Modelling and Optimization

INF170

#6:Transportation and Assignment Problem

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AGENDA

• Transportation Model

• Assignment Model

A single good is to be shipped from several origins to several destinations at minimum overall cost.

TRANSPORTATION PROBLEM

Suppose that we have decided to produce steel coils at three mill locations, in the following amounts:

SUM		6900
PITT	Pittsburgh, Pennsylvania	2900
CLEV	Cleveland, Ohio	2600
GARY	Gary, Indiana	1400

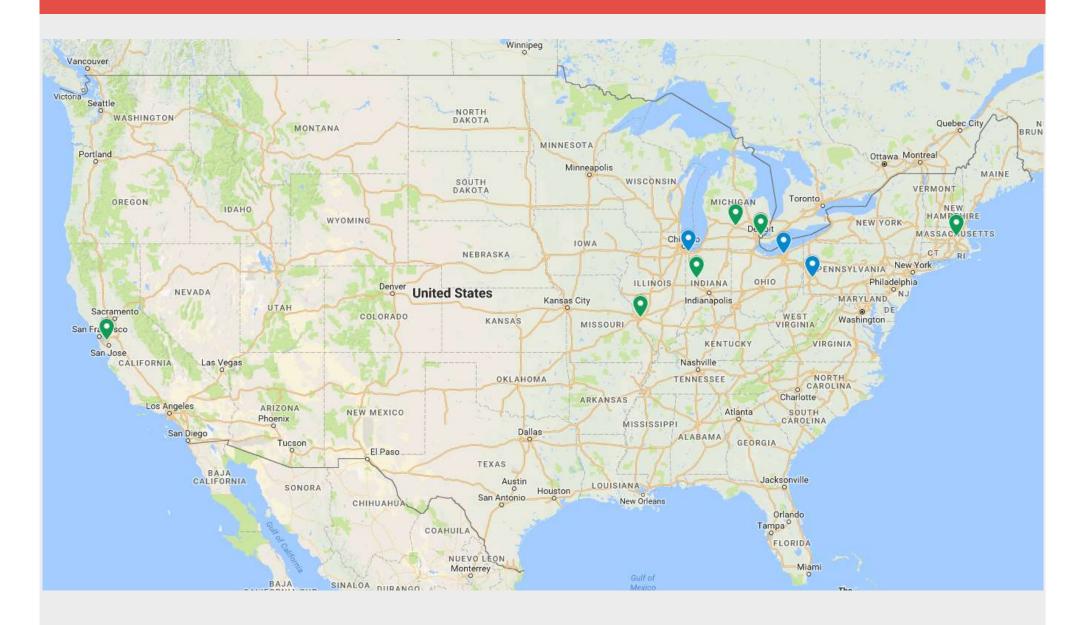
The total of 6,900 tons must be shipped in various amounts to meet orders at seven locations of automobile factories:

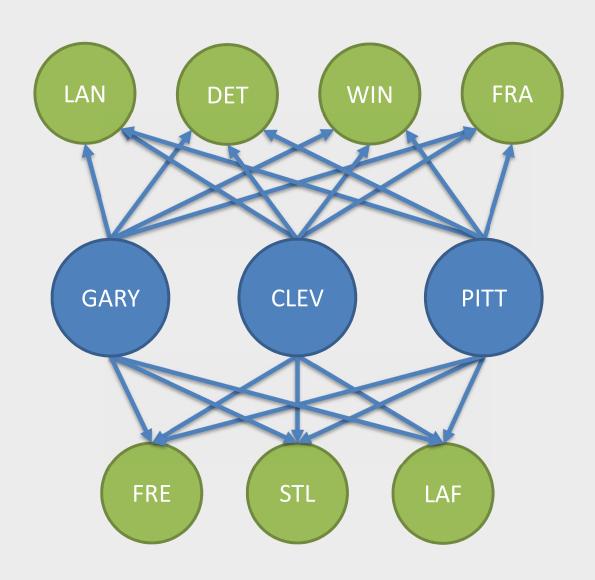
FRA	Framingham, Massachusetts	900
DET	Detroit, Michigan	1200
LAN	Lansing, Michigan	600
WIN	Windsor, Ontario	400
STL	St. Louis, Missouri	1700
FRE	Fremont, California	1100
LAF	Lafayette, Indiana	1000
SUM		6900

Table of shipping costs per ton:

	GARY	CLEV	PITT
FRA	39	27	24
DET	14	9	14
LAN	11	12	17
WIN	14	9	13
STL	16	26	28
FRE	82	95	99
LAF	8	17	20

What is the least expensive plan for shipping the coils from mills to plants?





