# Modelling and Optimization

**INF170** 

#13:Heurisitcs for TSP

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Dept. of Informatics
University of Bergen

Fall Semester 2018



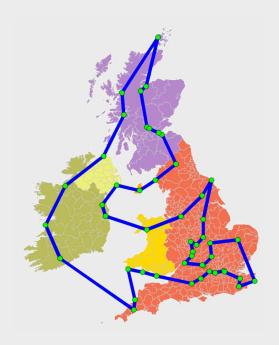
### AGENDA

Construction heuristics

• Improvement heuristics (neighborhood search)

#### Travelling Salesman Problem

• Starting from an origin node, find the minimum distance required to visit each node once and only once and return to the origin.



### TSP - IMPLEMENTATION

20 nodes:

MTZ gain(Svestka)

Steps(Dantzig)

nvars 400 7600

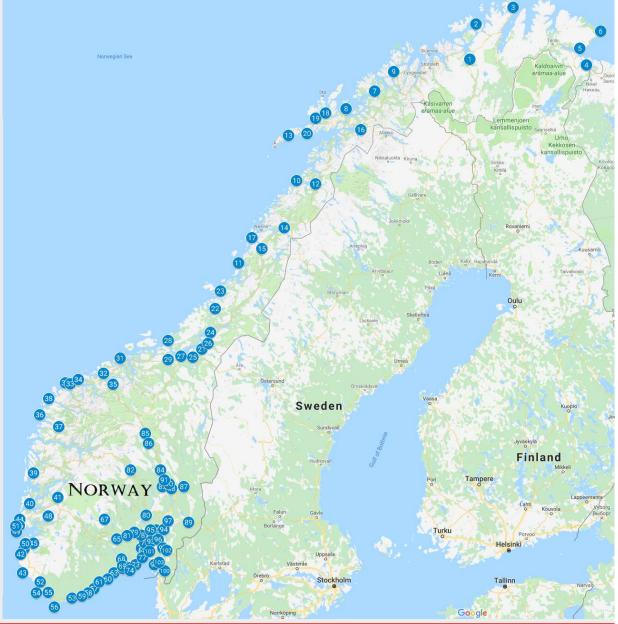
ncons 420 459 420

Time 0.14 0.23 7.4

# TSP - Implementation

# 103 cities in Norway





### TSP - IMPLEMENTATION

103 cities in Norway:

MTZ gain(Svestka) Steps(Dantzig)

nvars 10609 21012 1082120

ncons 10712 10917 10712

Time ? ?

- Why do we use a heuristic method to solve a TSP?
  - The problem is difficult (known to be NP-Hard)
  - No polynomial time algorithm for solving it to optimality
  - Exponential in the number of cities
  - We must solve relatively "large" instances of the problem
- Heuristics aims to efficiently generate very good solutions. They do
  not find the optimal solution, or at least do not guarantee the
  optimality of the found solutions.

#### Type of heuristics for TSP

- Construction heuristics:
  - builds a solution from scratch (starting with nothing).
- Improvement heuristics (neighborhood search):
  - starts with a solution, and then tries to improve the solution, usually by making small changes in the current solution.
- Metaheuristics
  - a high-level problem-independent algorithmic framework that combines
     operators/heuristics intelligently and provide a sufficiently good solution!

- Construction heuristics:
  - Nearest Neighbor Heuristic
  - Greedy Heuristic
  - Insertion Heuristics
    - Nearest insertion of arbitrary city
    - Nearest insertion
    - Farthest Insertion
    - Cheapest insertion
  - Christofides algorithm

• Improvement heuristics:

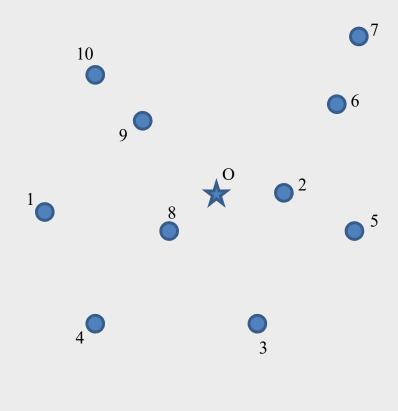
- 2opt
- 3opt
- k-opt

• Starting from an origin node, find the minimum distance required to visit each node once and only once and return to the origin.

