

# Algorithm in Computational Biology :Assignment 1

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## Question 2

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Since we need to check whether sequence X and Y come from the same linear sequence, we will need to create a sequence which has all possible single rotations, this can be done by adding the sequence to itself one more time.

Let's take an example,

- X= "GTACA" and Y= "CAGTA",  
We will consider Pattern = X = "GTACA"
- We will create a text string with all possible single rotations, Text= Y+Y = CAGTA +CAGTA = "**CAGTACAGTA**". This string contains all possible rotations of CAGTA in the form of substring, which are CAGTA, AGTAC, **GTACA**, TACAG, ACAGT, CAGTA. Of all the rotations, our pattern X matches the substring 3.
- Given Pattern = "GTACA" and Text = "CAGTACAGTA" , we can visibly notice that our text contains the pattern at index starting from 2 till index 7.
- Now for the z algorithm the input pattern is "GTACA" and the input text is "CAGTACAGTA", let's add the "\$" symbol between them to process the algorithm "GTACA\$CAGTACAGTA".
- Now if we run the algorithm to calculate the z values we get z[]= [x,0,0,0,0,0,0,5,0,0,0,0,3].
- After which we check in the z array, is there a value with length of the input pattern, which is equal to 5 in our example.
- Finally we can display from what index our pattern matches.

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### Question 3

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Given:

P = ATCATCT and S = TCATCATGATGATCATCT

Let :

x be the pattern Index iterator

y be the Text index iterator

i = 1 2 3 4 5 6 7

P[i] = A T C A T C T

lps[i] = 0 0 0 1 2 3 0

lps'[i] = 0 0 0 0 0 3 0

X = 0 1 2 3 4 5 6

S = T C A T C A T G A T G A T C A T C T

Y = 0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17

KMP Algo:

```
M = len(P)
```

```
N = len(S)
```

```
y = 0 # index for txt
```

```
x = 0 # index for pat
```

```
while (N - y) >= (M - x):
```

```
    if pat[x] == txt[y]:
```

```
        x += 1
```

```
        y += 1
```

```
    if x == M:
```

```
        result.append(y - x + 1)
```

```
        x = lps[x - 1]
```

```
    elif y < N and pat[x] != txt[y]:
```

```
        if x != 0:
```

```
            x = lps[x - 1]
```

```
        else:
```

```
            y += 1
```

### Manually running lps:

Index x = 0 1 2 3 4 5 6  
lps[x] = 0 0 0 1 2 3 0  
P[x] = A T C A T C T

| Steps    | Pattern               | Text   | Inference |
|----------|-----------------------|--------|-----------|
| <u>1</u> | x=0 A                 | y=0 T  | No equal  |
| <u>2</u> | x=0 A                 | y=1 C  | No equal  |
| 3        | x=0 A                 | y=2 A  | Equal     |
| 4        | x=1 T                 | y=3 T  | Equal     |
| 5        | x=2 C                 | y=4 C  | Equal     |
| 6        | X=3 A                 | y=5 A  | Equal     |
| 7        | x=4 T                 | y=6 T  | Equal     |
| 8        | x=5 C                 | y=7 G  | Not Equal |
| 9        | x=lps[5-1]=lps[4]=2 C | y=7 G  | Not Equal |
| 10       | x=lps[2-1]=lps[1]=0 A | y=7 G  | Not Equal |
| 11       | x=0 A                 | y=8 A  | Equal     |
| 12       | x=1 T                 | y=9 T  | Equal     |
| 13       | x=2 C                 | y=10 G | Not Equal |
| 14       | x=lps[2-1]=lps[1]=0 A | y=10 G | Not Equal |
| 15       | x=0 A                 | y=11 A | Equal     |
| 16       | x=1 T                 | y=12 T | Equal     |
| 17       | x=2 C                 | y=13 C | Equal     |
| 18       | X=3 A                 | y=14 A | Equal     |
| 19       | x=4 T                 | y=15 T | Equal     |
| 20       | x=5 C                 | y=16 C | Equal     |
| 21       | x=6 T                 | y=17 T | Equal     |
| Stop     | x=7=len(P)            |        | Print     |

