

Day 3 (Morning)

- ① As far from land as possible
- ② Mother vertex
- ③ Kosaraju Algorithm

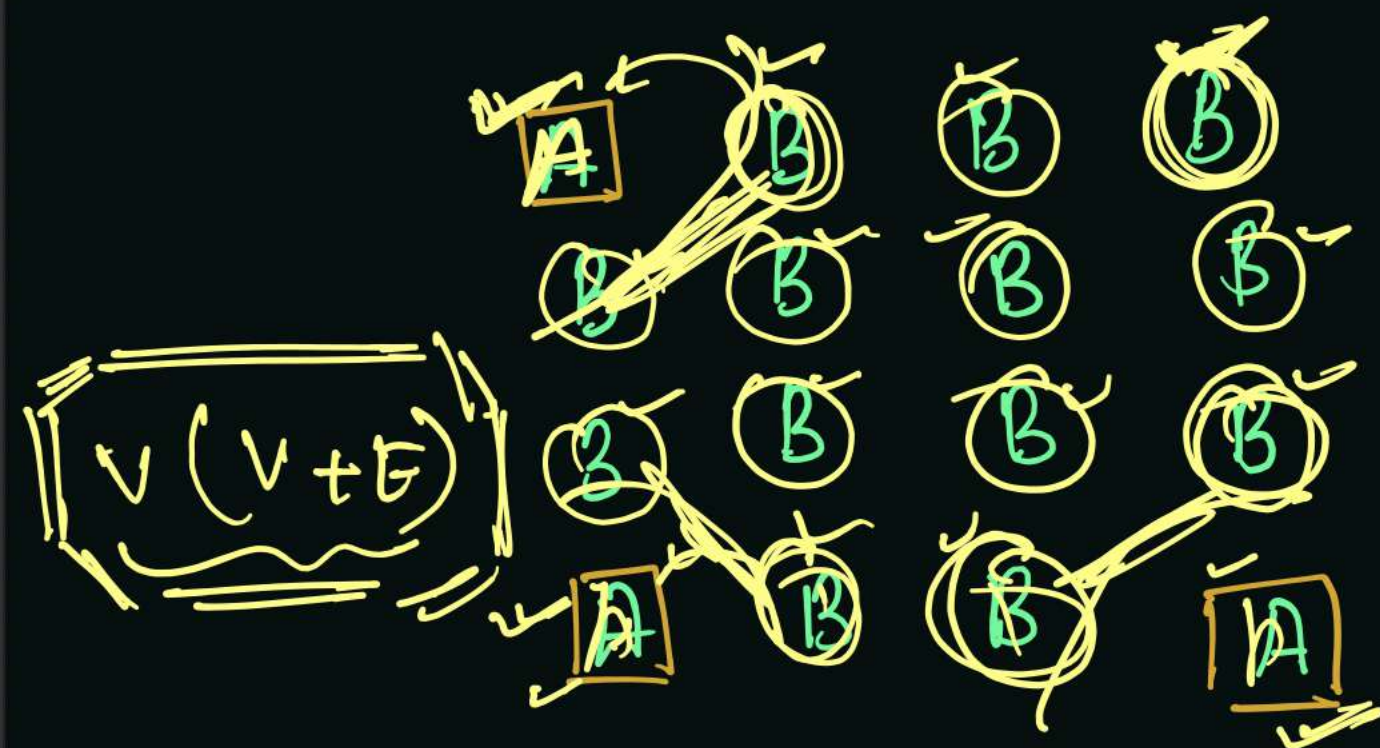
Day 4 (Evening)

- ① Shortest bridge
- ② Articulation point and bridge
- ③ Critical connection

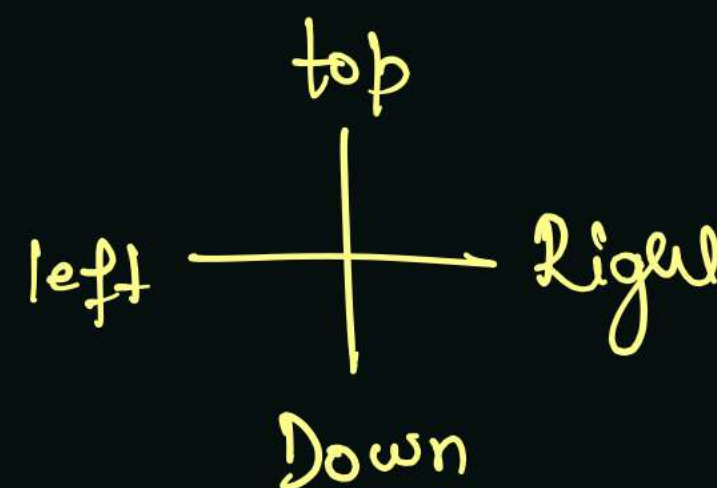
As Far From Land As Possible

Saturday, 25 September 2021 7:04 PM

A \Rightarrow Land, B \Rightarrow water (Manhattan dist.)



Maximise min distance from B to A.



put all A's at same level
Simultaneously BFS from land

Min distance from B to A \rightarrow

A	1	2	3
1	2	3	2
1	2	2	1
A	1	1	A

Result = 3

or

take all coordinate at level-0
(initial A's coordinate)

pair = level

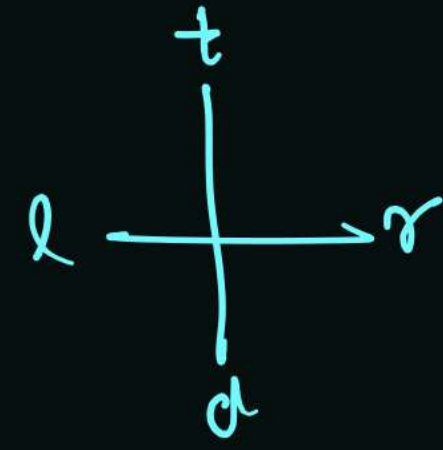
while(



pair = level
while() (V+E)

