Longest Palindromic Subsequence

Ex. a b c b c b → b c c b

String → longest Pal. Subseq. LPS

a b c d

lps(abcd) → lps(String) / or lps(Str)

lps(Str) → S(Str) → set of seq. of Str.

S(abcd) = a S(bc) d (a: option, d: opt)

S(abcd) →

- - - -
- - - d
- - c -
- - c d
- b - -
- b - d
- bc -
- bc d
a - - -
a - - d
a - c -
a - c d
a b - -
a b - d
a bc -
a bc d

S(bc) → a & d

d
- - -
- c -
- b -
- bc -

a & d

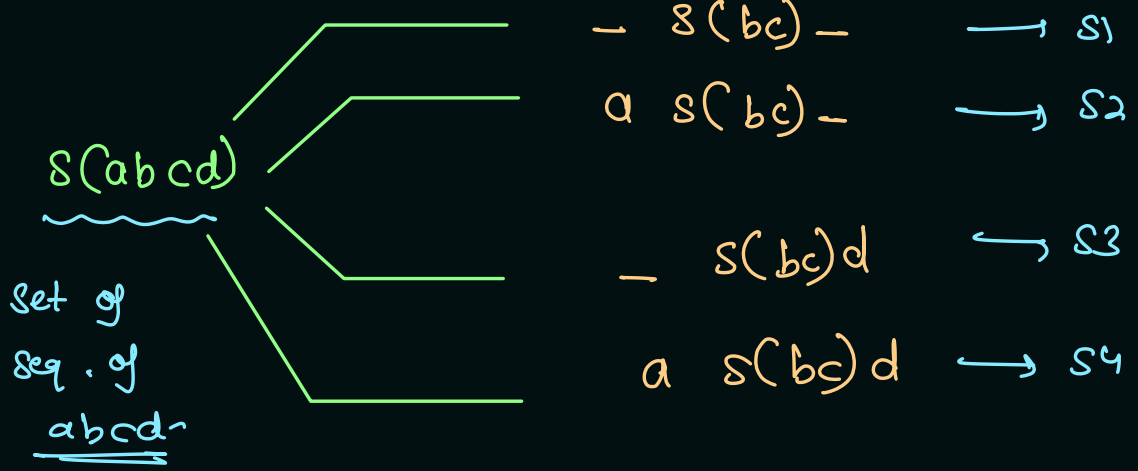
d
- - d
- c d
- b d
- bc d

'd' & d

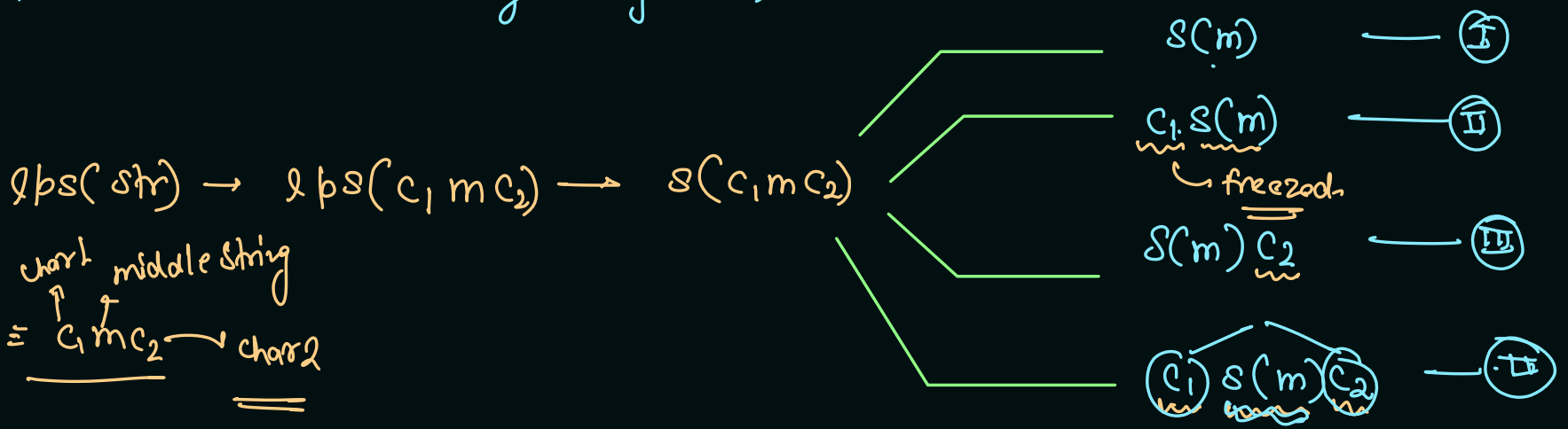
a
- - -
- c -
b - -
bc -

a, a

a
- - d
- c d
b - d
bc d



Now move toward generic ity →



$c_1 == c_2$	$c_1 \neq c_2$
[I, II, III] ✗	IV ✗ Result
IV ✓	[I, II, III] ✓
l(Str) = 2 + l(m)	longest from I, II & IV