Variables:

GARY:FRA → the number of tons to be shipped from GARY to FRA

... (21 decision variables in all! One for each combination of mill and factory)

Objective function: Minimize

	GARY	CLEV	PITT
FRA	39	27	24
DET	14	9	14
LAN	11	12	17
WIN	14	9	13
STL	16	26	28
FRE	82	95	99
LAF	8	17	20

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What about supplying the factories from the mill that can ship most cheaply to it!

	GARY	CLEV	PITT		GARY	CLEV	PITT
FRA	39	27	24	FRA	0	0	900
DET	14	9	14	DET	0	1200	0
LAN	11	12	17	LAN	600	0	0
WIN	14	9	13	WIN	0	400	0
STL	16	26	28	STL	1700	0	0
FRE	82	95	99	FRE	1100	0	0
LAF	8	17	20	LAF	1000	0	0
				SUM	4400	1600	900

Shipments from GARY to the seven factories is equal to the production level of 1400

GARY:FRA + GARY:DET + GARY:LAN + GARY:WIN +

GARY:STL + GARY:FRE + GARY:LAF = 1400

GARY:FRA + GARY:DET + GARY:LAN + GARY:WIN +

GARY:STL + GARY:FRE + GARY:LAF = 1400

CLEV:FRA + CLEV:DET + CLEV:LAN + CLEV:WIN +

CLEV:STL + CLEV:FRE + CLEV:LAF = 2600

PITT:FRA + PITT:DET + PITT:LAN + PITT:WIN +

PITT:STL + PITT:FRE + PITT:LAF = 2900



At FRA, the sum of the shipments received from the

three mills must equal the 900 tons ordered:

GARY:FRA + CLEV:FRA + PITT:FRA = 900

GARY:FRA + CLEV:FRA + PITT:FRA = 900

GARY:DET + CLEV:DET + PITT:DET = 1200

GARY:LAN + CLEV:LAN + PITT:LAN = 600

GARY:WIN + CLEV:WIN + PITT:WIN = 400

GARY:STL + CLEV:STL + PITT:STL = 1700

GARY:FRE + CLEV:FRE + PITT:FRE = 1100

GARY:LAF + CLEV:LAF + PITT:LAF = 1000

Two fundamental sets of objects:

- ✓ The sources or origins (mills, in our example)
- ✓ The destinations (factories, in our example)

Declaration of these two sets in AMPL model:

```
set ORIG;
set DEST;
```

There is

- a supply of something at each origin (tons of steel coils produced, in our case)
- > a demand for the same thing at each destination (tons of coils ordered).

```
param supply {ORIG} >= 0;
param demand {DEST} >= 0;
```

```
check: sum {i in ORIG} supply[i] = sum {j in DEST} demand[j];
```

For each combination of an origin and a destination, there is a transportation cost and a variable representing the amount transported.

```
param cost {ORIG, DEST} >= 0;
var Trans {ORIG, DEST} >= 0;
```

For a particular origin i and destination j,

We ship Trans[i,j] units from i to j, at a cost of cost[i,j] per unit

> The total cost for this pair is

```
cost[i,j] * Trans[i,j]
```

Adding over all pairs, we have the objective function:

```
minimize Total_Cost:
sum {i in ORIG, j in DEST} cost[i,j] * Trans[i,j];
```

These are the same:

```
sum {i in ORIG, j in DEST} cost[i,j] * Trans[i,j];
```

```
sum {j in DEST, i in ORIG} cost[i,j] * Trans[i,j];
```

```
sum {i in ORIG} sum {j in DEST} cost[i,j] * Trans[i,j];
```

Constraints:

Declaration in AMPL model:

```
subject to Supply {i in ORIG}: ...
subject to Demand {j in DEST}: ...
```

- > Note that the names supply and Supply are unrelated
- > AMPL distinguishes upper and lower case.