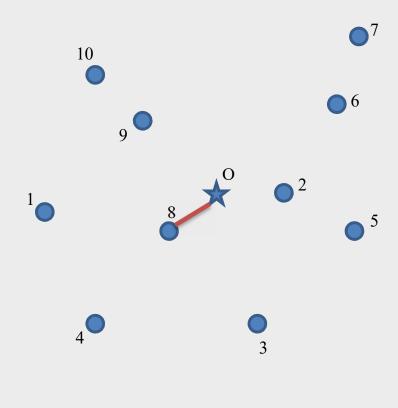
# Nearest Neighbor Heuristic:

- 1. Select any node to be the active node
- 2. Connect the active node to the closest unconnected node, make that the new active node.
- 3. If there are more unconnected nodes go to step 2, otherwise connect to the starting node and end.

Dis.	1	2	3	4	5	6	7	8	9	10	О
1	0	10	10.8	6.3	13	13.4	16.4	5.1	6.4	8.2	7.1
2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
7	16.4	9.5	16.5	19.4	11	4.1	0	13.6	10.3	11.2	10.8
8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0

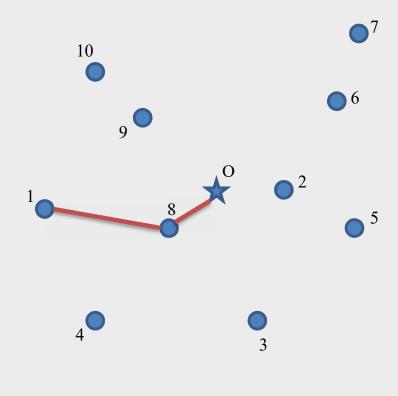


Dis.	1	2	3	4	5	6	7	8	9	10	0
1	0	10	10.8	6.3	13	13.4	16.4	5.1	6.4	8.2	7.1
2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
7	16.4	9.5	16.5	19.4	11	4.1	0	13.6	10.3	11.2	10.8
8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0



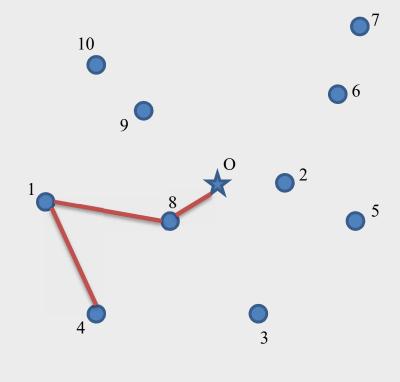
Tour: O-8

Dis.	1	2	3	4	5	6	7	8	9	10	0
1	0	10	10.8	6.3	13	13.4	16.4	5.1	6.4	8.2	7.1
2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
7	16.4	9.5	16.5	19.4	11	4.1	0	13.6	10.3	11.2	10.8
8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0



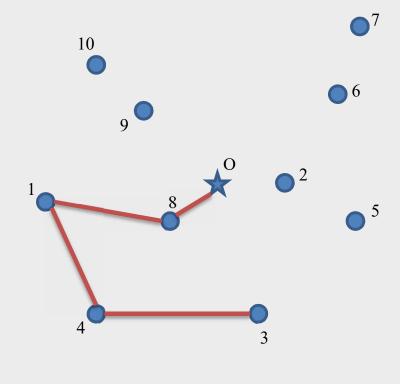
Tour: 0-8-1

Dis.	1	2	3	4	5	6	7	8	9	10	0
1	0	10	10.8	6.3	13	13.4	16.4	5.1	6.4	8.2	7.1
2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
7	16.4	9.5	16.5	19.4	11	4.1	0	13.6	10.3	11.2	10.8
8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0



