

1 Introduction

A table of frequently used symbols:

Symbol	Representation
g	the input graph
v	the set of nodes in g
e	the set of edges in g
n	a node in v
s	the start node in v
t	the target node in v
f	the number of atoms to conserve (flow to move)
G	the constructed metagraph
V	the set of metanodes in G
E	the set of metaedges in G
N	a metanode in V of the form $\{n_i : x_i\}$ where $n_i \in v$ and x_i is the number of atoms at n_i

Each $N \in V$ is composed of i many nodes from v . A metanode takes the form $\{n_i : x_i\}$ where $n_i \in v$ and x_i is the number of atoms at n_i . The $\sum_0^i x_i = k$ at any given N ensuring that all atoms are accounted for at each metanode.

The terms *state* and *metanode* are used interchangeably. Each metanode is simply a representation of where flow is distributed. The state $\{s : f\}$ indicates that there is f amount of flow at node s . Similarly, $\{a : 3, b : 2\}$ indicates that there are 3 atoms at node a and 2 atoms at node b .

The goal is to move all f atoms (flow) from node s to node t . In other words, we seek a path from metanode $\{s : f\}$ to metanode $\{t : f\}$ in the metagraph.

2 Overview

Algorithm 1: Metagraph Generation

Input: g, s, t, f
Output: G

```
1  $S = \{s : f\}$  /* Create the start state */
2  $V \leftarrow V \cup S$  /* Add start state to the metagraph */
3  $stack.push(S)$ ;
4 while  $|stack| \neq 0$  do
5    $current \leftarrow stack.pop()$ ;
6    $partials \leftarrow$  initialize empty list;
7   for each inner node,  $n$ , in  $current$  do
8      $innerPartials \leftarrow$  list of all unique ways to move the flow residing at  $n$ ;
9      $partials.append(innerPartials)$ ;
10   $nbrs \leftarrow$  all complete states, where each complete state is formed by merging together one partial
    state from each list in  $partials$ ;
11  for  $nbr \in nbrs$  do
12    if  $nbr$  is valid and  $nbr \notin deadset$  then
13       $V \leftarrow V \cup nbr$ ;
14       $E \leftarrow E \cup (current, nbr)$ ;
15       $nbrCount \leftarrow nbrCount + 1$ ;
16       $stack.push(nbr)$ ;
17 return  $G$ ;
```

3 Algorithms

Algorithm 2: ConstructMetaGraph

Input: $g = (v, e)$: the input graph, $\{s, t\} \in v$: the start and target compounds, k : flow/number of atoms to conserve

Output: $G = (V, E)$: the meta-graph

```

1  $MG \leftarrow (MV = \emptyset, ME = \emptyset)$  /* initialize new metagraph */
2  $stack \leftarrow$  initialize empty stack;
3  $start \leftarrow \{s : k\}$  /* Create a metanode with all  $k$  atoms at the start compound */
4  $V \leftarrow V \cup start$  /* Add the start state to the set of metanodes */
5  $stack.push(start)$  /* Add the start state to the stack to find its neighbors */
6  $G \leftarrow \text{PopulateMetaGraph}(G, stack, target)$  /* Find neighboring metanodes to build  $G$  */
7 return  $G$ 
```

Algorithm 3: PopulateMetaGraph

Input: $G = (V, E)$: the metagraph, $stack$: stack of metanodes that need to be explored, $target$: the metanode state with all k atoms at node t

Output: $G = (V, E)$: metagraph with any newly found metanodes added in

```

1 while  $|stack| > 0$  do
2    $current \leftarrow stack.pop()$  /* Pop off a metanode to explore */
3   if  $current = target$  then
4     continue /* If we've reached the target, no need to find nbrs */
5   else
6      $nbrCount \leftarrow \text{IterativeFindMetaNbrs}(current)$ ;
7     if  $nbrCount = 0$  and  $current \neq target$  then
8        $G \leftarrow \text{Prune}(current)$  /* This metanode is a terminus, so remove it from the
          metagraph */
9 return  $G$ 
```

Algorithm 4: IterativeFindMetaNbrs

Input: *current*: the metanode to find nbrs for, $G = (V, E)$: the metagraph so far

Output: $G = (V, E)$: metagraph with any newly found metanodes added in

```
1 partialStates  $\leftarrow$  initialize new list;
2 for  $n \in N$  do
3    $\lfloor$  partialStates.append(RecursiveNbrSearch(n, flow(n), nbrs(n)));
4 metanbrs  $\leftarrow$  generateCompleteStates(partialStates);
5 for  $nbr \in \text{metanbrs}$  do
6   if isValid(nbr) and  $nbr \notin \text{deadset}$  then
7      $V \leftarrow V \cup nbr$ ;
8      $E \leftarrow E \cup (current, nbr)$ ;
9     nbrCount  $\leftarrow$  nbrCount + 1;
10     $\lfloor$  stack.push(nbr);
11 return  $G$ 
```

Algorithm 5: RecursiveNbrSearch

Input: *n*: the node we're trying to move flow away from, *flow*: the amount of flow residing at *n* that needs to be moved, *nbrs*: the unexplored neighbors of *n*

Output: *partialStates*: list of partial states (each state represents a unique move of the flow at *n* to its nbrs)

```
1 states  $\leftarrow$  initialize new list;
2 if |nbrs|  $\neq 1$  then
3    $nbr \leftarrow \text{nbrs}[0]$ ;
4   for  $i = 0; i \leq \min(\text{flow}, \text{capacity}(\text{parent}, nbr)); i++$  do
5     partialState  $\leftarrow$  initialize mapping of nodes (string) to flow (number);
6     if  $i \neq 0$  then
7        $\lfloor$  partialState.put(nbr, i);
8       remaining  $\leftarrow$  recursiveNbrSearch(n, flow-i, nbrs-nbr)
9       for  $state \in \text{remaining}$  do
10         $\lfloor$  states.add(mergePartialStates(partialState, state))
11 if |nbrs| = 0 and flow  $\neq 0$  and current = target then
12   /* If there is flow remaining, but the current node is the target node, the flow
13      can remain stationary */
14    $\lfloor$  state.put(current, flow); states.add(state);
15 else
16   /* If there are no more nbrs and flow remains, no more valid states can be added,
17      so return */
18 return states
```
