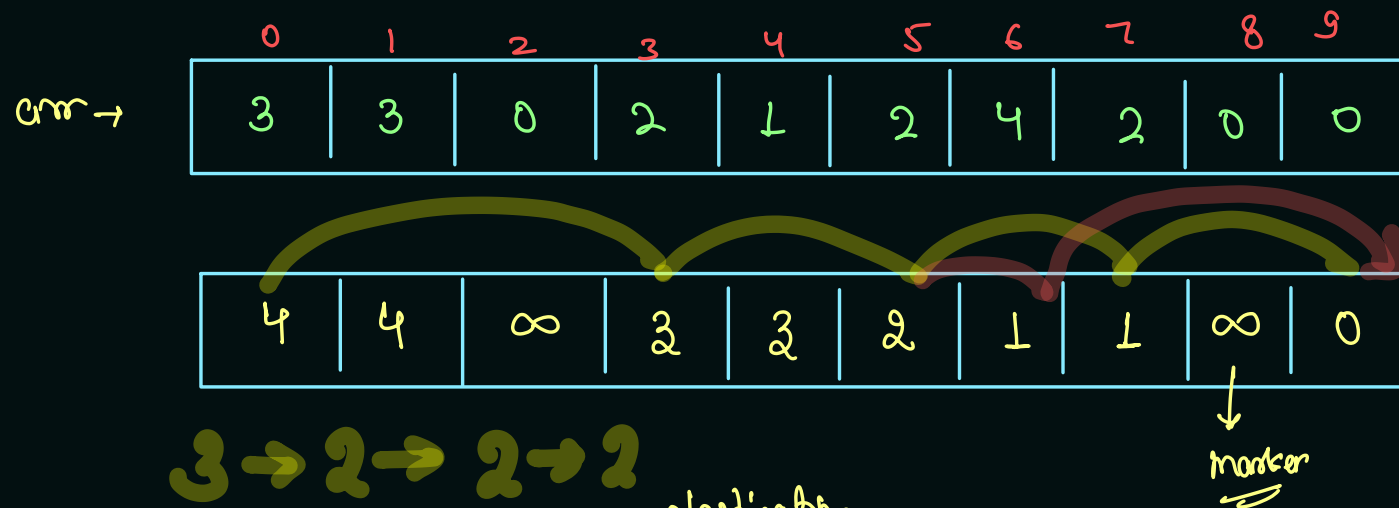


Print all path with minimum Jumps: →



arr[i] → no. of Jumps allowed from ith index

0 to 9 → Min Jumps

Jump
vs
path

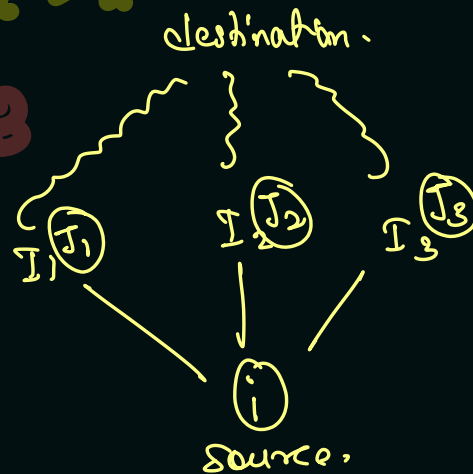
3 → 2 → 2 → 2

8 → 2 → 1 → 3

Smallest → 9 to 9 = 0
↳ no. of jumps
problem