Constraints:

Supply constraint for origin *i*:

 \triangleright The sum of all shipments out of i is equal to the supply available.

```
subject to Supply {i in ORIG}:
sum {j in DEST} Trans[i,j] = supply[i];
```

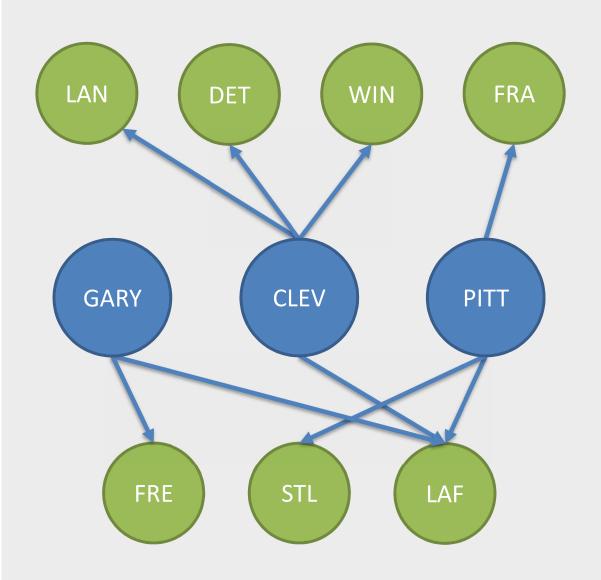
Constraints:

Demand constraint for destination *j*:

The sum of all shipments into *j* is equal to the demand available.

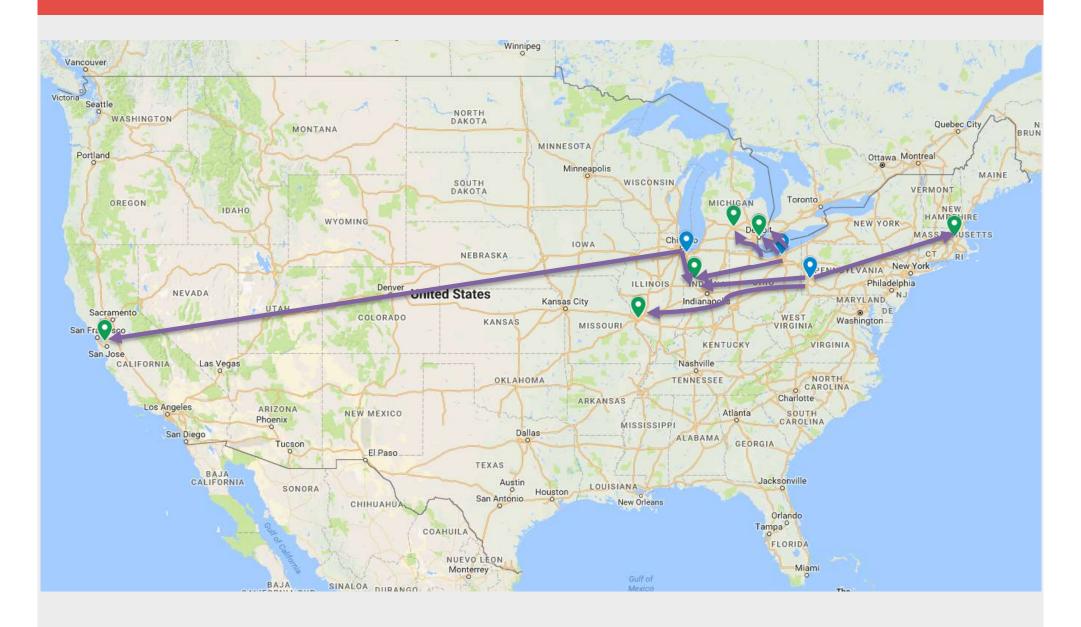
```
subject to Demand {j in DEST}:
sum {i in ORIG} Trans[i,j] = demand[j];
```