	0	1	2	3
0	1	0	0	0
1	0	0	0	0
2	0	0	0	0
3	1	0	0	1

① add initial ones



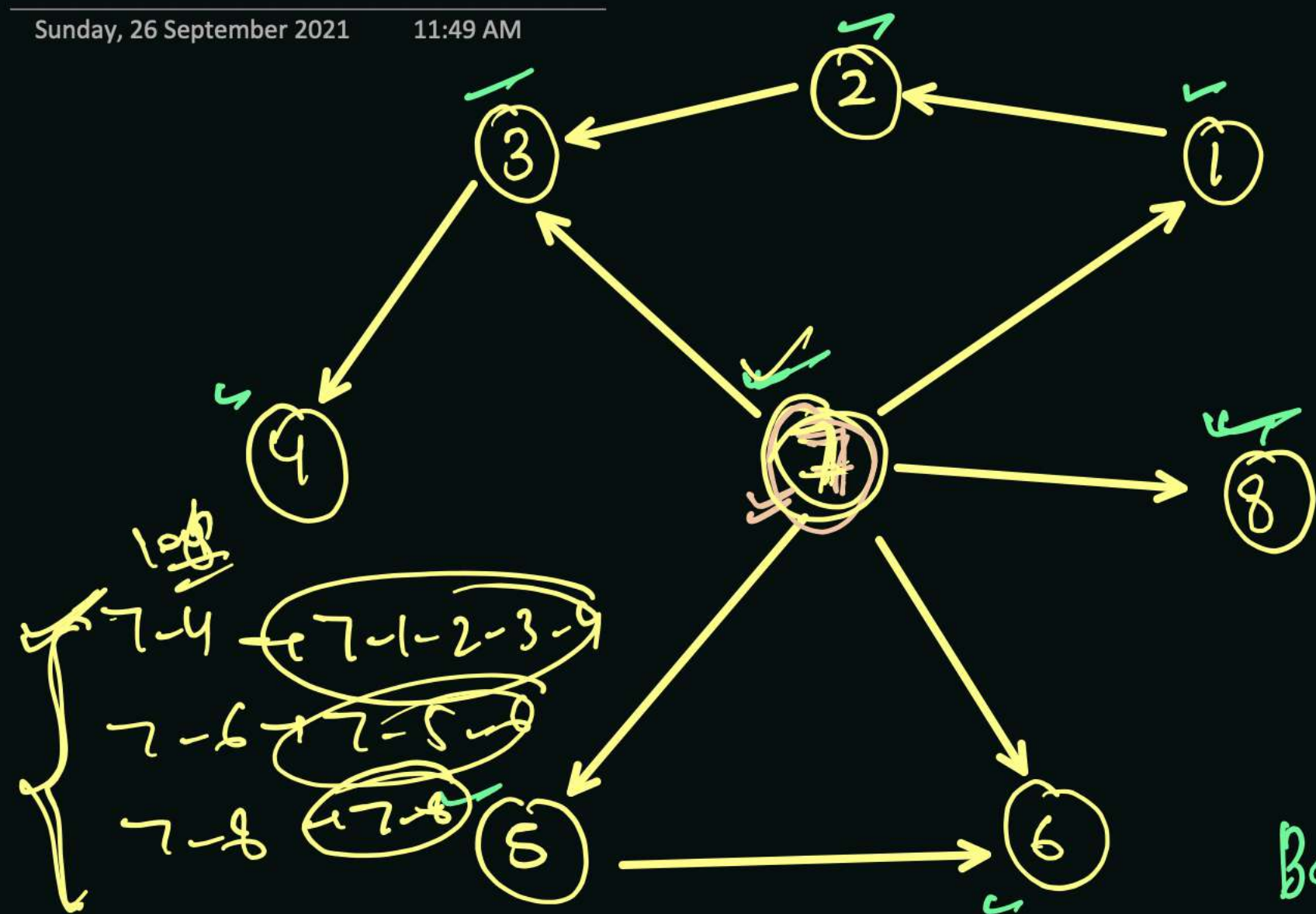
level = ~~1~~ ~~0~~ ~~1~~ ~~2~~ 3

~~(0,0)~~ | ~~(3,0)~~ | ~~(2,0)~~ | ~~(1,0)~~ | ~~(0,0)~~ | ~~(2,0)~~ | ~~(3,1)~~ | ~~(2,1)~~ | ~~(3,2)~~ | ~~(1,1)~~ | ~~(0,2)~~ | ~~(2,1)~~ | ~~(1,3)~~ | ~~(2,2)~~

Size = ~~2~~ ~~6~~ ~~5~~ 2

~~(1,2)~~ | ~~(0,3)~~ |

level = 3 → Max distance
↳ Manhattan distance



We can visit all vertex from 7, so

'7' is mother vertex

① Brute force
→ Try

Dijkstra's

Shortest path in
terms of weights

Topological
Sort

Back track → stack.push(vertices)