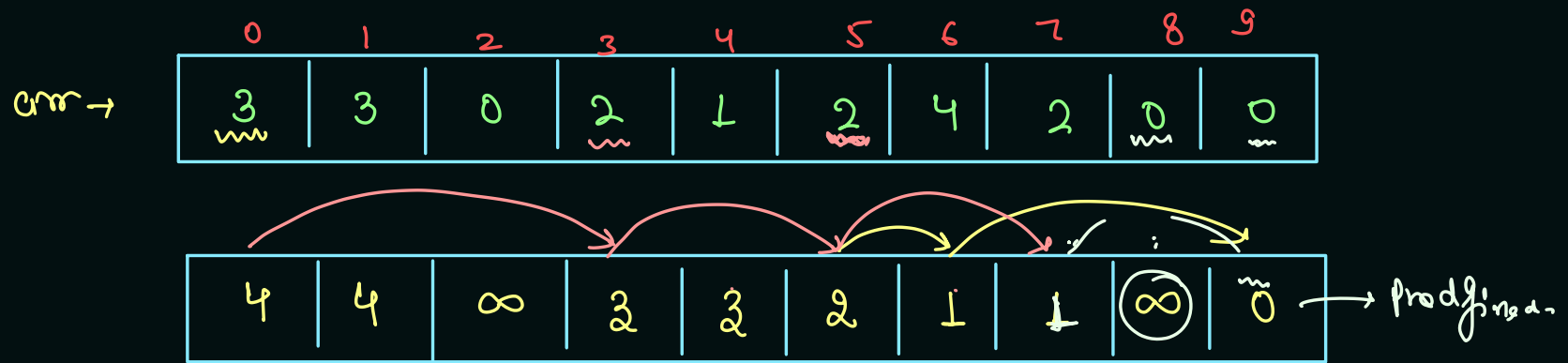
9 to 9 → no. of paths
↳ 1
i.e. don't move.

Size
vs
Jump Count

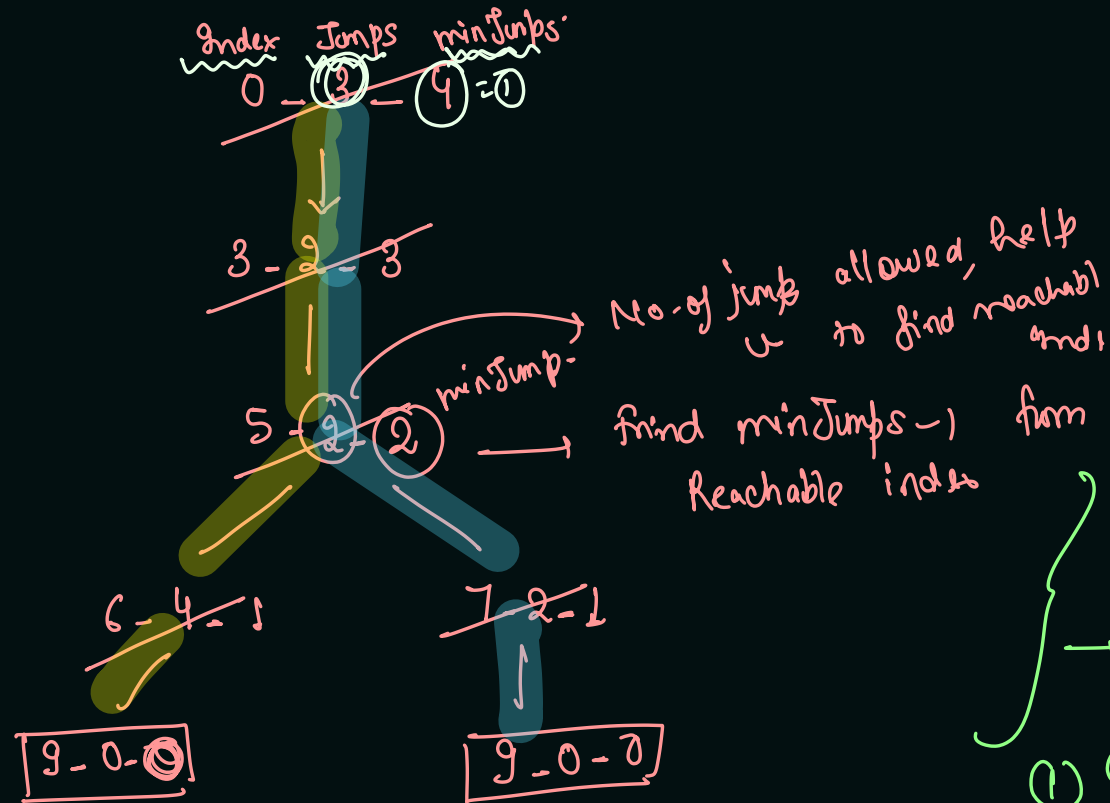


i to destination →
= min(J₁, J₂, J₃) + 1 → from source to intermediate

Result from arr[0] to arr[9]
we can reach in 4 Jumps
path ??