lps(str) → l(c1 m c2) → s(c1 m c2) → s(m), c1 s(m), s(m) c2, c1 s(m) c2

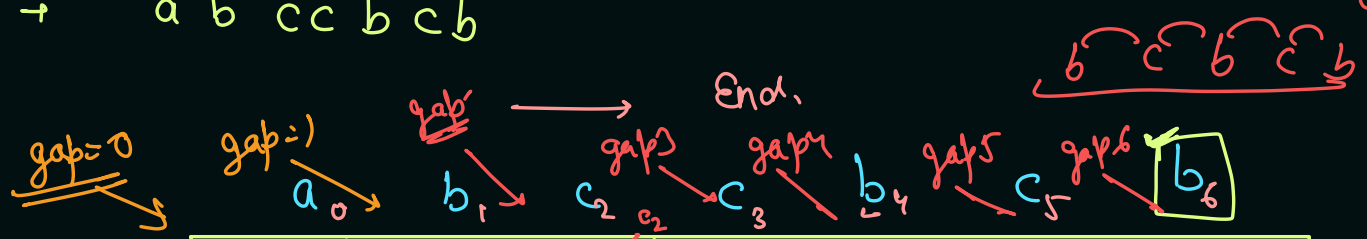
l(c1 m) → s(c1 m) → s(m), c1 s(m)

l(m c2) → s(m c2) → s(m), s(m) c2

a b c c b c b

$c_1 == c_2$	$c_1 \neq c_2$
④ ↳ [①, ②, ③] ×	④ × [①, ②, ③] ↳
$l(str) = 2 + l(m)$	$l(str) = \max(l(c_1 m), l(m c_2))$
$l(c_1 m c_2) = 2 + l(m)$	$l(c_1 m c_2) = \max(l(c_1 m), l(m c_2))$

string → a b cc b cb



	gap=0	gap=1	gap=2	gap=3	gap=4	gap=5	gap=6
a_0	1 a	1 ab	1 abc	2 abcc	4 abccb	4 abccbc	5 abccbcb
b_1	x	1 b	1 bc	2 bcc	4 bccb	4 bccbc	5 bccbcb
c_2	x	x	1 c	2 cc	2 ccb	3 (1,c) cbc	3 cbcb
c_3	x	x	x	1 c	1 (x+1,c) cb	3 cbc	3 cbcb
b_4	x	x	x	x	1 b	1 bc	2 (4,c) bcb
c_5	x	x	x	x	x	1 c	1 cb
b_6	x	x	x	x	x	x	1 b

abc → $c_1 m c_2$

$c_1 m$ → ab

$m c_2$ → bc

a b c c

m

c_1 c_2

$c_1 m$ → abc

$m c_2$ → bcc

Count all palindromic subseq. from given string :

count(S) → C[S(str)] → C[S(ch1 m ch2)]

C[S(m)] → C1 (Count)

C[ch1 S(m)] → C3

C[S(m) ch2] → C2

C[ch1 S(m) ch2] → C4

bc

m
S(bc) → (q) - - - (q)
q - c q
b - q
q bc q

Palindrome → (2)

