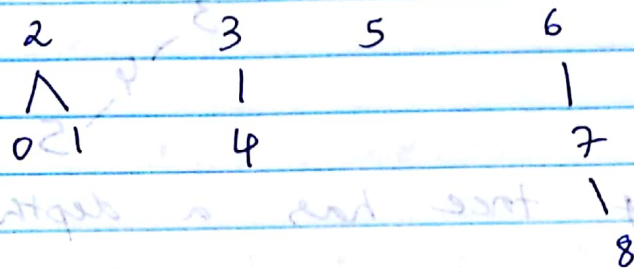


# Assignment - 8

1.

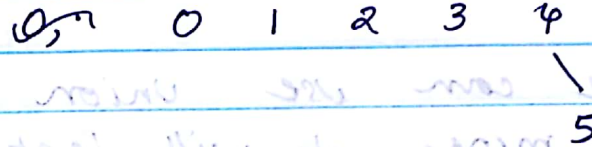
2	2	-1	-1	3	-1	-1	6	7
0	1	2	3	4	5	6	7	8



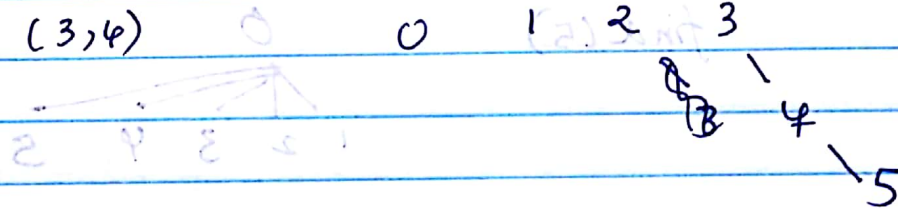
(1) 2.

0, 1, 2, 3, 4, 5

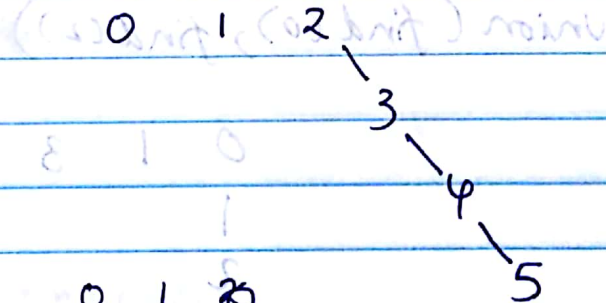
Union (4, 5)



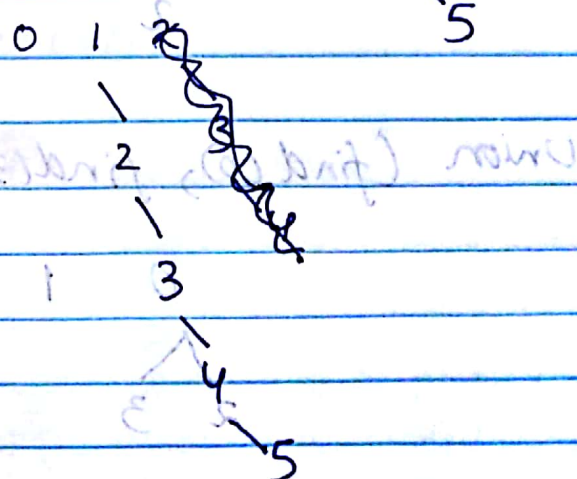
Union (3, 4)



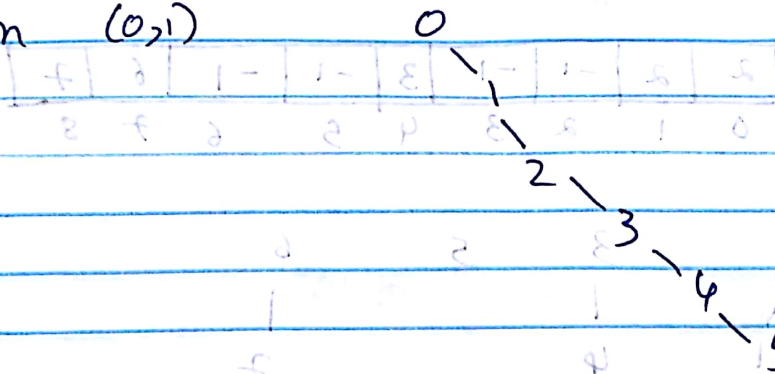
Union (2, 3)



Union (1, 2)



Union (0,1)

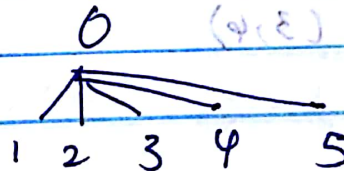


The resulting tree has a depth of 5 i.e.  $(N-1)$ .

Although '5' can be accessed at  $O(1)$  time,  $\text{find}(5)$  in the resulting tree will take  $O(N)$  time.

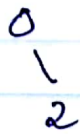
We can use union function to make a merge, it will look as following:

$\text{find}(5)$



3

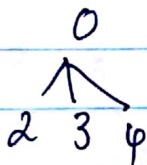
union (find(0), find(2))



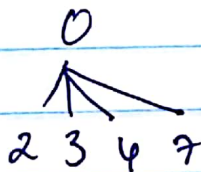
union (find(0), find(3))



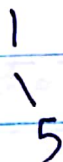
union (find(0), find(4))



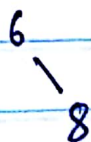
union (find(0), find(7))



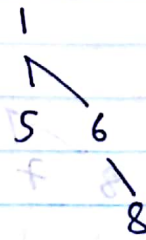
union (find(1), find(5))



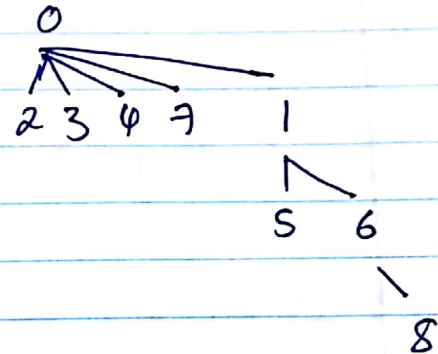
union (find(6), find(8))



union (find(5), find(8))



union (find(7), find(8))



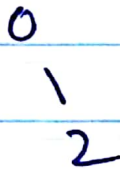


4.