mark ① → DFS

→ mark

→ 4

→ 3

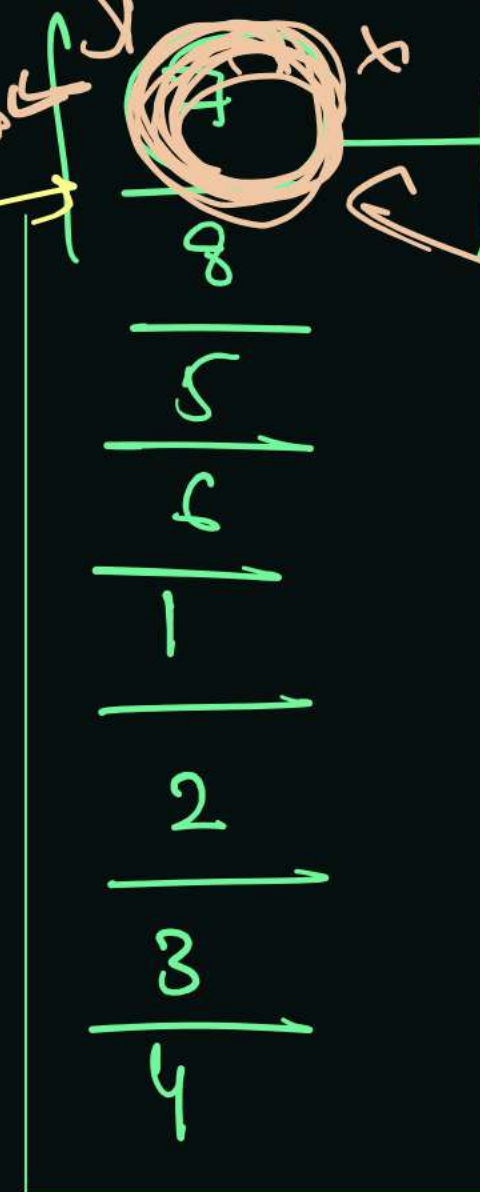
→ ⑤ → DFS

→ 6 → mark

→ 7 → mark

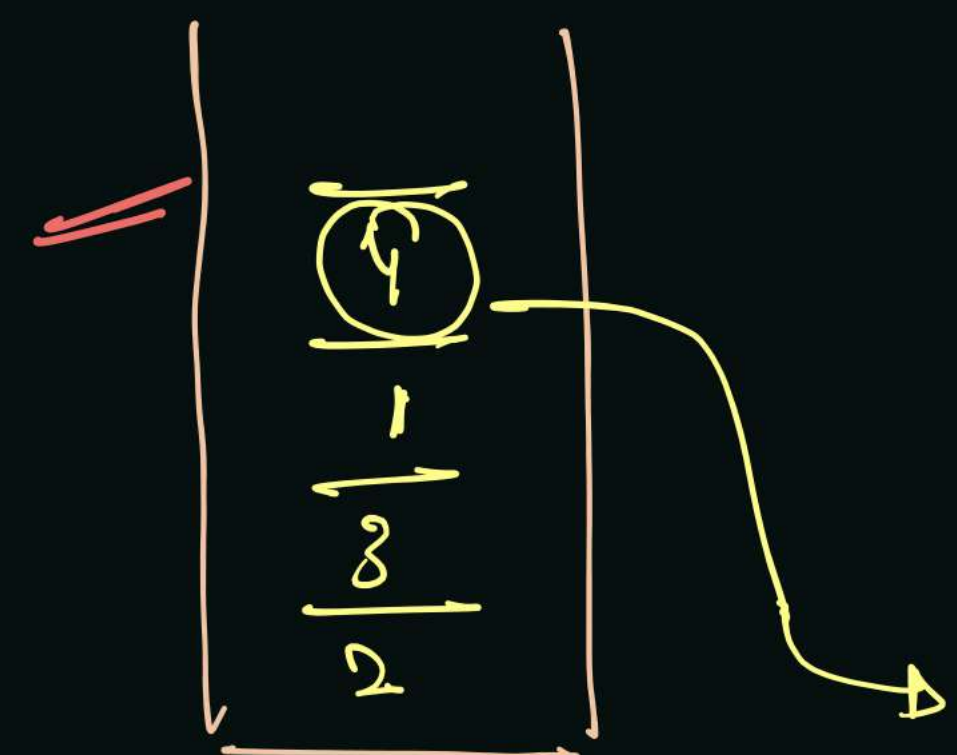