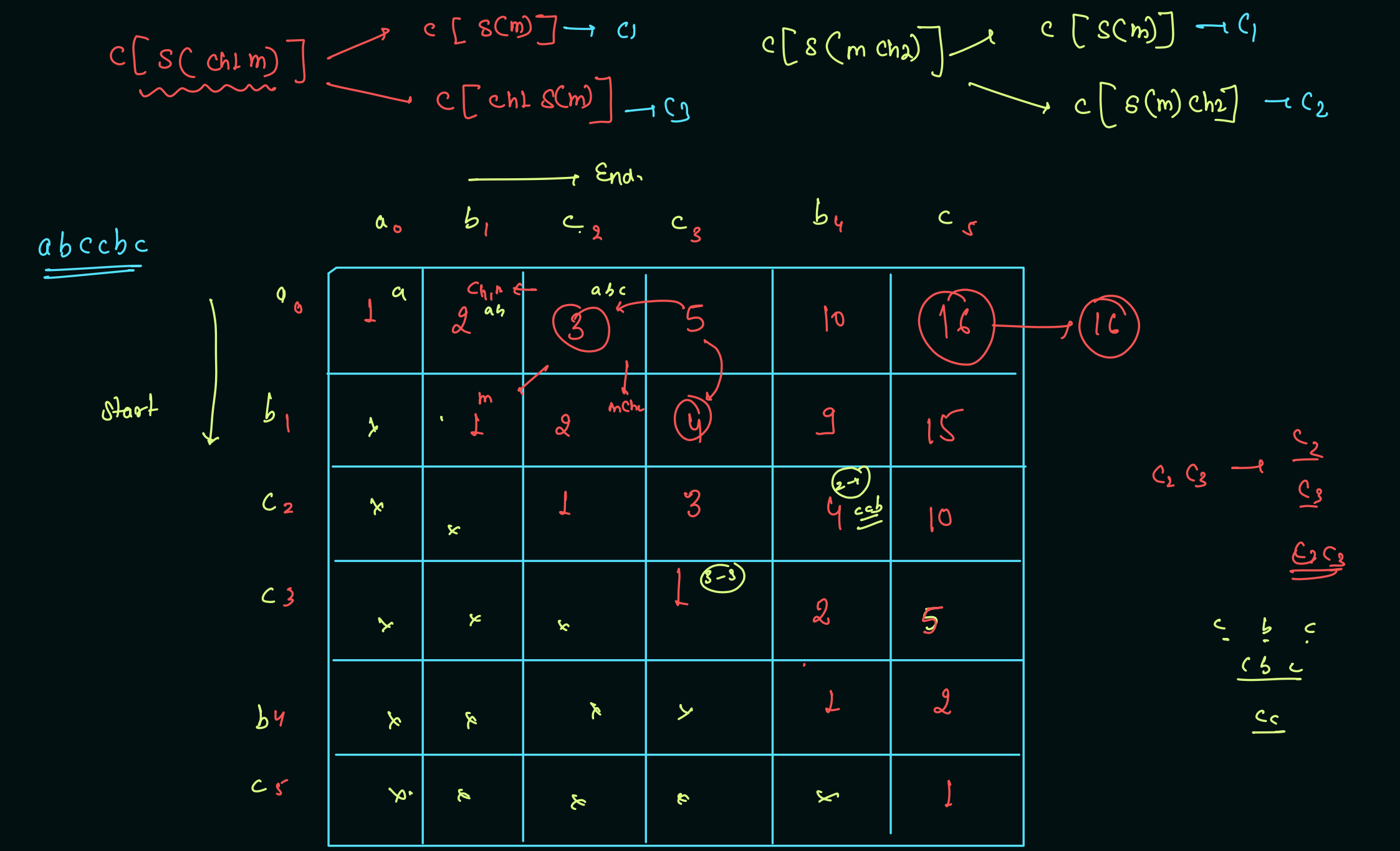
S(m) → (2)

@S(m)@ = S(m) + 1

Total count [without impact] → $C_1 + C_2 + C_3 + C_4$.

$ch1 == ch2$	$ch1 \neq ch2$
Total count = $C_1 + C_2 + C_3 + C_4$	Total Count = $C_1 + C_2 + C_3 + C_4$ → [0]
we know that $ch1 == ch2$	= $C_1 + C_2 + C_3 + C_1 - C_1$
$\Rightarrow C_4 = C_1 + 1$	= $(C_1 + C_2) + (C_1 + C_3) - C_1$
Total Count = $C_1 + C_2 + C_3 + C_4 + 1$	Total Cut = $C[S(m\,ch_2)] + C[S(ch_1\,m)] - C[S(m)]$
= $(C_1 + C_2) + (C_1 + C_3) + 1$	

Total Cout = $C[S(m\,ch_2)] + C[S(ch_1\,m)] + 1$


C[S(ch1 m)]
C[S(m)] → C1
C[ch1 S(m)] → C2
C[S(m ch2)]
C[S(m)] → C1
C[S(m) ch2] → C2
End.
a0 b1 c2 c3 b4 c5
abccbc
Start
q0
b1
c2
c3
b4
c5
a
ch1 m
ab
abc
m
mCh1
2-1
ccab
3-3
c2 c3 → c2 c3
c b c
cb c
cc

Ugly Numbers →

0 based → 0 1 2 3 4 5 6 7 8 9 10 wrong.
indx 1 2 3 4 5 6 8 9 10 12 14

2 → 2 4 6 8 10 12 14
3 → 3 6 9 12 15
5 → 5 10 15

val * indx

ugly Number Seq

2,3,5

1 2 2 3 4 5 6 8 9 10 12 15 16 ---
0 1 2 3 4 5 6 7 8 9 10 11 12

sp

indx

2p → 0 1 2 3 4 5 6
3p → 0 1 2 3 4 5
5p → 0 1 2 3

2,3,5

2×3 ✓

2×4 = 8
3×3 = 9
5×2 = 10

2×5 = 10
3×3 = 9
5×2 = 10

2×5 = 10
3×4 = 12
5×2 = 10

5×3 = 15
3×4 = 12
2×6 = 12

3×5 = 15
5×3 = 15
2×8 = 16

2×8 = 16
3×6 = 18
5×4 = 20