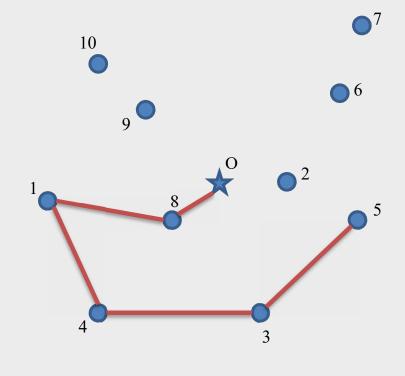
Tour: O-8-1-4

Dis.	1	2	3	4	5	6	7	8	9	10	0
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2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
7	16.4	9.5	16.5	19.4	11	4.1	0	13.6	10.3	11.2	10.8
8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0



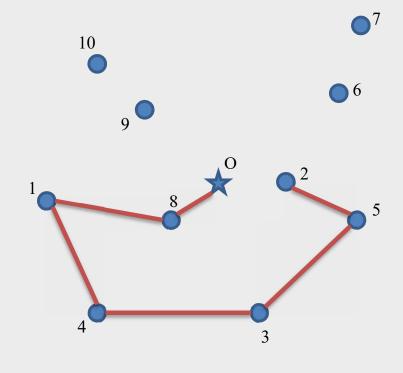
Tour: O-8-1-4-3

Dis.	1	2	3	4	5	6	7	8	9	10	0
1	0	10	10.8	6.3	13	13.4	16.4	5.1	6.4	8.2	7.1
2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
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8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0



Tour: O-8-1-4-3-5

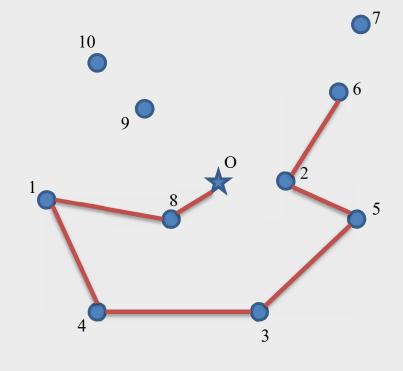
Dis.	1	2	3	4	5	6	7	8	9	10	0
1	0	10	10.8	6.3	13	13.4	16.4	5.1	6.4	8.2	7.1
2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
7	16.4	9.5	16.5	19.4	11	4.1	0	13.6	10.3	11.2	10.8
8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0



Tour: O-8-1-4-3-5-2

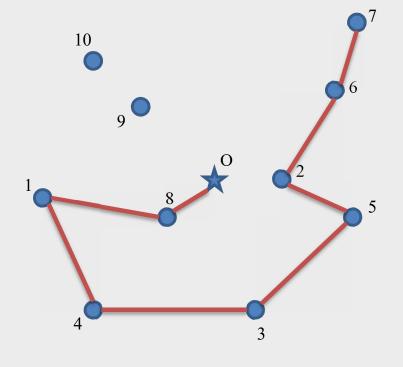
Ahmad Hemmati

Dis.	1	2	3	4	5	6	7	8	9	10	0
1	0	10	10.8	6.3	13	13.4	16.4	5.1	6.4	8.2	7.1
2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
7	16.4	9.5	16.5	19.4	11	4.1	0	13.6	10.3	11.2	10.8
8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0



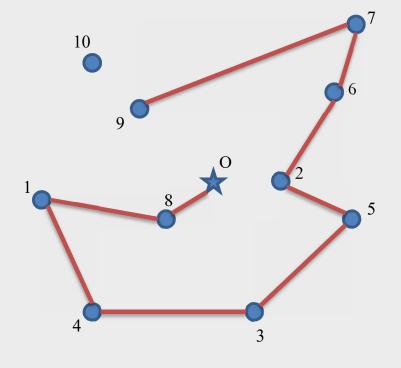
Tour: O-8-1-4-3-5-2-6

Dis.	1	2	3	4	5	6	7	8	9	10	0
1	0	10	10.8	6.3	13	13.4	16.4	5.1	6.4	8.2	7.1
2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
7	16.4	9.5	16.5	19.4	11	4.1	0	13.6	10.3	11.2	10.8
8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0