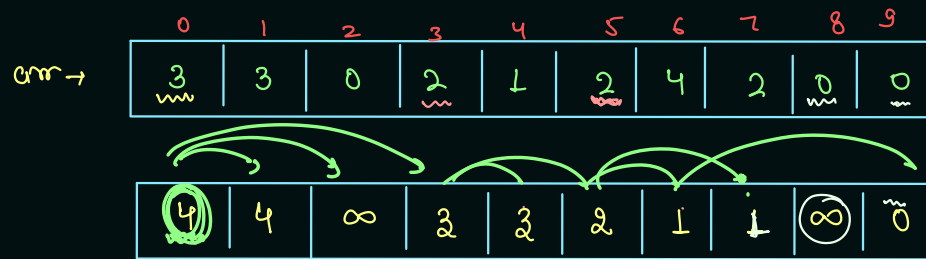
0 → 3 → 5 → 6 → 9 .
 3 → 2 → 1 → 3 .
 0 → 3 → 5 → 7 → 9 .
 3 → 2 → 2 → 2 .



Reverse Engineering

DP:

- ① Storage
- ② Assign one array total
- ③ Fill prerequisite
 → BFS
 Analyse Smaller
- ④ Get + Remove problem
- ⑤ make and direct work to fill DP
- ⑥ Add unvisited nb.
- ⑦ Travel & Fill



$n=10$

```

Queue<MJHelper> que = new LinkedList<>();
que.add(new MJHelper(0, arr[0], dp[0], "0"));

while(que.size() > 0) {
    MJHelper rem = que.remove();
    if(rem.indx == n - 1) {
        System.out.println(rem.psf + " .");
        continue;
    }
    for(int jump = 1; jump <= rem.jumps && jump + rem.indx < n; jump++) {
        if(dp[rem.indx + jump] == rem.minJumps - 1) {
            int nindx = rem.indx + jump;
            que.add(new MJHelper(nindx, arr[nindx], rem.minJumps - 1, rem.psf + " -> " + nindx));
        }
    }
}

```

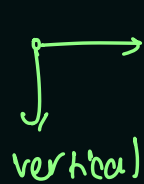
0 → 3 → 5 → 6 → 9
0 → 3 → 5 → 7 → 9

indx jump minJumps psf
~~0, 3, 4, 0~~
 |
~~3, 2, 2, 0 → 3~~
 |
~~5, 2, 2, 0 → 3 → 5~~
 / \
~~6, 4, 1, 0 → 3 → 5 → 6~~ ~~7, 2, 1, 0 → 3 → 5 → 7~~
 | |
~~9, 0, 0, 0 → 3 → 5 → 6 → 9~~ ~~9, 0, 0, 0 → 3 → 5 → 7 → 9~~

Print all path for min cost

	0	1	2	3	4	5
0	0	1	4	2	8	2
1	4	3	6	5	0	4
2	1	2	4	1	4	6
3	2	0	7	3	2	2
4	3	1	5	9	2	4
5	2	7	0	8	5	1

	0	1	2	3	4	5
0	23	23	24	20	21	19
1	24	22	23	18	13	17
2	20	19	17	18	13	12
3	21	19	19	12	9	7
4	23	20	19	16	11	8
5	23	21	14	14	6	1


 Horizontal
 vertical

with two moves i.e. H, & v,
 what is the min cost to reach
 at Bottom Right from top
 left?

✓ H V V H H V H V H V
 ✓ H V V H H V H H V V

Print All path for max Gold %

Goldmine $\begin{matrix} \nearrow D_1 \\ \rightarrow D_2 \\ \searrow D_3 \end{matrix}$

$x, y \begin{matrix} \nearrow x-1, y+1 \\ \rightarrow x, y+1 \\ \searrow x+1, y+1 \end{matrix}$

last Colun

first Row

1

2

3

4

5

last

	0	1	2	3	4	5
0	0	1	4	2	8	2
1	4	3	6	5	0	4
2	1	2	4	1	4	6
3	2	0	7	3	2	2
4	3	1	5	9	2	4
5	2	7	0	8	5	1

0

33-1

33-2

3

33-4

5

33

4

0	1	2	3	4	5
26	24	21	14	12	2
31	26	23	17	6	4
28	27	21	11	10	6
29	25	25	13	8	2
32	30	18	17	9	1

