



**Vidyavardhini's College of Engineering and Technology**

**Department of Artificial Intelligence & Data Science**

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Experiment No. 7
Kruskal's Algorithm
Date of Performance:
Date of Submission:



## **Experiment No. 7**

**Title:** Kruskal's Algorithm.

**Aim:** To study and implement Kruskal's Minimum Cost Spanning Tree Algorithm.

**Objective:** To introduce Greedy based algorithms

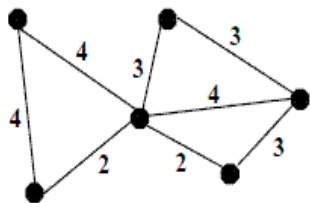

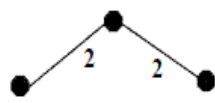
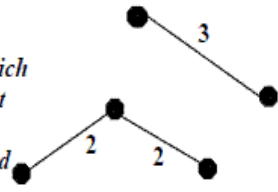
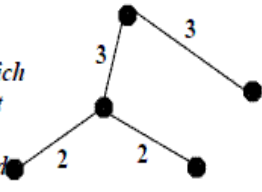
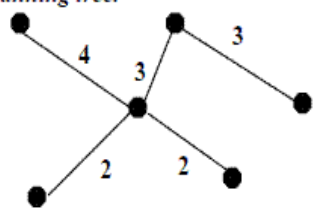
### **Theory:**

Kruskal's algorithm finds a minimum spanning forest of an undirected edge-weighted graph. If the graph is connected, it finds a minimum spanning tree. (A minimum spanning tree of a connected graph is a subset of the edges that forms a tree that includes every vertex, where the sum of the weights of all the edges in the tree is minimized. For a disconnected graph, a minimum spanning forest is composed of a minimum spanning tree for each connected component.) It is a greedy algorithm in graph theory as in each step it adds the next lowest-weight edge that will not form a cycle to the minimum spanning forest.

### **Example:**



## Kruskal's Algorithm

<p>1 Given a network.....</p> 	<p>2 Choose the shortest edge (if there is more than one, choose any of the shortest).....</p> 	<p>3 Choose the next shortest edge and add it.....</p> 
<p>4 Choose the next shortest edge which wouldn't create a cycle and add it.</p> 	<p>5 Choose the next shortest edge which wouldn't create a cycle and add it.</p> 	<p>6 Repeat until you have a minimal spanning tree.</p> 

Algorithm and Complexity:



```
1  Algorithm Kruskal( $E, cost, n, t$ )
2  //  $E$  is the set of edges in  $G$ .  $G$  has  $n$  vertices.  $cost[u, v]$  is the
3  // cost of edge  $(u, v)$ .  $t$  is the set of edges in the minimum-cost
4  // spanning tree. The final cost is returned.
5  {
6      Construct a heap out of the edge costs using Heapify;
7      for  $i := 1$  to  $n$  do  $parent[i] := -1$ ;
8      // Each vertex is in a different set.
9       $i := 0$ ;  $mincost := 0.0$ ;
10     while  $((i < n - 1)$  and  $(\text{heap not empty}))$  do
11     {
12         Delete a minimum cost edge  $(u, v)$  from the heap
13         and reheapify using Adjust;
14          $j := \text{Find}(u)$ ;  $k := \text{Find}(v)$ ;
15         if  $(j \neq k)$  then
16         {
17              $i := i + 1$ ;
18              $t[i, 1] := u$ ;  $t[i, 2] := v$ ;
19              $mincost := mincost + cost[u, v]$ ;
20             Union $(j, k)$ ;
21         }
22     }
23     if  $(i \neq n - 1)$  then write ("No spanning tree");
24     else return  $mincost$ ;
25 }
```

Time Complexity is  $O(n \log n)$ , Where,  $n$  = number of Edges

### Implementation:

```
#include <stdio.h>
```

```
#include <stdlib.h>
```

```
#define MAX_EDGES 1000
```

```
typedef struct Edge {
```

```
    int src, dest, weight;
```

```
} Edge;
```

```
typedef struct Graph {
```



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```
int V, E;

Edge edges[MAX_EDGES];

} Graph;

typedef struct Subset {

    int parent, rank;

} Subset;

Graph* createGraph(int V, int E) {

    Graph* graph = (Graph*) malloc(sizeof(Graph));

    graph->V = V;

    graph->E = E;

    return graph;

}

int find(Subset subsets[], int i) {

    if (subsets[i].parent != i) {

        subsets[i].parent = find(subsets, subsets[i].parent);

    }

    return subsets[i].parent;

}

void Union(Subset subsets[], int x, int y) {

    int xroot = find(subsets, x);

    int yroot = find(subsets, y);

    if (subsets[xroot].rank < subsets[yroot].rank) {

        subsets[xroot].parent = yroot;

    } else if (subsets[xroot].rank > subsets[yroot].rank) {

        subsets[yroot].parent = xroot;

    } else {

        subsets[yroot].parent = xroot;
```



```
    subsets[xroot].rank++;
}
}

int compare(const void* a, const void* b) {
    Edge* a_edge = (Edge*) a;
    Edge* b_edge = (Edge*) b;
    return a_edge->weight - b_edge->weight;
}

void kruskalMST(Graph* graph) {
    Edge mst[graph->V];
    int e = 0, i = 0;
    qsort(graph->edges, graph->E, sizeof(Edge), compare);
    Subset* subsets = (Subset*) malloc(graph->V * sizeof(Subset));
    for (int v = 0; v < graph->V; ++v) {
        subsets[v].parent = v;
        subsets[v].rank = 0;
    }
    while (e < graph->V - 1 && i < graph->E) {
        Edge next_edge = graph->edges[i++];
        int x = find(subsets, next_edge.src);
        int y = find(subsets, next_edge.dest);
        if (x != y) {
            mst[e++] = next_edge;
            Union(subsets, x, y);
        }
    }
    printf("Minimum Spanning Tree:\n");
```



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```
for (i = 0; i < e; ++i) {  
    printf("(%d, %d) -> %d\n", mst[i].src, mst[i].dest, mst[i].weight);  
}  
}  
  
int main() {  
    int V, E;  
  
    printf("Enter number of vertices and edges: ");  
  
    scanf("%d %d", &V, &E);  
  
    Graph* graph = createGraph(V, E);  
  
    printf("Enter edges and their weights:\n");  
  
    for (int i = 0; i < E; ++i) {  
        scanf("%d %d %d", &graph->edges[i].src, &graph->edges[i].dest, &graph->edges[i].weight);  
    }  
  
    kruskalMST(graph);  
  
    return 0;  
}
```



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### Output:

```
Enter number of vertices and edges: 5 7
Enter edges and their weights:
0 1 2
0 3 6
1 3 8
3 4 9
1 4 5
1 2 3
2 4 7
Minimum Spanning Tree:
(0, 1) -> 2
(1, 2) -> 3
(1, 4) -> 5
(0, 3) -> 6

...Program finished with exit code 0
Press ENTER to exit console. □
```

**Conclusion:** Implementing Kruskal's algorithm proved effective in finding the minimum spanning tree of a given graph. Its simplicity and efficiency make it a valuable tool for solving graph optimization problems, demonstrating its practical applicability in various domains.