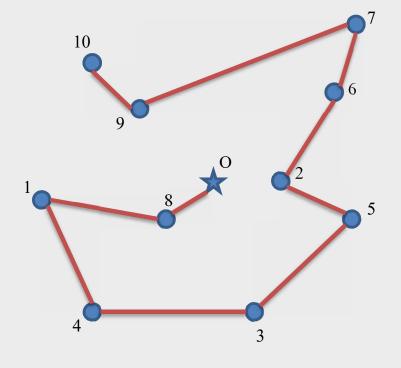
Tour: O-8-1-4-3-5-2-6-7

Dis.	1	2	3	4	5	6	7	8	9	10	0
1	0	10	10.8	6.3	13	13.4	16.4	5.1	6.4	8.2	7.1
2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
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8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0



Tour: O-8-1-4-3-5-2-6-7-9

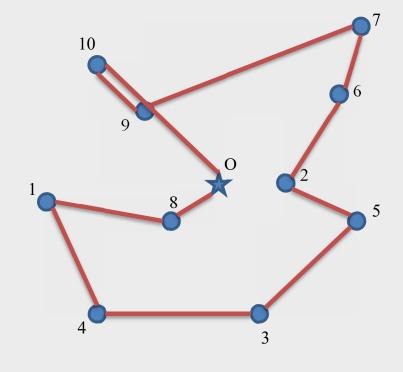
Dis.	1	2	3	4	5	6	7	8	9	10	0
1	0	10	10.8	6.3	13	13.4	16.4	5.1	6.4	8.2	7.1
2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
7	16.4	9.5	16.5	19.4	11	4.1	0	13.6	10.3	11.2	10.8
8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0



Tour: O-8-1-4-3-5-2-6-7-9-10-O

# TSP - NEAREST NEIGHBOR HEURISTIC

Dis.	1	2	3	4	5	6	7	8	9	10	0
1	0	10	10.8	6.3	13	13.4	16.4	5.1	6.4	8.2	7.1
2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
7	16.4	9.5	16.5	19.4	11	4.1	0	13.6	10.3	11.2	10.8
8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0



Tour: O-8-1-4-3-5-2-6-7-9-10-O

Length: 63.2

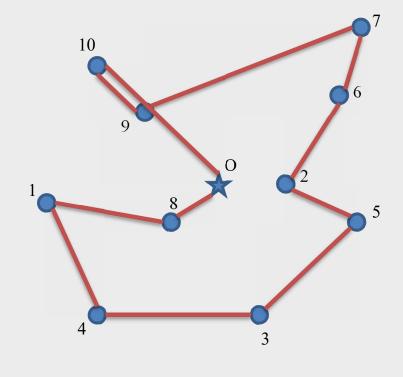
# TSP - 2-Opt

# Improvement Heuristic: 2-Opt

- 1. Identify pairs of arcs (*i-j* and *k-l*), where d(ij) + d(kl) > d(ik) + d(jl) (usually where they cross)
- 2. Select the pair with the largest difference, and re-connect the arcs (i-k and j-l)
- 3. Continue until there are no more crossed arcs.

# TSP - 2-Opt

Dis.	1	2	3	4	5	6	7	8	9	10	0
1	0	10	10.8	6.3	13	13.4	16.4	5.1	6.4	8.2	7.1
2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
7	16.4	9.5	16.5	19.4	11	4.1	0	13.6	10.3	11.2	10.8
8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0