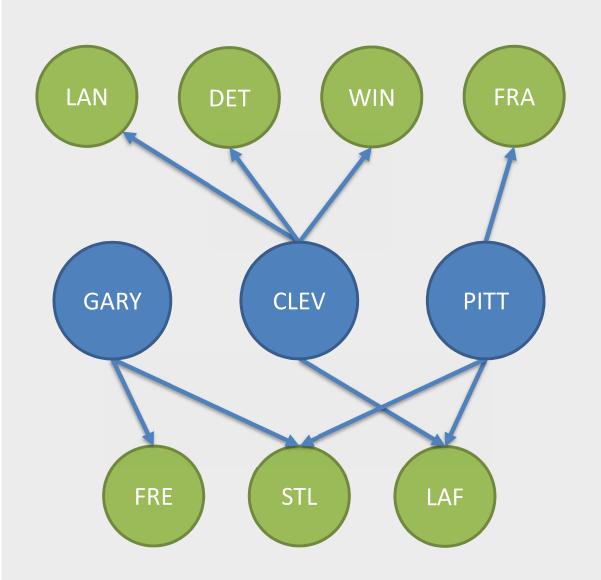
```
set ORIG; # origins
set DEST; # destinations
param supply {ORIG} >= 0; # amounts available at origins
param demand {DEST} >= 0; # amounts required at destinations
check: sum {i in ORIG} supply[i] = sum {j in DEST} demand[j];
param cost {ORIG, DEST} >= 0; # shipment costs per unit
var Trans {ORIG, DEST} >= 0; # units to be shipped
minimize Total Cost:
sum {i in ORIG, j in DEST} cost[i,j] * Trans[i,j];
subject to Supply {i in ORIG}:
sum {j in DEST} Trans[i,j] = supply[i];
subject to Demand {j in DEST}:
sum {i in ORIG} Trans[i,j] = demand[j];
```



	CLEV	GARY	PITT
DET	1200	0	0
FRA	0	0	900
FRE	0	1100	0
LAF	400	300	300
LAN	600	0	0
STL	0	0	1700
WIN	400	0	0

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	CLEV	GARY	PITT
DET	1200	0	0
FRA	0	0	900
FRE	0	1100	0
LAF	400	0	600
LAN	600	0	0
STL	0	300	1400
WIN	400	0	0

$$minimize \sum_{i \in N_1} \sum_{j \in N_2} c_{i,j} x_{i,j}$$

subject to

$$\sum_{j \in N_2} x_{i,j} = s_i$$

for all $i \in N_1$

(Supply)

$$\sum_{i \in N_1} x_{i,j} = d_j$$

for all $j \in N_2$

(Demand)

$$x_{i,j} \geq 0$$

ASSIGNMENT PROBLEM

Consider a department that needs to assign some number of people to an equal number of offices.

- > The origins
- **Destinations**

Consider a department that needs to assign some number of people to an equal number of offices.

- > The origins now represent individual people
- > Destinations represent individual offices

- Since each person is assigned one office, and each office is occupied by one person,
 - \triangleright supply[i] = 1
 - \rightarrow demand[j] = 1
- Interpret *Trans*[*i,j*] as the "amount" of person *i* that is assigned to office j
 - \succ Trans[i,j] is 1 then person i will occupy office j
 - \succ *Trans*[*i,j*] is 0 then person *i* will not occupy office *j*.

Objective function?

- ✓ Ask people to rank the offices, giving their first choice, second choice, and so forth. Then let cost[i,j] be the rank that person i gives to office j.
- ✓ Assignment that will please a lot of people!
- ✓ Same model as transportation model but new data!

ASSIGNMENT PROBLEM

```
set ORIG := Coullard Daskin Hazen Hopp Iravani Linetsky Mehrotra Nelson
Smilowitz Tamhane White ;
set DEST := C118 C138 C140 C246 C250 C251 D237 D239 D241 M233 M239;
param supply default 1;
param demand default 1;
param cost:
          C118 C138 C140 C246 C250 C251 D237 D239 D241 M233 M239 :=
Coullard 6
                                    10
                                               5
                                                    3
                               11
                                          4
Daskin
                                    10
Hazen
               10
                     11
                                    6
                     8
         11
                          10
                                    5
Hopp
                     8
                          9
Travani
                               10
                                    11
                     10
Linetsky
         11
Mehrotra
               11
                     10
                          9
                     4
                                               9
                                                    10
Nelson
          11
Smilowitz 11
                                               3
                     10
                                                    4
Tamhane
                                               10
                                                    11
White
                          4
                                    5
                                               10
          11
                                                              1 ;
```

ASSIGNMENT PROBLEM

Coullard	C118	6
Daskin	D241	4
Hazen	C246	1
Норр	D237	1
Iravani	C138	2
Linetsky	C250	3
Mehrotra	D239	2
Nelson	C140	4
Smilowitz	M233	1
Tamhane	C251	3
White	M239	1

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$$minimize \sum_{(i,j)\in A} c_{i,j} x_{i,j}$$

subject to

$$\sum_{(i,k)\in A} x_{i,k} = 1$$

for all $i \in N_1$

(Supply)

$$\sum_{(k,j)\in A} x_{k,j} = 1$$

for all $j \in N_2$

(Demand)

$$x_{i,j} \geq 0$$

for all $(i, j) \in A$ (Nonnegativity)

But how did we know that every Trans[i,j] would equal either 0 or 1 in the optimal solution, rather than, say, 0.5?