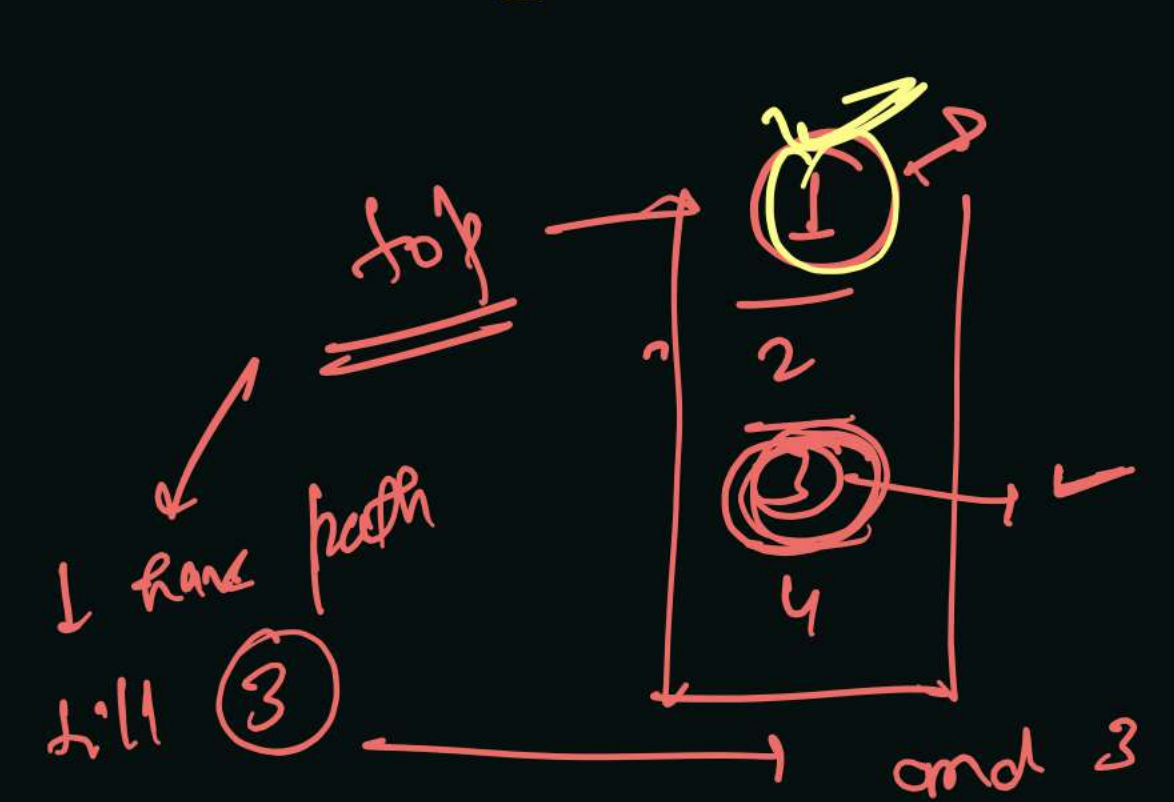
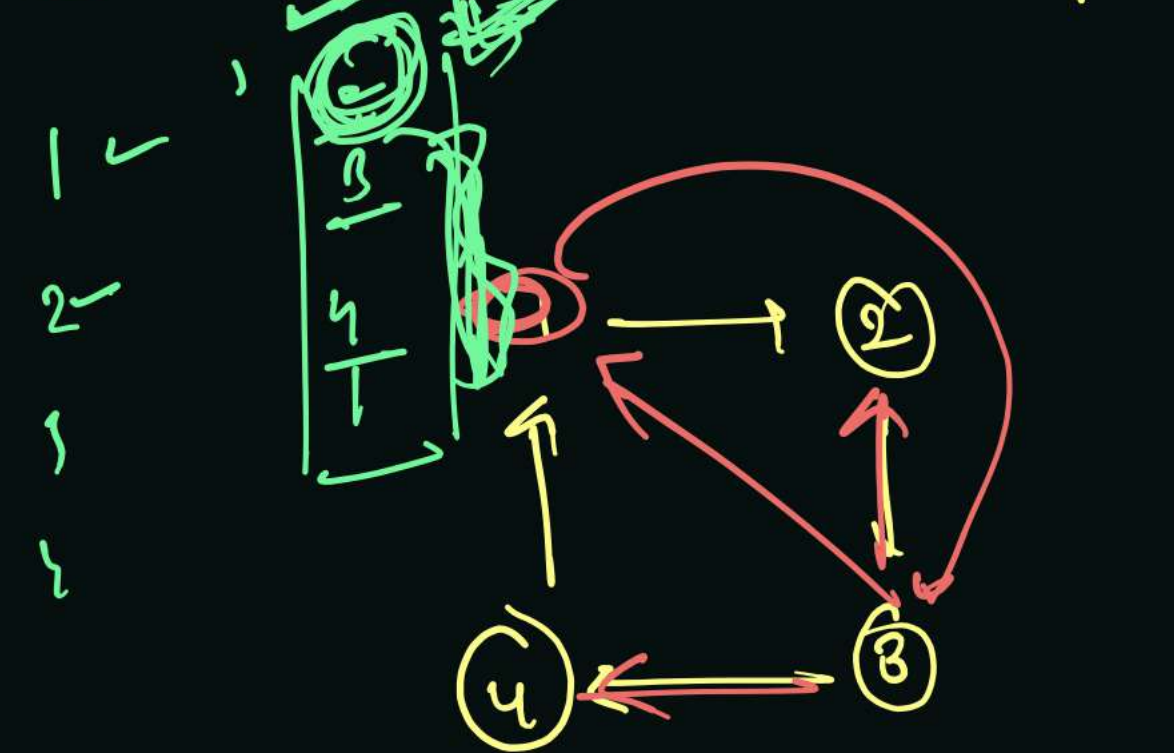
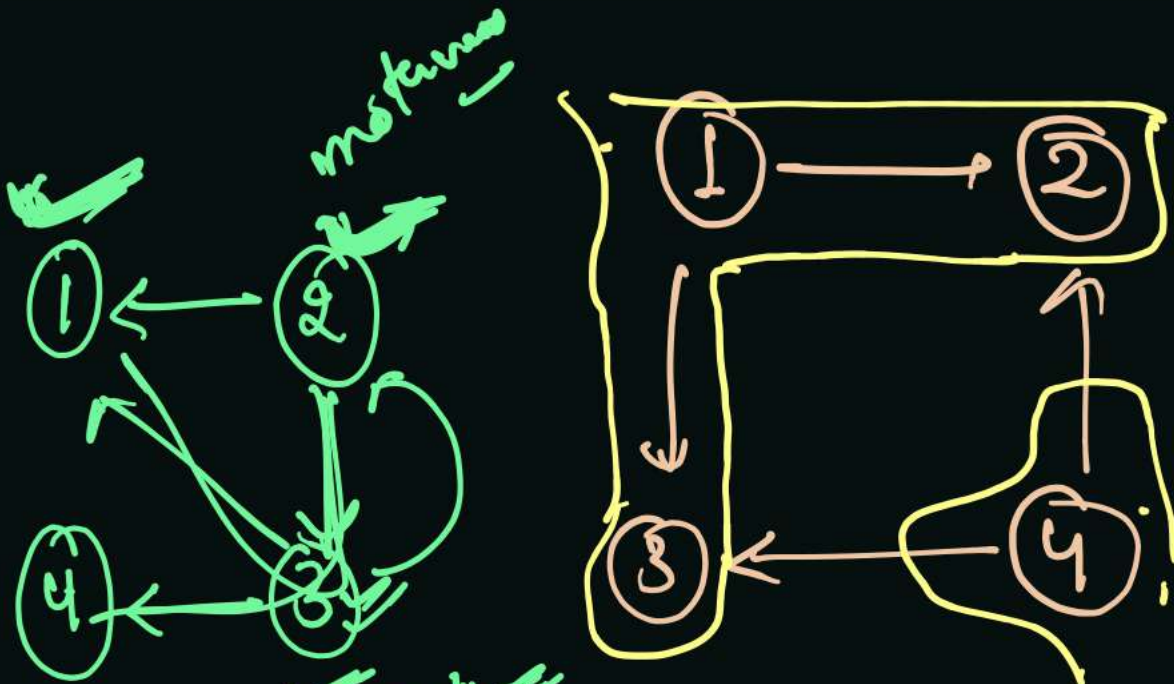
→ 8 → "