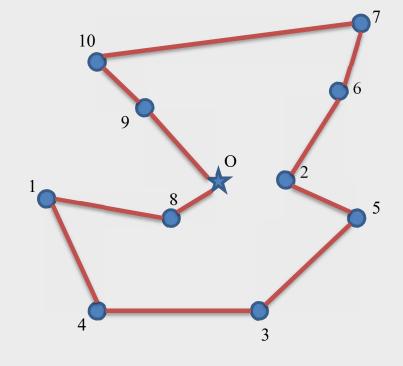


- Arcs 7-9 and 10-O cross
- d(79) + d(10-0) = 18.9 > d(7-10) + d(9-0) = 16.2
- Re-connect arcs 7-10 and 9-O



# TSP - 2-Opt

Dis.	1	2	3	4	5	6	7	8	9	10	0
1	0	10	10.8	6.3	13	13.4	16.4	5.1	6.4	8.2	7.1
2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
7	16.4	9.5	16.5	19.4	11	4.1	0	13.6	10.3	11.2	10.8
8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0



Tour: O-8-1-4-3-5-2-6-7-10-9-O

Tour length reduces from 63.2 to 60.5

Length: 60.5

### TSP - NEAREST NEIGHBOR HEURISTIC

- Initialization Start with a partial tour with just one city i,
   randomly chosen;
- 2. Selection Let (1, ..., k) be the current partial tour (k < n). Find city k + 1 that is not yet in the tour and that is closer to k.
- 3. Insertion Insert k + 1 at the end of the partial tour.
- 4. If all cities are inserted then STOP, else go back to 2.

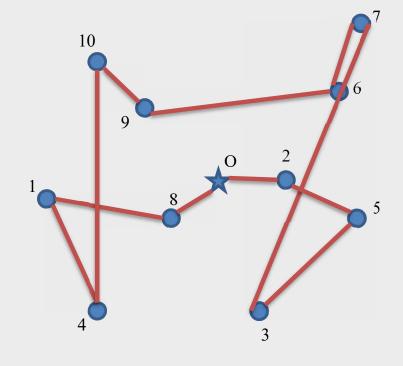
#### TSP - GREEDY HEURISTIC

- 1. Sort all edges.
- 2. Select the shortest edge and add it to our tour if it doesn't violate any of the above constraints.

3. Do we have *N* edges in our tour? If no, repeat step 2.

# TSP - GREEDY HEURISTIC

Dis.	1	2	3	4	5	6	7	8	9	10	0
1	0	10	10.8	6.3	13	13.4	16.4	5.1	6.4	8.2	7.1
2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
7	16.4	9.5	16.5	19.4	11	4.1	0	13.6	10.3	11.2	10.8
8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0



Tour: O-8-1-4-10-9-6-7-3-5-2-O

Length: 73.5

#### TSP - Nearest Insertion of Arbitrary City

- 1. Initialization Start with a partial tour with just one city i, randomly chosen; find the city j for which  $c_{ij}$  (distance or cost from i to j) is minimum and build the partial tour (i, j).
- 2. Selection Given a partial tour, arbitrary select a city k that is not yet in the partial tour.
- 3. Insertion Find the edge  $\{i, j\}$ , belonging to the partial tour, that minimizes  $c_{ik} + c_{kj} c_{ij}$ . Insert k between i and j.
- 4. If all cities are inserted then STOP, else go back to 2.



### TSP - NEAREST INSERTION

- 1. Initialization Start with a partial tour with just one city i, randomly chosen; find the city j for which  $c_{ij}$  (distance or cost from i to j) is minimum and build the partial tour (i, j).
- 2. Selection Find cities k and j (j belonging to the partial tour and k not belonging) for which  $c_{kj}$  is minimized.
- 3. Insertion Find the edge  $\{i, j\}$ , belonging to the partial tour, that minimizes  $c_{ik} + c_{kj} c_{ij}$ . Insert k between i and j.
- 4. If all cities are inserted then STOP, else go back to 2.

#### TSP - FARTHEST INSERTION

- 1. Initialization Start with a partial tour with just one city i, randomly chosen; find the city j for which  $c_{ij}$  (distance or cost from i to j) is minimum and build the partial tour (i, j).
- 2. Selection Find cities k not belonging to the partial tour that is farthest from any of the cities belonging to the partial tour.
- 3. Insertion Find the edge  $\{i, j\}$ , belonging to the partial tour, that minimizes  $c_{ik} + c_{kj} c_{ij}$ . Insert k between i and j.
- 4. If all cities are inserted then STOP, else go back to 2.