### Manually running lps':

Index x = 0 1 2 3 4 5 6  
lps'[x] = 0 0 0 0 0 3 0  
P[x] = A T C A T C T

| Steps | Pattern               | Text   | Inference |
|-------|-----------------------|--------|-----------|
| 1     | x=0 A                 | y=0 T  | No equal  |
| 2     | x=0 A                 | y=1 C  | No equal  |
| 3     | x=0 A                 | y=2 A  | Equal     |
| 4     | x=1 T                 | y=3 T  | Equal     |
| 5     | x=2 C                 | y=4 C  | Equal     |
| 6     | X=3 A                 | y=5 A  | Equal     |
| 7     | x=4 T                 | y=6 T  | Equal     |
| 8     | x=5 C                 | y=7 G  | Not Equal |
| 9     | x=lps[5-1]=lps[4]=0 A | y=7 G  | Not Equal |
| 10    | x=0 A                 | y=8 A  | Equal     |
| 11    | x=1 T                 | y=9 T  | Equal     |
| 12    | x=2 C                 | y=10 G | Not Equal |
| 13    | x=lps[2-1]=lps[1]=0 A | y=10 G | Not Equal |
| 14    | x=0 A                 | y=11 A | Equal     |
| 15    | x=1 T                 | y=12 T | Equal     |
| 16    | x=2 C                 | y=13 C | Equal     |
| 17    | X=3 A                 | y=14 A | Equal     |
| 18    | x=4 T                 | y=15 T | Equal     |
| 19    | x=5 C                 | y=16 C | Equal     |
| 20    | x=6 T                 | y=17 T | Equal     |
| Stop  | x=7=len(P)            |        | Print     |

From the above calculation I can see that **lps takes 21 steps** which is greater than **lps' which takes only 20 steps**. For the above example I can confirm that Dr.wiz solution is valid, but however one example cannot be used to assure the validity of the algorithm in all cases. A valid proof would be

required to prove that it works for all cases or an example which breaks the algorithm would be required to disprove it.

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## Question 4

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Let us take a example 1 Input = **CAGTACAGTA**

i            = 1 2 3 4 5 6 7 8 9 10  
input [i] = C A G T A C A G T A

- Lets calculate the z values: which store the length of the longest substring starting from input[i] if it is the prefix of input[1..n-1].

z[] = [x,0,0,0,0,5,0,0,0,0]

- Lets calculate the lps values: which stores length of the longest non trivial suffix of input[1..i] that matches a prefix of input.

lps[i] = [ 0,0,0,0,0,1,2,3,4,5]

- Lets calculate the lps' values: which stores length of the longest non trivial suffix of input[1..i] that matches a prefix of input such that input[lps[i] + 1] != input[i + 1].

lps[i] = [ 0,0,0,0,0,0,0,0,0,5]

Let us take a example 2 Input = **CAGTCCAGTC**

i            = 1 2 3 4 5 6 7 8 9 10  
input [i] = C A G T C C A G T C

- Lets calculate the z values: which store the length of the longest substring starting from input[i] if it is the prefix of input[1..n-1].

z[] = [x,0,0,0,1,5,0,0,0,1]

- Lets calculate the lps values: Which stores length of the longest non trivial suffix of input[1..i] that matches a prefix of input.

lps[i] = [ 0,0,0,0,1,1,2,3,4,5]

- Lets calculate the lps' values: which stores length of the longest non trivial suffix of input[1..i] that matches a prefix of input such that input[lps[i] + 1] != input[i + 1].

```
lps[i]= [ 0,0,0,0,1,0,0,0,0,5]
```

From both the above examples, we can notice that the z values store the length of the longest substring at the **start index** of the match, whereas lps' stores at the **last index** of the match.

So we can calculate the lps's values for each index of z values by assigning:

$\text{lps}' [\text{current z index} + \text{current z value} - 1] \Rightarrow \text{current z value}$ , given lps' is in initial state (0)

Code:

#### **Def calculateLpsDash (zvalues):**

```
# Initialize lps' with 0 for length of zValues array
```

```
lpsDash=[0] * len(zValues)
```

```
# Iterate over the zValues array from index 1, as first value is x
```

```
For i in range(1,len(zvalues)):
```

```
    # If zValue contains value & lps' is in the initial state (0) than assign the value to lps'
```

```
    If zValues[i]>0 and lpsDash[i + zValues[i] -1] == 0 :
```

```
        lpsDash[i + zValues[i] -1]= zValues[i]
```

#### **Main function:**

```
# Input
```

```
S= "CAGTACAGTA"
```

```
zValues=z[]= [x,0,0,0,0,5,0,0,0,0]
```

```
# Call the function
```

```
result = calculateLpsDash ( zValues)
```

```
# Display the function
```

```
Print (result)
```