Stack

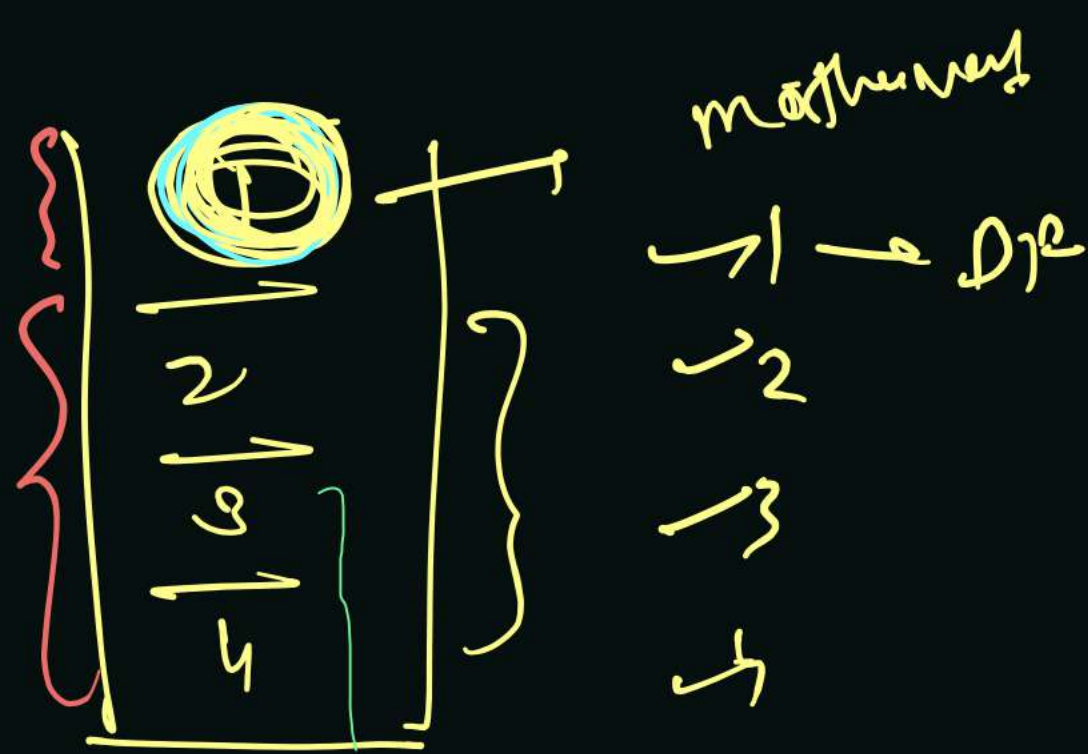


$${}^nC_2 = \frac{n(n-1)(n-2)\dots(1)}{(2)!}$$

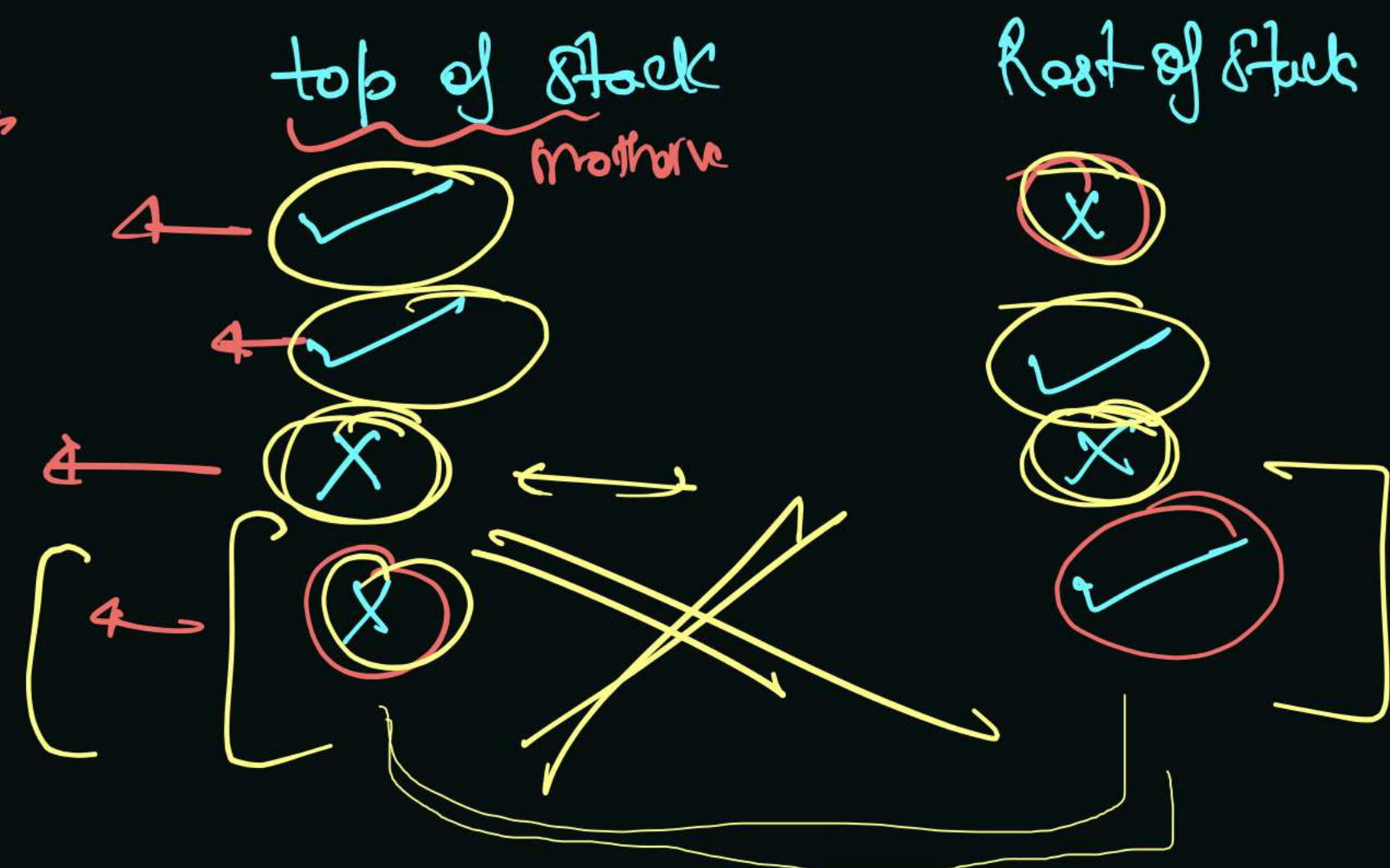
$$= \underline{\underline{V^2}} \times \underline{\underline{DFS}}$$



✓ 1 → DP
 ✓ 2 → mother
 ✓ 3 → mother
 ✓ 4 →
 4 is at top but it is not mother vertex

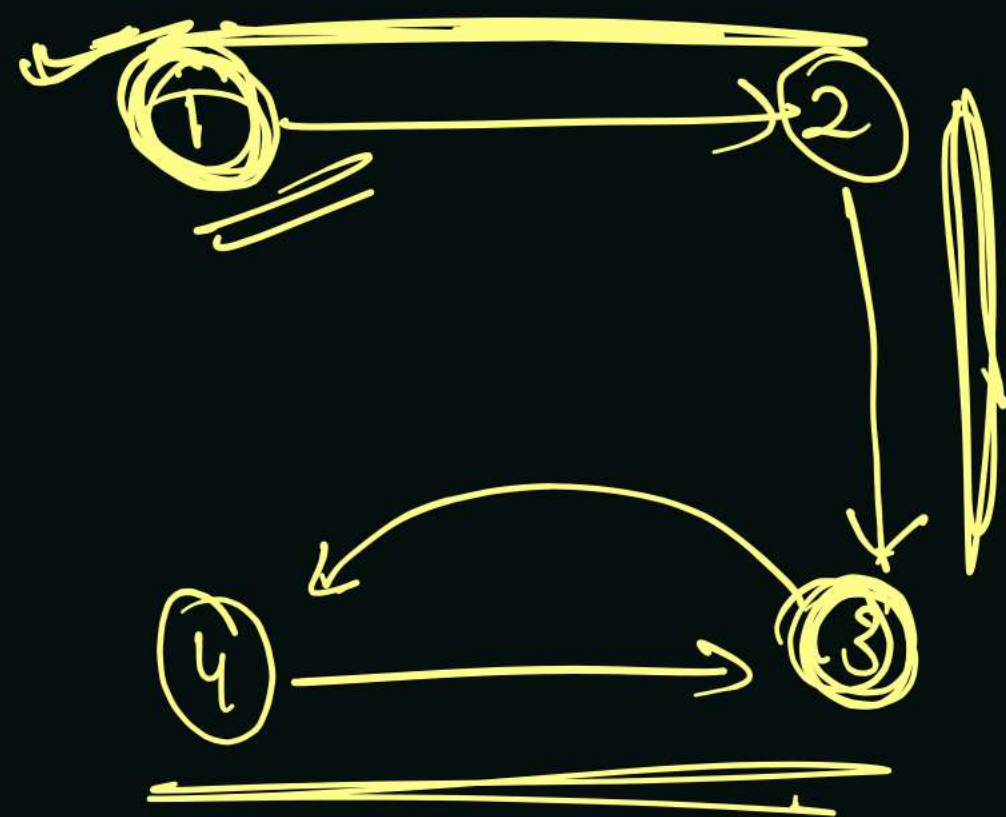


4
 4



definitely it is mothermost
 is [Mothermost]