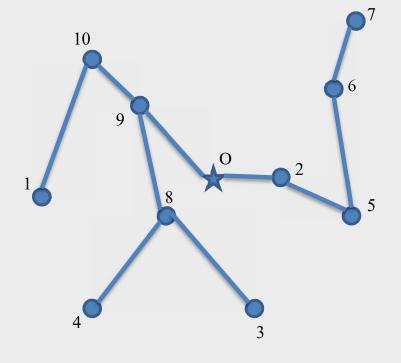
#### TSP - CHEAPEST INSERTION

- 1. Initialization Start with a partial tour with just one city i, randomly chosen; find the city j for which  $c_{ij}$  (distance or cost from i to j) is minimum and build the partial tour (i, j).
- 2. Selection Find cities k, i and j (i and j being the extremes of an edge belonging to the partial tour and k not belonging to that tour) for which  $c_{ik} + c_{kj} c_{ij}$  is minimized.
- 3. Insertion Insert k between i and j.
- 4. If all cities are inserted then STOP, else go back to 2.

# SPANNING TREE

• A tree which includes all of the vertices of a graph

Dis.	1	2	3	4	5	6	7	8	9	10	0
1	0	10	10.8	6.3	13	13.4	16.4	5.1	6.4	8.2	7.1
2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
7	16.4	9.5	16.5	19.4	11	4.1	0	13.6	10.3	11.2	10.8
8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0



Length: 52.9

#### MINIMUM SPANNING TREE

Objective: Find the minimum distance such that all nodes

are visited once (i.e. no cycles).

The following algorithms run in polynomial time

- Kruskal's algorithm
- *Prim*'s algorithm

#### MINIMUM SPANNING TREE - ALGORITHMS

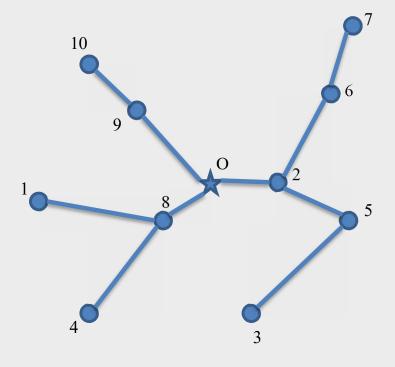
### Kruskal's algorithm

- 1. Select the shortest edge in a network
- 2. Select the next shortest edge which does not create a cycle
- 3. Repeat step 2 until all vertices have been connected

#### Prim's algorithm

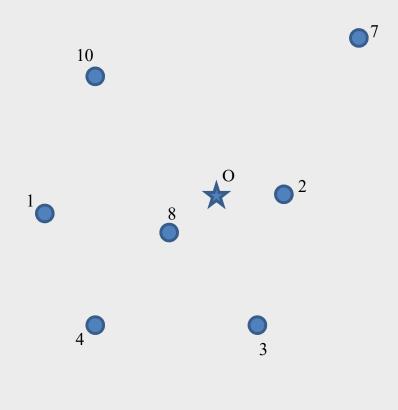
- 1. Select any vertex
- 2. Select the shortest edge connected to that vertex
- 3. Select the shortest edge connected to any vertex already connected
- 4. Repeat step 3 until all vertices have been connected

Dis.	1	2	3	4	5	6	7	8	9	10	0
1	0	10	10.8	6.3	13	13.4	16.4	5.1	6.4	8.2	7.1
2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
7	16.4	9.5	16.5	19.4	11	4.1	0	13.6	10.3	11.2	10.8
8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0



