.26.52.117 - - [24/Apr/2013:03:35:49 - 0500] \ "GET / img/welcome/regionals.jpg
```

- Pretty similar, huh?
- \bullet I heard indexes make stuff faster \to indexed sequence of strings
- Rank query: Number of requests for /img/welcome/corner.png?
- Select query: Position of *i*-th occurrence of /img/welcome/corner.png
- We can do prefix operations too

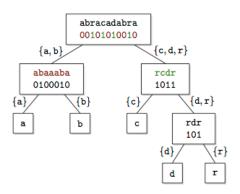
Patricia Trie

- Trie is ordered tree data structure used to store a dynamic set
- Space-efficient trie.
- Node always has at least two children.



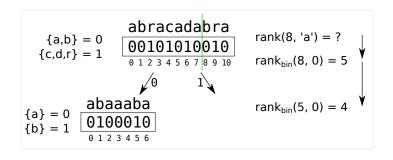
Wavelet Tree

- Organizes a string into a balanced binary tree of bit vectors.
- Alphabet $\Sigma = \{a, b, c, d, r\} \rightarrow \{0, 0, 1, 1, 1\}$
- At root, we have ambiguity, reducing ambiguity towards leaves
- Only the binary vectors are stored!



Efficient Computation of Rank in Wavelet Tree

- rank(8,a) = how many a's before position 8
- rank_{bin}(pos, s) binary rank, # of occurences of s before pos



E.g.: # of requests to /img/welcome/corner.png before April 10th.

Mutable & Compressed Indexed Sequences

- Sequences can change over time:
 Insert(s, pos), Append(s), Delete(pos)
- Alphabet not always known in advance
- Traditional approach, store explicitly (e.g. array), make an index
- Space-inefficient, we want to query the compressed representation
- Succint data structure uses space close to lower information-theoretic bound

Wavelet Trie to rescue!

Wavelet Trie

- Wavelet Tree + Patricia Trie
- Compressed data structure
- Can support Insert, Append, Delete, and dynamic alphabet
- Variants: Static, Append-only, Fully-dynamic
- Note $O(|s| + h_s)$ is very fast for balanced trees

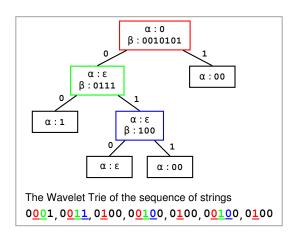
	Query	Append	Insert	Delete	Space
Static	$O(s +h_s)$	-	-	-	$LB + o(\tilde{h}n)$
Append-only	$O(s +h_s)$	$O(s +h_s)$	-	-	$LB + PT + o(\tilde{h}n)$
Fully-dynamic	$O(s + h_s \log n)$	$LB + PT + O(nH_0)$			

 h_s – number of nodes traversed while searching for s is Patricia Tree \tilde{h} – average height of the Wavelet Trie |s| – length of query string

LB – lower bound on storing the unique elements of a sequence

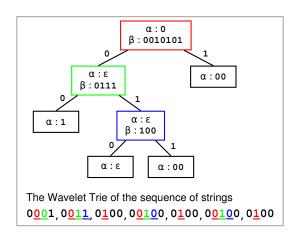
Wavelet Trie - Example (Static)

- $\Sigma = \{0001, 0011, 0100, 00100\}$, α common prefix, β bit-vector
- Splitting based on common prefix like in Patricia Trie, not halving an alphabet like in Wavelet Tree!



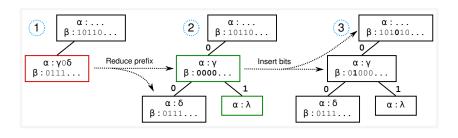
Wavelet Trie - Example (Static)

- $Rank(6,0100) \rightarrow Rank(6,1) = 2$
- Notice how ambiguity decreases towards the leaves



Dynamic Wavelet Tries

- Bitvectors must support insertion and deletion
- We also want dynamic alphabet, e.g. we want to add "0111"



Inserting new $s = ... \gamma 1\lambda$ at pos = 3; γ – common prefix, λ – suffix

- Original Wavelet Trie
- After splitting red node, adding internal node and leaf node (green)
- Inserting bits in the root-to-leaf path nodes

Summary

- Wavelet Trie = Wavelet Tree + Patricia Trie
- Compressed sequences of strings
- Very space efficient (no need for data+index)
- Applications in column databases, access log aggregation, ...
- Quite fast, querying in $O(|s| + h_s)$
- Can support prefix searches

Thank you.

Resources

- Grossi, Roberto, and Giuseppe Ottaviano. "The wavelet trie: Maintaining an indexed sequence of strings in compressed space."
 In Proceedings of the 31st symposium on Principles of Database Systems, pp. 203-214. ACM, 2012.
- http://alexbowe.com/wavelet-trees/
- http://en.wikipedia.org/wiki/Wavelet_Tree
- http://siganakis.com/ challenge-design-a-data-structure-thats-small