Example →



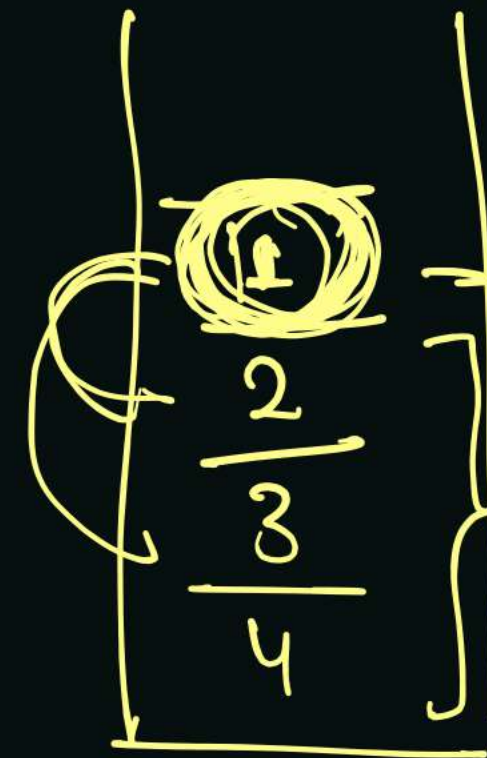
✓ 3 → DFS

✓ 2 → DFS

- 1 -

✓ 4 → already
vis

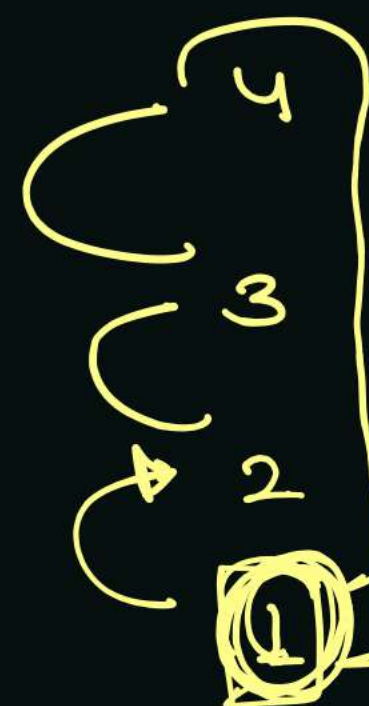
Stack



Mother vertex

Valid candidate
for mother

Vertices

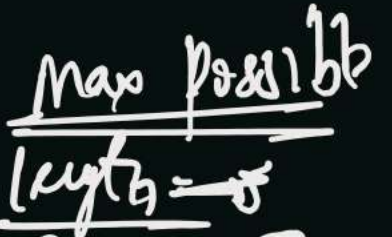
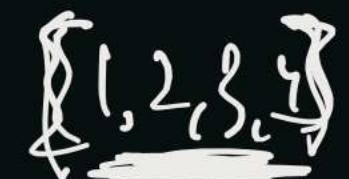