Special property of transportation models, which guarantees that

- ✓ as long as all supply and demand values are integers, and
- ✓ as long as all all lower and upper bounds on the variables are integers, there will be an optimal solution that is entirely integral.

But don't let this favorable result mislead you into assuming that integrality can be assured in all other circumstances; even in examples that seem to be much like the transportation model, finding integral solutions can require a special solver, and a lot more work.

$$minimize \sum_{i \in N_1} \sum_{j \in N_2} c_{i,j} x_{i,j}$$

subject to

$$\sum_{j \in N_2} x_{i,j} = s_i$$

for all $i \in N_1$

(Supply)

$$\sum_{i \in N_1} x_{i,j} = d_j$$

for all $j \in N_2$

(Demand)

$$x_{i,j} \geq 0$$

$$\sum s_i \geq \sum d_j$$

$$minimize \sum_{i \in N_1} \sum_{j \in N_2} c_{i,j} x_{i,j}$$

subject to

$$\sum_{j \in N_2} x_{i,j} \le s_i$$

for all $i \in N_1$

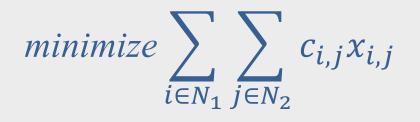
(Supply)

$$\sum_{i \in N_1} x_{i,j} \ge d_j$$

for all $j \in N_2$

(Demand)

$$x_{i,j} \geq 0$$



subject to

$$\sum_{j \in N_2} x_{i,j} \le s_i$$

 $\sum_{i \in N_1} x_{i,j} \le d_i$

$$x_{i,j} \geq 0$$

for all $i \in N_1 \cap V$

(Supply)

for all $j \in N_2$

(Demand)

$$\sum s_i \leq \sum d_j$$

$$minimize \sum_{i \in N_1} \sum_{j \in N_2} c_{i,j} x_{i,j}$$

subject to

$$\sum_{j \in N_2} x_{i,j} \ge s_i$$

for all $i \in N_1$

(Supply)

$$\sum_{i \in N_1} x_{i,j} \le d_j$$

for all $j \in N_2$

(Demand)

$$x_{i,j} \geq 0$$

Use of Dummy Sources!

- ✓ Add a dummy Origin in case of shortage in supply!
 - With capacity equal to shortage and new links with high costs
- ✓ Add a dummy destination in case of shortage in demand!
 - With capacity equal to shortage and new links with high costs

Use of Penalty!

$$minimize \sum_{i \in N_1} \sum_{j \in N_2} c_{i,j} x_{i,j} + P \left[\sum_{j \in N_2} \left(d_j - \sum_{i \in N_1} x_{i,j} \right) \right]$$

subject to

$$\sum_{j \in N_2} x_{i,j} \le s_i$$

for all $i \in N_1$

(Supply)

$$\sum_{i \in N_1} x_{i,j} \le d_j$$

for all $j \in N_2$

(Demand)

$$x_{i,j} \geq 0$$

$$minimize \sum_{i \in N_1} \sum_{j \in N_2} c_{i,j} x_{i,j} + \sum_{j \in N_2} P_j \left(d_j - \sum_{i \in N_1} x_{i,j} \right)$$

subject to

$$\sum_{j \in N_2} x_{i,j} \le s_i$$

for all $i \in N_1$

(Supply)

$$\sum_{i \in N_1} x_{i,j} \le d_j$$

for all $j \in N_2$

(Demand)

$$x_{i,j} \geq 0$$

ASSIGNMENT

ASSIGNMENT #3:

AMPL BOOK
CHAPTER 3. EXERCISES(1-4)

IN CLASS ASSIGNMENT (TEST)

IN CLASS ASSIGNMENT (TEST) #1:

- ✓ Closed-book
- ✓ Based on slides
- ✓ No need of AMPL

NEXT LECTURE

LECTURE #7:

NETWORK OPTIMIZATION