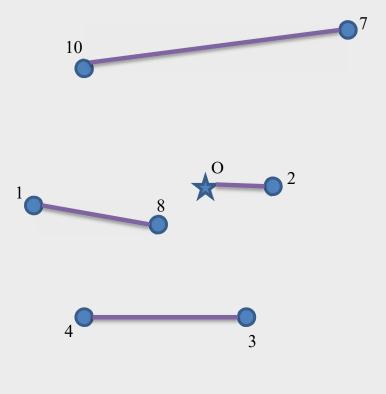
Ahmad Hemmati Page 3'

Dis.	1	2	3	4	5	6	7	8	9	10	0
1	0	10	10.8	6.3	13	13.4	16.4	5.1	6.4	8.2	7.1
2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
7	16.4	9.5	16.5	19.4	11	4.1	0	13.6	10.3	11.2	10.8
8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0



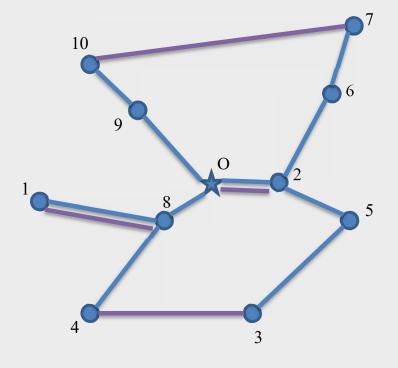
Ahmad Hemmati Page

Dis.	1	2	3	4	5	6	7	8	9	10	0
1	0	10	10.8	6.3	13	13.4	16.4	5.1	6.4	8.2	7.1
2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
7	16.4	9.5	16.5	19.4	11	4.1	0	13.6	10.3	11.2	10.8
8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0



Ahmad Hemmati Page 3

Dis.	1	2	3	4	5	6	7	8	9	10	0
1	0	10	10.8	6.3	13	13.4	16.4	5.1	6.4	8.2	7.1
2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
7	16.4	9.5	16.5	19.4	11	4.1	0	13.6	10.3	11.2	10.8
8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0

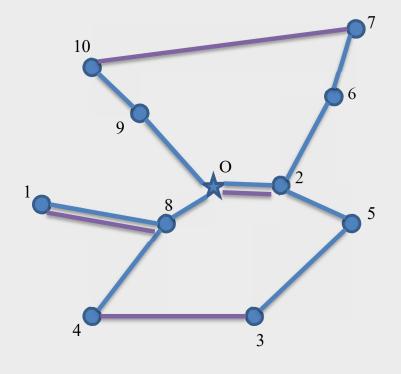


- It is an approximation algorithm that guarantees that its solutions will be within a factor of 3/2 of the optimal solution length
- As of 2017, this is the best approximation ratio that has been proven for the traveling salesman problem on general metric spaces, although better approximations are known for some special cases.

- 1. Create a minimum spanning tree *T* of *G*.
- 2. Let *O* be the set of vertices with odd degree in *T*.
- 3. Find a minimum-weight perfect matching *M* in the induced subgraph given by the vertices from *O*.
- 4. Combine the edges of *M* and *T* to form a connected multigraph *H* in which each vertex has even degree.
- 5. Form an Eulerian circuit in *H*.
- 6. Make the circuit found in previous step into a Hamiltonian circuit by skipping repeated vertices (shortcutting).

# EULER CYCLE

Dis.	1	2	3	4	5	6	7	8	9	10	0
1	0	10	10.8	6.3	13	13.4	16.4	5.1	6.4	8.2	7.1
2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
7	16.4	9.5	16.5	19.4	11	4.1	0	13.6	10.3	11.2	10.8
8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0



O-9-10-7-6-2-O-2-5-3-4-8-1-8-O

Dis.	1	2	3	4	5	6	7	8	9	10	0
1	0	10	10.8	6.3	13	13.4	16.4	5.1	6.4	8.2	7.1
2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
7	16.4	9.5	16.5	19.4	11	4.1	0	13.6	10.3	11.2	10.8
8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0