if exist just it is
mother vertex
otherwise

Travel all element

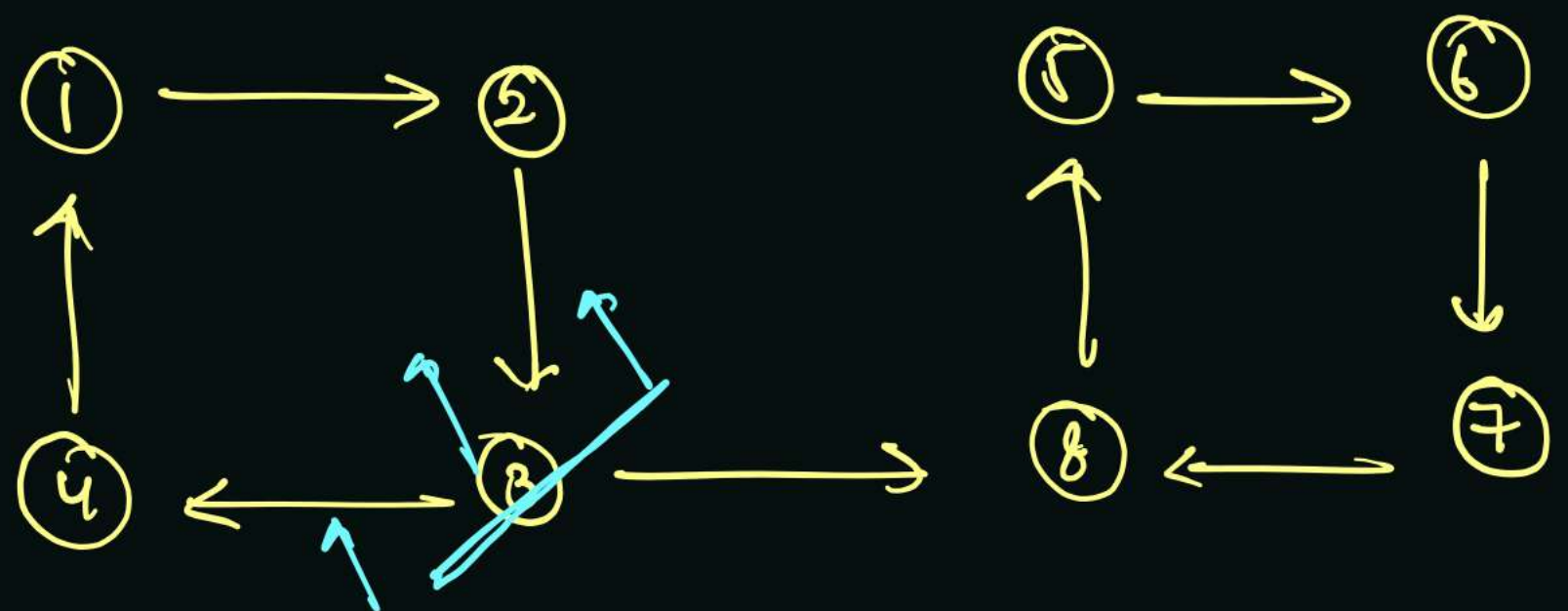
no mother vertex is
here

1-8


$$\{5, 1, 7, 8\}$$

$$[(1, 2, 3, 4) \quad (5, 6, 2, 8)]$$

is cyclic \mathbb{Z}

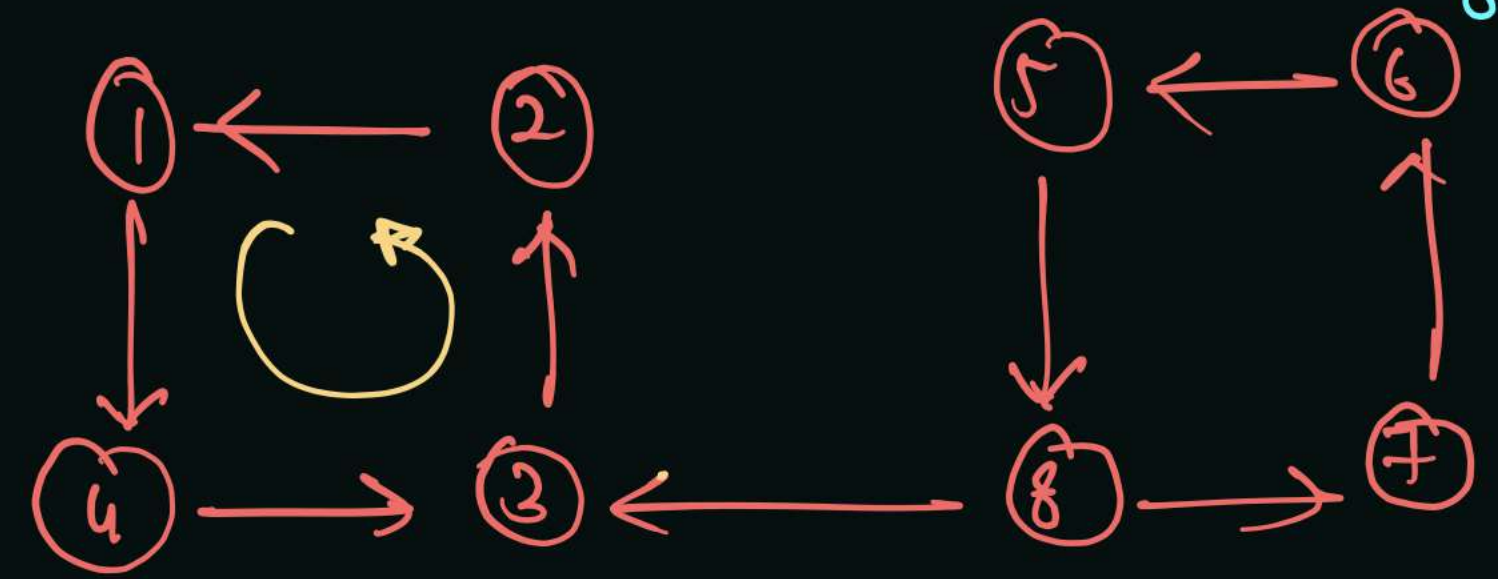
~~No. of cycle is no. of strongly connected component it have~~



Step ① → travel on vertex and call DFS from unvisited vertex and add vertex in stack in post order

~~1~~ → top element

1	vis
2	vis
3	vis
4	vis
5	vis
6	vis
7	vis
8	vis



Step ② → Reverse all Edge

Step ③ → Stack DFS in order of stack's element and count No. of calls

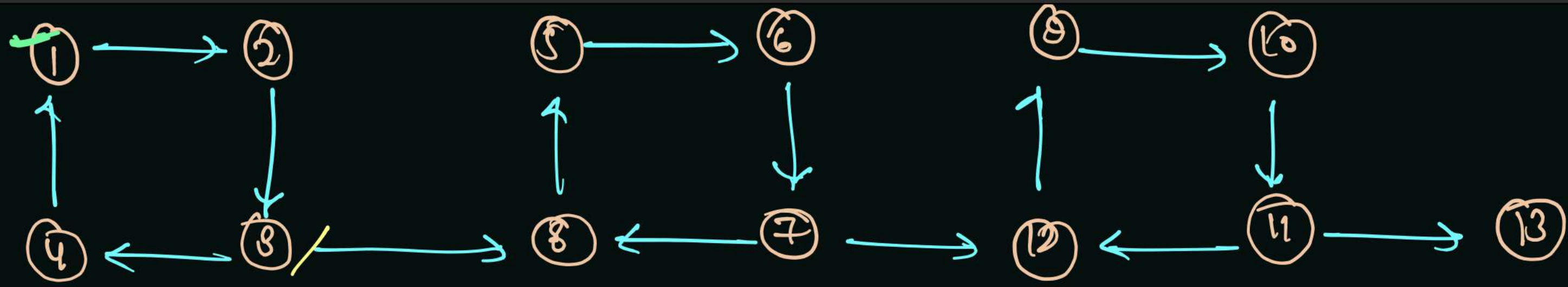
- Mark
- ✓ 1
 - ✓ 2
 - ✓ 3
 - ✓ 4
 - ✓ 5
 - ✓ 6
 - ✓ 7
 - ✓ 8

1 4 3 2
call 1

8 7 6 5
call 2

Count

$[[1, 4, 3, 2], [8, 7, 6, 5]]$



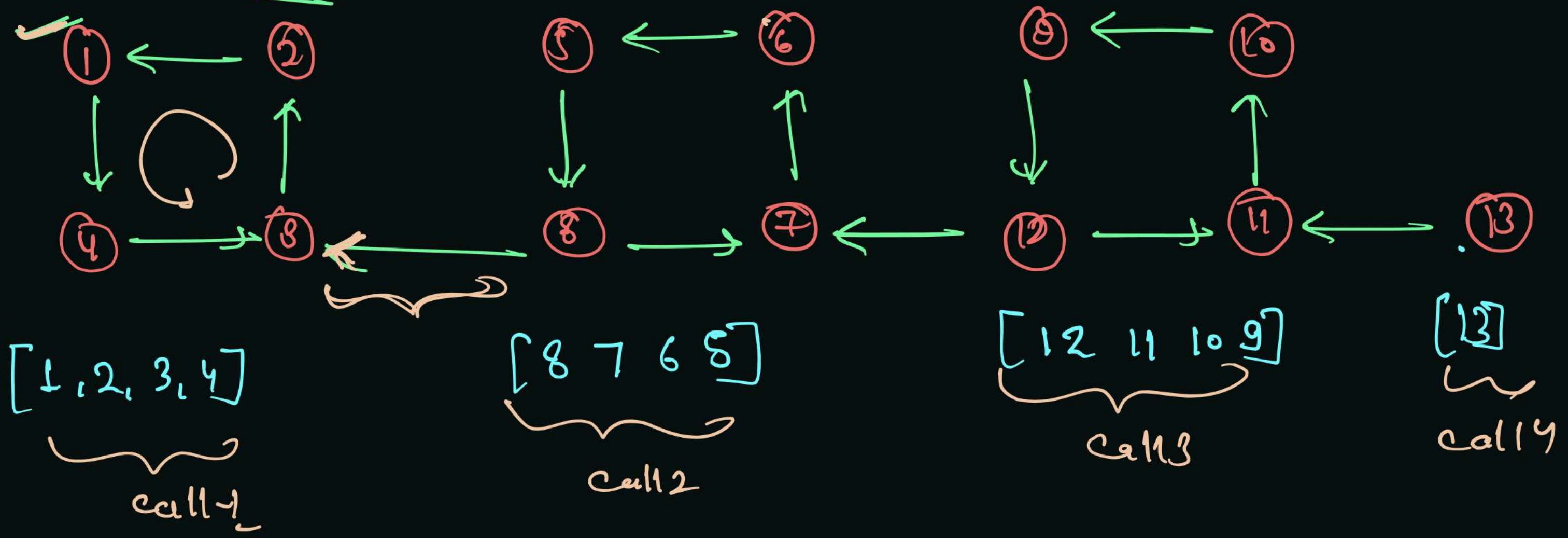
Step 1 - Stack fill
 Step 2 - Reverse Edge