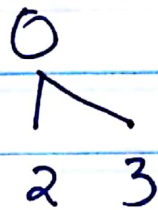
-9	0	0	0	0	1	0	0	6
0	1	2	3	4	5	6	7	8

5.

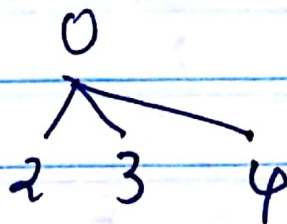
Union (find(0), find(2))



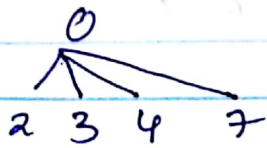
Union (find(0), find(3))



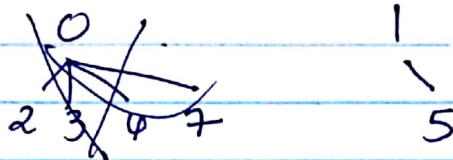
Union (find(0), find(4))



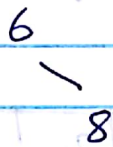
union (find(0), find(7))



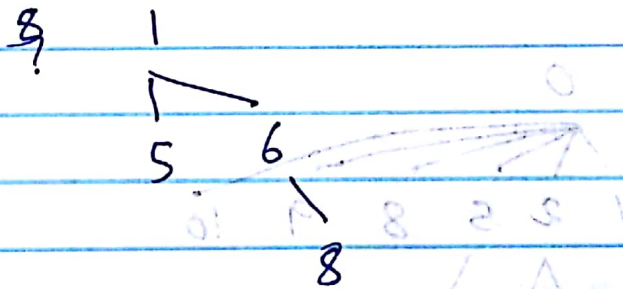
union (find(1), find(5))



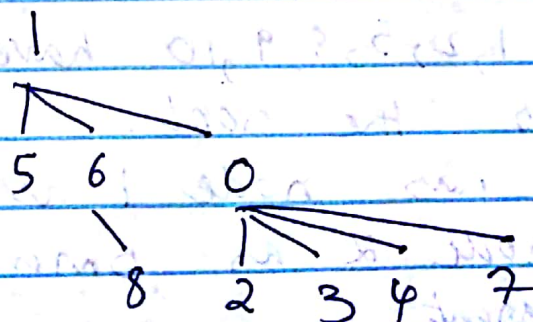
union (find(6), find(8))



union (find(5), find(8))



union (find(7), find(8))



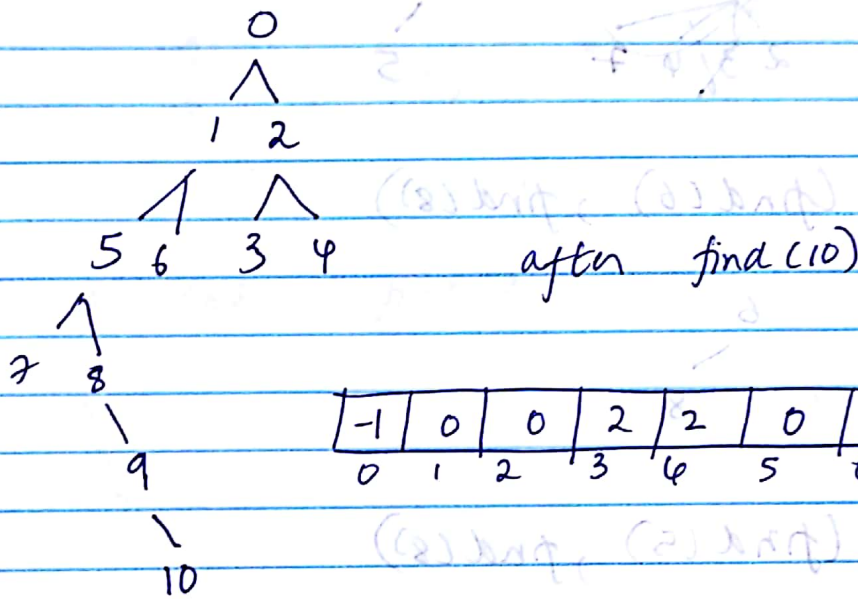


6.

1	-3	0	0	0	1	1	0	6
0	1	2	3	4	5	6	7	8

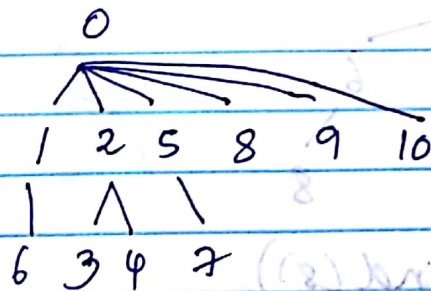
7.

-1	0	0	2	2	1	1	5	5	8	9
0	1	2	3	4	5	6	7	8	9	10



-1	0	0	2	2	0	1	5	0	0	0
0	1	2	3	4	5	6	7	8	9	10

8.



Nodes 1, 2, 5, 8, 9, 10 have 0 as parent, 0 is also the root.

Node 6 has node 1 as parent, Nodes 3, 4 have node 2 as parent, node 7 has 5 as parent.