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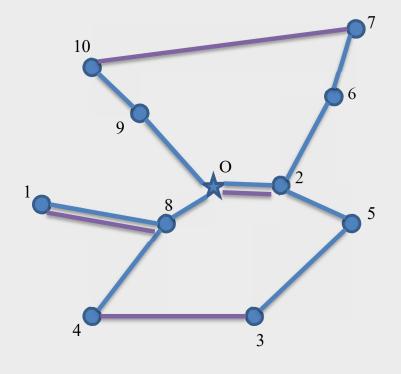
Tour: O-9-10-7-6-2-5-3-4-8-1-O

Length: 64.3

Ahmad Hemmati Page

# EULER CYCLE

Dis.	1	2	3	4	5	6	7	8	9	10	0
1	0	10	10.8	6.3	13	13.4	16.4	5.1	6.4	8.2	7.1
2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
7	16.4	9.5	16.5	19.4	11	4.1	0	13.6	10.3	11.2	10.8
8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0



O-8-1-8-4-3-5-2-O-2-6-7-10-9-O

Dis.	1	2	3	4	5	6	7	8	9	10	0
1	0	10	10.8	6.3	13	13.4	16.4	5.1	6.4	8.2	7.1
2	10	0	7.1	10.6	3.6	5.4	9.5	5.4	7.2	10.6	3
3	10.8	7.1	0	7	6.4	12.4	16.5	6.4	12.1	15.7	7.3
4	6.3	10.6	7	0	12.1	15.6	19.4	5.8	11.2	14	8.6
5	13	3.6	6.4	12.1	0	7.1	11	8	10.8	14.2	6.3
6	13.4	5.4	12.4	15.6	7.1	0	4.1	9.9	8.1	10.2	7.1
7	16.4	9.5	16.5	19.4	11	4.1	0	13.6	10.3	11.2	10.8
8	5.1	5.4	6.4	5.8	8	9.9	13.6	0	6.1	9.5	2.8
9	6.4	7.2	12.1	11.2	10.8	8.1	10.3	6.1	0	3.6	5
10	8.2	10.6	15.7	14	14.2	10.2	11.2	9.5	3.6	0	8.6
0	7.1	3	7.3	8.6	6.3	7.1	10.8	2.8	5	8.6	0

9 8 O 2 5

10

Tour: O-8-1-4-3-5-2-6-7-10-9-O

Length: 60.5

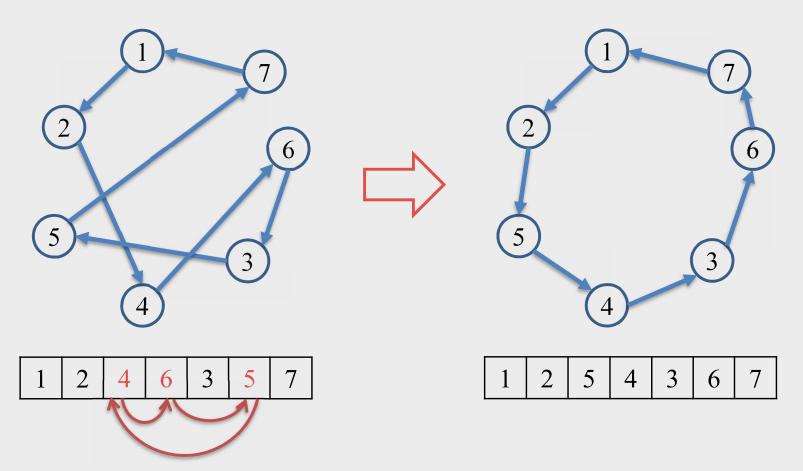
# TSP - 2-Opt

- 2-opt
  - Note: a finite sequence of "swap (2-exchange)" can generate any tour (for TSP), including the optimum tour

- <u>Strategy</u>:
  - Select the best swap among N \* (N-1) / 2 possible swap
  - Repeat this process until no improvement can be made)

# TSP - 3-Opt

• 3-exchange  $\rightarrow$  3-opt





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#### NEXT LECTURE

# LECTURE #14:

# VEHICLE ROUTING PROBLEM