- ✓✓ 1 → DFS
- ✓✓ 2 ✓
- ✓✓ 3 ✗
- ✓✓ 4 ✗
- ✓✓ 5 ✗
- ✓✓ 6 ✗
- ✓✓ 7 ✗
- ✓✓ 8 ✗
- ✓✓ 9 ✗
- ✓✓ 10 ✗
- ✓✓ 11 ✗
- ✓✓ 12 ✗
- ✓✓ 13 ✗

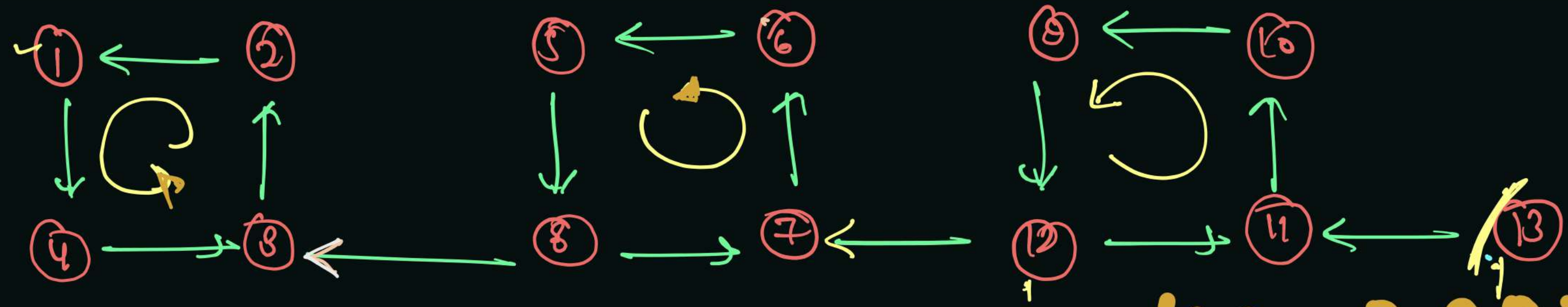
1	DFS
2	vis
3	vis
4	DFS
5	vis
6	vis
7	vis
8	DFS
9	vis
10	vis
11	vis
12	DFS
13	vis
4	vis

stack

Reverse Edge

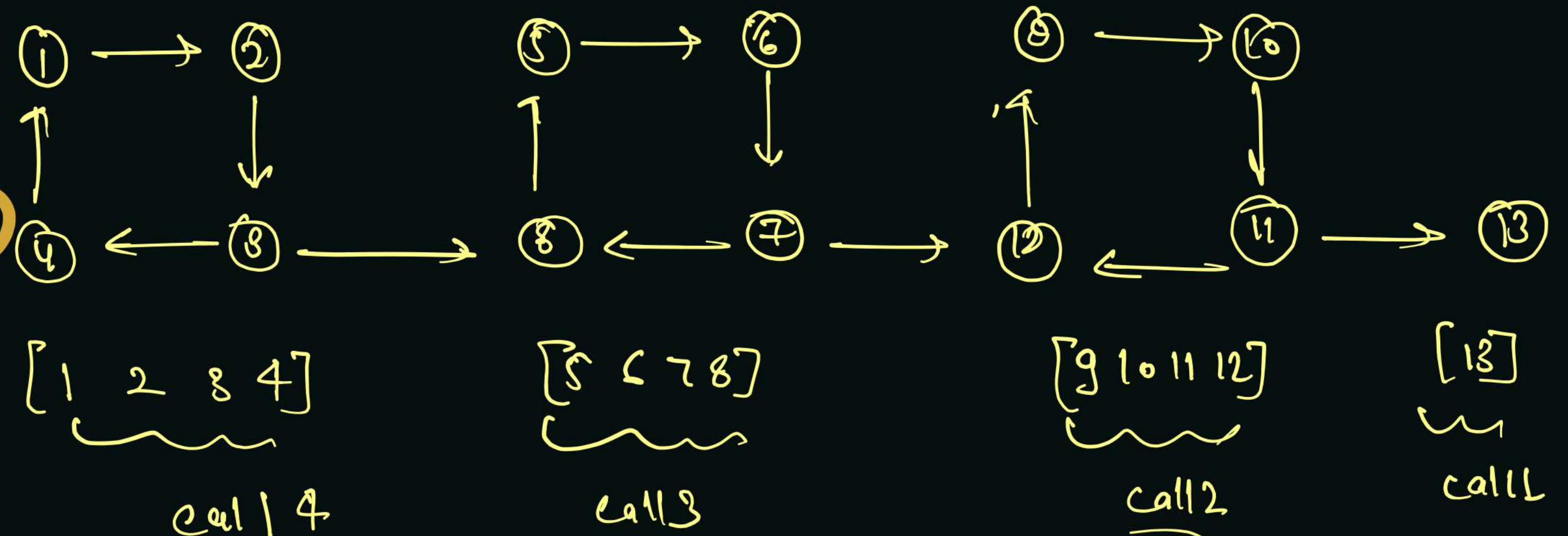


Strongly connected components.



KOSARAJU ALGO.

Reverse Edge



Total = 4 calls = 4 strongly connected comp.

- ✓ 1 → DFS
- ✓ 2 → vis
- ✓ 3 → vis
- ✓ 4 → vis
- ✓ 5 → DFS
- ✓ 6 → vis
- ✓ 7 → vis
- ✓ 8 → vis
- ✓ 9 → DFS
- ✓ 10 → vis
- ✓ 11 → vis
- ✓ 12 → vis
- ✓ 13

- Stack
- ✓ 13 → DFS
 - ✓ 9 → DFS
 - 12 vis
 - 11 vis
 - 10 vis
 - 5 → DFS
 - 8 vis
 - 7 vis
 - 6 vis
 - 1 → DFS
 - 4 vis
 - 3 vis
 - 2 vis
- Stack