$D_3 \rightarrow D_1 \rightarrow D_2 \rightarrow D_3 \rightarrow D_1$

Direction of Mining \rightarrow left to Right

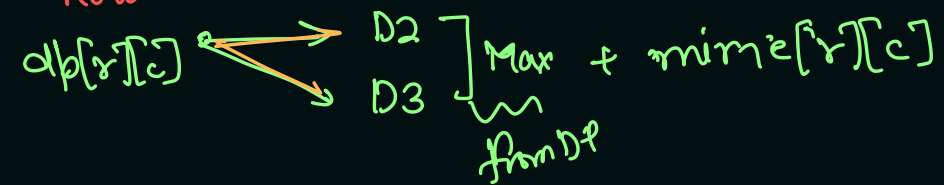
Direction of Solving problem \rightarrow Right to left

last column

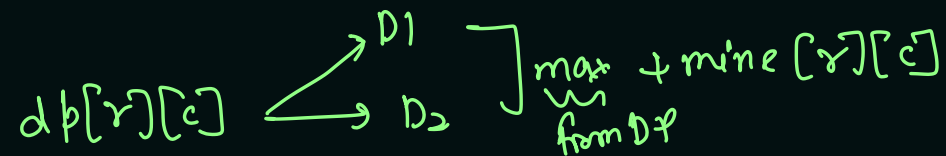
→ same as min

$$dp[r][c] = mine[r][c]$$

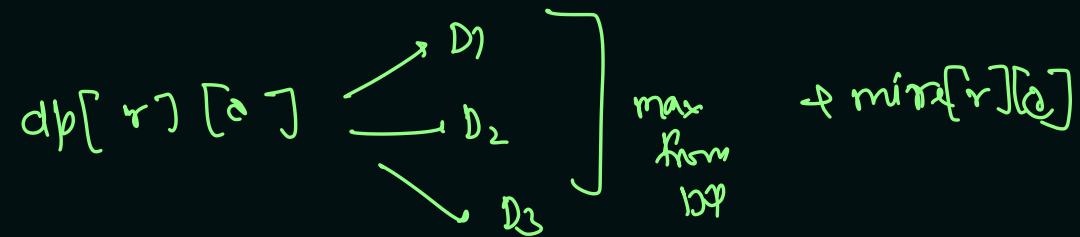
First Row



Last Row



Middle



	0	1	2	3	4	5
0	26	24	21	14	12	2
1	31	26	28	17	6	4
2	28	27	21	11	10	6
3	29	25	25	13	8	2
4	33	26	23	18	6	4
5	32	30	18	17	9	4

if $C == 5$

Destination is
arrived

4 D3 D1 D2 D3 D1

① Travel in zeroth column and add max
Coordinate in queue of pairs.

② Then solve BFS:

$\frac{x, y, \text{psf}}{i, j, i}$

~~4, 0, "4"~~

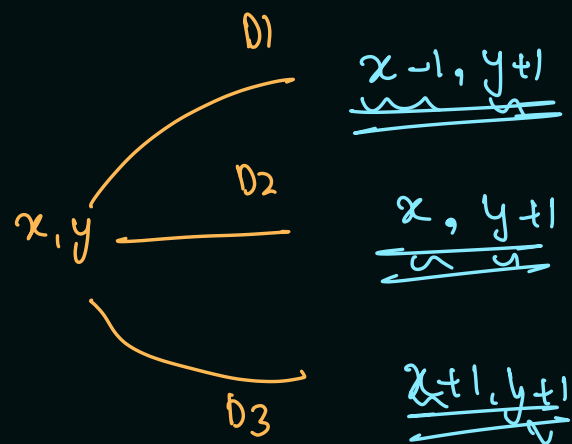
~~5, 1, "4 D3"~~

~~4, 2, "4 D3 D1"~~

~~4, 3, "4 D3 D1 D2"~~

~~5, 4, "4 D3 D1 D2 D3"~~

~~4, 5, "4 D3 D1 D2 D3 D1"~~ 5, 5, "4 D3 D1 D2 D3 D2"



$$\begin{array}{ccc} x \rightarrow & \textcircled{1} & 0 & 1 \\ y \rightarrow & \textcircled{1} & 1 & 1 \end{array}$$

$$\underline{\underline{D = 1}}$$

$$\underline{\underline{D = 2}}$$

$$\underline{\underline{D = 3}}$$