Algorithm in Computational Biology: Assignment 1

Name : Preethi Raksha Thorali Gobinath ID: 1229592977

Question 2

Since we need to check whether sequence X and Y come from the same linear sequence, we will need to create a sequence which has all possible single rotations, this can be done by adding the sequence to itself one more time.

Let's take an example,

- X= "GTACA" and Y= "CAGTA",
 We will consider Pattern = X = "GTACA"
- We will create a text string with all possible single rotations, Text= Y+Y = CAGTA +CAGTA =
 "CAGTACAGTA". This string contains all possible rotations of CAGTA in the form of
 substring, which are CAGTA, AGTAC, GTACA, TACAG, ACAGT, CAGTA. Of all the rotations,
 our pattern X matches the substring 3.
- Given Pattern = "GTACA" and Text = "CAGTACAGTA", we can visibly notice that our text contains the pattern at index starting from 2 till index 7.
- Now for the z algorithm the input pattern is "GTACA" and the input text is "CAGTACAGTA", let's add the "\$" symbol between them to process the algorithm "GTACA\$CAGTACAGTA".
- Now if we run the algorithm to calculate the z values we get z[]=[x,0,0,0,0,0,0,0,0,0,0,0,0,0].
- After which we check in the z array, is there a value with length of the input pattern, which is
 equal to 5 in our example. If the length of the input pattern is present in the z array, then we
 output "both the sequences are from the same circular sequence" else we output they are
 different sequences.

Question 3

```
Given:
```

P = ATCATCT and S = TCATCATGATGATCATCT

<u>Let:</u>

x be the pattern Index iterator y be the Text index iterator

S=TCATCATGAT G A T C A T C T Y=01 234 56 78 9 10 11 12 13 14 15 16 17

KMP Algo:

```
M = len(P)
N = len(S)
y = 0 # index for txt
x = 0 # index for pat
while (N - y) >= (M - x):
       if pat[x] == txt[y]:
              x += 1
              y += 1
       if x == M:
              result.append(y - x + 1)
              x = lps[x - 1]
       elif y < N and pat[x] != txt[y]:
              if x != 0:
                      x = lps[x - 1]
              else:
                      y += 1
```

Manually running lps:

Steps	Pattern	Text	Inference
1	x=0 A	y=0 T	No equal
2	x=0 A	y=1 C	No equal
3	x=0 A	y=2 A	Equal
4	x=1 T	y=3 T	Equal
5	x=2 C	y=4 C	Equal
6	X=3 A	y=5 A	Equal
7	x=4 T	y=6 T	Equal
8	x=5 C	y=7 G	Not Equal
9	x=lps[5-1]=lps[4]=2 C	y=7 G	Not Equal
10	x=lps[2-1]=lps[1]=0 A	y=7 G	Not Equal
11	x=0 A	y=8 A	Equal
12	x=1 T	y=9 T	Equal
13	x=2 C	y=10 G	Not Equal
14	x=lps[2-1]=lps[1]=0 A	y=10 G	Not Equal
15	x=0 A	y=11 A	Equal
16	x=1 T	y=12 T	Equal
17	x=2 C	y=13 C	Equal
18	X=3 A	y=14 A	Equal
19	x=4 T	y=15 T	Equal
20	x=5 C	y=16 C	Equal
21	x=6 T	y=17 T	Equal
Stop	x=7=len(P)		Print

Manually running lps':

Steps	Pattern	Text	Inference
1	x=0 A	y=0 T	No equal
2	x=0 A	y=1 C	No equal
3	x=0 A	y=2 A	Equal
4	x=1 T	y=3 T	Equal
5	x=2 C	y=4 C	Equal
6	X=3 A	y=5 A	Equal
7	x=4 T	y=6 T	Equal
8	x=5 C	y=7 G	Not Equal
9	x=lps[5-1]=lps[4]=0 A	y=7 G	Not Equal
10	x=0 A	y=8 A	Equal
11	x=1 T	y=9 T	Equal
12	x=2 C	y=10 G	Not Equal
13	x=lps[2-1]=lps[1]=0 A	y=10 G	Not Equal
14	x=0 A	y=11 A	Equal
15	x=1 T	y=12 T	Equal
16	x=2 C	y=13 C	Equal
17	X=3 A	y=14 A	Equal
18	x=4 T	y=15 T	Equal
19	x=5 C	y=16 C	Equal
20	x=6 T	y=17 T	Equal
Stop	x=7=len(P)		Print

From the above calculation I can see that **Ips takes 21 steps** which is greater than **Ips' which takes only 20 steps**. For the above example I can confirm that Dr.wiz solution is valid for this example, but however one example cannot be used to assure the validity of the algorithm in all cases. A valid proof

would be required to prove that it works for all cases or an example which breaks the algorithm would be required to disprove it.

Question 4

Let us take a example 1 Input = CAGTACAGTA

• Lets calculate the z values: which store the length of the longest substring starting from input[i] if it is the prefix of input[1..n-1].

$$z[]=[x,0,0,0,0,5,0,0,0,0]$$

• Lets calculate the lps values: hich stores length of the longest non trivial suffix of input[1..i] that matches a prefix of input.

$$lps[i] = [0,0,0,0,0,1,2,3,4,5]$$

• Lets calculate the lps' values: which stores length of the longest non trivial suffix of input[1..i] that matches a prefix of input such that input[lps[i] + 1]! = input[i + 1].

$$lps[i] = [0,0,0,0,0,0,0,0,0,5]$$

Let us take a example 2 Input = CAGTCCAGTC

• Lets calculate the z values: which store the length of the longest substring starting from input[i] if it is the prefix of input[1..n-1].

$$z[]=[x,0,0,0,1,5,0,0,0,1]$$

• Lets calculate the lps values: Which stores length of the longest non trivial suffix of input[1..i] that matches a prefix of input.

$$lps[i] = [0,0,0,0,1,1,2,3,4,5]$$

• Lets calculate the lps' values: which stores length of the longest non trivial suffix of input[1..i] that matches a prefix of input such that input[lps[i] + 1]! = input[i + 1].

```
lps[i] = [0,0,0,0,1,0,0,0,0,5]
```

From both the above examples, we can notice that the z values store the length of the longest substring at the **start index** of the match, whereas lps' stores at the **last index** of the match.

So we can calculate the lps's values for each index of z values by assigning:

```
lps' [ current z index + current z value -1 ] => current z value, given lps' is in initial state (0)
```

Code:

```
Def calculateLpsDash (zvalues):
```

```
# Initialize lps' with 0 for length of zValues array
lpsDash=[0] * len(zValues)

# Iterate over the zValues array from index 1, as first value is x
For i in range(1,len(zvalues)):
```

Main function:

```
# Input
S= "CAGTACAGTA"
zValues=z[]= [x,0,0,0,0,5,0,0,0,0]

# Call the function
result = calculateLpsDash ( zValues)

# Display the function
Print (result)
```